- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: $2.6 \mathbf{2}$.5 2.4 2.3 2.2 $2.1 \mathbf{2} \mathbf{2}$


## Table of Contents

pgRouting extends the PostGIS/PostgreSQL geospatial database to provide geospatial routing and other network analysis functionality.

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## General

- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: $2.6 \mathbf{2}$.5 $2.42 .3 \mathbf{2}$.2 $2.1 \mathbf{2} \mathbf{2}$


## Introduction

pgRouting is an extension of PostGIS and PostgreSQL geospatial database and adds routing and other network analysis functionality. A predecessor of pgRouting - pgDijkstra, written by Sylvain Pasche from Camptocamp, was later extended by Orkney and renamed to pgRouting. The project is now supported and maintained by Georepublic, Paragon Corporation and a broad user community. pgRouting is part of OSGeo Community Projects from the OSGeo Foundation and included on OSGeoLive.

## Licensing

The following licenses can be found in pgRouting:

| License |  |  |
| :--- | :--- | :--- |
| GNU General Public License v2.0 or | Most features of pgRouting are available underGNU General Public License <br> later | v2.0 or later. |
| Boost Software License - Version 1.0 | Some Boost extensions are available under Boost Software License - Version <br>  <br> MIT-X License | Some code contributed by iMaptools.com is available under MIT-X license. |
| Creative Commons Attribution-Share | The pgRouting Manual is licensed under a Creative Commons Attribution- <br> Alike 3.0 License | Share Alike 3.0 License. |

In general license information should be included in the header of each source file.

Contributors
This Release Contributors

Individuals in this release (in alphabetical order)

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And all the people that give us a little of their time making comments, finding issues, making pull requests etc. in any of our products: osm2pgrouting, pgRouting, pgRoutingLayer, workshop.

Corporate Sponsors in this release (in alphabetical order)
These are corporate entities that have contributed developer time, hosting, or direct monetary funding to the pgRouting project:

- Google Summer of Code
- Paragon Corporation


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- Camptocamp
- CSIS (University of Tokyo)
- Georepublic
- Google Summer of Code
- iMaptools
- Leopark
- Orkney
- Paragon Corporation
- Versaterm Inc.


## More Information

- The latest software, documentation and news items are available at the pgRouting web sitehttps://pgrouting.org.
- PostgreSQL database server at the PostgreSQL main sitehttps://www.postgresql.org.
- PostGIS extension at the PostGIS project web sitehttps://postgis.net.
- Boost C++ source libraries athttps://www.boost.org.
- The Migration guide from 2.6 can be found athttps://github.com/pgRouting/pgrouting/wiki/Migration-Guide.
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: $2.6 \mathbf{2}$.5 2.4 2.3 2.2 2.12 .0


## Installation

## Table of Contents

- Short Version
- Get the sources
- Enabling and upgrading in the database
- Dependencies
- Configuring
- Building
- Testing

Instructions for downloading and installing binaries for different operating systems, additional notes and corrections not included in this documentation can be found in Installation wiki

To use pgRouting PostGIS needs to be installed, please read the information about installation in thislnstall Guide

## Short Version

Extracting the tar ball

```
tar xvfz pgrouting-3.3.4.tar.gz
cd pgrouting-3.3.4
```

To compile assuming you have all the dependencies in your search path:

```
mkdir build
cd build
cmake
make
sudo make install
```

Once pgRouting is installed, it needs to be enabled in each individual database you want to use it in.

```
createdb routing
psql routing -c 'CREATE EXTENSION PostGIS'
psql routing -c 'CREATE EXTENSION pgRouting'
```


## Get the sources

The pgRouting latest release can be found inhttps://github.com/pgRouting/pgrouting/releases/latest

## wget

To download this release:

```
wget -O pgrouting-3.3.4.tar.gz https://github.com/pgRouting/pgrouting/archive/v3.3.4.tar.gz
```

Go to Short Version for more instructions on extracting tar ball and compiling pgRouting.

## git

To download the repository

```
git clone git://github.com/pgRouting/pgrouting.git
cd pgrouting
git checkout v3.3.4
```

Go to Short Version for more instructions on compiling pgRouting (there is no tar ball involved while downloading pgRouting repository from GitHub).

Enabling and upgrading in the database

## Enabling the database

pgRouting is a PostgreSQL extension and depends on PostGIS to provide functionalities to end user. Below given code demonstrates enabling PostGIS and pgRouting in the database.

```
CREATE EXTENSION postgis;
CREATE EXTENSION pgrouting;
```

Checking PostGIS and pgRouting version after enabling them in the database.

```
SELECT PostGIS_full_version();
SELECT * FROM pgr_version();
```


## Upgrading the database

To upgrade pgRouting in the database to version 3.3.4 use the following command:

```
ALTER EXTENSION pgrouting UPDATE TO "3.3.4";
```


## More information can be found inhttps://www.postgresql.org/docs/current/sql-createextension.html

## Dependencies

## Compilation Dependencies

To be able to compile pgRouting, make sure that the following dependencies are met:

- C and C++0x compilers
- Compiling with Boost 1.56 up to Boost 1.74 requires $C++$ Compiler with $C++03$ or $C++11$ standard support
- Compiling with Boost 1.75 requires $C++$ Compiler with $C++14$ standard support
- Postgresql version $=$ Supported versions by PostgreSQL
- The Boost Graph Library (BGL). Version >= 1.56
- CMake $>=3.2$


## optional dependencies

For user's documentation

- Sphinx >= 1.1
- Latex

For developer's documentation

- Doxygen >=1.7

For testing

- pgtap
- pg_prove

For using:

- PostGIS version $>=2.2$


## Example: Installing dependencies on linux

Installing the compilation dependencies

## Database dependencies

```
sudo apt install postgresql-14
sudo apt install postgresql-server-dev-14
sudo apt install postgresql-14-postgis
```


## Configuring PostgreSQL

Entering psql console

```
sudo systemctl start postgresql.service
sudo -i -u postgres
psq|
```


## To exit psql console

9

Entering psql console directly without switching roles can be done by the following commands

```
sudo -u postgres psq
```

Then use the above given method to exit out of the psql console
Checking PostgreSQL version

```
psql --version
```

or
Enter the psql console using above given method and then enter

```
SELECT VERSION();
```

Creating PostgreSQL role

```
sudo -i -u postgres
createuser --interactive
```

or

Default role provided by PostgreSQL is postgres. To create new roles you can use the above provided commands. The prompt will ask the user to type name of the role and then provide affirmation. Proceed with the steps and you will succeed in creating PostgreSQL role successfully.

To add password to the role or change previously created password of the role use the following commands
ALTER USER <role name> PASSWORD <password>

To get additional details on the flags associated withcreateuser below given command can be used

```
man createuser
```


## Creating Database in PostgreSQL

```
sudo -i -u postgres
createdb <database name>
```

or
sudo -u postgres createdb <database name>

## Connecting to a PostgreSQL Database

Enter the psql console and type the following commands

```
connect <database name>
```


## Build dependencies

```
sudo apt install cmake
sudo apt install g++
sudo apt install libboost-graph-dev
```


## Optional dependencies

## For documentation and testing

```
pip install sphinx
pip install sphinx-bootstrap-theme
sudo apt install texlive
sudo apt install doxygen
sudo apt install libtap-parser-sourcehandler-pgtap-perl
sudo apt install postgresql-14-pgtap
```


## Configuring

pgRouting uses the cmake system to do the configuration.
The build directory is different from the source directory
Create the build directory
\$ mkdir build

## Configurable variables

## To see the variables that can be configured

```
$ cd build
$ cmake -L ..
```


## Configuring The Documentation

Most of the effort of the documentation has been on the HTML files. Some variables for building documentation:

| Variable | Default | Comment |
| :--- | :--- | :--- |
| WITH_DOC | BOOL=OFF | Turn on/off building the documentation |
| BUILD_HTML | BOOL=ON | If ON, turn on/off building HTML for user's documentation |


| Variable | Default | Comment |
| :--- | :--- | :--- |
| BUILD_DOXY | BOOL=ON | If ON, turn on/off building HTML for developer's documentation |
| BUILD_LATEX | BOOL=OFF | If ON, turn on/off building PDF |
| BUILD_MAN | BOOL=OFF | If ON, turn on/off building MAN pages |
| DOC_USE_BOOTSTRAP | BOOL=OFF | If ON, use sphinx-bootstrap for HTML pages of the users |
| documentation |  |  |

Configuring cmake to create documentation before building pgRouting
\$ cmake -DWITH_DOC=ON -DDOC_USE_BOOTSTRAP=ON .

## Note

Most of the effort of the documentation has been on the html files.

## Building

Using make to build the code and the documentation
The following instructions start from path/to/pgrouting/build

```
$ make # build the code but not the documentation
$ make doc # build only the user's documentation
$ make all doc # build both the code and the user's documentation
$ make doxy # build only the developer's documentation
```

We have tested on several platforms, For installing or reinstalling all the steps are needed.

## Warning

The sql signatures are configured and build in thecmake command.

## MinGW on Windows

```
$ mkdir build
$ cd build
$ cmake -G"MSYS Makefiles" ..
$ make
$ make install
```


## Linux

The following instructions start from path/to/pgrouting

```
mkdir build
cd build
cmake
make
sudo make install
```

To remove the build when the configuration changes, use the following code:

```
rm -rf build
```

and start the build process as mentioned previously.

## Testing

Currently there is no make test and testing is done as follows
The following instructions start from path/to/pgrouting/

```
tools/testers/doc_queries_generator.pl
createdb -U <user> ___pgr___test___
sh ./tools/testers/pg_prove_tests.sh <user>
dropdb -U <user> __pgr___test
```


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2 2.1 2.0


## Support

pgRouting community support is available through the pgRouting website, documentation, tutorials, mailing lists and others. If you're looking for commercial support, find below a list of companies providing pgRouting development and consulting services.

## Reporting Problems

Bugs are reported and managed in an issue tracker. Please follow these steps:

1. Search the tickets to see if your problem has already been reported. If so, add any extra context you might have found, or at least indicate that you too are having the problem. This will help us prioritize common issues.
2. If your problem is unreported, create a new issue for it.
3. In your report include explicit instructions to replicate your issue. The best tickets include the exact SQL necessary to replicate a problem.
4. If you can test older versions of PostGIS for your problem, please do. On your ticket, note the earliest version the problem appears.
5. For the versions where you can replicate the problem, note the operating system and version of pgRouting, PostGIS and PostgreSQL.
6. It is recommended to use the following wrapper on the problem to pin point the step that is causing the problem.

SET client_min_messages TO debug; <your code>
SET client_min_messages TO notice;

## Mailing List and GIS StackExchange

There are two mailing lists for pgRouting hosted on OSGeo mailing list server:

- User mailing list: https://lists.osgeo.org/mailman/listinfo/pgrouting-users
- Developer mailing list: https://lists.osgeo.org/mailman/listinfo/pgrouting-dev

For general questions and topics about how to use pgRouting, please write to the user mailing list.
You can also ask at GIS StackExchange and tag the question withpgrouting. Find all questions tagged with pgrouting under https://gis.stackexchange.com/questions/tagged/pgrouting or subscribe to the pgRouting questions feed.

Commercial Support

For users who require professional support, development and consulting services, consider contacting any of the following organizations, which have significantly contributed to the development of pgRouting:

| Company | Offices in | Website |
| :--- | :--- | :--- |
| Georepublic | Germany, Japan | https://georepublic.info |
| Paragon Corporation | United States | https://www.paragoncorporation.com |
| Netlab | Capranica, Italy | https://www.osgeo.org/service- <br> providers/netlab/ |

- Sample Data that is used in the examples of this manual.
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.6 2.5 2.42 .3 2.2 2.12 .0


## Sample Data

The documentation provides very simple example queries based on a small sample network that resembles a city. To be able to execute the mayority of the examples queries, follow the instructions bellow.

```
- Main graph
```

- Edges
- Edges data
- Vertices
- Vertices data
- The topology
- Topology data
- Points outside the graph
- Points of interest
- Points of interest fillup
- Points of interest geometry
- Points of interest data
- Support tables
- Combinations
- Combinations data
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- Directed graph with cost and reverse_cost
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- Directed graph with cost
- Undirected graph with cost
- Pick \& Deliver Data
- The vehicles
- The original orders
- The orders

Main graph

A graph consists of a set of edges and a set of vertices
The following city is to be inserted into the database:


Information known at this point is the geometry of the edges, cost values, cpacity values, category values and some locations that are not in the graph.

The process to have working topology starts by inserting the edges. After that everything else is calculated.

## Edges

The database design for the documentation of pgRouting, keeps in the same row 2 segments, one in the direction of the geometry and the second in the oposite direction. Therfore some information has the reverse_ prefix which corresponds to the segment on the oposite direction of the geometry.

| Column | Description |  |
| :--- | :--- | :--- |
| id | A unique identifier. |  |
| source | Identifier of the starting vertex of the geometrygeom. |  |
| target | Identifier of the ending vertex of the geometrygeom |  |
| cost | Cost to traverse from source to target. |  |
| reverse_cost | Cost to traverse from target to source. |  |
| capacity | Flow capacity from source to target. |  |
| reverse_capacity | Flow capacity from target to source. |  |
| category | Flow capacity from target to source. |  |
| reverse_category | Flow capacity from target to source. |  |
| x1 | $\backslash(x \backslash)$ coordinate of the starting vertex of the geometry. |  |


| Column | Description |
| :--- | :--- |
| y2 | $\backslash(\mathrm{y} \backslash)$ coordinate of the ending vertex of the geometry. |
|  | - For convinience it is saved on the table but can be calculated as  <br>  ST_Y(ST_EndPoint(geom)). |
| The geometry of the segments. |  |

```
CREATE TABLE edges (
    id BIGSERIAL PRIMARY KEY,
    source BIGINT,
    target BIGINT,
    cost FLOAT,
    reverse_cost FLOAT,
    capacity BIGINT,
    reverse_capacity BIGINT,
    x1 FLOAT,
    y1 FLOAT,
    x2 FLOAT,
    y2 FLOAT,
    geom geometry
);
CREATE TABLE
```

Starting on PostgreSQL 12:

```
x1 FLOAT GENERATED ALWAYS AS (ST_X(ST_StartPoint(geom))) STORED,
y1 FLOAT GENERATED ALWAYS AS (ST_Y(ST_StartPoint(geom))) STORED,
x1 FLOAT GENERATED ALWAYS AS (ST_X(ST_EndPoint(geom))) STORED,
y1 FLOAT GENERATED ALWAYS AS (ST_Y(ST_EndPoint(geom))) STORED,
```

Optionally indexes on different columns can be created. The recomendation is to have

- id indexed.
- source and target columns indexed to speed up pgRouting queries.
- geom indexed to speed up gemetry processes that might be needed in the front end.

For this small example the indexes are skipped, except forid

Edges data

Inserting into the database the information of the edges:

```
INSERT INTO edges (
    cost, reverse_cost,
    capacity, reverse_capacity, geom) VALUES
(1, 1, 80, 130, ST_MakeLine(ST_POINT(2, 0), ST_POINT(2, 1))),
(-1, 1, -1,100, ST_MakeLine(ST_POINT(2, 1), ST_POINT(3, 1))),
(-1, 1, -1, 130, ST_MakeLine(ST_POINT(3, 1), ST_POINT(4, 1))),
(1, 1, 100, 50, ST_MakeLine(ST_POINT(2, 1), ST_POINT(2, 2))),
(1,-1,130,-1, ST_MakeLine(ST_POINT(3,1),ST_POINT(3, 2))),
(1, 1, 50, 100, ST_MakeLine(ST_POINT(0, 2), ST_POINT(1, 2))),
(1, 1, 50,130, ST_MakeLine(ST_POINT(1, 2), ST_POINT(2, 2))),
( 1, 1, 100, 130, ST_MakeLine(ST_POINT(2, 2), ST_POINT(3, 2))),
(1, 1, 130, 80, ST_MakeLine(ST_POINT(3, 2),ST_POINT(4, 2))),
(1, 1, 130, 50, ST_MakeLine(ST_POINT(2, 2), ST_POINT(2, 3))),
(1,-1,130, -1, ST_MakeLine(ST_POINT(3, 2), ST_POINT(3, 3))),
( 1,-1, 100, -1, ST_MakeLine(ST_POINT(2, 3), ST_POINT(3, 3))),
(1,-1,100,-1, ST_MakeLine(ST_POINT(3,3), ST_POINT(4,3))),
(1, 1, 80,130, ST_MakeLine(ST_POINT(2, 3), ST_POINT(2, 4))),
(1, 1, 80, 50, ST_MakeLine(ST_POINT(4, 2), ST_POINT(4, 3))),
(1, 1, 80, 80, ST_MakeLine(ST_POINT(4, 1), ST_POINT(4, 2))),
(1, 1, 130, 100, ST_MakeLine(ST_POINT(0.5, 3.5), ST_POINT(1.999999999999, 3.5))),
(1, 1, 50, 130, ST_MakeLine(ST_POINT(3.5, 2.3), ST_POINT(3.5, 4)));
INSERT 0 }1
```

Negative values on the cost, capacity and category means that the edge do not exist.

## Vertices

The vertex information is calculated based on the identifier of the edge and the geometry and saved on a table. Saving all the information provided by pgr_extractVertices - Proposed:

```
SELECT * INTO vertices
FROM pgr_extractVertices('SELECT id, geom FROM edges ORDER BY id');
SELECT 17
```

In this case the because the CREATE statement was not used, the definition of an index on the table is needed.

## CREATE SEQUENCE vertices_id_seq <br> CREATE SEQUENCE

ALTER TABLE vertices ALTER COLUMN id SET DEFAULT nextval('vertices_id_seq');
ALTER TABLE
ALTER SEQUENCE vertices id seq OWNED BY vertices.id;
ALTER SEQUENCE
SELECT setval('vertices_id_seq', (SELECT coalesce(max(id))) FROM vertices));
setval
17
(1 row)

The structure of the table is:
$\square$

## Vertices data

The saved information of the vertices is:


Here is where adding more columns to the vertices table can be done. Additional columns names and types will depend on the application.

## The topology

This queries based on the vertices data create a topology by filling thesource and target columns in the edges table.

```
/* -- set the source information */
UPDATE edges AS e
SET source = v.id, x1 = x, y1 = y
FROM vertices AS v
WHERE ST_StartPoint(e.geom) = v.geom;
UPDATE 18
/* -- set the target information */
UPDATE edges AS e
SET target = v.id, x2 = x, y2 = y
FROM vertices AS v
WHERE ST_EndPoint(e.geom) = v.geom
UPDATE }1
```

```
SELECT id, source, target
FROM edges ORDER BY id
id | source | target
1| 5| 6
2 6 6 10
3| 10| 15
4| 6| 7
5| 10| 11
6| 1| 3
7| 3| 7
8| 7| 11
| 11| 16
0| 7 8
11 11| 12
2| 8| 12
3| 12| 17
4| 8| 9
5| 16| 17
16| 15| 16
17| 2| 4
18| 13| 14
(18 rows)
```


## Points outside the graph

Points of interest

Some times the applications work "on the fly" starting from a location that is not a vertex in the graph. Those locations, in pgRrouting are called points of interest

The information needed in the points of interest ispid, edge_id, side, fraction.

On this documentation there will be some 6 fixed points of interest and they will be stored on a table.

| Column | Description |
| :--- | :--- |
| pid | A unique identifier. |
| edge_id | Identifier of the edge nearest edge that allows an arrival to the <br> point. |
| side | Is it on the left, right or both sides of the segmentedge_id |
| fraction | Where in the segment is the point located. |
| geom | The geometry of the points. |
| newPoint | The geometry of the points moved on top of the segment. |

```
CREATE TABLE pointsOfInterest
    pid BIGSERIAL,
    x FLOAT,
    y FLOAT,
    edge_id BIGINT,
    side CHAR,
    fraction FLOAT,
    geom geometry,
    newPoint geometry
);
CREATE TABLE
```

Points of interest fillup

Inserting the data of the points of interest:

```
INSERT INTO pointsOfInterest ( }\textrm{x},\textrm{y}\mathrm{ y, edge_id, side, fraction) VALUES
(1.8, 0.4, 1, 'I', 0.4),
(4.2, 2.4, 15, 'r', 0.4),
(2.6,3.2, 12, 'l', 0.6),
(0.3, 1.8, 6, 'r', 0.3),
(2.9, 1.8, 5, 'l', 0.8),
(2.2, 1.7, 4, 'b', 0.7)
INSERT 0 6
```


## Points of interest geometry

Calculating for visual purposes the points over the graph.

```
UPDATE pointsOfInterest SET geom = st_makePoint(x,y);
UPDATE }
UPDATE pointsOfInterest
    SET newPoint = ST LineInterpolatePoint(e.geom, fraction)
    FROM edges AS e WWHERE edge_id = id;
UPDATE }
```


## oints of interest data

```
SELECT pid, edge_id, side, fraction,
    ST_AsText(geom), ST_AsText(newPoint)
FROM pointsOfInterest
ORDER BY pid;
pid | edge_id | side | fraction | st_astext | st_astext
    1| | 0.4|POINT(1.8 0.4)|POINT(2 0.4)
    15|r | 0.4|POINT(4.2 2.4)|POINT(4 2.4)
    12|| | 0.6 | POINT(2.6 3.2)| POINT(2.6 3)
    6 | | | 0.3 | POINT(0.3 1.8) | POINT(0.3 2)
    5 | | 0.8|POINT(2.9 1.8)|POINT(3 1.8)
    4|b | 0.7 |POINT(2.2 1.7)|POINT(2 1.7)
(6 rows)
```


## Support tables

## Combinations

Many functions can be used with a combinations of(source, target) pairs when wanting a route from source to target.

For convinence of this documentations, some combinations will be stored on a table:

```
CREATE TABLE combinations (
    source BIGINT,
    target BIGINT
);
CREATE TABLE
```

Inserting the data:

```
INSERT INTO combinations (
    source, target) VALUES
(5,6),
(5, 10),
(6,5),
(6, 15),
(6, 14);
INSERT 0 5
```


## Combinations data

```
SELECT * FROM combinations;
source | target
    5| 6
    5| 10
    6| 5
    6| 15
    6| 14
(5 rows)
```


## Restrictions

Some functions accept soft restrictions about the segments.
The creation of the restrictions table

```
CREATE TABLE restrictions (
    rid BIGINT NOT NULL,
    to_cost FLOAT,
    target_id BIGINT
    from_edge BIGINT,
    via_path TEXT
);
CREATE TABLE
```

Adding the restrictions

## Restrictions used on pgr_turnRestrictedPath - Experimental

```
CREATE TABLE new_restrictions (
    id SERIAL PRIMARY KEY,
    path BIGINT[],
    cost float
);
CREATE TABLE
INSERT INTO new_restrictions (path, cost) VALUES
(ARRAY[4, 7], 100),
(ARRAY[8, 11], 100),
(ARRAY[7, 10], 100),
(ARRAY[3, 5, 9], 4),
(ARRAY[9, 16], 100);
INSERT 0 5
```


## Images

- Red arrows correspond when cost >0 in the edge table.
- Blue arrows correspond when reverse_cost $>0$ in the edge table.
- Points are outside the graph.
- Click on the graph to enlarge.

Directed graph with cost and reverse_cost
When working with city networks, this is recommended for point of view of vehicles.


Directed, with cost and reverse_cost

Undirected graph with cost and reverse_cost

When working with city networks, this is recommended for point of view of pedestrians.


Undirected, with cost and reverse cost

Directed graph with cost


Directed, with cost

## Undirected graph with cost



Undirected, with cost

## Pick \& Deliver Data

This data example lc101 is from data published at https://www.sintef.no/projectweb/top/pdptw/li-lim-benchmark/ The vehicles

There are 25 vehciles in the problem all with the same characteristics.

```
CREATE TABLE v Ic101(
    id BIGINT NOT NULL primary key,
    capacity BIGINT DEFAULT 200,
    start x FLOAT DEFAULT 30,
    start y FLOAT DEFAULT 50,
    start_open INTEGER DEFAULT 0,
start_close INTEGER DEFAULT 1236);
CREATE TABLE
/* create 25 vehciles *
INSERT INTO v_lc101 (id)
(SELECT * FROM generate_series(1, 25));
INSERT 0 25
```


## The original orders

The data comes in different rows for the pickup and the delivery of the same order.

```
CREATE table Ic101_c(
    id BIGINT not null primary key,
    x DOUBLE PRECISION,
    y DOUBLE PRECISION,
    demand INTEGER,
    open INTEGER,
    close INTEGER,
    service INTEGER,
    pindex BIGINT,
    dindex BIGINT
);
CREATE TABLE
/* the original data */
INSERT INTO Ic101_c
    id, x, y, demand, open, close, service, pindex, dindex) VALUES
( 1, 45, 68, -10, 912, 967, 90, 11, 0),
( 2, 45, 70, -20, 825, 870, 90, 6, 0),
( 3, 42, 66, 10, 65, 146, 90, 0, 75),
(4, 42, 68, -10, 727, 782, 90, 9, 0),
( 5, 42, 65, 10, 15, 67, 90, 0, 7),
( 6, 40, 69, 20, 621, 702, 90, 0, 2),
( 7, 40, 66, -10, 170, 225, 90, 5, 0),
( 8, 38, 68, 20, 255, 324, 90, 0, 10),
(9, 38, 70, 10, 534, 605, 90, 0, 4),
(10, 35, 66, -20, 357, 410, 90, 8, 0),
(11, 35, 69, 10, 448, 505, 90, 0, 1),
(12, 25, 85, -20, 652, 721, 90, 18, 0),
(13, 22, 75, 30, 30, 92, 90, 0, 17),
(14, 22, 85, -40, 567, 620, 90, 16, 0),
(15, 20, 80, -10, 384, 429, 90, 19, 0),
(16, 20, 85, 40, 475, 528, 90, 0, 14),
(17, 18, 75, -30, 99, 148, 90, 13, 0),
(18, 15, 75, 20, 179, 254, 90, 0, 12),
(19, 15, 80, 10, 278, 345, 90, 0, 15),
(20, 30, 50, 10, 10, 73, 90, 0, 24),
(21, 30, 52, -10, 914, 965, 90, 30, 0),
(22, 28, 52, -20, 812, 883, 90, 28, 0),
(23, 28, 55, 10, 732, 777, 0, 0, 103),
(24, 25, 50, -10, 65, 144, 90, 20, 0),
(25, 25, 52, 40, 169, 224, 90, 0, 27),
(26, 25, 55, -10, 622, 701, 90, 29, 0),
(27, 23, 52, -40, 261, 316, 90, 25, 0),
(28, 23, 55, 20, 546, 593, 90, 0, 22),
(29, 20, 50, 10, 358, 405, 90, 0, 26),
(30, 20, 55, 10, 449, 504, 90, 0, 21),
(31, 10, 35, -30, 200, 237, 90, 32, 0),
(32, 10, 40, 30, 31, 100, 90, 0, 31),
(33, 8, 40, 40, 87, 158, 90, 0, 37),
(34, 8, 45, -30, 751, 816, 90, 38, 0),
(35, 5, 35, 10, 283, 344, 90, 0, 39),
(36, 5, 45, 10, 665, 716, 0, 0, 105),
(37, 2, 40, -40, 383, 434, 90, 33, 0),
(38, 0, 40, 30, 479, 522, 90, 0, 34),
39, 0, 45, -10, 567, 624, 90, 35, 0),
(40, 35, 30, -20, 264, 321, 90, 42, 0),
(41, 35, 32, -10, 166, 235, 90, 43, 0),
(42, 33, 32, 20, 68, 149, 90, 0, 40),
(43, 33, 35, 10, 16, 80, 90, 0, 41),
(44, 32, 30, 10, 359, 412, 90, 0, 46),
(45, 30, 30, 10, 541, 600, 90, 0, 48),
(46, 30, 32, -10, 448, 509, 90, 44, 0),
(47, 30, 35, -10, 1054, 1127, 90, 49, 0),
(48, 28, 30, -10, 632, 693, 90, 45, 0),
(49, 28, 35, 10, 1001, 1066, 90, 0, 47),
(50, 26, 32, 10, 815, 880, 90, 0, 52),
(51, 25, 30, 10, 725, 786, 0, 0, 101),
(52, 25, 35, -10, 912, 969, 90, 50, 0),
(53, 44, 5, 20, 286, 347, 90, 0, 58),
(54, 42, 10, 40, 186, 257, 90, 0, 60),
(55, 42, 15, -40, 95, 158, 90, 57, 0),
(56, 40, 5, 30, 385, 436, 90, 0, 59),
(57, 40, 15, 40, 35, 87, 90, 0, 55),
(58, 38, 5, -20, 471, 534, 90, 53, 0),
(59, 38, 15, -30, 651, 740, 90, 56, 0),
(60, 35, 5, -40, 562, 629, 90, 54, 0),
```



The orders

The original data needs to be converted to an appropiate table:

```
WITH deliveries AS (SELECT * FROM Ic101_c WHERE dindex = 0)
SELECT
row_number() over() AS id, p.demand,
p.id as p_node_id, p.x AS p_x, p.y AS p_y, p.open AS p_open, p.close as p_close, p.service as p_service,
d.id as d_node_id, d.x AS d_x, d.y AS d_y, d.open AS d_open, d.close as d_close, d.service as d_service
INTO c_Ic101
FROM deliveries as d JOIN Ic101_c as p ON (d.pindex = p.id);
SELECT 53
SELECT * FROM c_Ic101 LIMIT 1;
id| demand | p_node_id |p_x | p_y | p_open | p_close | p_service | d_node_id |d_x|d_y | d_open | d_close|d_service
1| 10| 3| 42| 66| 65| 146| 90| 75| 45| 65| 997|-------------------------------------------------------------------------------------------
(1 row)
```

Pgrouting Concepts
. Supported versions: Latest (3.3) 3.2 3.1 3.0
-
Unsupported versions: 2.6 2.5 2.4 2.3 2.2 $2.1 \mathbf{2 . 0}$

## pgRouting Concepts

This is a simple guide that go through some of the steps for getting started with pgRouting. This guide covers:

```
- Graphs
Graphs without geometries
Graphs with geometries
Check the Routing Topology
Function's structure
- Function's overloads
```


## Graphs

```
- Graph definition
- Graph with cost
- Graph with cost and reverse_cost
```


## Graph definition

A graph is an ordered pair $\backslash(G=(V, E) \backslash)$ where:

- $\backslash(\mathrm{V} \backslash)$ is a set of vertices, also called nodes.
- $\backslash(E \backslash s u b s e t e q \backslash\{(u, v) \backslash m i d u$, $v \backslash i n ~ V \backslash\} \backslash)$

There are different kinds of graphs:

- Undirected graph
- $\backslash(E \backslash$ subseteq $\backslash\{(u, v) \backslash$ mid $u, v \backslash i n ~ V \backslash\} \backslash)$
- Undirected simple graph
- $\backslash(E$ \subseteq $\backslash\{(\mathrm{u}, \mathrm{v})$ \mid $u$, v \in V , u \neq v $\backslash\} \backslash)$
- Directed graph
- $\backslash(E \backslash$ subseteq $\backslash\{(u, v) \backslash m i d(u, v) \backslash i n(V X V) \backslash\} \backslash)$
- Directed simple graph
- <br>(E \subseteq <br>{( u, v ) \mid (u, v) \in (V X V), u \neq v<br>}<br>)


## Graphs:

- Do not have geometries.
- Some graph theory problems require graphs to have weights, calledcost in pgRouting.

In pgRouting there are several ways to represent a graph on the database:

- With cost
- (id, source, target, cost)
- With cost and reverse_cost
- (id, source, target, cost, reverse_cost)

Where:

| Column | Description |
| :---: | :---: |
| id | Identifier of the edge. Requirement to use the database in a consistent. manner. |
| source | Identifier of a vertex. |
| target | Identifier of a vertex. |
| cost | Weight of the edge (source, target): |
|  | When negative the edge (source, target) do not exist on the graph. cost must exist in the query. |
| reverse_cost | Weight of the edge (target, source) |
|  | - When negative the edge (target, source) do not exist on the graph. |

The decision of the graph to bedirected or undirected is done when executing a pgRouting algorithm.

## Graph with cost

The weighted directed graph, $\backslash\left(G \_d(V, E) \backslash\right)$ :

- Graph data is obtained with a query

SELECT id, source, target, cost FROM edges

- the set of edges $\backslash(E \backslash)$
- $\backslash(E=\backslash\{($ source_\{id $\}$, target_\{id\}, cost_\{id\}) \text\{ when \} cost_\{id\} \ge $0 \backslash\} \backslash)$
- Edges where cost is non negative are part of the graph.
- the set of vertices $\backslash(\mathrm{V} \backslash)$
- $\backslash\left(\mathrm{V}=\backslash\{\right.$ source_\{id $\left.\left.\} \backslash c u p ~ t a r g e t \_\{i d\} \backslash\right\} \backslash\right)$
- All vertices in source and target are part of the graph.


## Directed graph

In a directed graph the edge <br>((source_\{id\}, target_\{id\}, cost_\{id\})<br>) has directionality: <br>(source_\{id\} \rightarrow target_\{id\}<br>)

For the following data:

```
SELECT*
FROM (VALUES (1, 1, 2, 5), (2, 1, 3, -3))
    AS t(id, source, target, cost);
id | source | target | cost
1| 1| 2| 5
2| 1| 3| -3
(2 rows)
```

Edge $\backslash(2 \backslash)(\backslash(1 \backslash$ rightarrow $3 \backslash))$ is not part of the graph.

The data is representing the following graph:


## Undirected graph

In an undirected graph the edge<br>((source_\{id\}, target_\{id\}, cost_\{id\})<br>) does not have directionality: <br>(source_\{id\} $\backslash f r a c\{\backslash ; \backslash \backslash ; ; ; \backslash ;\}\}$ target_\{id\}<br>)

- In terms of a directed graph is like having two edges:<br>(source_\{id\} \leftrightarrow target_\{id\}<br>)

For the following data:

```
SELECT
FROM (VALUES (1, 1, 2, 5), (2, 1, 3, -3))
    AS t(id, source, target, cost);
id | source | target | cost
1|}10|\mp@code{2|}
(2 rows)
```

Edge $\backslash(2 \backslash)(\backslash(1 \backslash \operatorname{frac}\{\backslash ; \backslash ; \backslash \backslash \backslash ; \backslash ;\}\} 3 \backslash))$ is not part of the graph.
The data is representing the following graph:


Graph with cost and reverse_cost
The weighted directed graph, $\backslash\left(\mathrm{G}_{-} \mathrm{d}(\mathrm{V}, \mathrm{E}) \backslash\right)$, is defined by:

- Graph data is obtained with a query

SELECT id, source, target, cost, reverse_cost FROM edges

- The set of edges $\backslash(\mathrm{E} \backslash)$ :
- $\backslash(E=\backslash$ begin $\{$ split $\} \backslash$ begin $\{$ align $\} \&\{\backslash\{$ (source_\{id\}, target_\{id\}, cost_\{id\}) \text\{ when $\}$ cost_\{id\} $>=0 \backslash\}\} \backslash \backslash \& \backslash c u p$
 lend\{split\}<br>)
- Edges <br>((source \rightarrow target)<br>) where cost is non negative are part of the graph.
- Edges $\backslash(($ target $\backslash r i g h t a r r o w ~ s o u r c e) \backslash) ~ w h e r e ~ r e v e r s e \_c o s t ~ i s ~ n o n ~ n e g a t i v e ~ a r e ~ p a r t ~ o f ~ t h e ~ g r a p h . ~$
- The set of vertices $\backslash(\mathrm{V} \backslash)$ :
- $\backslash(\mathrm{V}=\backslash\{$ source_\{id $\} \backslash$ cup target_\{id $\} \backslash\} \backslash)$
- All vertices in source and target are part of the graph.


## Directed graph

In a directed graph both edges have directionality

- edge <br>((source_\{id\}, target_\{id\}, cost_\{id\})<br>) has directionality: <br>(source_\{id\} \rightarrow target_\{id\}<br>)
- edge <br>((target_\{id\}, source_\{id\}, reverse\_cost_\{id\})<br>) has directionality: <br>(target_\{id\} \rightarrow source_\{id\}<br>)

For the following data:

```
SELECT *
FROM (VALUES (1, 1, 2, 5, 2), (2, 1, 3, -3, 4), (3, 2, 3, 7, -1))
    AS t(id, source, target, cost, reverse_cost);
id | source | target | cost | reverse_cost
1| 1| 2| 5| 2
2|
(3 rows)
```

Edges not part of the graph:

- $\backslash(2 \backslash)(\backslash(1$ \rightarrow $3 \backslash))$
- $\quad \backslash(3 \backslash)(\(3 \backslash$ rightarrow $2 \backslash))$

The data is representing the following graph:


## Undirected graph

In a directed graph both edges do not have directionality

- Edge $\backslash(($ source_\{id $\}$, target_\{id\}, cost_\{id\})<br>) is $\backslash($ source_\{id $\} \backslash f r a c\{\backslash ; \backslash ; ; ; \backslash ; \backslash ;\}\}$ target_\{id $\} \backslash)$

- In terms of a directed graph is like having four edges:
- <br>(source_i \leftrightarrow target_i<br>)
- <br>(target_i \leftrightarrow source_i<br>)

For the following data:

```
SELECT*
```

FROM (VALUES (1, 1, 2, 5, 2), (2, 1, 3, -3, 4), (3, 2, 3, 7, -1))
AS t(id, source, target, cost, reverse_cost);
id | source | target | cost | reverse_cost

| $1 \mid$ | $1 \mid$ | $2 \mid$ | $5 \mid$ | 2 |
| ---: | ---: | ---: | ---: | ---: |
| $2 \mid$ | $1 \mid$ | $3 \mid$ | $-3 \mid$ | 4 |
| $3 \mid$ | $2 \mid$ | $3 \mid$ | $7 \mid$ | -1 |
| (3 rows) |  |  |  |  |

Edges not part of the graph:


- <br>(3<br>) (<br>(3 \frac\{i;i;i;i;i;\}\{\} 2\))

The data is representing the following graph:


## Graphs without geometries

Personal relationships, genealogy, file dependency problems can be solved using pgRouting. Those problems, normally, do not come with geometries associated with the graph.

```
- Wiki example
```

- Prepare the database
- Create a table
- Insert the data
- Find the shortest path
- Vertex information


## Wiki example

Solve the example problem taken from wikipedia):


Where:

- Problem is to find the shortest path from $\backslash(1)$ to $\backslash(5 \backslash)$.
- Is an undirected graph.
- Although visually looks like to have geometries, the drawing is not to scale.
- No geometries associated to the vertices or edges
- Has 6 vertices $\backslash(\backslash\{1,2,3,4,5,6 \backslash\} \backslash)$
- Has 9 edges:
$\backslash(\backslash$ begin $\{$ split $\} \backslash$ begin $\{$ align $\} E=\& \backslash\{(1,2,7),(1,3,9),(1,6,14), \backslash \ \&(2,3,10),(2,4,13), \backslash \backslash(3,4,11),(3,6,2), \backslash \backslash \&(4,5,6), \backslash \backslash \&$ $(5,6,9) \backslash\}$ \end \{align\} \end\{split\}\) }
- The graph can be represented in many ways for example:


Prepare the database

Create a database for the example, access the database and install pgRouting:

```
$ createdb wiki
$ psql wiki
wiki =# CREATE EXTENSION pgRouting CASCADE;
```


## Create a table

The basic elements needed to perform basic routing on an undirected graph are:

| Column | Type | Description |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER | Identifier of the second end point vertex of the <br> edge. |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Using this table design for this example:

```
CREATE TABLE wiki (
    id SERIAL,
    source INTEGER,
    target INTEGER,
    cost INTEGER);
CREATE TABLE
```


## Insert the data

```
INSERT INTO wiki (source, target, cost) VALUES
(1, 2, 7), (1, 3, 9), (1, 6, 14),
(2, 3, 10), (2, 4, 15),
(3, 6, 2), (3, 4, 11),
(4, 5, 6),
(5, 6, 9);
INSERT09
```


## Find the shortest path

To solve this example pgr_dijkstra is used:

```
SELECT * FROM pgr_dijkstra(
    'SELECT id, source, target, cost FROM wiki',
    1,5,false);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 1| 2| 9| 0
2| 2| 3| 6| 2| 9
3| 3| 6| 9| 9| 11
(4 rows)
```

To go from $\backslash(1 \backslash)$ to $\backslash(5 \backslash)$ the path goes thru the following vertices: $\backslash(1$ rightarrow 3 \rightarrow 6 \rightarrow 5<br>)


Vertex information

To obtain the vertices information, use pgr_extractVertices - Proposed

```
SELECT id, in_edges,out_edges
FROM pgr_extractVertices('SELECT id, source, target FROM wiki');
id | in_edges | out_edges
3|{2,4} |{6,7}
5|{8} |{9}
4|{5,7} |{8}
2|{1} |{4,5}
1| |{1,2,3}
6 |{3,6,9}
(6 rows)
```


## Graphs with geometries

```
- Create a routing Database
- Load Data
- Build a routing topology
- Adjust costs
    - Update costs to length of geometry
    - Update costs based on codes
```


## Create a routing Database

The first step is to create a database and load pgRouting in the database.

Typically create a database for each project.
Once having the database to work in, load your data and build the routing application in that database.

```
createdb sampledata
psql sampledata -c "CREATE EXTENSION pgrouting CASCADE"
```


## Load Data

There are several ways to load your data into pgRouting.

- Manually creating a database.
- Graphs without geometries
- Sample Data: a small graph used on the documentation examples
- Using osm2pgrouting

There are various open source tools that can help, like:

```
shp2pgsql:
    - postgresql shapefile loader
ogr2ogr:
    - vector data conversion utility
osm2pgsq|:
    - load OSM data into postgresql topology needs to be adjusted.
- Breakup a segments on each segment-segment intersection
- When missing, add columns and assign values tosource, target, cost, reverse_cost.
- Connect a disconnected graph.
- Create the complete graph topology
- Create one or more graphs based on the application to be developed.
```

Please note that these tools will not import the data in a structure compatible with pgRouting and when this happens the

- Create a contracted graph for the high speed roads
- Create graphs per state/country

In few words:

Prepare the graph
What and how to prepare the graph, will depend on the application and/or on the quality of the data and/or on how close the information is to have a topology usable by pgRouting and/or some other factors not mentioned.

The steps to prepare the graph involve geometry operations usingPostGIS and some others involve graph operations like pgr_contraction to contract a graph.

The workshop has a step by step on how to prepare a graph using Open Street Map data, for a small application.
The use of indexes on the database design in general:

- Have the geometries indexed.
- Have the identifiers columns indexed.

Please consult the PostgreSQL documentation and the PostGIS documentation.

## Build a routing topology

The basic information to use the majority of the pgRouting functionsid, source, target, cost, [reverse_cost] is what in pgRouting is called the routing topology.
reverse_cost is optional but strongly recommended to have in order to reduce the size of the database due to the size of the geometry columns. Having said that, in this documentation reverse_cost is used in this documentation.

When the data comes with geometries and there is no routing topology, then this step is needed.
All the start and end vertices of the geometries need an identifier that is to be stored in isource and target columns of the table of the data. Likewise, cost and reverse_cost need to have the value of traversing the edge in both directions.

If the columns do not exist they need to be added to the table in question. (seeALTER TABLE)
The function pgr_extractVertices - Proposed is used to create a vertices table based on the edge identifier and the geometry of the edge of the graph.

Finally using the data stored on the vertices tables the source and target are filled up.
See Sample Data for an example for building a topology.
Data coming from OSM and usingosm2pgrouting as an import tool, comes with the routing topology. See an example of using osm2pgrouting on the workshop.

## Adjust costs

For this example the cost and reverse_cost values are going to be the double of the length of the geometry.

```
Update costs to length of geometry
```

Suppose that cost and reverse_cost columns in the sample data represent:

- $\backslash(1 \backslash)$ when the edge exists in the graph
- $\backslash(-1 \backslash)$ when the edge does not exist in the graph

Using that information updating to the length of the geometries:

```
UPDATE edges SET
cost = sign(cost) * ST_length(geom) * 2,
reverse_cost = sign(reverse_cost) * ST_length(geom) * 2;
UPDATE }1
```


## Which gives the following results:



Note that to be able to follow the documentation examples, everything is based on the original graph.
Returning to the original data:

```
UPDATE edges SET
cost = sign(cost)
reverse_cost = sign(reverse_cost);
UPDATE }1
```


## Update costs based on codes

Other datasets, can have a column with values like

- FT vehicle flow on the direction of the geometry
- TF vehicle flow opposite of the direction of the geometry
- B vehicle flow on both directions

Preparing a code column for the example:

```
ALTER TABLE edges ADD COLUMN direction TEXT;
ALTER TABLE
UPDATE edges SET
direction = CASE WHEN (cost>0 AND reverse_cost>0) THEN 'B'
    WHEN (cost>0 AND reverse_cost<0) THEN 'FT'
    WHEN (cost<0 AND reverse_cost>0) THEN 'TF
    ELSE " END;
UPDATE }1
```

Adjusting the costs based on the codes:

```
UPDATE edges SET
cost = CASE WHEN (direction = 'B' OR direction = 'FT')
    THEN ST_length(geom) * 2
    ELSE -1 END,
reverse_cost = CASE WHEN (direction = 'B' OR direction = 'TF')
    THEN ST_length(geom) * 2
    ELSE -1 END;
UPDATE 18
```

Which gives the following results:

```
SELECT id, cost, reverse_cost FROM edges;
id| cost | reverse_cos
\begin{tabular}{ccc}
\(6 \mid\) & \(2 \mid\) & 2 \\
\(7 \mid\) & \(2 \mid\) & 2 \\
\(4 \mid\) & \(2 \mid\) & 2 \\
\(5 \mid\) & \(2 \mid\) & -1 \\
\(8 \mid\) & \(2 \mid\) & 2 \\
\(12 \mid\) & \(2 \mid\) & -1 \\
\(11 \mid\) & \(2 \mid\) & -1 \\
\(10 \mid\) & \(2 \mid\) &
\end{tabular}
    2.999999999998| 2.999999999998
|| 2| 2
| 3.40000000000000004 | 3.40000000000000004
| 2 | \(\quad-1\)
\(6 |\)\begin{tabular}{ll}
\(2 \mid\) & 2
\end{tabular}
|) 2|
|i|
2|
(18 rows)
```

Returning to the original data:

```
UPDATE edges SET
cost = sign(cost),
reverse_cost = sign(reverse_cost);
UPDATE 18
ALTER TABLE edges DROP COLUMN direction;
ALTER TABLE
```


## Check the Routing Topology

```
- Crossing edges
```

    - Adding split edges
    - Adding new vertices
    - Updating edges topology
    - Removing the surplus edges
    - Updating vertices topology
    - Checking for crossing edges
    - Disconnected graphs
    - Prepare storage for connection information
    - Save the vertices connection information
    - Save the edges connection information
    - Get the closest vertex
    - Connecting components
    - Checking components
- Contraction of a graph
    - Dead ends
    - Linear edges

There are lots of possible problems in a graph.

- The data used may not have been designed with routing in mind.
- A graph has some very specific requirements.
- The graph is disconnected.
- There are unwanted intersections.
- The graph is too large and needs to be contracted
- A sub graph is needed for the application.
- and many other problems that the pgRouting user, that is the application developer might encounter.


## Crossing edges

To get the crossing edges:

```
SELECT a.id, b.id
FROM edges AS a, edges AS b
WHERE a.id < b.id AND st_crosses(a.geom, b.geom);
id id
--------
13|18
(1 row)
```



That information is correct, for example, when in terms of vehicles, is it a tunnel or bride crossing over another road.
It might be incorrect, for example:

1. When it is actually an intersection of roads, where vehicles can make turns.
2. When in terms of electrical lines, the electrical line is able to switch roads even on a tunnel or bridge.

When it is incorrect, it needs fixing:

1. For vehicles and pedestrians

- If the data comes from OSM and was imported to the database usingosm2pgrouting, the fix needs to be done in theOSM portal and the data imported again.
- In general when the data comes from a supplier that has the data prepared for routing vehicles, and there is a problem, the data is to be fixed from the supplier

2. For very specific applications

- The data is correct when from the point of view of routing vehicles or pedestrians.
- The data needs a local fix for the specific application.

Once analyzed one by one the crossings, for the ones that need a local fix, the edges need to besplit.

```
SELECT ST_AsText((ST_Dump(ST_Split(a.geom, b.geom))).geom)
FROM edges AS a, edges AS b
WHERE a.id = 13 AND b.id = 18
UNION
SELECT ST_AsText((ST_Dump(ST_Split(b.geom, a.geom))).geom)
FROM edges AS a, edges AS b
WHERE a.id = 13 AND b.id = 18;
    st_astext
LINESTRING(3.5 2.3,3.5 3)
LINESTRING(3 3,3.5 3)
LINESTRING(3.5 3,4 3)
LINESTRING(3.5 3,3.5 4)
(4 rows)
```

The new edges need to be added to the edges table, the rest of the attributes need to be updated in the new edges, the old edges need to be removed and the routing topology needs to be updated.

Adding split edges

For each pair of crossing edges a process similar to this one must be performed.
The columns inserted and the way are calculated are based on the application. For example, if the edges have a trainame, then that column is to be copied.

## For pgRouting calculations

- factor based on the position of the intersection of the edges can be used to adjust thecost and reverse_cost columns.
- Capacity information, used on the Flow - Family of functions functions does not need to change when splitting edges.

```
WITH
first_edge AS (
    SELECT (ST_Dump(ST_Split(a.geom, b.geom))).path[1],
    (ST_Dump(ST_Split(a.geom, b.geom))).geom,
    ST_LineLocatePoint(a.geom,ST Intersection(a.geom,b.geom)) AS factor
    FROM edges AS a, edges AS b
WHERE a.id = 13 AND b.id = 18),
first_segments AS (
    SELECT path, first_edge.geom,
    capacity, reverse_capacity,
    CASE WHEN path=1 THEN factor * cost
        ELSE (1 - factor) * cost END AS cost,
    CASE WHEN path=1 THEN factor * reverse_cost
        ELSE (1 - factor) * reverse_cost END AS reverse_cost
    FROM first_edge , edges WHERE id = 13),
second_edge AS 
    SELECT (ST_Dump(ST_Split(b.geom, a.geom))).path[1],
    (ST_Dump(ST_Split(b.geom, a.geom))).geom,
    ST_LineLocatePoint(b.geom,ST_Intersection(a.geom,b.geom)) AS factor
    FROM edges AS a, edges AS b
WHERE a.id = 13 AND b.id = 18),
second_segments AS (
    SELECT path, second_edge.geom,
    capacity, reverse_capacity,
    CASE WHEN path=1 THEN factor * cost
            ELSE (1 - factor) * cost END AS cost,
    CASE WHEN path=1 THEN factor * reverse_cost
        ELSE (1 - factor) * reverse_cost END AS reverse_cost
    FROM second_edge, edges WHERE id=18),
all_segments A\overline{S}
    SELECT * FROM first_segments
    UNION
    SELECT * FROM second_segments)
INSERT INTO edges
    (capacity, reverse_capacity,
    cost, reverse cost,
    x1, y1, x2, y2,
    geom)
(SELECT capacity, reverse_capacity, cost, reverse_cost,
ST X(ST_StartPoint(geom)), ST_Y(ST_StartPoint(geom)),
ST_X(ST_EndPoint(geom)), ST_Y(ST_EndPoint(geom)),
    geom
    FROM all_segments)
INSERT 0 4
```


## Adding new vertices

After adding all the split edges required by the application, the newly created vertices need to be added to the vertices table.

```
INSERT INTO vertices (in_edges, out_edges, x, y, geom)
(SELECT nv.in_edges, nv.out_edges, nv.x, nv.y, nv.geom
FROM pgr_extractVertices('SELECT id, geom FROM edges') AS nv
LEFT JOIN vertices AS v USING(geom) WHERE v.geom IS NULL);
```

INSERT 01

## Updating edges topology

```
/* -- set the source information */
UPDATE edges AS e
SET source = v.id
FROM vertices AS v
WHERE source IS NULL AND ST_StartPoint(e.geom) = v.geom;
UPDATE 4
/* -- set the target information */
UPDATE edges AS e
SET target = v.id
FROM vertices AS v
WHERE target IS NULL AND ST_EndPoint(e.geom) = v.geom;
UPDATE }
```


## Removing the surplus edges

Once all significant information needed by the application has been transported to the new edges, then the crossing edges can be deleted.

```
DELETE FROM edges WHERE id IN (13, 18);
DELETE 2
```

There are other options to do this task, like creating a view, or a materialized view.

```
UPDATE vertices AS v SET
in_edges = nv.in_edges, out_edges = nv.out_edges
FROM (SELECT * FROM pgr_extractVertices('SELECT id, geom FROM edges')) AS nv
WHERE v.geom = nv.geom
UPDATE 18
```


## Checking for crossing edges

There are no crossing edges on the graph.

```
SELECT a.id, b.id
FROM edges AS a, edges AS b
WHERE a.id < b.id AND st_crosses(a.geom, b.geom);
id | id
---+----
```


## Disconnected graphs

To get the graph connectivity:


In this example, the component $\backslash(2 \backslash)$ consists of vertices $\backslash(\backslash\{2,4 \backslash\} \backslash)$ and both vertices are also part of the dead end result set.

This graph needs to be connected.

## Note

With the original graph of this documentation, there would be 3 components as the crossing edge in this graph is a different component.

## Prepare storage for connection information

```
ALTER TABLE vertices ADD COLUMN component BIGINT
ALTER TABLE
ALTER TABLE edges ADD COLUMN component BIGINT;
ALTER TABLE
```


## Save the vertices connection information

```
UPDATE vertices SET component = c.component
FROM (SELECT * FROM pgr_connectedComponents(
    'SELECT id, source, target, cost, reverse_cost FROM edges'
)) AS c
WHERE id = node;
UPDATE }1
```

```
UPDATE edges SET component = v.component
FROM (SELECT id, component FROM vertices) AS v
WHERE source = v.id;
UPDATE 20
```


## Get the closest vertex

The closest vertex to component $\backslash(1 \backslash)$ is vertex $\backslash(4 \backslash)$. And the closest edge to vertex $\backslash(4 \backslash)$ is edge $\backslash(14 \backslash)$.

```
WITH
edges_sql AS (SELECT id, geom FROM edges WHERE component = 1),
point_sql AS (SELECT geom AS point FROM vertices WHERE component = 2),
results AS (
    SELECT
    id::BIGINT AS edge id,
    ST_LineLocatePoint(geom, point) AS fraction,
    CA\overline{SE WHEN ST_Intersects(ST_Buffer(geom, 2, 'side=right endcap=flat'), point)}
        THEN 'r'
        ELSE 'I' END::CHAR AS side,
    geom <-> point AS distance,
    point,
    ST_MakeLine(point, ST_ClosestPoint(geom, point)) AS new_line
    FROM edges_sql, point_sq
    WHERE ST_DWithin(geom, point, 2)
    ORDER BY geom <-> point),
prepare_cap AS (
    SELECT row number() OVER (PARTITION BY point ORDER BY point, distance) AS rn, *
    FROM results)
cap AS (
    SELECT edge_id, fraction, side, distance, point, new_line
    FROM prepare_cap
    WHERE rn <= 1
)
SELECT edge_id, fraction, side, distance, point AS geom, new_line AS edge, id AS closest_vertex
INTO closest
FROM cap JOIN vertices ON (point = geom) ORDER BY distance LIMIT 1;
SELECT }
```

The edge can be used to connect the components, using thefraction information about the edge $\backslash(14 \backslash)$ to split the connecting edge.

Connecting components
There are three basic ways to connect the components

- From the vertex to the starting point of the edge
- From the vertex to the ending point of the edge
- From the vertex to the closest vertex on the edge - This solution requires the edge to be split.

The following query shows the three ways to connect the components:

```
WITH
info AS (
    SELECT
    edge_id, fraction, side, distance, ce.geom, edge, v.id AS closest,
    source, target, capacity, reverse_capacity, e.geom AS e_geom
    FROM closest AS ce
    JOIN vertices AS v USING (geom)
    JOIN edges AS e ON (edge_id = e.id)
ORDER BY distance LIMIT 1),
three_options AS (
    SELECT
    closest AS source, target, 0 AS cost, 0 AS reverse_cost,
    capacity, reverse_capacity,
    ST_X(geom) AS x1, ST_Y(geom) AS y1,
    ST_X(ST_EndPoint(e_geom)) AS x2, ST_Y(ST_EndPoint(e_geom)) AS y2,
    ST_MakeLine(geom, ST_EndPoint(e_geom)) AS geom
    FROM info
    UNION
    SELECT closest, source, 0, 0, capacity, reverse_capacity,
    ST_X(geom) AS x1, ST_Y(geom) AS y1,
    ST_X(ST_StartPoint(e_geom)) AS x2, ST_Y(ST_StartPoint(e_geom)) AS y2,
    ST_MakeLine(info.geom, ST_StartPoint(e_geom))
FROM info
/*
UNION
    -- This option requires splitting the edge
    SELECT closest, NULL, 0, 0, capacity, reverse_capacity,
    ST_X(geom) AS x1, ST_Y(geom) AS y1,
    ST_X(ST_EndPoint(edge)) AS x2, ST_Y(ST_EndPoint(edge)) AS y2,
    edge
    FROM info *
)
INSERT INTO edges
    (source, target,
    cost, reverse_cost,
    capacity, reverse_capacity,
    x1, y1, x2, y2,
    geom)
(SELECT
    source, target, cost, reverse_cost, capacity, reverse_capacity,
    x1,y1, x2, y2, geom
FROM three_options);
INSERT 0 2
```


## Checking components

Ignoring the edge that requires further work. The graph is now fully connected as there is only one component.

```
SELECT * FROM pgr_connectedComponents(
    'SELECT id, source, target, cost, reverse_cost FROM edges'
);
seq | component | node
| 1
1|}
1| 3
1| 4
1| 5
1| 6
1| 7
1| 8
1| 9
1| }1
1| 11
1| }1
1| 13
1| 14
1| 15
1| 16
1| 17
1| 18
(18 rows)
```


## Contraction of a graph

The graph can be reduced in size usingContraction - Family of functions
When to contract will depend on the size of the graph, processing times, correctness of the data, on the final application, or any other factor not mentioned.

A fairly good method of finding out if contraction can be useful is because of the number of dead ends and/or the number of linear edges.

A complete method on how to contract and how to use the contracted graph is described onContraction - Family of

## functions

Dead ends

To get the dead ends:

```
SELECT id FROM vertices
WHERE array_length(in_edges || out_edges, 1) = 1;
id
----
5
5
13
14
2
4
(7 rows)
```

That information is correct, for example, when the dead end is on the limit of the imported graph.
Visually node $\backslash(4 \backslash)$ looks to be as start/ending of 3 edges, but it is not.

## Is that correct?

- Is there such a small curb:
- That does not allow a vehicle to use that visual intersection?
- Is the application for pedestrians and therefore the pedestrian can easily walk on the small curb?
- Is the application for the electricity and the electrical lines than can easily be extended on top of the small curb?
- Is there a big cliff and from eagles view look like the dead end is close to the segment?

When there are many dead ends, to speed up, theContraction - Family of functions functions can be used to divide the problem.

## Linear edges

To get the linear edges:

```
SELECT id FROM vertices
WHERE array_length(in_edges || out_edges, 1) = 2;
id
3
15
17
(3 rows)
```

This information is correct, for example, when the application is taking into account speed bumps, stop signals.
When there are many linear edges, to speed up, theContraction - Family of functions functions can be used to divide the problem.

## Function's structure

Once the graph preparation work has been done above, it is time to use a
The general form of a pgRouting function call is:
pgr_<name>(Inner queries, parameters, [ Optional parameters)

## Where:

- Inner queries: Are compulsory parameters that areTEXT strings containing SQL queries.
- parameters: Additional compulsory parameters needed by the function.
- Optional parameters: Are non compulsory named parameters that have a default value when omitted.

The compulsory parameters are positional parameters, the optional parameters are named parameters.
For example, for this pgr_dijkstra signature:
pgr_dijkstra(Edges SQL, start vid, end vid [, directed])

- Edges SQL:
- Is the first parameter.
- It is compulsory.
- It is an inner query.
- It has no name, so Edges SQL gives an idea of what kind of inner query needs to be used
- start vid:
- Is the second parameter.
- It is compulsory.
- It has no name, so start vid gives an idea of what the second parameter's value should contain.
- end vid
- Is the third parameter.
- It is compulsory.
- It has no name, so end vid gives an idea of what the third parameter's value should contain
- directed
- Is the fourth parameter.
- It is optional.
- It has a name.

The full description of the parameters are found on theParameters section of each function.

## Function's overloads

A function might have different overloads. The most common are called:

- One to One
- One to Many
- Many to One
- Many to Many
- Combinations

Depending on the overload the parameters types change.

- One: ANY-INTEGER
- Many: ARRAY [ANY-INTEGER]

Depending of the function the overloads may vary. But the concept of parameter type change remains the same.

## One to One

When routing from:

- From one starting vertex
- to one ending vertex


## One to Many

When routing from:

- From one starting vertex
- to many ending vertices


## Many to One

When routing from:

- From many starting vertices
- to one ending vertex

Many to Many
When routing from:

- From many starting vertices
- to many ending vertices


## Combinations

When routing from:

- From many different starting vertices
- to many different ending vertices
- Every tuple specifies a pair of a start vertex and an end vertex
- Users can define the combinations as desired.


## Inner Queries

```
- Edges SQL
    - General
    - General without id
    - General with (X,Y)
    - Flow
- Combinations SQL
- Restrictions SQL
- Points SQL
```

There are several kinds of valid inner queries and also the columns returned are depending of the function. Which kind of inner query will depend on the function(s) requirements. To simplify variety of types, ANY-INTEGER and ANY-NUMERICAL is used.

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Edges SQL

General

## Edges SQL for

- Dijkstra - Family of functions
- withPoints - Family of functions
- Bidirectional Dijkstra - Family of functions
- Components - Family of functions
- Kruskal - Family of functions
- Prim - Family of functions
- Some uncategorised functions

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL -1 | Weight of the edge (target, source) |
|  |  |  |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

General without id

## Edges SQL for

All Pairs - Family of Functions

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |  |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |  |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |  |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |
|  |  |  | When negative: edge (target, source) does not exist, therefore it's <br>  |
|  |  |  | not part of the graph. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

General with (X,Y)

## Edges SQL for

- $A^{*}$ - Family of functions
- Bidirectional $A^{*}$ - Family of functions

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER |  | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) <br> - When negative: edge (source, target) does not exist, therefore it's not part of the graph. |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source), <br> - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |
| x1 | ANY-NUMERICAL |  | $X$ coordinate of source vertex. |
| y1 | ANY-NUMERICAL |  | $Y$ coordinate of source vertex. |
| x2 | ANY-NUMERICAL |  | X coordinate of target vertex. |
| y2 | ANY-NUMERICAL |  | Y coordinate of target vertex. |

Where:

```
ANY-INTEGER:
SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT
```

Flow

Edges SQL for Flow - Family of functions

## Edges SQL for

pgr_pushRelabel
pgr_edmondsKarp

- pgr_boykovKolmogorov

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER |  | Identifier of the second end point vertex of the edge. |
| capacity | ANY-INTEGER |  | Weight of the edge (source, target) |
| reverse_capacity | ANY-INTEGER | -1 | Weight of the edge (target, source) |

- When negative: edge (target, source) does not exist, therefore it's not part of the graph.

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Edges SQL for the following functions of Flow - Family of functions

- pgr_maxFlowMinCost - Experimental

| Column | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER |  | Identifier of the second end point vertex of the edge. |
| capacity | ANY-INTEGER |  | Capacity of the edge (source, target) |
|  |  |  | - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |
| reverse_capacity | ANY-INTEGER | -1 | Capacity of the edge (target, source) |
|  |  |  | - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) if it exist |
| reverse_cost | ANY-NUMERICAL | $\(-1 \backslash)$ | Weight of the edge (target, source) if it exist |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT
Combinations SQL

## Used on combination signatures

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- | Identifier of the departure vertex. |
|  | INTEGER |  |
| target | ANY- <br> INTEGER | Identifier of the arrival vertex. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

Restrictions SQL

| Column | Type | Description |
| :--- | :--- | :--- |
| path | ARRAY [ANY-INTEGER] | Sequence of edge identifiers that form a path that is not allowed to be taken. - <br> Empty arrays or NULL arrays are ignored. - Arrays that have aNULL element will <br> raise an exception. |
| Cost | ANY-NUMERICAL | Cost of taking the forbidden path. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Points SQL

## Points SQL for

- withPoints - Family of functions

| Parameter | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| pid | ANY-INTEGER | value | Identifier of the point. |

- Use with positive value, as internally will be converted to negative value
- If column is present, it can not be NULL.
- If column is not present, a sequential negativevalue will be given automatically.

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| fraction | ANY-NUMERICAL |  | Value in $\langle 0,1\rangle$ that indicates the relative postition from the first end point of the edge. |
| side | CHAR | b | Value in [b, r, l, NULL] indicating if the point is: |
|  |  |  | - In the right $r$, <br> - In the left I, <br> - In both sides b, NULL |

## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Parameters

The main parameter of the majority of the pgRouting functions is a query that selects the edges of the graph.

| Parameter | Type | Description |  |
| :--- | :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |  |

Depending on the family or category of a function it will have additional parameters, some of them are compulsory and some are optional.

The compulsory parameters are nameless and must be given in the required order. The optional parameters are named parameters and will have a default value.

## Parameters for the Via functions

- pgr_dijkstraVia - Proposed

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| Edges SQL | TEXT |  | SQL query as described. |
| via vertices | ARRAY [ANY-INTEGER] |  | Array of ordered vertices identifiers that are going to be visited. |
| directed | BOOLEAN | true | When true Graph is considered Directed <br> - When false the graph is considered as Undirected. |
| strict | BOOLEAN | false | When true if a path is missing stops and returnsEMPTY SET <br> When false ignores missing paths returning all paths found |
| U_turn_on_edge | BOOLEAN | true | When true departing from a visited vertex will not try to avoid using the edge used to reach it. In other words, U turn using the edge with same identifier is allowed. <br> - When false when a departing from a visited vertex tries to avoid using the edge used to reach it. In other words, $U$ turn using the edge with same identifier is used when no other path is found. |

## For the TRSP functions

- pgr_trsp - Turn Restriction Shortest Path (TRSP)

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | SQL query as described. |
| Restrictions SQL | TEXT | SQL query as described. |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | ANY-INTEGER | Identifier of the departure vertex. |
| start vids | ARRAY [ANY-INTEGER] | Array of identifiers of destination vertices. |
| end vid | ANY-INTEGER | Identifier of the departure vertex. |
| end vids | ARRAY [ANY-INTEGER] | Array of identifiers of destination vertices. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Return columns

## - Return columns for a path

- Multiple paths
- Selective for multiple paths
- Non selective for multiple paths
- Return columns for cost functions
- Return columns for flow functions
- Return columns for spanning tree functions

There are several kinds of columns returned are depending of the function.

Return columns for a path

## Used on functions that return one path solution

Returns set of (seq, path_seq [, start_vid] [, end_vid], node, edge, cost, agg_cost)

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. |
| path_seq | INTEGER | Relative position in the path. Has value1 for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex. Returned when multiple starting vetrices are in the query. Many to One Many to Many |
| end_vid | BIGINT | Identifier of the ending vertex. Returned when multiple ending vertices are in the query. <br> One to Many <br> Many to Many |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for the last node of the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | FLOAT | Aggregate cost from start_vid to node. |

## Used on functions the following:

- pgr_withPoints - Proposed

Returns set of (seq, path_seq [, start_pid] [, end_pid], node, edge, cost, agg_cost)

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. |
| path_seq | INTEGER | Relative position in the path. <br> 1 For the first row of the path. |
| start_pid | BIGINT | Identifier of a starting vertex/point of the path. <br> - When positive is the identifier of the starting vertex. <br> - When negative is the identifier of the starting point. <br> - Returned on Many to One and Many to Many |
| end_pid | BIGINT | Identifier of an ending vertex/point of the path. <br> - When positive is the identifier of the ending vertex. <br> - When negative is the identifier of the ending point. <br> - Returned on One to Many and Many to Many |
| node | BIGINT | Identifier of the node in the path fromstart_pid to end_pid. <br> - When positive is the identifier of the a vertex. <br> - When negative is the identifier of the a point. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence. <br> -1 for the last row of the path. |


| Column | Type | Description |
| :--- | :--- | :--- |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | FLOAT | $\mathbf{0}$ For the first row of the path. |
|  |  | Aggregate cost from start_vid to node. |
|  | $0 \quad \mathbf{0}$ For the first row of the path. |  |

## Used on functions the following:

- pgr_dijkstraNear - Proposed

Returns (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Sequential value starting from $\mathbf{1}$. |
| path_seq | INTEGER | Relative position in the path. Has value $\mathbf{1}$ for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex of the current path. |
| end_vid | BIGINT | Identifier of the ending vertex of the current path. |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for <br> the last node of the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | FLOAT | Aggregate cost from start_vid to node. |

Multiple paths

Selective for multiple paths.

The columns depend on the function call.
Set of (seq, path_id, path_seq [, start_vid] [, end_vid], node, edge, cost, agg_cost)

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. |
| path_id | INTEGER |  |
|  |  | - Has value $\mathbf{1}$ for the first of a path fromstart_vid to end_vid. |
| path_seq | INTEGER | Relative position in the path. Has value 1 for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex. Returned when multiple starting vetrices are in the query. |
|  |  | Many to One <br> Many to Many <br> Combinations |
| end_vid | BIGINT | Identifier of the ending vertex. Returned when multiple ending vertices are in the query. |
|  |  | - One to Many <br> - Many to Many <br> - Combinations |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for the last node of the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | FLOAT | Aggregate cost from start_vid to node. |

Non selective for multiple paths

Regardless of the call, al the columns are returned.

- pgr_trsp - Turn Restriction Shortest Path (TRSP)

Returns set of (seq, path_id, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Sequential value starting from $\mathbf{1}$. |


| Column | Type | Description |
| :--- | :--- | :--- |
| path_id | INTEGER | Path identifier. |
|  |  | Has value $\mathbf{1}$ for the first of a path fromstart_vid to end_vid. |
| path_seq | INTEGER | Relative position in the path. Has value $\mathbf{1}$ for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence. - 1 for <br> the last node of the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | Aggregate cost from start_vid to node. |  |

Return columns for cost functions

## Used in the following

- Cost - Category
- Cost Matrix - Category
- All Pairs - Family of Functions

Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

Note
When start_vid or end_vid columns have negative values, the identifier is for a Point.

Return columns for flow functions

## Edges SQL for the following

Flow - Family of functions

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INT | Sequential value starting from $\mathbf{1}$. |
| edge | BIGINT | Identifier of the edge in the original query (edges_sql). |
| start_vid | BIGINT | Identifier of the first end point vertex of the edge. |
| end_vid | BIGINT | Identifier of the second end point vertex of the edge. |
| flow | BIGINT | Flow through the edge in the direction (start_vid, end_vid). |
| residual_capacity | BIGINT | Residual capacity of the edge in the direction (start_vid, <br> end_vid). |

## Edges SQL for the following functions of Flow - Family of functions

- pgr_maxFlowMinCost - Experimental

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INT | Sequential value starting from $\mathbf{1}$. |
| edge | BIGINT | Identifier of the edge in the original query (edges_sql). |
| source | BIGINT | Identifier of the first end point vertex of the edge. |
| target | BIGINT | Identifier of the second end point vertex of the edge. |
| flow | BIGINT | Flow through the edge in the direction (source, target). |
| residual_capacity | BIGINT | Residual capacity of the edge in the direction (source, target). |
| cost | FLOAT | The cost of sending this flow through the edge in the direction (source, |
|  |  | target). |
| agg_cost | FLOAT | The aggregate cost. |

## Edges SQL for the following

- pgr_prim
- pgr_kruskal

Returns SET OF (edge, cost)

| Column | Type | Description |
| :--- | :--- | :--- | :--- |
| edge | BIGINT | Identifier of the edge. |
| cost | FLOAT | Cost to traverse the <br> edge. |

## Performance Tips

## - For the Routing functions

For the Routing functions

To get faster results bound the queries to an area of interest of routing.
In this example Use an inner query SQL that does not include some edges in the routing function and is within the area of the results.

```
SELECT * FROM pgr_dijkstra($$
    SELECT id, source, target, cost, reverse_cost from edges
    WHERE geom && (SELECT st_buffer(geom, 1) AS myarea
    FROM edges WHERE id = 2)$$,
1, 2);
seq | path_seq | node | edge | cost | agg_cost
(0 rows)
```


## How to contribute

## Wiki

- Edit an existing pgRouting Wiki page.
- Or create a new Wiki page
- Create a page on the pgRouting Wiki
- Give the title an appropriate name
- Example


## Adding Functionaity to pgRouting

Consult the developer's documentation

## Indices and tables

- Index
- Search Page


## Function Families

- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2 $2.1 \mathbf{2 . 0}$


## Function Families

## All Pairs - Family of Functions

- pgr_floydWarshall - Floyd-Warshall's algorithm.
- pgr_johnson - Johnson's algorithm


## A* - Family of functions

pgr_aStar - A* algorithm for the shortest path.

- pgr_aStarCost - Get the aggregate cost of the shortest paths.
- pgr_aStarCostMatrix - Get the cost matrix of the shortest paths.


## Bidirectional A* - Family of functions

- pgr_bdAstar - Bidirectional A* algorithm for obtaining paths.
- pgr_bdAstarCost - Bidirectional A* algorithm to calculate the cost of the paths.
- pgr_bdAstarCostMatrix - Bidirectional $A^{*}$ algorithm to calculate a cost matrix of paths.


## Bidirectional Dijkstra - Family of functions

- pgr_bdDijkstra - Bidirectional Dijkstra algorithm for the shortest paths.
- pgr_bdDijkstraCost - Bidirectional Dijkstra to calculate the cost of the shortest paths
- pgr_bdDijkstraCostMatrix - Bidirectional Dijkstra algorithm to create a matrix of costs of the shortest paths.


## Components - Family of functions

pgr_connectedComponents - Connected components of an undirected graph.
pgr_strongComponents - Strongly connected components of a directed graph.

- pgr_biconnectedComponents - Biconnected components of an undirected graph.
- pgr_articulationPoints - Articulation points of an undirected graph.
- pgr_bridges - Bridges of an undirected graph.


## Contraction - Family of functions

- pgr_contraction


## Dijkstra - Family of functions

pgr_dijkstra - Dijkstra's algorithm for the shortest paths.
pgr_dijkstraCost - Get the aggregate cost of the shortest paths.
pgr_dijkstraCostMatrix - Use pgr_dijkstra to create a costs matrix. pgr_drivingDistance - Use pgr_dijkstra to calculate catchament information. pgr_KSP - Use Yen algorithm with pgr_dijkstra to get the K shortest paths.

## Flow - Family of functions

pgr_maxFlow - Only the Max flow calculation using Push and Relabel algorithm. pgr_boykovKolmogorov - Boykov and Kolmogorov with details of flow on edges. pgr_edmondsKarp - Edmonds and Karp algorithm with details of flow on edges. pgr_pushRelabel - Push and relabel algorithm with details of flow on edges.
Applications

- pgr_edgeDisjointPaths - Calculates edge disjoint paths between two groups of vertices.
- pgr_maxCardinalityMatch - Calculates a maximum cardinality matching in a graph.


## Kruskal - Family of functions

```
pgr_kruskal
pgr_kruskalBFS
pgr_kruskaIDD
pgr_kruskaIDFS
```


## Prim - Family of functions

- pgr_prim
- pgr_primBFS
- pgr_primDD
- pgr_primDFS


## Reference

- pgr_version
- pgr_full_version


## Topology - Family of Functions

The following functions modify the database directly therefore the user must have special permissions given by the administrators to use them.

- pgr_createTopology - create a topology based on the geometry.
- pgr_createVerticesTable - reconstruct the vertices table based on the source and target information.
- pgr_analyzeGraph - to analyze the edges and vertices of the edge table.
- pgr_analyzeOneWay - to analyze directionality of the edges.
- pgr_nodeNetwork -to create nodes to a not noded edge table.


## Traveling Sales Person - Family of functions

- pgr_TSP - When input is given as matrix cell information.
- pgr_TSPeuclidean - When input are coordinates.
pgr_trsp - Turn Restriction Shortest Path (TRSP) - Turn Restriction Shortest Path (TRSP)

Functions by categories

## Cost - Category

. pgr_aStarCost

- pgr_bdAstarCost
- pgr_dijkstraCost
- pgr_bdDijkstraCost
- pgr_dijkstraNearCost - Proposed


## Cost Matrix - Category

pgr_aStarCostMatrix

- pgr_bdAstarCostMatrix
- pgr_bdDijkstraCostMatrix
- pgr_dijkstraCostMatrix
- pgr_bdDijkstraCostMatrix

Driving Distance - Category

- pgr_drivingDistance - Driving Distance based on Dijkstra's algorithm
- pgr_primDD - Driving Distance based on Prim's algorithm
- pgr_kruskaIDD - Driving Distance based on Kruskal's algorithm
- Post pocessing
- pgr_alphaShape - Alpha shape computation


## K shortest paths - Category

- pgr_KSP - Yen's algorithm based on pgr_dijkstra

Spanning Tree - Category

- Kruskal - Family of functions
- Prim - Family of functions

```
BFS - Category
```

- pgr_kruskalBFS
- pgr_primBFS

DFS - Category

- pgr_kruskaIDFS
- pgr_primDFS
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: $2.6 \mathbf{2 . 5} 2.42 .32 .2$

All Pairs - Family of Functions
The following functions work on all vertices pair combinations

- pgr_floydWarshall - Floyd-Warshall's algorithm.
- pgr_johnson-Johnson's algorithm
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: $2.6 \mathbf{2}$.5 2.4 2.3 2.2 2.12 .0

```
pgr_floydWarshall
```

pgr_floydWarshall - Returns the sum of the costs of the shortest path for each pair of nodes in the graph using Floyd-Warshall algorithm.

## Availability

- Version 2.2.0
- Signature change
- Old signature no longer supported
- Version 2.0.0
- Official function


## Description

The Floyd-Warshall algorithm, also known as Floyd's algorithm, is a good choice to calculate the sum of the costs of the shortest path for each pair of nodes in the graph, for dense graphs. We use Boost's implementation which runs in (\Theta(V^3)<br>) time,

The main characteristics are:

- It does not return a path.
- Returns the sum of the costs of the shortest path for each pair of nodes in the graph.
- Process is done only on edges with positive costs.
- Boost returns a $\backslash(\mathrm{V} \backslash$ times $\mathrm{V} \backslash$ ) matrix, where the infinity values. Represent the distance between vertices for which there is no path.
- We return only the non infinity values in form of a set of(start_vid, end_vid, agg_cost).
- Let be the case the values returned are stored in a table, so the unique index would be the pair:(start_vid, end_vid).
- For the undirected graph, the results are symmetric.
- The agg_cost of $(u, v)$ is the same as for $(v, u)$.
- When start_vid = end_vid, the agg_cost $=0$.
- Recommended, use a bounding box of no more than 3500 edges.


## Signatures

## Summary

```
pgr_floydWarshall(Edges SQL, [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET
```


## Example:

For a directed subgraph with edges $\backslash(\backslash\{1,2,3,4 \backslash\} \backslash)$.

```
SELECT * FROM pgr_floydWarshall(
    SELECT id, source, target, cost, reverse_cost
    FROM edges where id < 5'
ORDER BY start_vid, end vid;
start_vid | end_vid | agg_cost
\begin{tabular}{|c|c|c|}
\hline 51 & 61 & 1 \\
\hline 51 & 71 & 2 \\
\hline 61 & 51 & 1 \\
\hline 61 & 71 & 1 \\
\hline 71 & 5 & 2 \\
\hline 71 & 61 & 1 \\
\hline \(10 \mid\) & 51 & 2 \\
\hline \(10 \mid\) & 61 & 1 \\
\hline \(10 \mid\) & 71 & 2 \\
\hline 15 | & 51 & 3 \\
\hline 15 | & 61 & 2 \\
\hline 15 | & 71 & 3 \\
\hline 15| & \(10 \mid\) & 1 \\
\hline (13 rows) & & \\
\hline
\end{tabular}
```


## Parameter

| Parameter | Type | Default | Description |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |  |  |
|  |  |  |  |  |



Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| source | ANY-INTEGER | Description |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Result Columns

Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

## See Also

- pgrjohnson
- Boost floyd-Warshall
- Queries uses the Sample Data network.


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.62 .52 .42 .32 .22 .12 .0
pgrjohnson
pgr_johnson - Returns the sum of the costs of the shortest path for each pair of nodes in the graph using Floyd-Warshall algorithm.

Boost Graph Inside

## Availability

- Version 2.2.0
- Signature change
- Old signature no longer supported
- Version 2.0.0
- Official function

The Johnson algorithm, is a good choice to calculate the sum of the costs of the shortest path for each pair of nodes in the graph, for sparse graphs. It usees the Boost's implementation which runs in $\backslash(\mathrm{O}(\mathrm{V} \mathrm{E} \backslash \mathrm{log} \mathrm{V}) \backslash)$ time,

The main characteristics are:

- It does not return a path.
- Returns the sum of the costs of the shortest path for each pair of nodes in the graph.
- Process is done only on edges with positive costs.
- Boost returns a $\backslash(\mathrm{V} \backslash$ times $\mathrm{V} \backslash)$ matrix, where the infinity values. Represent the distance between vertices for which there is no path.
- We return only the non infinity values in form of a set of(start_vid, end_vid, agg_cost).
- Let be the case the values returned are stored in a table, so the unique index would be the pair:(start_vid, end_vid).
- For the undirected graph, the results are symmetric.
- The agg_cost of $(u, v)$ is the same as for $(v, u)$.
- When start_vid = end_vid, the agg_cost $=0$.
- Recommended, use a bounding box of no more than $\mathbf{3 5 0 0}$ edges.

Signatures

## Summary

```
pgr johnson(Edges SQL, [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET
```


## Example:

For a directed subgraph with edges $\backslash(\backslash\{1,2,3,4 \backslash\} \backslash)$.

```
SELECT * FROM pgr johnson(
    'SELECT source, target, cost FROM edges
    WHERE id < 5'
) ORDER BY start_vid, end_vid;
start vid | end vid | agg cost
    5| 6| 1
    5| 7| 2
(3 rows)
```

Parameters

| Parameter | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |  |

## Optional parameters

| Column | Type | Default | Description |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| directed | BOOLEAN | true |  | When true the graph is considered Directed |  |
|  |  |  |  | Whenfalse the graph is considered Undirected. | as |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| source | ANY-INTEGER | Description |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Column

Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

## See Also

- pgr_floydWarshall
- Boost Johnson
- Queries uses the Sample Data network.


## Indices and tables

- Index
- Search Page


## Introduction

The main characteristics are:

- It does not return a path.
- Returns the sum of the costs of the shortest path for each pair of nodes in the graph.
- Process is done only on edges with positive costs.
- Boost returns a $\backslash(\mathrm{V} \backslash$ times $\mathrm{V} \backslash)$ matrix, where the infinity values. Represent the distance between vertices for which there is no path.
- We return only the non infinity values in form of a set of(start_vid, end_vid, agg_cost).
- Let be the case the values returned are stored in a table, so the unique index would be the pair:(start_vid, end_vid).
- For the undirected graph, the results are symmetric.
- The agg_cost of $(u, v)$ is the same as for $(v, u)$.
- When start_vid = end_vid, the agg_cost $=0$.
- Recommended, use a bounding box of no more than $\mathbf{3 5 0 0}$ edges.


## Parameters

| Parameter | Type | Default | Description |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |  |  |  |

Optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true | When true the graph is considered Directed |  |
|  |  |  When false the graph is considered as <br>  Undirected. |  |

Inner Queries

Edges SQL

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |  |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |  |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |
|  |  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |  |
|  |  |  |  |
|  |  |  |  |
|  |  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Column

## Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

## Performance

The following tests:

- non server computer
- with AMD 64 CPU
- 4G memory
- trusty
- posgreSQL version 9.3


## Data

The following data was used

```
BBOX="-122.8,45.4,-122.5,45.6"
wget --progress=dot:mega -O "sampledata.osm" "https://www.overpass-api.de/api/xapi?*[bbox=][@meta]"
```

Data processing was done with osm2pgrouting-alpha

```
createdb portland
psql -c "create extension postgis" portland
psql -c "create extension pgrouting" portland
osm2pgrouting -f sampledata.osm -d portland -s 0
```


## Results

## Test:

One
This test is not with a bounding box The density of the passed graph is extremely low. For each <SIZE> 30 tests were executed to get the average The tested query is:

```
SELECT count(*) FROM pgr_floydWarshall(
    'SELECT gid as id, source, target, cost, reverse_cost
    FROM ways where id <= <SIZE>');
SELECT count(*) FROM pgr_johnson(
    'SELECT gid as id, source, target, cost, reverse_cost
    FROM ways where id <= <SIZE>');
```

The results of this tests are presented as:

## SIZE:

is the number of edges given as input.

## EDGES:

is the total number of records in the query.

## DENSITY:

is the density of the data <br>(\dfrac $\{\mathrm{E}\}\{\mathrm{V}$ \times (V-1) $\} \backslash$ ).

## OUT ROWS:

is the number of records returned by the queries.

## Floyd-Warshall:

is the average execution time in seconds of pgr_floydWarshall.

## Johnson:

is the average execution time in seconds of pgr johnson.

| SIZE | EDGES | DENSITY | OUT ROWS | Floyd-Warshall | Johnson |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 500 | 500 | $0.18 \mathrm{E}-7$ | 1346 | 0.14 | 0.13 |
| 1000 | 1000 | $0.36 \mathrm{E}-7$ | 2655 | 0.23 | 0.18 |


| SIZE | EDGES | DENSITY | OUT ROWS | Floyd-Warshall | Johnson |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1500 | 1500 | $0.55 \mathrm{E}-7$ | 4110 | 0.37 | 0.34 |
| 2000 | 2000 | $0.73 \mathrm{E}-7$ | 5676 | 0.56 | 0.37 |
| 2500 | 2500 | $0.89 \mathrm{E}-7$ | 7177 | 0.84 | 0.51 |
| 3000 | 3000 | $1.07 \mathrm{E}-7$ | 8778 | 1.28 | 0.68 |
| 3500 | 3500 | $1.24 \mathrm{E}-7$ | 10526 | 2.08 | 0.95 |
| 4000 | 4000 | $1.41 \mathrm{E}-7$ | 12484 | 3.16 | 1.24 |
| 4500 | 4500 | $1.58 \mathrm{E}-7$ | 14354 | 4.49 | 1.47 |
| 5000 | 5000 | $1.76 \mathrm{E}-7$ | 16503 | 6.05 | 1.78 |
| 5500 | 5500 | $1.93 \mathrm{E}-7$ | 18623 | 7.53 | 2.03 |
| 6000 | 6000 | $2.11 \mathrm{E}-7$ | 20710 | 8.47 | 2.37 |
| 6500 | 6500 | $2.28 \mathrm{E}-7$ | 22752 | 9.99 | 2.68 |
| 7000 | 7000 | $2.46 \mathrm{E}-7$ | 24687 | 11.82 | 3.12 |
| 7500 | 7500 | $2.64 \mathrm{E}-7$ | 26861 | 13.94 | 3.60 |
| 8000 | 8000 | $2.83 \mathrm{E}-7$ | 29050 | 15.61 | 4.09 |
| 8500 | 8500 | $3.01 \mathrm{E}-7$ | 31693 | 17.43 | 4.63 |
| 9000 | 9000 | $3.17 \mathrm{E}-7$ | 33879 | 19.19 | 5.34 |
| 9500 | 9500 | $3.35 \mathrm{E}-7$ | 36287 | 20.77 | 6.24 |
| 10000 | 10000 | $3.52 \mathrm{E}-7$ | 38491 | 23.26 | 6.51 |

## Test:

Two
This test is with a bounding box The density of the passed graph higher than of the Test One. For each <SIZE> 30 tests were executed to get the average The tested edge query is:

```
WITH
buffer AS (
    SELECT ST_Buffer(ST_Centroid(ST_Extent(the_geom)), SIZE) AS geom
    FROM ways),
bbox AS (
    SELECT ST_Envelope(ST_Extent(geom)) as box FROM buffer)
SELECT gid as id, source, target, cost, reverse_cost
FROM ways where the_geom && (SELECT box from bbox);
```

The tested queries

```
SELECT count(*) FROM pgr_floydWarshall(<edge query>)
SELECT count(*) FROM pgr_johnson(<edge query>)
```

The results of this tests are presented as:

## SIZE:

is the size of the bounding box.
EDGES:
is the total number of records in the query.

## DENSITY:

is the density of the data <br>(\dfrac\{E\}\{V \times (V-1) $\} \backslash$ ).
OUT ROWS:
is the number of records returned by the queries.

## Floyd-Warshall:

is the average execution time in seconds of pgr_floydWarshall.

## Johnson:

is the average execution time in seconds of pgrjohnson.

| SIZE | EDGES | DENSITY | OUT ROWS | Floyd-Warshall | Johnson |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 0.001 | 44 | 0.0608 | 1197 | 0.10 | 0.10 |
| 0.002 | 99 | 0.0251 | 4330 | 0.10 | 0.10 |
| 0.003 | 223 | 0.0122 | 18849 | 0.12 | 0.12 |
| 0.004 | 358 | 0.0085 | 71834 | 0.16 | 0.16 |
| 0.005 | 470 | 0.0070 | 116290 | 0.22 | 0.19 |
| 0.006 | 639 | 0.0055 | 207030 | 0.37 | 0.27 |
| 0.007 | 843 | 0.0043 | 346930 | 0.64 | 0.38 |
| 0.008 | 996 | 0.0037 | 469936 | 0.90 | 0.49 |
| 0.009 | 1146 | 0.0032 | 613135 | 1.26 | 0.62 |
| 0.010 | 1360 | 0.0027 | 849304 | 1.87 | 0.82 |
| 0.011 | 1573 | 0.0024 | 1147101 | 2.65 | 1.04 |


| SIZE | EDGES | DENSITY | OUT ROWS | Floyd-Warshall | Johnson |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 0.012 | 1789 | 0.0021 | 1483629 | 3.72 | 1.35 |
| 0.013 | 1975 | 0.0019 | 1846897 | 4.86 | 1.68 |
| 0.014 | 2281 | 0.0017 | 2438298 | 7.08 | 2.28 |
| 0.015 | 2588 | 0.0015 | 3156007 | 10.28 | 2.80 |
| 0.016 | 2958 | 0.0013 | 4090618 | 14.67 | 3.76 |
| 0.017 | 3247 | 0.0012 | 4868919 | 18.12 | 4.48 |

See Also

- pgr_johnson
- pgr_floydWarshall
- Boost floyd-Warshall


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: $2.6 \mathbf{2 . 5} 2.4$

A* - Family of functions
The A* (pronounced "A Star") algorithm is based on Dijkstra's algorithm with a heuristic that allow it to solve most shortest path problems by evaluation only a sub-set of the overall graph.

- pgr_aStar - A* algorithm for the shortest path.
- pgr_aStarCost - Get the aggregate cost of the shortest paths.
- pgr_aStarCostMatrix - Get the cost matrix of the shortest paths.
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.6 2.5 2.42 .32 .22 .12 .0
pgr_astar
pgr_aStar - Shortest path using the A* algorithm.

Boost Graph Inside

## Availability

- Version 3.2.0
- New proposed signature:
- pgr_aStar (Combinations)
- Version 3.0.0
- Official function
- Version 2.4.0
- New Proposed signatures:
- pgr_aStar (One to Many)
- pgr_aStar (Many to One)
- pgr_aStar (Many to Many)
- Version 2.3.0
- Signature change on pgr_astar (One to One)
- Old signature no longer supported
- Version 2.0.0
- Official pgr_aStar (One to One)


## Description

- Process works for directed and undirected graphs.
- Ordering is:
- first by start_vid (if exists)
- then by end_vid
- Values are returned when there is a path.
- Let $\backslash(v \backslash)$ and $\backslash(u \backslash)$ be nodes on the graph:
- If there is no path from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ :
- no corresponding row is returned
- agg_cost from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ is $\backslash(\backslash i n f t y \backslash)$
- There is no path when $\backslash(v=u \backslash)$ therefore
- no corresponding row is returned
- agg_cost from $v$ to $u$ is $\backslash(0 \backslash)$
- When $\backslash((x, y) \backslash)$ coordinates for the same vertex identifier differ:
- A random selection of the vertex's $\backslash((x, y) \backslash)$ coordinates is used.
- Running time: $\backslash(\mathrm{O}((\mathrm{E}+\mathrm{V}) * \backslash \log \mathrm{~V}) \backslash)$
- The results are equivalent to the union of the results of thepgr_aStar( One to One ) on the:
- pgr_aStar(One to Many)
- pgr_aStar(Many to One )
- pgr_aStar( Many to Many )
- start_vid and end_vid in the result is used to distinguish to which path it belongs.


## Signatures

## Summary

```
pgr_aStar(Edges SQL, start vid, end vid, [options])
pgr_aStar(Edges SQL, start vid, end vids, [options])
pgr_aStar(Edges SQL, start vids, end vid, [options])
pgr_aStar(Edges SQL, start vids, end vids, [options])
pgr_aStar(Edges SQL, Combinations SQL, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (seq, path_seq, [start_vid], [end_vid], node, edge, cost, agg_cost)
OR EMPTY SET
```

Optional parameters are named parameters and have a default value.

```
One to One
```

```
pgr_aStar(Edges SQL, start vid, end vid, [options])
```

options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET

## Example:

From vertex $\backslash(6 \backslash)$ to vertex $\backslash(12 \backslash)$ on a directed graph with heuristic $\backslash(2 \backslash)$

```
SELECT * FROM pgr_aStar(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    6, 12,
    directed => true, heuristic => 2);
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{lllccl}
\(1 \mid\) & \(1 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(7 \mid\) & \(10 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(3 \mid\) & \(8 \mid\) & \(12 \mid\) & \(1 \mid\) & 2 \\
\(4 \mid\) & \(4 \mid\) & \(12 \mid\) & \(-1 \mid\) & \(0 \mid\) & 3 \\
(4 rows) & & & & &
\end{tabular}
```


## One to Many

```
pgr_aStar(Edges SQL, start vid, end vids, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (seq, path_seq, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(6 \backslash)$ to vertices $\backslash(\backslash\{10,12 \backslash\} \backslash)$ on a directed graph with heuristic $\backslash(3 \backslash)$ and factor $\backslash(3.5 \backslash)$

```
SELECT * FROM pgr_aStar(
    SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    6, ARRAY[10, 12]
    heuristic => 3, factor := 3.5);
seq | path_seq | end_vid | node | edge | cost | agg_cost
\begin{tabular}{ccccc}
\(1 \mid\) & \(10 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(10 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(10 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(10 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) \\
\(5 \mid\) & 3 \\
\(5 \mid\) & \(10 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(1 \mid\) & \(12 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(12 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(12 \mid\) & \(11 \mid\) & \(11 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(12 \mid\) & \(12 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
4 & 3
\end{tabular}
(10 rows)
```

Many to One

## pgr_aStar(Edges SQL, start vids, end vid, [options])

options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (seq, path_seq, start_vid, node, edge, cost, agg_cost)
OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,8 \backslash\} \backslash)$ to vertex $\backslash(10 \backslash)$ on an undirected graph with heuristic $\backslash(4 \backslash)$

```
SELECT * FROM pgr_aStar(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges'
    ARRAY[6, 8], 10,
    false, heuristic => 4);
seq | path_seq | start_vid | node | edge | cost | agg_cost
1| 1| 6| 6| 2| 1| 0
2| 2| 6| 10| -1| 0| 1
3| 1| 8| 8| 12| 1| 0
4| 2| 8| 12| 11| 1| 1
5| 3| 8| 11| 5| 1| 2
6| 4| 8| 10| -1| 0| 3
(6 rows)
```


## Many to Many

pgr_aStar(Edges SQL, start vids, end vids, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost) OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,8 \backslash\} \backslash)$ to vertices $\backslash(\backslash\{10,12 \backslash\} \backslash)$ on a directed graph with factor $\backslash(0.5 \backslash)$

## SELECT * FROM pgr aStar

'SELECT id, source, target, cost, reverse_cost, $\mathrm{x} 1, \mathrm{y} 1, \mathrm{x} 2, \mathrm{y} 2$
FROM edges',
ARRAY[6, 8], ARRAY[10, 12
factor $=>0.5$ );
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $6 \mid$ | $10 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $6 \mid$ | $10 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 1 |
| $3 \mid$ | $3 \mid$ | $6 \mid$ | $10 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 2 |
| $4 \mid$ | $4 \mid$ | $6 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 3 |
| $5 \mid$ | $5 \mid$ | $6 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ | 4 |
| $6 \mid$ | $6 \mid$ | $6 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 5 |
| $7 \mid$ | $1 \mid$ | $6 \mid$ | $12 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| $8 \mid$ | $2 \mid$ | $6 \mid$ | $12 \mid$ | $7 \mid$ | $10 \mid$ | $1 \mid$ | 1 |
| $9 \mid$ | $3 \mid$ | $6 \mid$ | $12 \mid$ | $8 \mid$ | $12 \mid$ | $1 \mid$ | 2 |
| $10 \mid$ | $4 \mid$ | $6 \mid$ | $12 \mid$ | $12 \mid$ | $-1 \mid$ | $0 \mid$ | 3 |
| $11 \mid$ | $1 \mid$ | $8 \mid$ | $10 \mid$ | $8 \mid$ | $10 \mid$ | $1 \mid$ | 0 |
| $12 \mid$ | $2 \mid$ | $8 \mid$ | $10 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 1 |
| $13 \mid$ | $3 \mid$ | $8 \mid$ | $10 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 2 |
| $14 \mid$ | $4 \mid$ | $8 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 3 |
| $15 \mid$ | $5 \mid$ | $8 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ | 4 |
| $16 \mid$ | $6 \mid$ | $8 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 5 |
| $17 \mid$ | $1 \mid$ | $8 \mid$ | $12 \mid$ | $8 \mid$ | $12 \mid$ | $1 \mid$ | 0 |
| $18 \mid$ | $2 \mid$ | $8 \mid$ | $12 \mid$ | $12 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |

## Combinations

## pgr_aStar(Edges SQL, Combinations SQL, [options]) <br> options: [directed, heuristic, factor, epsilon]

RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost) OR EMPTY SET

## Example:

Using a combinations table on a directed graph with factor $\backslash(0.5 \backslash)$.
The combinations table:

```
SELECT * FROM combinations
source | target
    5| 6
    5| 10
    6 5
    6| 15
    6 | 14
(5 rows)
```

The query:


| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | BIGINT | Identifier of the starting vertex of the path. |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting vertices. |
| end vid | BIGINT | Identifier of the ending vertex of the path. |
| end vids | ARRAY[BIGINT] | Array of identifiers of ending vertices. |

Optional parameters

aStar optional Parameters

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| heuristic | INTEGER | 5 | ```Heuristic number. Current valid values 0~5. 0:\(h(v) = O\) (Use this value to compare with pgr_dijkstra) 1: \(h(v) = abs(max(\Delta x,\Delta y))\) 2: \(h(v) = abs(min(\Delta x, \Delta y))\) 3: \(h(v) = \Delta x * \Delta x + \Delta y * \Delta y\) 4: \(h(v) = sqrt(\Delta x * \Delta x + \Delta y * \Delta y)\) 5: \(h(v) = abs(\Delta x) + abs(\Delta y)\)``` |
| factor | FLOAT | 1 | For units manipulation. $\backslash($ factor $>0 \backslash$ ). |
| epsilon | FLOAT | 1 | For less restricted results. (epsilon >= 1). |

See heuristics available and factor handling.

Inner Queries

Edges SQL

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER |  | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
|  |  |  | - When negative: edge (source, target) does not exist, therefore it's not part of the graph. |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source), |
|  |  |  | When negative: edge (target, source) does not exist, therefore it's not part of the graph. |
| $\times 1$ | ANY-NUMERICAL |  | $X$ coordinate of source vertex. |
| y1 | ANY-NUMERICAL |  | Y coordinate of source vertex. |
| $\times 2$ | ANY-NUMERICAL |  | X coordinate of target vertex. |
| y 2 | ANY-NUMERICAL |  | Y coordinate of target vertex. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Combinations SQL

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- | Identifier of the departure vertex. |
|  | INTEGER |  |


| Parameter | Type | Description |
| :--- | :--- | :--- |
| target | ANY- <br> INTEGER | Identifier of the arrival vertex. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

Result Columns

Returns set of (seq, path_seq [, start_vid] [, end_vid], node, edge, cost, agg_cost)

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. |
| path_seq | INTEGER | Relative position in the path. Has value $\mathbf{1}$ for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex. Returned when multiple starting vetrices are in the query. Many to One Many to Many |
| end_vid | BIGINT | Identifier of the ending vertex. Returned when multiple ending vertices are in the query. <br> One to Many <br> Many to Many |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for the last node of the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | FLOAT | Aggregate cost from start_vid to node. |

## Additional Examples

## Example 1:

Demonstration of repeated values are ignored, and result is sorted.

```
SELECT * FROM pgr_aStar(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    ARRAY[7, 10, 15, 10, 10, 15], ARRAY[10, 7, 10, 15]);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
\begin{tabular}{ccccccc|c}
\(1 \mid\) & \(1 \mid\) & \(7 \mid\) & \(10 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(7 \mid\) & \(10 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(3 \mid\) & \(7 \mid\) & \(10 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) & 2 \\
\(4 \mid\) & \(4 \mid\) & \(7 \mid\) & \(10 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) & 3 \\
\(5 \mid\) & \(5 \mid\) & \(7 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 4 \\
\(6 \mid\) & \(1 \mid\) & \(7 \mid\) & \(15 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 0 \\
\(7 \mid\) & \(2 \mid\) & \(7 \mid\) & \(15 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) & 1 \\
\(8 \mid\) & \(3 \mid\) & \(7 \mid\) & \(15 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) & 2 \\
\(9 \mid\) & \(4 \mid\) & \(7 \mid\) & \(15 \mid\) & \(15 \mid\) & \(-1 \mid\) & \(0 \mid\) & 3 \\
\(10 \mid\) & \(1 \mid\) & \(10 \mid\) & \(7 \mid\) & \(10 \mid\) & \(5 \mid\) & \(1 \mid\) & 0 \\
\(11 \mid\) & \(2 \mid\) & \(10 \mid\) & \(7 \mid\) & \(11 \mid\) & \(8 \mid\) & \(1 \mid\) & 1 \\
\(12 \mid\) & \(3 \mid\) & \(10 \mid\) & \(7 \mid\) & \(7 \mid\) & \(-1 \mid\) & \(0 \mid\) & 2 \\
\(13 \mid\) & \(1 \mid\) & \(10 \mid\) & \(15 \mid\) & \(10 \mid\) & \(5 \mid\) & \(1 \mid\) & 0 \\
\(14 \mid\) & \(2 \mid\) & \(10 \mid\) & \(15 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) & 1 \\
\(15 \mid\) & \(3 \mid\) & \(10 \mid\) & \(15 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) & 2 \\
\(16 \mid\) & \(4 \mid\) & \(10 \mid\) & \(15 \mid\) & \(15 \mid\) & \(-1 \mid\) & \(0 \mid\) & 3 \\
\(17 \mid\) & \(1 \mid\) & \(15 \mid\) & \(7 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) & 0 \\
\(18 \mid\) & \(2 \mid\) & \(15 \mid\) & \(7 \mid\) & \(10 \mid\) & \(2 \mid\) & \(1 \mid\) & 1 \\
\(19 \mid\) & \(3 \mid\) & \(15 \mid\) & \(7 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) & 2 \\
\(20 \mid\) & \(4 \mid\) & \(15 \mid\) & \(7 \mid\) & \(7 \mid\) & \(-1 \mid\) & \(0 \mid\) & 3 \\
\(21 \mid\) & \(1 \mid\) & \(15 \mid\) & \(10 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) & 0 \\
\(22 \mid\) & \(2 \mid\) & \(15 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 1
\end{tabular}
```


## Example 2:

Making start vids the same as end vids.

SELECT * FROM pgr_aStar(
'SELECT id, source, target, cost, reverse_cost, x 1 , y 1 , x 2 , y2
FROM edges',
ARRAY[7, 10, 15], ARRAY[7, 10, 15]);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1 \| | 1 \| | 71 | $10 \mid$ | 71 | 81 | $1 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | 21 | 71 | 10\| | 11\| | 9\| | 1\| | 1 |
| 3 | 3 | 71 | 10\| | 16\| | 16\| | 1 \| | 2 |
| $4 \mid$ | 4 | 71 | 10\| | 15\| | 3\| | 1\| | 3 |
| 5 | 51 | 71 | 10\| | 10\| | -1\| | 0 \| | 4 |
| 6 | 1\| | 71 | 15\| | 71 | 8। | 1\| | 0 |
| 71 | 21 | 71 | 15\| | 11\| | 9\| | 1\| | 1 |
| 8 | 31 | 71 | 15\| | 16\| | 16\| | $1 \mid$ | 2 |
| 9 | 4\| | 71 | 15\| | 15 | -1\| | $0 \mid$ | 3 |
| $10 \mid$ | 1 \| | $10 \mid$ | 71 | 10\| | 5\| | 1 \| | 0 |
| 11\| | $2 \mid$ | 10\| | 71 | 11\| | 8\| | 1 \| | 1 |
| $12 \mid$ | 31 | $10 \mid$ | 71 | 71 | -1\| | 01 | 2 |
| $13 \mid$ | 1 \| | $10 \mid$ | 15 \| | $10 \mid$ | 51 | 1 \| | 0 |
| 14 \| | $2 \mid$ | $10 \mid$ | 15 \| | 11\| | 91 | 1 |  |
| 15 \| | 31 | $10 \mid$ | 15 \| | $16 \mid$ | $16 \mid$ | 1 |  |
| $16 \mid$ | 4 \| | $10 \mid$ | 15 \| | 15 | -1\| | 01 |  |
| 17\| | 1 \| | 15 \| | 71 | 15\| | 3\| | 1 \| | 0 |
| 18 \| | $2 \mid$ | 15 \| | 71 | 10\| | $2 \mid$ | 1 \| | 1 |
| 19 \| | 31 | 15 \| | 71 | 61 | 4\| | 1\| | 2 |
| 201 | 4 \| | 15 \| | 71 | 71 | -1\| | 01 | 3 |
| 21\| | 1 \| | 15\| | $10 \mid$ | 15 | 31 | 1 | 0 |
| $22 \mid$ | $2 \mid$ | 15 \| | 10 \| | $10 \mid$ | -1\| | $0 \mid$ |  |
| (22 ro |  |  |  |  |  |  |  |

## Example 3:

Manually assigned vertex combinations.

```
SELECT * FROM pgr_aStar(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    'SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)');
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
\begin{tabular}{rccc|c|c|c}
\(1 \mid\) & \(1 \mid\) & \(6 \mid\) & \(7 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(2 \mid\) & \(6 \mid\) & \(7 \mid\) & \(7 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(3 \mid\) & \(1 \mid\) & \(6 \mid\) & \(10 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(2 \mid\) & \(6 \mid\) & \(10 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(5 \mid\) & \(3 \mid\) & \(6 \mid\) & \(10 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(4 \mid\) & \(6 \mid\) & \(10 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) \\
\(7 \mid\) & \(5 \mid\) & \(6 \mid\) & \(10 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) \\
\(8 \mid\) & \(6 \mid\) & \(6 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & 0 \\
\(9 \mid\) & \(1 \mid\) & \(12 \mid\) & \(10 \mid\) & \(12 \mid\) & \(13 \mid\) & \(1 \mid\) \\
\(10 \mid\) & \(2 \mid\) & \(12 \mid\) & \(10 \mid\) & \(17 \mid\) & \(15 \mid\) & \(1 \mid\) \\
\(11 \mid\) & \(3 \mid\) & \(12 \mid\) & \(10 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) \\
\(12 \mid\) & \(4 \mid\) & \(12 \mid\) & \(10 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) \\
\(12 \mid\) & 3 \\
\(13 \mid\) & \(5 \mid\) & \(12 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\((13\) rows) & & & & & 4
\end{tabular}
```

See Also

- $A^{*}$ - Family of functions
- Bidirectional $A^{*}$ - Family of functions
- Sample Data
- https://www.boost.org/libs/graph/doc/astar_search.html
- https://en.wikipedia.org/wiki/A*_search_algorithm


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.6 2.5 2.4
pgr_aStarCost
pgr_aStarCost - Total cost of the shortest path(s) using the A* algorithm.


## Availability

- Version 3.2.0
- New proposed signature:
- pgr_aStarCost (Combinations)
- Version 3.0.0
- Official function
- Version 2.4.0
- New proposed function


## Description

The pgr_aStarCost function sumarizes of the cost of the shortest path(s) using the A* algorithm.

## The main characteristics are:

- Process works for directed and undirected graphs.
- Ordering is:
- first by start_vid (if exists)
- then by end_vid
- Values are returned when there is a path.
- Let $\backslash(v \backslash)$ and $\backslash(u \backslash)$ be nodes on the graph:
- If there is no path from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ :
- no corresponding row is returned
- agg_cost from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ is $\backslash(\backslash i n f t y \backslash)$
- There is no path when $\backslash(v=u \backslash)$ therefore - no corresponding row is returned - agg_cost from $v$ to $u$ is $\backslash(0 \backslash)$
- When $\backslash((x, y) \backslash)$ coordinates for the same vertex identifier differ:
- A random selection of the vertex's $\backslash((x, y) \backslash)$ coordinates is used.
- Running time: $\backslash(\mathrm{O}((\mathrm{E}+\mathrm{V}) * \backslash \log \mathrm{~V}) \backslash)$
- It does not return a path.
- Returns the sum of the costs of the shortest path of each pair combination of nodes requested.
- Let be the case the values returned are stored in a table, so the unique index would be the pair:(start_vid, end_vid)
- For undirected graphs, the results are symmetric.
- The agg_cost of $(u, v)$ is the same as for $(v, u)$.
- The returned values are ordered in ascending order:
- start_vid ascending
- end_vid ascending


## Signatures

## Summary

```
pgr_aStarCost(Edges SQL, start vid, end vid, [options])
pgr_aStarCost(Edges SQL, start vid, end vids, [options])
pgr_aStarCost(Edges SQL, start vids, end vid, [options])
pgr_aStarCost(Edges SQL, start vids, end vids, [options])
pgr_aStarCost(Edges SQL, Combinations SQL, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET
```


## One to One

```
pgr_aStarCost(Edges SQL, start vid, end vid, [options])
    options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(6 \backslash)$ to vertex $\backslash(12 \backslash)$ on a directed graph with heuristic $\backslash(2 \backslash)$

```
SELECT * FROM pgr_aStarCost(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    6,12,
directed => true, heuristic => 2);
start_vid | end_vid | agg_cost
    6| 12| 3
(1 row)
```


## One to Many

```
pgr_aStarCost(Edges SQL, start vid, end vids, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(6 \backslash)$ to vertices $\backslash(\backslash\{10,12 \backslash\} \backslash)$ on a directed graph with heuristic $\backslash(3 \backslash)$ and factor $\backslash(3.5 \backslash)$

```
SELECT * FROM pgr_aStarCost(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    6, ARRAY[10, 12],
    heuristic => 3, factor => 3.5);
start_vid | end_vid | agg_cost
    6 | 10| 5
    6| 12| 3
(2 rows)
```


## Many to One

```
pgr_aStarCost(Edges SQL, start vids, end vid, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET
```


## Example:

From vertices $\backslash(\backslash\{6,8 \backslash\} \backslash)$ to vertex $\backslash(10 \backslash)$ on an undirected graph with heuristic $\backslash(4 \backslash)$

```
SELECT * FROM pgr_aStarCost(
    SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    ARRAY[6,8], 10
    false, heuristic => 4);
start_vid | end_vid | agg_cost
    6| 10| 1
    8| 10| 3
(2 rows)
```


## Many to Many

pgr_aStarCost(Edges SQL, start vids, end vids, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,8 \backslash\} \backslash)$ to vertices $\backslash(\backslash\{10,12 \backslash\} \backslash)$ on a directed graph with factor $\backslash(0.5 \backslash)$

```
SELECT * FROM pgr_aStarCost
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    ARRAY[6, 8], ARRAY[10, 12],
factor => 0.5);
start_vid | end_vid | agg_cost
    6| 10| 5
    6| 12| 3
    8| 10| 5
    8| 12| 1
(4 rows)
```


## Combinations

```
pgr_aStarCost(Edges SQL, Combinations SQL, [options])
```

options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

## Example:

Using a combinations table on a directed graph with factor $\backslash(0.5 \backslash)$.

The combinations table:

```
SELECT * FROM combinations;
source | target
    5| 6
    5| 10
    6| 5
    6| 15
    6| 14
(5 rows)
```

The query:

```
SELECT * FROM pgr_aStarCost(
    'SELECT id, source, target, cost, reverse cost, x1, y1, x2, y2
    FROM edges',
    'SELECT * FROM combinations',
    factor => 0.5);
start_vid | end_vid | agg_cost
    5| 6| 1
    5| 10| 6
    6| 5| 1
(4 rows)
```

Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | BIGINT | Identifier of the starting vertex of the path. |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting vertices. |
| end vid | BIGINT | Identifier of the ending vertex of the path. |
| end vids | ARRAY[BIGINT] | Array of identifiers of ending vertices. |

## Optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true | When true the graph is considered Directed |  |
|  |  | When false the graph is considered as <br>  <br>  |  |


| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| heuristic | INTEGER | 5 | Heuristic number. Current valid values 0~5. |
| factor | FLOAT | 1 | For units manipulation. $\backslash($ factor $>0 \backslash$ ). |
| epsilon | FLOAT | 1 | For less restricted results. |
| (epsilon >= 1 ${ }^{\text {) }}$. |  |  |  |

See heuristics available and factor handling.

Inner Queries

Edges SQL

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER |  | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
|  |  |  | - When negative: edge (source, target) does not exist, therefore it's not part of the graph. |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source), |
|  |  |  | - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |
| x1 | ANY-NUMERICAL |  | $X$ coordinate of source vertex. |
| y1 | ANY-NUMERICAL |  | Y coordinate of source vertex. |
| $\times 2$ | ANY-NUMERICAL |  | $X$ coordinate of target vertex. |
| y2 | ANY-NUMERICAL |  | Y coordinate of target vertex. |

Where:
ANY-INTEGER:
SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Combinations SQL

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- | Identifier of the departure vertex. |
|  | INTEGER |  |
| target | ANY- | Identifier of the arrival vertex. |
|  | INTEGER |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Result Columns

Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

## Example 1:

```
SELECT * FROM pgr_aStarCost(
    FROM edges',
start_vid | end_vid | agg_cost
\begin{tabular}{ccc}
\(7 \mid\) & \(10 \mid\) & 4 \\
\(7 \mid\) & \(15 \mid\) & 3 \\
\(10 \mid\) & \(7 \mid\) & 2 \\
\(10 \mid\) & \(15 \mid\) & 3 \\
\(15 \mid\) & \(7 \mid\) & 3 \\
\(15 \mid\) & \(10 \mid\) & 1
\end{tabular}
```

    'SELECT id, source, target, cost, reverse_cost, \(x 1, y 1, x 2, y 2\)
    ARRAY[7, 10, 15, 10, 10, 15], ARRAY[10, 7, 10, 15]);

## Example 2:

Making start vids the same as end vids.

```
SELECT * FROM pgr_aStarCost(
'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
ARRAY[7, 10, 15], ARRAY[7, 10, 15]);
start_vid | end_vid | agg_cost
\begin{tabular}{rr}
\(7|------------+-------10|\) & 4 \\
\(7 \mid\) & \(15 \mid\) \\
\(10 \mid\) & \(7 \mid\) \\
10 \\
\(10 \mid\) & \(15 \mid\) \\
\(15 \mid\) & \(7 \mid\) \\
\(15 \mid\) & 3 \\
\(10 \mid\) & 1 \\
6 rows) & \\
\hline
\end{tabular}
```


## Example 3:

Manually assigned vertex combinations.

```
SELECT * FROM pgr_aStarCost(
    SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)');
start_vid | end_vid | agg_cost
    6| 7| 1
    6| 10| 5
    12| 10| 4
(3 rows)
```

See Also

- $A^{*}$ - Family of functions
- Cost-Category
- Sample Data


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: $2.6 \mathbf{2}$.5 2.4
pgr_aStarCostMatrix
pgr_aStarCostMatrix - Calculates the a cost matrix using pgr_aStar.

Boost Graph Inside

## Availability

- Official function
- Version 2.4.0
- New proposed function


## Description

## The main characteristics are:

- Using internaly the pgr_aStar algorithm
- Returns a cost matrix.
- No ordering is performed
- let $v$ and $u$ are nodes on the graph:
- when there is no path from $v$ to $u$ :
- no corresponding row is returned
- cost from $v$ to $u$ is $\backslash(\backslash i n f()$
- when $\backslash(v=u \backslash)$ then
- no corresponding row is returned
- cost from $v$ to $u$ is $\backslash(0 \backslash)$
- When the graph is undirected the cost matrix is symmetric


## Signatures

## Summary

```
pgr_aStarCostMatrix(Edges SQL, start vids, [options])
```

options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

## Example:

Symmetric cost matrix for vertices $\backslash(\backslash\{5,6,10,15 \backslash\} \backslash)$ on an undirected graph using heuristic $\backslash(2 \backslash)$

```
SELECT * FROM pgr_aStarCostMatrix(
'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2 FROM edges',
(SELECT array_agg(id) FROM vertices WHERE id IN (5, 6, 10, 15)),
directed => false, heuristic => 2);
start_vid | end_vid | agg_cost
    5| 6| 1
    5| 10| 2
    5| 15| 3
    6| 5| 1
    6| 10| 1
    6| 15| 2
    10| 5| 2
    10| 6| 1
    10| 15| 1
    15| 5| 3
    15| 6| 2
    15| 10| 1
(12 rows)
```


## Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting <br> vertices. |

Optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true |  When true the graph is considered Directed <br>   <br>   <br>   <br>   <br>   <br>  Undirected. n false the graph is considered as |  |


| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| heuristic | INTEGER | 5 | ```Heuristic number. Current valid values 0~5. 0:\(h(v) = 0\) (Use this value to compare with pgr_dijkstra) 1: \(h(v) = abs(max(\Delta x,\Delta y))\) 2: \(h(v) = abs(min(\Delta x,\Delta y))\) 3: \(h(v) = \Delta x * \Delta x + \Delta y * \Delta y\) 4: \(h(v) = sqrt(\Delta x * \Delta x + \Delta y * \Delta y)\) 5: \(h(v) = abs(\Delta x) + abs(\Delta y)\)``` |
| factor | FLOAT | 1 | For units manipulation. (factor > $0 \backslash$ ).  \hline epsilon & FLOAT & 1 & For less restricted results. (epsilon >=1). |

See heuristics available and factor handling.

Inner Queries

Edges SQL

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER |  | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
|  |  |  | - When negative: edge (source, target) does not exist, therefore it's not part of the graph. |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source), |
|  |  |  | - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |
| x1 | ANY-NUMERICAL |  | $X$ coordinate of source vertex. |
| y1 | ANY-NUMERICAL |  | Y coordinate of source vertex. |
| $\times 2$ | ANY-NUMERICAL |  | X coordinate of target vertex. |
| y2 | ANY-NUMERICAL |  | Y coordinate of target vertex. |

Where:
ANY-INTEGER:
SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Result Columns

Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

## Additional Examples

## Example:

Use with pgr_TSP

```
EELECT * FROM pgr_TSP
$$
SELECT * FROM pgr_aStarCostMatrix(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2 FROM edges',
    (SELECT array_agg(id) FROM vertices WHERE id IN (5, 6, 10, 15)),
    directed=> false, heuristic => 2)
$$,
randomize => false
);
NOTICE: pgr_TSP no longer solving with simulated annaeling
HINT: Ignoring annaeling parameters
seq | node | cost | agg cost
1| 5| 0| 0
2| 6| 1| 1
3| 10| 1| 2
4| 15| 1| 3
5| 5| 3| 6
(5 rows)
```


## See Also

A* - Family of functions

- Cost Matrix - Category
- Traveling Sales Person - Family of functions
- Sample Data


## Indices and tables

## Index

- Search Page


## Description

## The main Characteristics are:

- Process works for directed and undirected graphs.
- Ordering is:
- first by start_vid (if exists)
- then by end_vid
- Values are returned when there is a path.
- Let $\backslash(v \backslash)$ and $\backslash(u \backslash)$ be nodes on the graph:
- If there is no path from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ :
- no corresponding row is returned
- agg_cost from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ is $\backslash(\backslash i n f t y \backslash)$
- There is no path when $\backslash(v=u \backslash)$ therefore
- no corresponding row is returned
- agg_cost from $v$ to $u$ is $\backslash(0 \backslash)$
- When $\backslash((x, y) \backslash)$ coordinates for the same vertex identifier differ:
- A random selection of the vertex's $\backslash((x, y) \backslash)$ coordinates is used.
- Running time: $\backslash(\mathrm{O}((\mathrm{E}+\mathrm{V}) * \backslash \log \mathrm{~V}) \backslash)$
aStar optional Parameters

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| heuristic | INTEGER | 5 | ```Heuristic number. Current valid values 0~5. 0:\(h(v) = O\) (Use this value to compare with pgr_dijkstra) 1:\(h(v) = abs(max(\Delta x,\Delta y))\) 2: \(h(v) = abs(min(\Delta x,\Delta y))\) 3: \(h(v) = \Delta x * \Delta x + \Delta y * \Delta y\) 4: \(h(v) = sqrt(\Delta x * \Delta x + \Delta y * \Delta y)\) 5: \(h(v) = abs(\Delta x) + abs(\Delta y)\)``` |
| factor | FLOAT | 1 | For units manipulation. |
| (factor > $0 \backslash$ ). |  |  |  |
| epsilon | FLOAT | 1 | For less restricted results. |
| (epsilon >= 1 |  |  |  |

See heuristics available and factor handling.

Currently the heuristic functions available are:

- $0: \backslash(h(v)=0 \backslash)$ (Use this value to compare with pgr_dijkstra)
- $1: \backslash(h(v)=\operatorname{abs}(\max (\backslash$ Delta $x, \backslash$ Delta $y)) \backslash)$
- 2: $\backslash(\mathrm{h}(\mathrm{v})=\operatorname{abs}(\min (\backslash$ Delta $\mathrm{x}, \backslash$ Delta y$)) \backslash)$
- 3: $\backslash(h(v)=\$ DeIta $x * \backslash$ Delta $x+\backslash$ Delta $y * \backslash$ Delta $y \backslash)$
- 4: $\backslash\left(h(v)=\operatorname{sqrt}\left(\backslash\right.\right.$ Delta $x^{*} \backslash$ Delta $x+\backslash$ Delta $y * \backslash$ Delta $\left.\left.y\right) \backslash\right)$
- 5: $\backslash(h(v)=\operatorname{abs}(\backslash$ Delta $x)+a b s(\backslash$ Delta $y) \backslash)$
where <br>(\Delta $\left.x=x \_1-x \_0 \backslash\right)$ and $\backslash\left(\backslash D e l t a ~ y=y \_1-y \_0 \backslash\right)$

Factor

## Analysis 1

Working with cost/reverse_cost as length in degrees, $x / y$ in lat/lon: Factor $=1$ (no need to change units)

## Analysis 2

Working with cost/reverse_cost as length in meters, $\mathrm{x} / \mathrm{y}$ in lat/lon: Factor $=$ would depend on the location of the points:

| Latitude | Conversion | Factor |
| :--- | :--- | :--- |
| 45 | 1 longitude degree is 78846.81 m | 78846 |
| 0 | 1 longitude degree is 111319.46 | 111319 |
|  | m |  |
|  |  |  |

## Analysis 3

Working with cost/reverse_cost as time in seconds, $x / y$ in lat/lon: Factor: would depend on the location of the points and on the average speed say $25 \mathrm{~m} / \mathrm{s}$ is the speed.

| Latitude | Conversion | Factor |  |  |
| :--- | :--- | :--- | :--- | :--- |
| 45 | 1 longitude degree is $(78846.81 \mathrm{~m}) /(25 \mathrm{~m} / \mathrm{s})$ | 3153 s |  |  |
| 0 | $1 \quad$ longitude degree is $(111319.46$ | 4452 s |  |  |
|  | $\mathrm{~m}) /(25 \mathrm{~m} / \mathrm{s})$ |  |  |  |

## See Also

- Bidirectional A* - Family of functions
- https://www.boost.org/libs/graph/doc/astar_search.html
- https://en.wikipedia.org/wiki/A*_search_algorithm


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: $2.5 \mathbf{2 . 6}$


## Bidirectional A* - Family of functions

The bidirectional A* (pronounced "A Star") algorithm is based on the A* algorithm.

- pgr_bdAstar - Bidirectional A* algorithm for obtaining paths.
- pgr_bdAstarCost - Bidirectional A* algorithm to calculate the cost of the paths.
- pgr_bdAstarCostMatrix - Bidirectional A* algorithm to calculate a cost matrix of paths.
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2 2.12 .0
pgr_bdAstar
pgr_bdAstar - Shortest path using the bidirectional A* algorithm.

Boost Graph Inside

## Availability

- Version 3.2.0
- New proposed signature:
- pgr_bdAstar (Combinations)
- Version 3.0.0
- Official function
- Version 2.5.0
- New Proposed signatures:
- pgr_bdAstar (One to Many)
- pgr_bdAstar (Many to One)
- pgr_bdAstar (Many to Many)
- Signature change on pgr_bdAstar (One to One)
- Old signature no longer supported
- Version 2.0.0
- Official pgr_bdAstar (One to One)


## Description

## The main characteristics are:

- Process works for directed and undirected graphs.
- Ordering is:
- first by start_vid (if exists)
- then by end_vid
- Values are returned when there is a path.
- Let $\backslash(v \backslash)$ and $\backslash(u \backslash)$ be nodes on the graph:
- If there is no path from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ :
- no corresponding row is returned
- agg_cost from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ is $\backslash(\backslash i n f t y \backslash)$
- There is no path when $\backslash(v=u \backslash)$ therefore
- no corresponding row is returned
- agg_cost from $v$ to $u$ is $\backslash(0 \backslash)$
- When $\backslash((x, y) \backslash)$ coordinates for the same vertex identifier differ:
- A random selection of the vertex's $\backslash((x, y) \backslash)$ coordinates is used.
- Running time: $\backslash(\mathrm{O}((\mathrm{E}+\mathrm{V}) * \backslash \log \mathrm{~V}) \backslash)$
- The results are equivalent to the union of the results of thepgr_bdAStar( One to One ) on the:
- pgr_bdAstar(One to Many)
- pgr_bdAstar(Many to One )
- pgr_bdAstar(Many to Many )
- start_vid and end_vid in the result is used to distinguish to which path it belongs.


## Signatures

## Summary

```
pgr_bdAstar(Edges SQL, start vid, end vid, [options])
pgr_bdAstar(Edges SQL, start vid, end vids, [options])
pgr_bdAstar(Edges SQL, start vids, end vid, [options])
pgr_bdAstar(Edges SQL, start vids, end vids, [options])
pgr_bdAstar(Edges SQL, Combinations SQL, [options])
options: [directed, heuristic, factor, epsilon]
```

RETURNS SET OF (seq, path_seq, [start_vid], [end_vid], node, edge, cost, agg_cost)
OR EMPTY SET

Optional parameters are named parameters and have a default value.
pgr_bdAstar(Edges SQL, start vid, end vid, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET

## Example:

From vertex $\backslash(6 \backslash)$ to vertex $\backslash(12 \backslash)$ on a directed graph with heuristic $\backslash(2 \backslash)$

```
SELECT * FROM pgr_bdAstar(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    6,12,
    directed => true, heuristic => 2
);
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{ccccc|}
\(1 \mid\) & \(1 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(2 \mid\) & \(7 \mid\) & \(10 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(3 \mid\) & \(8 \mid\) & \(12 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(4 \mid\) & \(12 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\hline
\end{tabular}
(4 rows)
```


## One to Many

```
pgr_bdAstar(Edges SQL, start vid, end vids, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (seq, path_seq, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(6 \backslash)$ to vertices $\backslash(\backslash\{10,12 \backslash\} \backslash)$ on a directed graph with heuristic $\backslash(3 \backslash)$ and factor $\backslash(3.5 \backslash)$

```
SELECT * FROM pgr_bdAstar(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
6, ARRAY[10, 12],
heuristic => 3, factor := 3.5
);
seq | path_seq | end_vid | node | edge | cost | agg_cost
    1| 10| 6| 4| 1| 0
```



```
3| 3| 10| 11| 9| 1| 2
4| 4| 10| 16| 16| 1| 3
5| 5| 10| 15| 3| 1| 4
6|}6|\quad10| 10| -1| 0| 5 
8| 2| 12| 7| 8| 1| 1
9| 3| 12| 11| 11| 1| 2
```



```
(10 rows)
```


## Many to One

pgr_bdAstar(Edges SQL, start vids, end vid, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (seq, path_seq, start_vid, node, edge, cost, agg_cost)
OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,8 \backslash\} \backslash)$ to vertex $\backslash(10 \backslash)$ on an undirected graph with heuristic $\backslash(4 \backslash)$

```
SELECT * FROM pgr_bdAstar(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    ARRAY[6, 8], 10,
    false, heuristic => 4
);
seq | path_seq | start_vid | node | edge | cost | agg_cost
    1| 1| 6| 6| 2| 1| 0
2| 2| 6| 10| -1| 0| 1
3| 1| 8| 8| 12| 1| 0
4| 2| 8| 12| 11| 1| 1
5| 3| 8| 11| 5| 1| 2
6| 4| 8| 10|-1| 0| 3
(6 rows)
```


## Many to Many

pgr_bdAstar(Edges SQL, start vids, end vids, [options])
options: [directed, heuristic, factor, epsilon]

RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,8 \backslash\} \backslash)$ to vertices $\backslash(\backslash\{10,12 \backslash\} \backslash)$ on a directed graph with factor $\backslash(0.5 \backslash)$

## SELECT * FROM pgr_bdAstar(

'SELECT id, source, target, cost, reverse_cost, $x 1, y 1, x 2, y 2$
FROM edges',
ARRAY[6, 8], ARRAY[10, 12]
factor $=>0.5$
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1\| | 1\| | 6 | $10 \mid$ | 6 | 4\| | 1\| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 21 | $2 \mid$ | 61 | $10 \mid$ | 7\| | 81 | 1\| | 1 |
| 31 | 31 | 61 | $10 \mid$ | 11 | 91 | 1 \| | 2 |
| 4\| | 4 \| | 61 | $10 \mid$ | $16 \mid$ | 16\| | 1 \| | 3 |
| 51 | 51 | 61 | $10 \mid$ | $15 \mid$ | 3\| | 1 \| | 4 |
| 61 | 61 | 61 | $10 \mid$ | $10 \mid$ | -1\| | $0 \mid$ | 5 |
| 71 | 1 \| | 61 | 12\| | 6 | 4\| | 1\| | 0 |
| 81 | $2 \mid$ | 61 | $12 \mid$ | 71 | $10 \mid$ | 1 \| | 1 |
| 91 | 31 | 61 | 12\| | 8 | 12\| | 1 \| | 2 |
| 10 | 4 \| | 61 | $12 \mid$ | 12 | -1\| | $0 \mid$ | 3 |
| $11 \mid$ | 1 \| | 81 | $10 \mid$ | 81 | $10 \mid$ | 1 \| | 0 |
| $12 \mid$ | $2 \mid$ | 81 | $10 \mid$ | 71 | 8\| | 1 \| | 1 |
| 131 | 31 | 81 | $10 \mid$ | 11 | 91 | 1\| | 2 |
| 14 | 4 \| | 81 | $10 \mid$ | 16 | 16\| | $1 \mid$ | 3 |
| 15 | 51 | 81 | $10 \mid$ | 15 | 31 | 1 \| | 4 |
| 16 | $6 \mid$ | 8 | $10 \mid$ | 10 | -1\| | 01 | 5 |
| 17 \| | 1 \| | 8 | 12 \| | 81 | 12\| | 1\| | 0 |
| 18 \| | $2 \mid$ | 81 | 12 \| | 12 | -1\| | 0 \| | 1 |

(18 rows)

## Combinations

pgr_bdAstar(Edges SQL, Combinations SQL, [options])
options: [directed, heuristic, factor, epsilon]

RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost) OR EMPTY SET

## Example:

Using a combinations table on a directed graph with factor <br>(0.5<br>).
The combinations table:

```
SELECT * FROM combinations;
source | target
    5| 6
    5) 10
    6| 5
    6| 15
    6| 14
(5 rows)
```

```
SELECT * FROM pgr_bdAstar(
FROM edges',
'SELECT * FROM combinations',
factor => 0.5
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|c|}
\hline 1| & 11 & 51 & 61 & 5 & 1। & 1| & \\
\hline \(2 \mid\) & 21 & 51 & 61 & 6 & -1 & 01 & 1 \\
\hline 31 & 1| & 51 & 10| & 5 & 1| & 1| & 0 \\
\hline 4 | & \(2 \mid\) & 51 & \(10 \mid\) & 6| & \(4 \mid\) & 1 | & 1 \\
\hline 5| & 31 & 51 & \(10 \mid\) & 7| & \(8 \mid\) & 1| & 2 \\
\hline 61 & \(4 \mid\) & 51 & \(10 \mid\) & 11| & 9 | & 1| & 3 \\
\hline 7| & 51 & 51 & \(10 \mid\) & 16| & 16| & \(1 \mid\) & 4 \\
\hline 8| & 61 & 51 & \(10 \mid\) & 15| & \(3 \mid\) & 1| & 5 \\
\hline 9| & 71 & 51 & 10| & 10| & -1| & 01 & 6 \\
\hline \(10 \mid\) & 1| & 61 & 51 & 61 & 1| & 1| & 0 \\
\hline 11| & \(2 \mid\) & \(6 \mid\) & 51 & & -1| & 0| & 1 \\
\hline \(12 \mid\) & 1 & 61 & 15 & 6 & 4 & 1 & 0 \\
\hline 13 | & \(2 \mid\) & \(6 \mid\) & 15| & 7| & 8। & & 1 \\
\hline 14 | & 31 & 6 | & 15| & 11| & 9| & & 2 \\
\hline 15 | & 4 | & 61 & 15| & \(16 \mid\) & 16| & 1 & 3 \\
\hline 16 & 5 & 61 & 15 & 15 & -1| & 0 & 4 \\
\hline
\end{tabular}
```


## Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | BIGINT | Identifier of the starting vertex of the path. |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting vertices. |
| end vid | BIGINT | Identifier of the ending vertex of the path. |
| end vids | ARRAY[BIGINT] | Array of identifiers of ending vertices. |

## Optional parameters

| Column | Type | Default | Description |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true | When true the graph is considered Directed <br>  | Whenfalse the graph is considered as <br> Undirected. |  |

aStar optional Parameters

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| heuristic | INTEGER | 5 | ```Heuristic number. Current valid values 0~5. - 0:\(h(v) = O\) (Use this value to compare with pgr_dijkstra) 1: \(h(v) = abs(max(\Delta x,\Delta y))\) 2: \(h(v) = abs(min(\Delta x, \Delta y))\) 3: \(h(v) = \Delta x * \Delta x + \Delta y * \Delta y\) 4: \(h(v) = sqrt(\Delta x * \Delta x + \Delta y * \Delta y)\) 5: \(h(v) = abs(\Delta x) + abs(\Delta y)\)``` |
| factor | FLOAT | 1 | For units manipulation. |
| (factor > $0 \backslash$ ). |  |  |  |
| epsilon | FLOAT | 1 | For less restricted results. |
| (epsilon $>=1 \backslash$ ). |  |  |  |

See heuristics available and factor handling.

Inner Queries

Edges SQL

| Parameter | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |


| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
|  |  |  | - When negative: edge (source, target) does not exist, therefore it's not part of the graph. |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source), |
|  |  |  | - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |
| $\times 1$ | ANY-NUMERICAL |  | X coordinate of source vertex. |
| y1 | ANY-NUMERICAL |  | Y coordinate of source vertex. |
| $\times 2$ | ANY-NUMERICAL |  | $X$ coordinate of target vertex. |
| y2 | ANY-NUMERICAL |  | Y coordinate of target vertex. |

Where:

ANY-INTEGER:<br>SMALLINT, INTEGER, BIGINT<br>ANY-NUMERICAL:<br>SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Combinations SQL

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- | Identifier of the departure vertex. |
|  | INTEGER |  |
| target | ANY- | Identifier of the arrival vertex. |
|  | INTEGER |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Result Columns

Returns set of (seq, path_seq[, start_vid] [, end_vid], node, edge, cost, agg_cost)

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. |
| path_seq | INTEGER | Relative position in the path. Has value $\mathbf{1}$ for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex. Returned when multiple starting vetrices are in the query. |
|  |  | - Many to One <br> - Many to Many |
| end_vid | BIGINT | Identifier of the ending vertex. Returned when multiple ending vertices are in the query. |
|  |  | - One to Many <br> - Many to Many |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for the last node of the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | FLOAT | Aggregate cost from start_vid to node. |

Additional Examples

## Example 1:

Demonstration of repeated values are ignored, and result is sorted.

SELECT * FROM pgr_bdAstar(
'SELECT id, source, target, cost, reverse_cost, $x 1, y 1, x 2, y 2$
FROM edges',
ARRAY[7, 10, 15, 10, 10, 15], ARRAY[10, 7, 10, 15]);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost


## Example 2:

Making start vids the same as end vids.

## SELECT * FROM pgr_bdAstar(

'SELECT id, source, target, cost, reverse_cost, $\mathrm{x} 1, \mathrm{y} 1, \mathrm{x} 2, \mathrm{y} 2$
FROM edges',
ARRAY[7, 10, 15], ARRAY[7, 10, 15]);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $7 \mid$ | $10 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 0 |
| ---: | :--- | :--- | :--- | :--- | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $7 \mid$ | $10 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 1 |
| $3 \mid$ | $3 \mid$ | $7 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 2 |
| $4 \mid$ | $4 \mid$ | $7 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ | 3 |
| $5 \mid$ | $5 \mid$ | $7 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 4 |
| $6 \mid$ | $1 \mid$ | $7 \mid$ | $15 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 0 |
| $7 \mid$ | $2 \mid$ | $7 \mid$ | $15 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 1 |
| $8 \mid$ | $3 \mid$ | $7 \mid$ | $15 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 2 |
| $9 \mid$ | $4 \mid$ | $7 \mid$ | $15 \mid$ | $15 \mid$ | $-1 \mid$ | $0 \mid$ | 3 |
| $10 \mid$ | $1 \mid$ | $10 \mid$ | $7 \mid$ | $10 \mid$ | $5 \mid$ | $1 \mid$ | 0 |
| $11 \mid$ | $2 \mid$ | $10 \mid$ | $7 \mid$ | $11 \mid$ | $8 \mid$ | $1 \mid$ | 1 |
| $12 \mid$ | $3 \mid$ | $10 \mid$ | $7 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ | 2 |
| $13 \mid$ | $1 \mid$ | $10 \mid$ | $15 \mid$ | $10 \mid$ | $5 \mid$ | $1 \mid$ | 0 |
| $14 \mid$ | $2 \mid$ | $10 \mid$ | $15 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 1 |
| $15 \mid$ | $3 \mid$ | $10 \mid$ | $15 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 2 |
| $16 \mid$ | $4 \mid$ | $10 \mid$ | $15 \mid$ | $15 \mid$ | $-1 \mid$ | $0 \mid$ | 3 |
| $17 \mid$ | $1 \mid$ | $15 \mid$ | $7 \mid$ | $15 \mid$ | $16 \mid$ | $1 \mid$ | 0 |
| $18 \mid$ | $2 \mid$ | $15 \mid$ | $7 \mid$ | $16 \mid$ | $9 \mid$ | $1 \mid$ | 1 |
| $19 \mid$ | $3 \mid$ | $15 \mid$ | $7 \mid$ | $11 \mid$ | $8 \mid$ | $1 \mid$ | 2 |
| $20 \mid$ | $4 \mid$ | $15 \mid$ | $7 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ | 3 |
| $21 \mid$ | $1 \mid$ | $15 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ | 0 |
| $22 \mid$ | $2 \mid$ | $15 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |

Example 3:
Manually assigned vertex combinations.

## SELECT * FROM pgr_bdAstar(

'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
FROM edges',
'SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)'); seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost


- $A^{*}$ - Family of functions
- Bidirectional $A^{*}$ - Family of functions
- Sample Data
- https://www.boost.org/libs/graph/doc/astar_search.html
- https://en.wikipedia.org/wiki/A*_search_algorithm


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) $3.2 \mathbf{3 . 1} 3.0$
- Unsupported versions: $2.6 \mathbf{2 . 5}$
pgr_bdAstarCost
pgr_bdAstarCost - Total cost of the shortest path(s) using the bidirectional A* algorithm.


## boost

Boost Graph Inside

## Availability

- Version 3.2.0
- New proposed signature:
- pgr_bdAstarCost (Combinations)
- Version 3.0.0
- Official function
- Version 2.4.0
- New proposed function


## Description

The pgr_bdAstarCost function sumarizes of the cost of the shortest path(s) using the bidirectional A* algorithm.

## The main characteristics are:

- Process works for directed and undirected graphs.
- Ordering is:
- first by start_vid (if exists)
- then by end_vid
- Values are returned when there is a path.
- Let $\backslash(v \backslash)$ and $\backslash(u \backslash)$ be nodes on the graph:
- If there is no path from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ :
- no corresponding row is returned
- agg_cost from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ is $\backslash(\backslash i n f t y \backslash)$
- There is no path when $\backslash(v=u \backslash)$ therefore
- no corresponding row is returned
- agg_cost from $v$ to $u$ is $\backslash(0 \backslash)$
- When $\backslash((x, y) \backslash)$ coordinates for the same vertex identifier differ:
- A random selection of the vertex's $\backslash((x, y) \backslash)$ coordinates is used.
- Running time: $\backslash(\mathrm{O}((\mathrm{E}+\mathrm{V}) * \backslash \log \mathrm{~V}) \backslash)$
- It does not return a path.
- Returns the sum of the costs of the shortest path of each pair combination of nodes requested.
- Let be the case the values returned are stored in a table, so the unique index would be the pair:(start_vid, end_vid)
- For undirected graphs, the results are symmetric.
- The agg_cost of $(u, v)$ is the same as for $(v, u)$.
- The returned values are ordered in ascending order:
- start_vid ascending
- end_vid ascending


## Summary

```
pgr_bdAstarCost(Edges SQL, start vid, end vid, [options])
pgr_bdAstarCost(Edges SQL, start vid, end vids, [options])
pgr_bdAstarCost(Edges SQL, start vids, end vid, [options])
pgr_bdAstarCost(Edges SQL, start vids, end vids, [options])
pgr_bdAstarCost(Edges SQL, Combinations SQL, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET
```


## One to One

pgr_bdAstarCost(Edges SQL, start vid, end vid, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

## Example:

From vertex $\backslash(6 \backslash)$ to vertex $\backslash(12 \backslash)$ on a directed graph with heuristic $\backslash(2 \backslash)$

```
SELECT * FROM pgr_bdAstarCost(
'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
FROM edges',
6, 12,
directed => true, heuristic => 2
);
start_vid | end_vid | agg_cost
6| 12| 3
(1 row)
```


## One to Many

```
pgr_bdAstarCost(Edges SQL, start vid, end vids, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(6 \backslash)$ to vertices $\backslash(\backslash\{10,12 \backslash\} \backslash)$ on a directed graph with heuristic $\backslash(3 \backslash)$ and factor $\backslash(3.5 \backslash)$

```
SELECT * FROM pgr_bdAstarCost(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    6, ARRAY[10, 12]
heuristic => 3, factor := 3.5
);
start_vid | end_vid | agg_cost
    6 1 10| 5
(2 rows)
```


## Many to One

```
pgr_bdAstarCost(Edges SQL, start vids, end vid, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET
```


## Example:

From vertices $\backslash(\backslash\{6,8 \backslash\} \backslash)$ to vertex $\backslash(10 \backslash)$ on an undirected graph with heuristic $\backslash(4 \backslash)$

```
SELECT * FROM pgr_bdAstarCost(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    ARRAY[6, 8], 10,
false, heuristic => 4
);
start_vid | end_vid | agg_cost
    6| 10| 1
    8| 10| 3
(2 rows)
```


## Many to Many

## pgr_bdAstarCost(Edges SQL, start vids, end vids, [options])

options: [directed, heuristic, factor, epsilon]

RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,8 \backslash\} \backslash)$ to vertices $\backslash(\backslash\{10,12 \backslash\} \backslash)$ on a directed graph with factor $\backslash(0.5 \backslash)$

```
SELECT * FROM pgr bdAstarCost(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    ARRAY[6, 8], ARRAY[10, 12]
factor => 0.5
);
start_vid | end_vid | agg_cost
    6| 10| 5
    6| 12| 3
    8| 10| 5
    8| 12| 1
(4 rows)
```

Combinations
pgr_bdAstarCost(Edges SQL, Combinations SQL, [options])
options: [directed, heuristic, factor, epsilon]

RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

## Example:

Using a combinations table on a directed graph with factor $\backslash(0.5 \backslash)$.

The combinations table:

```
SELECT * FROM combinations;
source | target
    5| 6
    5| 10
    6| 5
    6| 15
    6| 14
(5 rows)
```

The query:

```
SELECT * FROM pgr bdAstarCost(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    'SELECT * FROM combinations',
factor => 0.5
);
start_vid | end_vid | agg_cost
5| 6| 1
    5| 10| 6
    6| 5| 1
    6| 15| 4
(4 rows)
```

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | BIGINT | Identifier of the starting vertex of the path. |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting vertices. |
| end vid | BIGINT | Identifier of the ending vertex of the path. |
| end vids | ARRAY[BIGINT] | Array of identifiers of ending vertices. |

Optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true | When true the graph is considered Directed <br>  |  When false the graph is considered as <br>  Undirected. |

aStar optional Parameters

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| heuristic | INTEGER | 5 | ```Heuristic number. Current valid values 0~5. 0 :\(h(v) = O\) (Use this value to compare with pgr_dijkstra) 1: \(h(v) = abs(max(\Delta x, \Delta y))\) 2: \(h(v) = abs(min(\Delta x, \Delta y)\\) 3: \(h(v) = \Delta x * \Delta x + \Delta y * \Delta y\) 4: \(h(v) = sqrt(\Delta x * \Delta x + \Delta y * \Delta y)\) 5: \(h(v) = abs(\Delta x) + abs(\Delta y)\)``` |
| factor | FLOAT | 1 | For units manipulation. |
| (factor $>0 \backslash$ ). |  |  |  |
| epsilon | FLOAT | 1 | For less restricted results. |
| (epsilon >= 1 ${ }^{\text {( }}$ ). |  |  |  |

See heuristics available and factor handling.

Inner Queries

Edges SQL

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER |  | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
|  |  |  | - When negative: edge (source, target) does not exist, therefore it's not part of the graph. |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source), |
|  |  |  | - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |
| $\times 1$ | ANY-NUMERICAL |  | $X$ coordinate of source vertex. |
| y1 | ANY-NUMERICAL |  | Y coordinate of source vertex. |
| $\times 2$ | ANY-NUMERICAL |  | X coordinate of target vertex. |
| y2 | ANY-NUMERICAL |  | Y coordinate of target vertex. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Combinations SQL

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- | Identifier of the departure vertex. |
|  | INTEGER |  |


| Parameter | Type | Description |
| :--- | :--- | :--- |
| target | ANY- | Identifier of the arrival vertex. |
|  | INTEGER |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Result Columns

Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

## Additional Examples

## Example 1:

Demonstration of repeated values are ignored, and result is sorted.

```
SELECT * FROM pgr_bdAstarCost(
    SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
ARRAY[7, 10, 15, 10, 10, 15], ARRAY[10, 7, 10, 15]);
start_vid | end_vid | agg_cost
    7| 10| 4
    7| 15| 3
    0) 15| 3
    15| 7| 3
    15| 10| 1
(6 rows)
```

Example 2:
Making start vids the same as end vids.

```
SELECT * FROM pgr_bdAstarCost(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    ARRAY[7, 10, 15], ARRAY[7, 10, 15]);
start_vid | end_vid | agg_cost
    7| 10| 4
    7| 15| 3
    10| 7| 2
    10| 15| 3
    15| 7| 3
    15| 10| 1
(6 rows)
```


## Example 3:

Manually assigned vertex combinations.

```
SELECT * FROM pgr_bdAstarCost(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2
    FROM edges',
    'SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)');
start_vid | end_vid | agg_cost
    6| 7| 1
    6| 10| 5
    12| 10| 4
(3 rows)
```

See Also

- Bidirectional $A^{*}$ - Family of functions
- Cost - Category
- Sample Data
- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: $2.6 \mathbf{2 . 5}$
pgr_bdAstarCostMatrix
pgr_bdAstarCostMatrix - Calculates the a cost matrix using pgr_aStar.


## boost

Boost Graph Inside

## Availability

- Version 3.0.0
- Official function
- Version 2.5.0
- New proposed function


## Description

## The main characteristics are:

- Using internaly the pgr_bdAstar algorithm
- Returns a cost matrix.
- No ordering is performed
- let $v$ and $u$ are nodes on the graph:
- when there is no path from $v$ to $u$ :
- no corresponding row is returned
- cost from $v$ to $u$ is <br>(\inf<br>)
- when $\backslash(v=u \backslash)$ then
- no corresponding row is returned
- cost from $v$ to $u$ is $\backslash(0 \backslash)$
- When the graph is undirected the cost matrix is symmetric

Signatures

## Summary

pgr_bdAstarCostMatrix(Edges SQL, start vids, [options])
options: [directed, heuristic, factor, epsilon]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

## Example:

Symmetric cost matrix for vertices $\backslash(\backslash\{5,6,10,15 \backslash\} \backslash)$ on an undirected graph using heuristic $\backslash(2 \backslash)$

```
SELECT * FROM pgr_bdAstarCostMatrix(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2 FROM edges',
    (SELECT array_agg(id) FROM vertices WHERE id IN (5, 6, 10, 15)),
directed => false, heuristic => 2
);
start_vid | end_vid | agg_cost
5| 6| 1
            5| 10| 2
            5| 15| 3
            6| 5| 1
            -15 |
            | 5| 2
            0| 6| 1
            | 15| 1
            5| 3
            5| 6| 2
            15| 10| 1
(12 rows)
```


## Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting <br> vertices. |

## Optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true |  When true the graph is considered Directed <br>   <br>   <br>   <br>  When n false the graph is considered as |  |

## aStar optional Parameters

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| heuristic | INTEGER | 5 | ```Heuristic number. Current valid values 0~5. 0 :\(h(v) = O\) (Use this value to compare with pgr_dijkstra) 1:\(h(v) = abs(max(\Delta x,\Delta y))\) 2: \(h(v) = abs(min(\Delta x, \Delta y))\) 3: \(h(v) = \Delta x * \Delta x + \Delta y * \Delta y\) 4: \(h(v) = sqrt(\Delta x * \Delta x + \Delta y * \Delta y)\) 5: \(h(v) = abs(\Delta x) + abs(\Delta y)\)``` |
| factor | FLOAT | 1 | For units manipulation. |
| (factor > $0 \backslash$ ). |  |  |  |
| epsilon | FLOAT | 1 | For less restricted results. |
| (epsilon $>=1 \backslash$ ). |  |  |  |

See heuristics available and factor handling.

Inner Queries

Edges SQL

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER |  | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
|  |  |  | - When negative: edge (source, target) does not exist, therefore it's not part of the graph. |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source), |
|  |  |  | - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |
| x1 | ANY-NUMERICAL |  | X coordinate of source vertex. |
| y1 | ANY-NUMERICAL |  | $Y$ coordinate of source vertex. |
| x2 | ANY-NUMERICAL |  | X coordinate of target vertex. |


| Parameter | Type | Default |
| :--- | :--- | :--- |
| y 2 | ANY-NUMERICAL |  |

## Where:

## ANY-INTEGER:

## SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

## Additional Examples

## Example:

Use with pgr_TSP

```
SELECT * FROM pgr_TSP(
    $$
    SELECT * FROM pgr_bdAstarCostMatrix(
    'SELECT id, source, target, cost, reverse_cost, x1, y1, x2, y2 FROM edges',
    (SELECT array_agg(id) FROM vertices WHERE id IN (5, 6, 10, 15)),
    directed=> false, heuristic => 2
)
$$,
randomize => false
);
NOTICE: pgr_TSP no longer solving with simulated annaeling
HINT: Ignoring annaeling parameters
seq | node | cost | agg_cost
----+----------+------
    2| 6| 1| 1
    3| 10| 1| 2
    4| 15| 1| 3
    5| 5| 3| 6
(5 rows)
```


## See Also

- Bidirectional A* - Family of functions
- Cost Matrix - Category
- Traveling Sales Person - Family of functions
- Sample Data


## Indices and tables

## - Index

- Search Page


## Description

Based on A* algorithm, the bidirectional search finds a shortest path from a starting vertex start_vid) to an ending vertex (end_vid). It runs two simultaneous searches: one forward from thestart_vid, and one backward from the end_vid, stopping when the two meet in the middle. This implementation can be used with a directed graph and an undirected graph.

The main Characteristics are:
Process works for directed and undirected graphs.

- Ordering is:
- first by start_vid (if exists)
- then by end_vid
- Values are returned when there is a path.
- Let $\backslash(v \backslash)$ and $\backslash(u)$ ) be nodes on the graph:
- If there is no path from $\backslash(\mathrm{v} \backslash)$ to $\backslash(\mathrm{u})$ ):
- no corresponding row is returned
- agg_cost from $\backslash(v \backslash)$ to $\backslash(u \backslash)$ is $\backslash(\backslash i n f t y \backslash)$
- There is no path when $\backslash(v=u \backslash)$ therefore
- no corresponding row is returned
- agg_cost from $v$ to $u$ is $\backslash(0 \backslash)$
- When $\backslash((x, y) \backslash)$ coordinates for the same vertex identifier differ:
- A random selection of the vertex's $\backslash((x, y) \backslash)$ coordinates is used.
- Running time: $\backslash(\mathrm{O}((\mathrm{E}+\mathrm{V}) * \backslash \log \mathrm{~V}) \backslash)$
- For large graphs where there is a path bewtween the starting vertex and ending vertex:
- It is expected to terminate faster than pgr_astar

See heuristics available and factor handling.

See Also

- $A^{*}$ - Family of functions
- https://www.boost.org/libs/graph/doc/astar_search.html
- https://en.wikipedia.org/wiki/A*_search_algorithm


## Indices and tables

- Index
- Search Page


## Previous versions of this page

- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: $2.6 \mathbf{2 . 5}$


## Bidirectional Dijkstra - Family of functions

- pgr_bdDijkstra - Bidirectional Dijkstra algorithm for the shortest paths.
- pgr_bdDijkstraCost - Bidirectional Dijkstra to calculate the cost of the shortest paths
- pgr_bdDijkstraCostMatrix - Bidirectional Dijkstra algorithm to create a matrix of costs of the shortest paths.
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2 2.12 .0
pgr_bdDijkstra
pgr_bdDijkstra - Returns the shortest path(s) using Bidirectional Dijkstra algorithm.


## boost

Boost Graph Inside

## Availability:

- Version 3.2.0
- New proposed signature:
pgr_bdDijkstra(Combinations)
- Version 3.0.0
- Official function
- Version 2.5.0
- New Proposed functions:
- pgr_bdDijkstra (One to Many)
- pgr_bdDijkstra (Many to One)
- pgr_bdDijkstra (Many to Many)
- Version 2.4.0
- Signature change on pgr_bdDijsktra (One to One)
- Old signature no longer supported
- Version 2.0.0
- Official pgr_bdDijkstra (One to One)


## The main characteristics are:

- Process is done only on edges with positive costs
- A negative value on a cost column is interpreted as the edge does not exist.
- Values are returned when there is a path.
- When there is no path:
- When the starting vertex and ending vertex are the same.
- The aggregate cost of the non included values $\backslash((v, v) \backslash)$ is $\backslash(0 \backslash)$
- When the starting vertex and ending vertex are the different and there is no path:
- The aggregate cost the non included values $\backslash((u, v) \backslash)$ is $\backslash(\backslash i n f t y \backslash)$
- For optimization purposes, any duplicated value in the starting vertices or on the ending vertices are ignored.
- Running time (worse case scenario): <br>( $O((V \backslash \log V+E)) \backslash)$
- For large graphs where there is a path bewtween the starting vertex and ending vertex:
- It is expected to terminate faster than pgr_dijkstra


## Signatures

## Summary

```
pgr_bdDijkstra(Edges SQL, start vid, end vid, [directed])
pgr_bdDijkstra(Edges SQL, start vid, end vids, [directed])
pgr_bdDijkstra(Edges SQL, start vids, end vid, [directed])
pgr_bdDijkstra(Edges SQL, start vids, end vids, [directed])
pgr_bdDijkstra(Edges SQL, Combinations SQL , [directed])
RETURNS SET OF (seq, path_seq, [start_vid], [end_vid], node, edge, cost, agg_cost)
OR EMPTY SET
```


## One to One

```
pgr_bdDijkstra(Edges SQL, start vid, end vid, [directed])
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(6 \backslash)$ to vertex $\backslash(10 \backslash)$ on a directed graph

```
SELECT * FROM pgr_bdDijkstra(
    'select id, source, target, cost, reverse_cost from edges',
6, 10, true);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 6| 4| 1| 0
    2| 7| 8| 1| 1
    3| 11| 9| 1| 2
    4| 16| 16| 1| 3
    5| 15| 3| 1| 4
    6| 10| -1| 0| 5
(6 rows)
```


## One to Many

```
pgr_bdDijkstra(Edges SQL, start vid, end vids, [directed])
RETURNS SET OF (seq, path_seq, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(6 \backslash)$ to vertices $\backslash(\backslash\{10,17 \backslash\} \backslash)$ on a directed graph

## SELECT * FROM pgr_bdDijkstra

'select id, source, target, cost, reverse_cost from edges'
6, ARRAY[10, 17]);
seq | path_seq | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $10 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| ---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $10 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ |
| $3 \mid$ | $3 \mid$ | $10 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ |
| $4 \mid$ | $4 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ |
| $5 \mid$ | $5 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ |
| $6 \mid$ | $6 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ |
| $7 \mid$ | $1 \mid$ | $17 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| $8 \mid$ | $2 \mid$ | $17 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ |
| $9 \mid$ | $3 \mid$ | $17 \mid$ | $11 \mid$ | $11 \mid$ | $1 \mid$ |
| 10 |  |  |  |  |  |
| $10 \mid$ | $4 \mid$ | $17 \mid$ | $12 \mid$ | $13 \mid$ | $1 \mid$ |
| $11 \mid$ | $5 \mid$ | $17 \mid$ | $17 \mid$ | $-1 \mid$ | $0 \mid$ |
| $(11$ rows) |  |  |  |  | 4 |

Many to One

```
pgr_bdDijkstra(Edges SQL, start vids, end vid, [directed])
RETURNS SET OF (seq, path_seq, start_vid, node, edge, cost, agg_cost)
```

OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,1 \backslash\} \backslash)$ to vertex $\backslash(17 \backslash)$ on a directed graph

```
SELECT * FROM pgr_bdDijkstra(
'select id, source, target, cost, reverse_cost from edges',
    ARRAY[6, 1], 17);
seq | path_seq | start_vid | node | edge | cost | agg_cost
```



Many to Many

```
pgr_bdDijkstra(Edges SQL, start vids, end vids, [directed])
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
``` OR EMPTY SET

\section*{Example:}

From vertices \(\backslash(\backslash\{6,1 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,17 \backslash\} \backslash)\) on an undirected graph

\section*{SELECT * FROM pgr_bdDijkstra(}
'select id, source, target, cost, reverse_cost from edges',
ARRAY[6, 1], ARRAY[10, 17],
directed => false);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
\begin{tabular}{ccccc|c|c|c}
\(1 \mid\) & \(1 \mid\) & \(1 \mid\) & \(10 \mid\) & \(1 \mid\) & \(6 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(1 \mid\) & \(10 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(3 \mid\) & \(1 \mid\) & \(10 \mid\) & \(7 \mid\) & \(4 \mid\) & \(1 \mid\) & 2 \\
\(4 \mid\) & \(4 \mid\) & \(1 \mid\) & \(10 \mid\) & \(6 \mid\) & \(2 \mid\) & \(1 \mid\) & 3 \\
\(5 \mid\) & \(5 \mid\) & \(1 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 4 \\
\(6 \mid\) & \(1 \mid\) & \(1 \mid\) & \(17 \mid\) & \(1 \mid\) & \(6 \mid\) & \(1 \mid\) & 0 \\
\(7 \mid\) & \(2 \mid\) & \(1 \mid\) & \(17 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) & 1 \\
\(8 \mid\) & \(3 \mid\) & \(1 \mid\) & \(17 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 2 \\
\(9 \mid\) & \(4 \mid\) & \(1 \mid\) & \(17 \mid\) & \(11 \mid\) & \(11 \mid\) & \(1 \mid\) & 3 \\
\(10 \mid\) & \(5 \mid\) & \(1 \mid\) & \(17|12|\) & \(13 \mid\) & \(1 \mid\) & 4 \\
\(11 \mid\) & \(6 \mid\) & \(1 \mid\) & \(17 \mid\) & \(17 \mid\) & \(-1 \mid\) & \(0 \mid\) & 5 \\
\(12 \mid\) & \(1 \mid\) & \(6 \mid\) & \(10 \mid\) & \(6 \mid\) & \(2 \mid\) & \(1 \mid\) & 0 \\
\(13 \mid\) & \(2 \mid\) & \(6 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 1 \\
\(14 \mid\) & \(1 \mid\) & \(6 \mid\) & \(17 \mid\) & \(6 \mid\) & \(2 \mid\) & \(1 \mid\) & 0 \\
\(15 \mid\) & \(2 \mid\) & \(6 \mid\) & \(17 \mid\) & \(10 \mid\) & \(5 \mid\) & \(1 \mid\) & 1 \\
\(16 \mid\) & \(3 \mid\) & \(6 \mid\) & \(17 \mid\) & \(11 \mid\) & \(11 \mid\) & \(1 \mid\) & 2 \\
\(17 \mid\) & \(4 \mid\) & \(6 \mid\) & \(17 \mid\) & \(12 \mid\) & \(13 \mid\) & \(1 \mid\) & 3 \\
\(18 \mid\) & \(5 \mid\) & \(6 \mid\) & \(17 \mid\) & \(17 \mid\) & \(-1 \mid\) & \(0 \mid\) & 4
\end{tabular}

\section*{Combinations}
```

pgr_bdDijkstra(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Using a combinations table on anundirected graph

The combinations table:
```

SELECT source, target FROM combinations;
source | target
5| 6
5| 10
6| 5
6| 15
6| 14
(5 rows)

```

The query:
```

SELECT * FROM pgr_bdDijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
SELECT source, target FROM combinations',
false);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $5 \mid$ | $6 \mid$ | $5 \mid$ | $1 \mid$ | $1 \mid$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $5 \mid$ | $6 \mid$ | $6 \mid$ | $-1 \mid$ | $0 \mid$ |
| $3 \mid$ | $1 \mid$ | $5 \mid$ | $10 \mid$ | $5 \mid$ | $1 \mid$ | $1 \mid$ |
| $4 \mid$ | $2 \mid$ | $5 \mid$ | $10 \mid$ | $6 \mid$ | $2 \mid$ | $1 \mid$ |
| $4 \mid$ | 1 |  |  |  |  |  |
| $5 \mid$ | $3 \mid$ | $5 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ |
| $6 \mid$ | $1 \mid$ | $6 \mid$ | $5 \mid$ | $6 \mid$ | $1 \mid$ | $1 \mid$ |
| $7 \mid$ | $2 \mid$ | $6 \mid$ | $5 \mid$ | $5 \mid$ | 0 |  |
| $8 \mid$ | $1 \mid$ | $0 \mid$ | 1 |  |  |  |
| $8 \mid$ | $6 \mid$ | $15 \mid$ | $6 \mid$ | $2 \mid$ | $1 \mid$ | 0 |
| $9 \mid$ | $2 \mid$ | $6 \mid$ | $15 \mid$ | $10 \mid$ | $3 \mid$ | $1 \mid$ |
| $10 \mid$ | $3 \mid$ | $6 \mid$ | $15 \mid$ | $15 \mid$ | $-1 \mid$ | $0 \mid$ |
| $10 \mid$ |  |  |  |  |  |  |

(10 rows)

```

\section*{Parameter}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline
\end{tabular}


Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & ANY- & Identifier of the departure vertex. \\
& INTEGER & \\
\hline target & ANY- & Identifier of the arrival vertex. \\
& INTEGER & \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Result Columns}

Returns set of (seq, path_seq [, start_vid] [, end_vid], node, edge, cost, agg_cost)
\begin{tabular}{|c|c|c|}
\hline Column & Type & Description \\
\hline seq & INTEGER & Sequential value starting from 1. \\
\hline path_seq & INTEGER & Relative position in the path. Has value 1 for the beginning of a path. \\
\hline start_vid & BIGINT & Identifier of the starting vertex. Returned when multiple starting vetrices are in the query.
Many to One
Many to Many \\
\hline end_vid & BIGINT & \begin{tabular}{l}
Identifier of the ending vertex. Returned when multiple ending vertices are in the query. \\
One to Many \\
- Many to Many
\end{tabular} \\
\hline node & BIGINT & Identifier of the node in the path fromstart_vid to end_vid. \\
\hline edge & BIGINT & Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for the last node of the path. \\
\hline cost & FLOAT & Cost to traverse from node using edge to the next node in the path sequence. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline
\end{tabular}

\section*{Example 1:}

Demonstration of repeated values are ignored, and result is sorted.
```

SELECT * FROM pgr_bdDijkstra(
'select id, source, target, cost, reverse cost from edges',
ARRAY[7, 10, 15, 10, 10, 15], ARRAY[10, 7, 10, 15]);
seq| path_seq |tart_vid | end_vid | node | edge | cost | agg_cost

| 1 \| | 1\| | 71 | 10\| | 71 | 8 | 1\| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | 7\| | 10\| | 11\| | 91 | $1 \mid$ | 1 |
| 31 | 31 | 7\| | 10\| | $16 \mid$ | 16\| | $1 \mid$ | 2 |
| 4 \| | 4 \| | 7\| | 10\| | 15\| | 31 | $1 \mid$ | 3 |
| 51 | 51 | 71 | 10\| | 10\| | -1\| | 01 | 4 |
| 61 | 1\| | 71 | 15\| | 7\| | 8\| | 1\| | 0 |
| 71 | 21 | 71 | 15 \| | 11\| | 91 | 1\| | 1 |
| 8 \| | 31 | 7\| | 15 \| | $16 \mid$ | 16\| | 1\| | 2 |
| 91 | 4\| | 71 | 151 | $15 \mid$ | -1\| | 01 | 3 |
| 10\| | 1 \| | 10\| | 71 | 10 | $2 \mid$ | 1 \| | 0 |
| 11\| | 21 | 10\| | 71 | 61 | 4। | 1 \| |  |
| 12\| | 31 | 10\| | 71 | 71 | -1\| | 01 | 2 |
| 13\| | 1 \| | $10 \mid$ | $15 \mid$ | $10 \mid$ | 51 | 1\| | 0 |
| 14\| | 21 | $10 \mid$ | $15 \mid$ | 11\| | 91 | 1 \| | 1 |
| 15\| | 31 | 10\| | $15 \mid$ | $16 \mid$ | $16 \mid$ | \| 1 | 2 |
| 16\| | 4 \| | 10\| | $15 \mid$ | 151 | -1\| | 01 | 3 |
| 17\| | 1 \| | 15\| | 7\| | 15\| | 31 | 1\| | 0 |
| 18\| | 21 | 15\| | 71 | $10 \mid$ | 21 | 1 \| | 1 |
| 19\| | 31 | 15\| | 71 | $6 \mid$ | 4\| | 1\| | 2 |
| $20 \mid$ | 4 \| | 15\| | 71 | 71 | -1\| | 01 | 3 |
| 21\| | 1 \| | 15\| | $10 \mid$ | 151 | 31 | 1\| | 0 |
| 221 | 21 | 15\| | $10 \mid$ | $10 \mid$ | -1\| | 01 | 1 |

```

\section*{Example 2:}

Making start vids the same as end vids.
```

SELECT FROM pgr bdDijkstra
'select id, source, target, cost, reverse cost from edges'
ARRAY[7, 10, 15], ARRAY[7, 10, 15]);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1 \| | 1 \| | 71 | 10\| | 71 | 8 | 1\| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2\| | 21 | 71 | 10\| | 11\| | 91 | 1\| | 1 |
| 3\| | 31 | 71 | $10 \mid$ | $16 \mid$ | 16\| | 1 \| | 2 |
| 4 \| | 4 \| | 71 | 10\| | 15\| | 3\| | 1 \| | 3 |
| 51 | 51 | 71 | $10 \mid$ | $10 \mid$ | -1\| | 01 | 4 |
| 61 | 11 | 71 | 15\| | 7\| | 8। | 1\| | 0 |
| 71 | 21 | 71 | 15\| | 11\| | 91 | $1 \mid$ | 1 |
| 8\| | 31 | 71 | 15\| | 16\| | 16\| | 1\| | 2 |
| 91 | 4 \| | 71 | 15\| | 15\| | -1\| | 01 | 3 |
| 10\| | 1 \| | $10 \mid$ | 71 | 10\| | $2 \mid$ | 1\| | 0 |
| 11\| | 2 \| | 101 | 71 | 61 | 4। | 1\| | 1 |
| 12\| | 31 | $10 \mid$ | 71 | $7 \mid$ | -1\| | 01 | 2 |
| 13\| | 1 \| | $10 \mid$ | 15 \| | $10 \mid$ | 51 | 1 \| | 0 |
| 14\| | 2 \| | $10 \mid$ | 15\| | 11\| | 91 | 1\| |  |
| 15\| | 31 | $10 \mid$ | 15 \| | $16 \mid$ | 16\| | 1 |  |
| 16\| | 4 \| | $10 \mid$ | 15\| | $15 \mid$ | -1\| | 01 |  |
| 17\| | 1 \| | 15\| | 71 | 15\| | 31 | 1\| | 0 |
| 18\| | 2 \| | 15\| | 71 | 10\| | $2 \mid$ | 1 \| | 1 |
| 19\| | 31 | 15\| | 71 | 61 | 4 \| | 1 \| | 2 |
| 20\| | $4 \mid$ | 15\| | 71 | 71 | -1\| | 01 | 3 |
| 21\| | 1\| | 15\| | 10\| | $15 \mid$ | 31 | 1 \| |  |
| $22 \mid$ | 21 | 15\| | $10 \mid$ | $10 \mid$ | -1\| | 01 |  |

(22 rows)

```

Example 3:
Manually assigned vertex combinations.

\section*{SELECT * FROM pgr_bdDijkstra}
'SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT * FROM (VALUES \((6,10),(6,7),(12,10))\) AS combinations (source, target)');
seq | path_seq \(\mid\) start_vid \(\mid\) end_vid \(\mid\) node \(\mid\) edge \(\mid\) cost | agg_cost


\section*{See Also}

\section*{- Bidirectional Dijkstra - Family of functions}
- Sample Data
- https://www.cs.princeton.edu/courses/archive/spr06/cos423/Handouts/EPP\%20shortest\%20path\%20algorithms.pdf
- https://en.wikipedia.org/wiki/Bidirectional_search

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: \(2.6 \mathbf{2 . 5}\)

\section*{pgr_bdDijkstraCost}
pgr_bdDijkstraCost — Returns the shortest path(s)'s cost using Bidirectional Dijkstra algorithm.

Boost Graph Inside

\section*{Availability}
- Version 3.2.0
- New proposed signature:
- pgr_bdDijkstraCost (Combinations)
- Version 3.0.0
- Official function
- Version 2.5.0
- New proposed function

\section*{Description}

The pgr_bdDijkstraCost function sumarizes of the cost of the shortest path using the bidirectional Dijkstra Algorithm.
- Process is done only on edges with positive costs.
- A negative value on a cost column is interpreted as the edge does not exist.
- Values are returned when there is a path.
- When there is no path:
- When the starting vertex and ending vertex are the same.
- The aggregate cost of the non included values \(\backslash((v, v) \backslash)\) is \(\backslash(0 \backslash)\)
- When the starting vertex and ending vertex are the different and there is no path:
- The aggregate cost the non included values \(\backslash((u, v) \backslash)\) is \(\backslash(\backslash i n f t y \backslash)\)
- For optimization purposes, any duplicated value in the starting vertices or on the ending vertices are ignored.
- Running time (worse case scenario): \\( \(O((V \backslash \log V+E)) \backslash)\)
- For large graphs where there is a path bewtween the starting vertex and ending vertex:
- It does not return a path.
- Returns the sum of the costs of the shortest path of each pair combination of nodes requested.
- Let be the case the values returned are stored in a table, so the unique index would be the pair:(start_vid, end_vid).
- Depending on the function and its parameters, the results can be symmetric.
- The aggregate cost of \(\backslash((u, v) \backslash)\) is the same as for \(\backslash((v, u) \backslash)\).
- Any duplicated value in the start or end vertex identifiers are ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending

\section*{Signatures}

\section*{Summary}
```

pgr_bdDijkstraCost(Edges SQL, start vid, end vid , [directed])
pgr_bdDijkstraCost(Edges SQL, start vid, end vids, [directed])
pgr_bdDijkstraCost(Edges SQL, start vids, end vid, [directed])
pgr_bdDijkstraCost(Edges SQL, start vids, end vids, [directed])
pgr_bdDijkstraCost(Edges SQL, Combinations SQL, [ directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{One to One}
```

pgr_bdDijkstraCost(Edges SQL, start vid, end vid , [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertex \(\backslash(6 \backslash)\) to vertex \(\backslash(10 \backslash)\) on a directed graph
```

SELECT * FROM pgr_bdDijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6,10, true);
start_vid | end_vid | agg_cost
6| 10| 5
(1 row)

```

\section*{One to Many}
```

pgr_bdDijkstraCost(Edges SQL, start vid, end vids, [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertex \(\backslash(6 \backslash)\) to vertices \(\backslash(\backslash\{10,17 \backslash\} \backslash)\) on a directed graph
```

SELECT * FROM pgr_bdDijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6, ARRAY[10, 17]);
start_vid | end_vid | agg_cost
6|---------------------
(2 rows)

```

\section*{Many to One}
```

pgr_bdDijkstraCost(Edges SQL, start vids, end vid , [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{6,1 \backslash\} \backslash)\) to vertex \(\backslash(17 \backslash)\) on a directed graph
```

SELECT * FROM pgr_bdDijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[6, 1], 17);
start_vid | end_vid | agg_cost
1| 17| 5
6| 17| 4
(2 rows)

```

\section*{Many to Many}
```

pgr_bdDijkstraCost(Edges SQL, start vids, end vids, [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{6,1 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,17 \backslash\} \backslash)\) on an undirected graph
```

SELECT * FROM pgr_bdDijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[6, 1], ARRAY[10, 17],
directed => false);
start_vid | end_vid | agg_cost

```
    1| \(10 \mid 4\)
    1| 17 | 5

(4 rows)

\section*{Combinations}
```

pgr_bdDijkstraCost(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Using a combinations table on anundirected graph
The combinations table:
```

SELECT source, target FROM combinations;
source | target
5| 6
5| 10
6| 5
6| 15
6| 14
(5 rows)

```

The query:
```

SELECT * FROM pgr_bdDijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT source, target FROM combinations',
false);
start_vid | end_vid | agg_cost
5| 6| 1
5| 10| 2
6| 5| 1
(4 rows)

```
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}

Optional parameters
\begin{tabular}{lllllll} 
Column & Type & Default & Description \\
\hline directed & BOOLEAN true & \begin{tabular}{l} 
When true the graph is consideredDirected \\
\\
\end{tabular} & \begin{tabular}{ll} 
Whe n false the graph is considered as \\
Undirected.
\end{tabular} & \\
\hline
\end{tabular}

\section*{Inner Queries}

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL -1 & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& & \\
& & \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & ANY- & Identifier of the departure vertex. \\
& INTEGER & \\
\hline target & \begin{tabular}{l} 
ANY- \\
\\
INTEGER
\end{tabular} & Identifier of the arrival vertex. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Result Columns}

Set of (start_vid, end_vid, agg_cost)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline start_vid & BIGINT & Identifier of the starting vertex. \\
\hline end_vid & BIGINT & Identifier of the ending vertex. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to end_vid. \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example 1:}

Demonstration of repeated values are ignored, and result is sorted.
\begin{tabular}{rcc}
\(7 \mid\) & \(10 \mid\) & 4 \\
\(7 \mid\) & \(15 \mid\) & 3 \\
\(10 \mid\) & \(7 \mid\) & 2 \\
\(10 \mid\) & \(15 \mid\) & 3 \\
\(15 \mid\) & \(7 \mid\) & 3 \\
\(15 \mid\) & \(10 \mid\) & 1
\end{tabular}

\section*{Example 2:}

Making start vids the same as end vids.
```

SELECT * FROM pgr_bdDijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[7, 10, 15], ARRAY[7, 10, 15]);
start_vid | end_vid | agg_cost
---------+---------------
7| 15| 3
10| 7| 2
10| 15| 3
15| 7| 3
15| 10| 1
(6 rows)

```

\section*{Example 3:}

Manually assigned vertex combinations.
```

SELECT * FROM pgr_bdDijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)');
start_vid | end_vid | agg_cost
6| 7| 1
6| 10| 5
12| 10| 4
(3 rows)

```

See Also

Bidirectional Dijkstra - Family of functions
- Sample Data
- https://www.cs.princeton.edu/courses/archive/spr06/cos423/Handouts/EPP\%20shortest\%20path\%20algorithms.pdf
- https://en.wikipedia.org/wiki/Bidirectional_search

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: \(2.6 \mathbf{2 . 5}\)
pgr_bdDijkstraCostMatrix
pgr_bdDijkstraCostMatrix - Calculates a cost matrix using pgr_bdDijkstra

Boost Graph Inside

\section*{Availability}

\footnotetext{
- Version 3.0.0
- Official function
- Version 2.5.0
}

\section*{Description}

Using bidirectional Dijkstra algorithm, calculate and return a cost matrix.
- Process is done only on edges with positive costs.
- A negative value on a cost column is interpreted as the edge does not exist.
- Values are returned when there is a path.
- When there is no path:
- When the starting vertex and ending vertex are the same.
- The aggregate cost of the non included values \(\backslash((v, v) \backslash)\) is \(\backslash(0 \backslash)\)
- When the starting vertex and ending vertex are the different and there is no path:
- The aggregate cost the non included values \(\backslash((u, v) \backslash)\) is \(\backslash(\backslash i n f t y \backslash)\)
- For optimization purposes, any duplicated value in the starting vertices or on the ending vertices are ignored.
- Running time (worse case scenario): \\(O((V \og V + E)) \\)
- For large graphs where there is a path bewtween the starting vertex and ending vertex:
- It is expected to terminate faster than pgr_dijkstra

The main Characteristics are:
- Can be used as input topgr_TSP.
- Use directly when the resulting matrix is symmetric and there is nd(\infty\\) value.
- It will be the users responsibility to make the matrix symmetric.
- By using geometric or harmonic average of the non symmetric values.
- By using max or min the non symmetric values.
- By setting the upper triangle to be the mirror image of the lower triangle.
- By setting the lower triangle to be the mirror image of the upper triangle.
- It is also the users responsibility to fix an\\(\infty\\) value.
- Each function works as part of the family it belongs to.
- It does not return a path.
- Returns the sum of the costs of the shortest path for pair combination of nodes in the graph.
- Process is done only on edges with positive costs.
- Values are returned when there is a path.
- When the starting vertex and ending vertex are the same, there is no path.
- The aggregate cost in the non included values \((v, v)\) is 0 .
- When the starting vertex and ending vertex are the different and there is no path. - The aggregate cost in the non included values(u,v) is \\(\infty\\).
- Let be the case the values returned are stored in a table:
- The unique index would be the pair: (start_vid, end_vid).
- Depending on the function and its parameters, the results can be symmetric.
- The aggregate cost of \((u, v)\) is the same as for \((v, u)\).
- Any duplicated value in the start vids are ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending

\section*{Signatures}

\section*{Summary}
```

pgr_bdDijkstraCostMatrix(Edges SQL, start vids, [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Symmetric cost matrix for vertices \(\backslash(\backslash\{5,6,10,15 \backslash\} \backslash)\) on an undirected graph
```

SELECT * FROM pgr bdDijkstraCostMatrix(
'SELECT id, source, target, cost, reverse_cost FROM edges',
(SELECT array_agg(id)
FROM vertices
WHERE id IN (5, 6, 10, 15)),
false);
start_vid | end_vid | agg_cost
5| 6| 1
5| 10| 2
5| 15| 3
5| 1
|0| 1
6| 15| 2
10| 5| 2
|| 6
| 15| 1
5| 5| 3
5| 6| 2
15| 10| 1
(12 rows)

```

Parameters
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline start vids & ARRAY[BIGINT] & \begin{tabular}{l} 
Array of identifiers of starting \\
vertices.
\end{tabular}
\end{tabular}

\section*{Optional parameters}


Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Set of (start_vid, end_vid, agg_cost)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline start_vid & BIGINT & Identifier of the starting vertex. \\
\hline end_vid & BIGINT & Identifier of the ending vertex. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to end_vid. \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Use with pgr_TSP.
```

SELECT * FROM pgr_TSP(

$$
SELECT * FROM pgr_bdDijkstraCostMatrix(
    'SELECT id, source, target, cost, reverse_cost FROM edges',
    (SELECT array_agg(id)
        FROM vertices
        WHERE id IN (5, 6, 10, 15)),
    false
$$);

NOTICE: pgr_TSP no longer solving with simulated annaeling
HINT: Ignoring annaeling parameters
seq | node | cost | agg_cost
1| 5| 0| 0
6| 1| 1
10| 1| 2
15| 1| 3
5| 5| 3| 6
(5 rows)

```

\section*{See Also}

\section*{- Bidirectional Dijkstra - Family of functions}
- Cost Matrix - Category
- Traveling Sales Person - Family of functions
- Sample Data

\section*{Indices and tables}
- Index
- Search Page

\section*{Synopsis}

Based on Dijkstra's algorithm, the bidirectional search finds a shortest path a starting vertex to an ending vertex.

It runs two simultaneous searches: one forward from the source, and one backward from the target, stopping when the two meet in the middle

This implementation can be used with a directed graph and an undirected graph.

\section*{Characteristics}

The main Characteristics are:
- Process is done only on edges with positive costs.
- A negative value on a cost column is interpreted as the edge does not exist.
- Values are returned when there is a path.
- When there is no path:
- When the starting vertex and ending vertex are the same.
- The aggregate cost of the non included values \(\backslash((v, v) \backslash)\) is \(\backslash(0 \backslash)\)
- When the starting vertex and ending vertex are the different and there is no path:
- The aggregate cost the non included values \(\backslash((u, v) \backslash)\) is \(\backslash(\backslash i n f t y \backslash)\)
- For optimization purposes, any duplicated value in the starting vertices or on the ending vertices are ignored.
- Running time (worse case scenario): \\(O((V \log V + E)) \\)
- For large graphs where there is a path bewtween the starting vertex and ending vertex:
- It is expected to terminate faster than pgr_dijkstra

See Also

\section*{Indices and tables}
```

Index
Search Page

```

Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: \(2.6 \mathbf{2 . 5}\)

\section*{Components - Family of functions}
pgr_connectedComponents - Connected components of an undirected graph.
- pgr_strongComponents - Strongly connected components of a directed graph.
- pgr_biconnectedComponents - Biconnected components of an undirected graph.
pgr_articulationPoints - Articulation points of an undirected graph.
- pgr_bridges - Bridges of an undirected graph.

\section*{Experimental}

\section*{Warning}

Possible server crash
- These functions might create a server crash

\section*{Warning}

Experimental functions
- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need \(c / c++\) coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting
- pgr_makeConnected - Experimental - Details of edges to make graph connected.
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: \(2.6 \mathbf{2 . 5}\)
pgr_connectedComponents
pgr_connectedComponents - Connected components of an undirected graph using a DFS-based approach.

\section*{boost}

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- Return columns change:
- n_seq is removed
- seq changed type to BIGINT
- Official function
- Version 2.5.0
- New experimental function

\section*{Description}

A connected component of an undirected graph is a set of vertices that are all reachable from each other.

\section*{The main characteristics are:}
- Works for undirected graphs.
- Components are described by vertices
- The returned values are ordered:
- component ascending
- node ascending
- Running time: \(\backslash(O(V+E) \backslash)\)
```

pgr_connectedComponents(Edges SQL)
RETURNS SET OF (seq, component, node)
OR EMPTY SET

```

\section*{Example:}

The connected components of the graph
```

SELECT * FROM pgr connectedComponents
'SELECT id, source, target, cost, reverse_cost FROM edges'
);
seq | component | node
1| 1| 1
2| 1| 3
3| 1| 5
4| 1/ 6
| 1| 7
1| 8
1|}
1| 10
1| 11
1| 12
1| 15
1) }1
1| 17
2| 2
2| 4
13| 13
13| }1
(17 rows)

```

\section*{Parameters}
\begin{tabular}{llll} 
Parameter & Type & Description \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below.
\end{tabular} \\
\hline
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
\hline
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns set of (seq, component, node)
\begin{tabular}{lcl} 
Column & Type & Description \\
\hline seq & BIGINT & Sequential value starting from 1. \\
\hline component & BIGINT & \begin{tabular}{c} 
Component identifier. \\
\\
\\
node
\end{tabular} \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Connecting disconnected components}

To get the graph connectivity:
```

SELECT * FROM pgr_connectedComponents
'SELECT id, source, target, cost, reverse_cost FROM edges'
);
seq | component | node
1| 1| 1
2| 1| 3
3| 1| 5
4| 1| 6
5| 1| 7
6| 1| 8
7| 1| 9
8| 1| 10
9| 1| 11
| 1| 12
10| 1/ 13
12 1 1 14
5| 1| 17
16| 1| 18
17| 2| 2
18| 2| 4
(18 rows)

```

In this example, the component \(\backslash(2 \backslash)\) consists of vertices \(\backslash(\backslash\{2,4 \backslash\} \backslash)\) and both vertices are also part of the dead end result set. This graph needs to be connected.

\section*{Note}

With the original graph of this documentation, there would be 3 components as the crossing edge in this graph is a different component.

\section*{Prepare storage for connection information}

ALTER TABLE vertices ADD COLUMN component BIGINT;
ALTER TABLE
ALTER TABLE edges ADD COLUMN component BIGINT;
ALTER TABLE

\section*{Save the vertices connection information}
```

UPDATE vertices SET component = c.component
FROM (SELECT * FROM pgr_connectedComponents(
'SELECT id, source, target, cost, reverse_cost FROM edges'
)) AS C
WHERE id = node;
UPDATE }1

```

\section*{Save the edges connection information}
```

UPDATE edges SET component = v.component
FROM (SELECT id, component FROM vertices) AS v
WHERE source = v.id;
UPDATE }2

```

The closest vertex to component \(\backslash(1 \backslash)\) is vertex \(\backslash(4 \backslash)\). And the closest edge to vertex \(\backslash(4 \backslash)\) is edge \(\backslash(14 \backslash)\).
```

WITH
edges_sql AS (SELECT id, geom FROM edges WHERE component = 1),
point_sql AS (SELECT geom AS point FROM vertices WHERE component = 2),
results AS (
SELECT
id::BIGINT AS edge_id,
ST_LineLocatePoint(geom, point) AS fraction,
CASE WHEN ST_Intersects(ST_Buffer(geom, 2, 'side=right endcap=flat'), point)
THEN 'r'
ELSE 'I' END::CHAR AS side,
geom <-> point AS distance,
point,
ST_MakeLine(point, ST_ClosestPoint(geom, point)) AS new_line
FRO-MM edges_sql, point_sql
WHERE ST_DWithin(geom, point, 2)
ORDER BY geom <-> point),
prepare_cap AS (
SELECT row_number() OVER (PARTITION BY point ORDER BY point, distance) AS rn, *
FROM results),
cap AS (
SELECT edge_id, fraction, side, distance, point, new_line
FROM prepare_cap
WHERE rn <= 1
)
SELECT edge_id, fraction, side, distance, point AS geom, new_line AS edge, id AS closest_vertex
INTO closest
FROM cap JOIN vertices ON (point = geom) ORDER BY distance LIMIT 1;
SELECT 1

```

The edge can be used to connect the components, using thefraction information about the edge \(\backslash(14 \backslash)\) to split the connecting edge.

\section*{Connecting components}

There are three basic ways to connect the components
- From the vertex to the starting point of the edge
- From the vertex to the ending point of the edge
- From the vertex to the closest vertex on the edge - This solution requires the edge to be split.

The following query shows the three ways to connect the components:
```

WITH
info AS (
SELECT
edge_id, fraction, side, distance, ce.geom, edge, v.id AS closest,
source, target, capacity, reverse_capacity, e.geom AS e_geom
FROM closest AS ce
JOIN vertices AS v USING (geom)
JOIN edges AS e ON (edge_id = e.id)
ORDER BY distance LIMIT 1),
three_options AS (
SELECT
closest AS source, target, 0 AS cost, 0 AS reverse_cost,
capacity, reverse_capacity,
ST_X(geom) AS x1, ST_Y(geom) AS y1,
ST_X(ST_EndPoint(e_geom)) AS x2, ST_Y(ST_EndPoint(e_geom)) AS y2,
ST_MakeLine(geom, ST_EndPoint(e_geom)) AS geom
FROM info
UNION
SELECT closest, source, 0, 0, capacity, reverse_capacity,
ST_X(geom) AS x1, ST_Y(geom) AS y1,
ST_X(ST_StartPoint(e_geom)) AS x2, ST_Y(ST_StartPoint(e_geom)) AS y2,
ST_MakeLine(info.geom, ST_StartPoint(e_geom))
FROM info
/*
UNION
- This option requires splitting the edge
SELECT closest, NULL, 0, 0, capacity, reverse_capacity,
ST_X(geom) AS x1, ST_Y(geom) AS y1,
ST_X(ST_EndPoint(edge)) AS x2, ST_Y(ST_EndPoint(edge)) AS y2,
edge
FROM info */
)
INSERT INTO edges
(source, target,
cost, reverse_cost,
capacity, reverse_capacity,
x1, y1, x2, y2,
geom)
(SELECT
source, target, cost, reverse_cost, capacity, reverse_capacity,
x1,y1, x2, y2, geom
FROM three_options);
INSERT02

```

\section*{Checking components}

Ignoring the edge that requires further work. The graph is now fully connected as there is only one component.
```

SELECT * FROM pgr_connectedComponents(
'SELECT id, source, target, cost, reverse_cost FROM edges'
);
seq | component | node

1) 2
|
1| 3
1/ 4
1| 5
1| 6
1| 7
1| 8
1| 9
1| 10
1| 11
1| 12
1| 13
1| 14
1| 15
1| 16
1| 17
1| 18
(18 rows)
```

See Also
- Components - Family of functions
- The queries use the Sample Data network.
- Boost: Connected components
- wikipedia: Connected component

\section*{Indices and tables}
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: \(2.6 \mathbf{2 . 5}\)
pgr_strongComponents
pgr_strongComponents - Strongly connected components of a directed graph using Tarjan's algorithm based on DFS.

\section*{boost}

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- Return columns change:
- n_seq is removed
- seq changed type to BIGINT
- Official function
- Version 2.5.0
- New experimental function

\section*{Description}

A strongly connected component of a directed graph is a set of vertices that are all reachable from each other.

\section*{The main characteristics are:}
- Works for directed graphs.
- Components are described by vertices identifiers.
- The returned values are ordered:
- component ascending
- node ascending
- Running time: \(\backslash(O(V+E) \backslash)\)

\section*{Signatures}
```

pgr_strongComponents(Edges SQL)
RETURNS SET OF (seq, component, node)

```
OR EMPTY SET

\section*{Example:}

The strong components of the graph
```

SELECT * FROM pgr_strongComponents(
'SELECT id, source, target, cost, reverse_cost FROM edges'
);
seq | component | node

| $1 \mid$ | $1 \mid c$ | 1 |
| :---: | :---: | :---: |
| $2 \mid$ | $1 \mid$ | 3 |
| $3 \mid$ | $1 \mid$ | 5 |
| $4 \mid$ | $1 \mid$ | 6 |
| $5 \mid$ | $1 \mid$ | 7 |
| $6 \mid$ | $1 \mid$ | 8 |
| $7 \mid$ | $1 \mid$ | 9 |
| $8 \mid$ | $1 \mid$ | 10 |
| $9 \mid$ | $1 \mid$ | 11 |
| $10 \mid$ | $1 \mid$ | 12 |
| $11 \mid$ | $1 \mid$ | 15 |
| $12 \mid$ | $1 \mid$ | 16 |
| $13 \mid$ | $1 \mid$ | 17 |
| $14 \mid$ | $2 \mid$ | 2 |
| $15 \mid$ | $2 \mid$ | 4 |
| $16 \mid$ | $13 \mid$ | 13 |
| $17 \mid$ | $13 \mid$ | 14 |
| $(17$ rows $)$ |  |  |

(17 rows)

```


Parameters
\begin{tabular}{lllll} 
Parameter & Type & Description \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below.
\end{tabular} \\
\hline
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& & \\
& & \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Result Columns

Returns set of (seq, component, node)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & BIGINT & Sequential value starting from 1. \\
\hline component & BIGINT & \begin{tabular}{c} 
Component identifier. \\
o
\end{tabular} \\
\hline Has the value of the minimum node identifier in the \\
nomponent.
\end{tabular}

\section*{See Also}
- Components - Family of functions
- The queries use the Sample Data network.
- Boost: Strong components
- wikipedia: Strongly connected component

\section*{Indices and tables}
- Index
- Search Page

\footnotetext{
- Supported versions: Latest (3.3) 3.23 .13 .0
}
- Unsupported versions: \(2.6 \mathbf{2 . 5}\)

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- Return columns change:
- n_seq is removed
- seq changed type to BIGINT
- Official function
- Version 2.5.0
- New experimental function

\section*{Description}

The biconnected components of an undirected graph are the maximal subsets of vertices such that the removal of a vertex from particular component will not disconnect the component. Unlike connected components, vertices may belong to multiple biconnected components. Vertices can be present in multiple biconnected components, but each edge can only be contained in a single biconnected component.

The main characteristics are:
- Works for undirected graphs.
- Components are described by edges.
- The returned values are ordered:
- component ascending.
- edge ascending.
- Running time: \(\backslash(O(V+E) \backslash)\)

\section*{Signatures}
```

pgr_biconnectedComponents(Edges SQL)
RETURNS SET OF (seq, component, edge)
OR EMPTY SET

```

\section*{Example:}

The biconnected components of the graph
```

SELECT * FROM pgr_biconnectedComponents(
'SELECT id, source, target, cost, reverse_cost FROM edges'
);
seq | component | edge
1| 1| 1
2| 2| 2
2| 3
2|4
2| 5
2| 8
2|}
2| 10
2| 11
2| 12
2| 13
2| 15
2| 16
| 6
7
74|}
17| 17
18|}1
(18 rows)

```


\section*{Parameters}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below.
\end{tabular}
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL -1 & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& & \\
& & \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns set of (seq, component, edge)
\begin{tabular}{|c|c|c|}
\hline Column & Type & Description \\
\hline seq & BIGINT & Sequential value starting from 1. \\
\hline component & BIGINT & \begin{tabular}{l}
Component identifier. \\
- Has the value of the minimum edge identifier in the component.
\end{tabular} \\
\hline edge & BIGINT & Identifier of the edge that belongs to thecomponent. \\
\hline
\end{tabular}

See Also
- Components - Family of functions
- The queries use the Sample Data network.
- Boost: Biconnected components
- wikipedia: Biconnected component

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: \(2.6 \mathbf{2 . 5}\)

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- Return columns change: seq is removed
- Official function
- Version 2.5.0
- New experimental function

\section*{Description}

Those vertices that belong to more than one biconnected component are called articulation points or, equivalently, cut vertices. Articulation points are vertices whose removal would increase the number of connected components in the graph. This implementation can only be used with an undirected graph.

\section*{The main characteristics are:}
- Works for undirected graphs.
- The returned values are ordered:
- node ascending
- Running time: \(\backslash(O(V+E) \backslash)\)

Signatures
pgr_articulationPoints(Edges SQL)
RETURNS SET OF (node)
OR EMPTY SET

\section*{Example:}

The articulation points of the graph
```

SELECT * FROM pgr_articulationPoints(
SELECT id, source, target, cost, reverse_cost FROM edges'
);
node
3
6
8
(4 rows)

```

Nodes in red are the articulation points.


Parameters
\begin{tabular}{lllll} 
Parameter & Type & Description \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below.
\end{tabular} & & \\
& & as
\end{tabular}

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Result Columns

Returns set of (node)
\begin{tabular}{llll} 
Column & Type & Description \\
\hline node & BIGINT & \\
& & \begin{tabular}{l} 
Identifier of the \\
vertex.
\end{tabular} & \\
\hline
\end{tabular}

\section*{See Also}
- Components - Family of functions
- The queries use the Sample Data network.
- Boost: Biconnected components \& articulation points
- wikipedia: Biconnected component

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: \(2.6 \mathbf{2 . 5}\)
pgr_bridges
pgr_bridges - Return the bridges of an undirected graph.

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- Return columns change: seq is removed
- Official function
- Version 2.5.0
- New experimental function

\section*{Description}

A bridge is an edge of an undirected graph whose deletion increases its number of connected components. This implementation can only be used with an undirected graph.

\section*{The main characteristics are:}
- Works for undirected graphs.
- The returned values are ordered:
- edge ascending
- Running time: \(\backslash(O(E *(V+E)) \backslash)\)

\section*{Signatures}
```

pgr_bridges(Edges SQL)
RETURNS SET OF (edge)
OR EMPTY SET

```

\section*{Example:}

The bridges of the graph
```

SELECT * FROM pgr_bridges(
'SELECT id, source, target, cost, reverse_cost FROM edges'
);
edge
1
6
7
14
18
(6 rows)

```

Parameters
\begin{tabular}{lllll} 
Parameter & Type & Description \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below.
\end{tabular} &
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL & -1 & Weight of the edge (target, source) \\
& & & When negative: edge (target, source) does not exist, therefore it's \\
& & & not part of the graph.
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT
\begin{tabular}{lll} 
Column & Type & Description \\
\hline edge & BIGINT & \begin{tabular}{l} 
Identifier of the edge that is a \\
bridge.
\end{tabular}
\end{tabular}

See Also
- https://en.wikipedia.org/wiki/Bridge_\%28graph_theory\%29
- The queries use the Sample Data network.

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2
pgr_makeConnected - Experimental
pgr_makeConnected - Set of edges that will connect the graph.


Boost Graph Inside

\section*{Warning}

Possible server crash
- These functions might create a server crash

\section*{Warning}

\section*{Experimental functions}
- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need \(c / c++\) coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting

\section*{Availability}

\section*{- Version 3.2.0}
- New experimental function

\section*{Description}

Adds the minimum number of edges needed to make the input graph connected. The algorithm first identifies all of the connected components in the graph, then adds edges to connect those components together in a path. For example, if a graph contains three connected components \(\mathrm{A}, \mathrm{B}\), and C , make_connected will add two edges. The two edges added might consist of one connecting a vertex in \(A\) with a vertex in \(B\) and one connecting a vertex in \(B\) with a vertex in \(C\).

The main characteristics are:

Works for undirected graphs.
- It will give a minimum list of all edges which are needed in the graph to make connect it.
- The algorithm does not considers traversal costs in the calculations.
- The algorithm does not considers geometric topology in the calculations.
- Running time: \(\backslash(O(V+E) \backslash)\)

Signatures
```

pgr_makeConnected(Edges SQL)
RETURNS SET OF (seq, start_vid, end_vid)
OR EMPTY SET

```

\section*{Example:}

Query done on Sample Data network gives the list of edges that are needed to connect the graph.
```

SELECT * FROM pgr_makeConnected(
'SELECT id, source, target, cost, reverse_cost FROM edges'
);
seq| start vid | end vid
1| 5| 2
2| 4| 13
(2 rows)

```

Parameters
\begin{tabular}{lllll} 
Parameter & Type & Description & \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below.
\end{tabular} & &
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL -1 & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& & \\
& & \\
& & \\
& &
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns set of (seq, start_vid, end_vid)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & BIGINT & Sequential value starting from \(\mathbf{1 .}\) \\
\hline start_vid & BIGINT & \begin{tabular}{l} 
Identifier of the first end point vertex of the edge. \\
\hline end_vid
\end{tabular} \\
& BIGINT & \begin{tabular}{l} 
Identifier of the second end point vertex of the \\
edge.
\end{tabular} \\
\hline
\end{tabular}

\footnotetext{
See Also
- https://www.boost.org/libs/graph/doc/make_connected.html
- The queries use the Sample Data network.
}
- Index
- Search Page

See Also

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: \(2.6 \mathbf{2}\).5 2.42 .32 .2

\section*{Contraction - Family of functions}
- pgr_contraction
- Supported versions: Latest (3.3) 3.2 \(3.1 \mathbf{3 . 0}\)
- Unsupported versions: \(2.6 \mathbf{2 . 5} 2.42 .3\)
pgr_contraction
pgr_contraction - Performs graph contraction and returns the contracted vertices and edges.

\section*{boost}

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- Return columns change: seq is removed
- Name change from pgr_contractGraph
- Bug fixes
- Official function
- Version 2.3.0
- New experimental function

\section*{Description}

Contraction reduces the size of the graph by removing some of the vertices and edges and, for example, might add edges that represent a sequence of original edges decreasing the total time and space used in graph algorithms.

The main Characteristics are:
- Process is done only on edges with positive costs.
- Does not return the full contracted graph
- Only changes on the graph are returned
- Currnetly there are two types of contraction methods
- Dead End Contraction
- Linear Contraction
- The returned values include
- the added edges by linear contraction.
- the modified vertices by dead end contraction.
- The returned values are ordered as follows:
- column id ascending when type is v
- column id descending when type ise

\section*{Signatures}

\section*{Summary}

The pgr_contraction function has the following signature:
options: [ max cycles, forbidden vertices, directed]
RETURNS SET OF (type, id, contracted_vertices, source, target, cost)

\section*{Example:}

Making a dead end and linear contraction in that order on an undirected graph.
```

SELECT * FROM pgr_contraction(
'SELECT id, source, target, cost, reverse cost FROM edges',
ARRAY[1, 2], directed => false);
type | id | contracted_vertices | source | target | cost
v | 4|{2} | -1| -1| -1
v | 7 |{1,3} | -1| -1| -1
v | 14|{13} | -1| -1| -1
e |-1|{5,6} | 7| 10| 2
e |-2|{8,9} | 7| 12| 2
e |-3|{17} | 12| 16| 2
e |-4|{15} | 10| 16| 2
(7 rows)

```

Parameters
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline Edges SQL & TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below.
\end{tabular} \\
\hline contraction Order & ARRAY[ ANY-INTEGER] & Ordered contraction operations. \\
& & o \(1=\) Dead end contraction \\
& & - \(2=\) Linear contraction
\end{tabular}

Optional Parameters
Column Type Default Description
\begin{tabular}{ll} 
directed BOOLEAN true & When true the graph is considered Directed \\
When false the graph is considered as \\
Undirected.
\end{tabular}

Contraction optional parameters
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline forbidden_vertices & ARRAY[ ANY-INTEGER] & Empty & Identifiers of vertices forbidden for contraction. \\
\hline max_cycles & INTEGER & \(\backslash(1 \backslash)\) & \begin{tabular}{l} 
Number of times the contraction operations on \\
contraction_order will be performed.
\end{tabular} \\
\hline
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & \\
& & Weight of the edge (target, source) \\
& & \\
& & not part of the graph.
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

The function returns a single row. The columns of the row are:
\begin{tabular}{|c|c|c|}
\hline Column & Type & Description \\
\hline type & TEXT & \begin{tabular}{l}
Type of the id. \\
- \(v\) when the row is a vertex. \\
- Column id has a positive value \\
- e when the row is an edge. \\
- Column id has a negative value
\end{tabular} \\
\hline id & BIGINT & \begin{tabular}{l}
All numbers on this column aredISTINCT \\
- When type \(=\) ' \(\mathbf{v}\) '. \\
- Identifier of the modified vertex. \\
- When type \(=\) ' \(\mathbf{e}\) '. \\
- Decreasing sequence starting from-1. \\
- Representing a pseudo id as is not incorporated in the set of original edges.
\end{tabular} \\
\hline contracted_vertices & ARRAY[BIGINT] & Array of contracted vertex identifiers. \\
\hline source & BIGINT & \begin{tabular}{l}
When type \(=\) ' \(\mathbf{v}\) ': \\(-1\\) \\
When type \(=\mathbf{~ ' ~} \mathbf{e}\) ': Identifier of the source vertex of the current edge source, target).
\end{tabular} \\
\hline target & BIGINT & \begin{tabular}{l}
- When type \(=\) ' \(\mathbf{v}\) ': \\(-1\\) \\
- When type \(=\) ' \(\mathbf{e}\) ': Identifier of the target vertex of the current edge source, target).
\end{tabular} \\
\hline cost & FLOAT & \begin{tabular}{l}
When type \(=\) ' \(\mathbf{v}\) ': \\(-1\\) \\
- When type = 'e': Weight of the current edge (source, target).
\end{tabular} \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Only dead end contraction
```

SELECT type, id, contracted_vertices FROM pgr_contraction(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[1]);
type | id | contracted_vertices
v | 4|{2}
v | 6|{5}
v | 7 |{1,3}
v | 8|{9}
v | 14|{13}
(5 rows)

```

\section*{Example:}

Only linear contraction
```

SELECT * FROM pgr_contraction(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[2]);
type | id | contracted_vertices | source | target | cost

| e | $\|-1\|\{3\}$ | $1 \mid$ | $7 \mid$ | 2 |
| :--- | :--- | :--- | :--- | :--- |
| e | $\|-2\|\{3\}$ |  | $7 \mid$ | $1 \mid$ |

(2 rows)

```

See Also
- Contraction - Family of functions

\section*{Indices and tables}
- Index
- Search Page

\section*{Introduction}

In large graphs, like the road graphs, or electric networks, graph contraction can be used to speed up some graph algorithms. Contraction reduces the size of the graph by removing some of the vertices and edges and, for example, might add edges that represent a sequence of original edges decreasing the total time and space used in graph algorithms.

This implementation gives a flexible framework for adding contraction algorithms in the future, currently, it supports two algorithms:
1. Dead end contraction
2. Linear contraction

Allowing the user to:
- Forbid contraction on a set of nodes.
- Decide the order of the contraction algorithms and set the maximum number of times they are to be executed.

\section*{Dead end contraction}

Contraction of the leaf nodes of the graph.

\section*{Dead end}

A node is considered a dead end node when
- On undirected graphs:
- The number of adjacent vertices is 1 .
- On directed graphs:
- The number of adjacent vertices is 1 .
- There are no outgoing edges and has at least one incoming edge.
- There are no incoming edges and has at least one outgoing edge.

When the conditions are true then the Operation: Dead End Contraction can be done.

\section*{Dead end vertex on undirected graph}
- The green nodes are dead end nodes
- The blue nodes have an unlimited number of edges.


Node Adjecent nodes Number of adjacent nodes
\begin{tabular}{lll}
\hline\(\backslash(a \backslash)\) & \(\backslash(\backslash\{u \backslash\} \backslash)\) & 1 \\
\hline\(\backslash(b \backslash)\) & \(\backslash(\backslash\{v \backslash\} \backslash)\) & 1 \\
\hline
\end{tabular}

\section*{Dead end vertex on directed graph}
- The green nodes are dead end nodes
- The blue nodes have an unlimited number of incoming and/or outgoing edges.


Node Adjecent nodes Number of adjacent nodes Number of incoming edges Number of outgoing edges
\begin{tabular}{lllll}
\(\backslash(a \backslash)\) & \(\backslash(\backslash\{u \backslash\} \backslash)\) & 1 & & \\
\hline\(\backslash(b \backslash)\) & \(\backslash(\backslash\{v \backslash\} \backslash)\) & 1 & 2 & 0 \\
\hline\(\backslash(c \backslash)\) & \(\backslash(\backslash\{v, w \backslash\} \backslash)\) & 2 & & \\
\hline\(\backslash(d \backslash)\) & \(\backslash(\backslash\{x \backslash\} \backslash)\) & 1 & 0 & 2 \\
\hline\(\backslash(e \backslash)\) & \(\backslash(\backslash\{x, y \backslash\} \backslash)\) & 2 & & \\
\hline
\end{tabular}

From above, nodes \(\backslash(\backslash\{a, b, d \backslash\} \backslash)\) are dead ends because the number of adjacent vertices is 1 . No further checks are needed for those nodes.

On the following table, nodes \(\backslash(\backslash\{c, e \backslash\} \backslash)\) because the even that the number of adjacent vertices is not 1 for
- \(\backslash(c \backslash)\)
- There are no outgoing edges and has at least one incoming edge.
- \(\backslash(e \backslash)\)
- There are no incoming edges and has at least one outgoing edge.

Operation: Dead End Contraction

The dead end contraction will stop until there are no more dead end nodes. For example from the following graph where( \(w \backslash\) ) is the dead end node:


After contracting \(\backslash(w \backslash)\), node \(\backslash(v \backslash)\) is now a dead end node and is contracted:


After contracting \(\backslash(v \backslash)\), stop. Node \(\backslash(u \backslash)\) has the information of nodes that were contrcted.

\section*{Rest of the Graph}

\section*{\(\mathrm{u}\{\mathrm{v}, \mathrm{w}\}\)}

Node \(\backslash(u \backslash)\) has the information of nodes that were contracted.

\section*{Linear contraction}

In the algorithm, linear contraction is represented by 2.

\section*{Linear}

In case of an undirected graph, a node is considered alinear node when
- The number of adjacent vertices is 2 .

In case of a directed graph, a node is considered alinear node when
- The number of adjacent vertices is 2 .
- Linearity is symmetrical

\section*{Linear vertex on undirected graph}
- The green nodes are linear nodes
- The blue nodes have an unlimited number of incoming and outgoing edges.

\section*{Undirected}

\begin{tabular}{lll}
\hline Node & Adjecent nodes & Number of adjacent nodes \\
\hline\(\backslash(v \backslash)\) & \(\backslash(\backslash\{u, w \backslash\} \backslash)\) & 2
\end{tabular}

\section*{Linear vertex on directed graph}
- The green nodes are linear nodes
- The blue nodes have an unlimited number of incoming and outgoing edges.
- The white node is not linear because the linearity is not symetrical.
- It is possible to go \\(y \rightarrow c \rightarrow \(\mathrm{z} \backslash\) )
- It's not possible to go\\(z \rightarrow c \rightarrow \(y \backslash\) )


Operation: Linear Contraction

The linear contraction will stop when there are no more linear nodes. For example from the following graph where( \(v \backslash\) ) and \(\backslash\) \((w \backslash)\) are linear nodes:

\section*{Rest of the Graph}


Contracting \(\backslash(w \backslash)\),
- The vertex \(\backslash(w \backslash)\) is removed from the graph
- The edges \(\backslash(v \backslash\) rightarrow \(w \backslash)\) and \(\backslash(w \backslash\) rightarrow \(z \backslash)\) are removed from the graph.
- A new edge \(\backslash(v\) \rightarrow \(z \backslash)\) is inserted represented with red color.


Contracting \(\backslash(\mathrm{v} \backslash)\) :
- The vertex \(\backslash(v \backslash)\) is removed from the graph
- The edges \(\backslash(u\) \rightarrow \(v \backslash)\) and \(\backslash(v \backslash\) rightarrow \(z \backslash)\) are removed from the graph.
- A new edge \(\backslash(u\) \rightarrow \(z \backslash)\) is inserted represented with red color.


Edge \(\backslash(u \backslash\) rightarrow \(z \backslash)\) has the information of nodes that were contracted.

The cycle

Contracting a graph, can be done with more than one operation. The order of the operations affect the resulting contracted graph, after applying one operation, the set of vertices that can be contracted by another operation changes.

This implementation, cycles max_cycles times through operations_order .
```

<input>
do max_cycles times {
for (operation in operations_order)
{ do operation }
}
<output>

```

\section*{Contracting sample data}

In this section, building and using a contracted graph will be shown by example.
- The Sample Data for an undirected graph is used
- Construction of the graph in the database
- Contraction results
- Add additional columns
- Store contraction information
- The vertex table update
- The edge table update
- The contracted graph
- Vertices that belong to the contracted graph.
- Edges that belong to the contracted graph.
- Contracted graph
- Using the contracted graph
- Case 1: Both source and target belong to the contracted graph.
- Case 2: Source and/or target belong to an edge subgraph.
- Case 3: Source and/or target belong to a vertex.

Construction of the graph in the database

\section*{Original Data}

The following query shows the original data involved in the contraction operation.

SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id;
id | source | target | cost | reverse_cost
\begin{tabular}{ccccc}
\(1 \mid\) & \(5 \mid\) & \(6 \mid\) & \(1 \mid\) & 1 \\
\(2 \mid\) & \(6 \mid\) & \(10 \mid\) & \(-1 \mid\) & 1 \\
\(3 \mid\) & \(10 \mid\) & \(15 \mid\) & \(-1 \mid\) & 1 \\
\(4 \mid\) & \(6 \mid\) & \(7 \mid\) & \(1 \mid\) & 1 \\
\(5 \mid\) & \(10 \mid\) & \(11 \mid\) & \(1 \mid\) & -1 \\
\(6 \mid\) & \(1 \mid\) & \(3 \mid\) & \(1 \mid\) & 1 \\
\(7 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) & 1 \\
\(8 \mid\) & \(7 \mid\) & \(11 \mid\) & \(1 \mid\) & 1 \\
\(9 \mid\) & \(11 \mid\) & \(16 \mid\) & \(1 \mid\) & 1 \\
\(10 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 1 \\
\(11 \mid\) & \(11 \mid\) & \(12 \mid\) & \(1 \mid\) & -1 \\
\(12 \mid\) & \(8 \mid\) & \(12 \mid\) & \(1 \mid\) & -1 \\
\(13 \mid\) & \(12 \mid\) & \(17 \mid\) & \(1 \mid\) & -1 \\
\(14 \mid\) & \(8 \mid\) & \(9 \mid\) & \(1 \mid\) & 1 \\
\(15 \mid\) & \(16 \mid\) & \(17 \mid\) & \(1 \mid\) & 1 \\
\(16 \mid\) & \(15 \mid\) & \(16 \mid\) & \(1 \mid\) & 1 \\
\(17 \mid\) & \(2 \mid\) & \(4 \mid\) & \(1 \mid\) & 1 \\
\(18 \mid\) & \(13 \mid\) & \(14 \mid\) & \(1 \mid\) & 1 \\
\((18\) rows) & & &
\end{tabular}

The original graph:


\section*{Contraction results}

The results do not represent the contracted graph. They represent the changes done to the graph after applying the contraction algorithm.

Observe that vertices, for example, \(\backslash(6 \backslash)\) do not appear in the results because it was not affected by the contraction algorithm.
\begin{tabular}{lll|r|r}
v & \(|4|\{2\}\) & \(-1 \mid\) & \(-1 \mid\) & -1 \\
v & \(7 \mid\{1,3\}\) & \(\mid\) & \(-1 \mid\) & \(-1 \mid\) \\
v & \(14 \mid\{13\}\) & -1 \\
e & \(|-1|\{5,6\}\) & \(-1 \mid\) & \(-1 \mid\) & -1 \\
e & \(|-2|\{8,9\}\) & \(7 \mid\) & \(10 \mid\) & 2 \\
e & \(|-3|\{17\}\) & \(7 \mid\) & \(12 \mid\) & 2 \\
e & \(|-4|\{15\}\) & \(12 \mid\) & \(16 \mid\) & 2 \\
\hline
\end{tabular}

After doing the dead end contraction operation:


After doing the linear contraction operation to the graph above:


The process to create the contraction graph on the database:

Add additional columns

Adding extra columns to the edge_table and edge_table_vertices_pgr tables, where:
\begin{tabular}{ll} 
Column & Description \\
\hline contracted_vertices & The vertices set belonging to the vertex/edge \\
\hline is_contracted & On the vertex table \\
& \begin{tabular}{l} 
O when true the vertex is contracted, its not part of the contracted graph. \\
is_new
\end{tabular} \\
& \begin{tabular}{l} 
when false the vertex is not contracted, its part of the contracted graph. \\
\\
\\
graph. \\
when false the edge is an original edge, might be or not part of the contracted graph.
\end{tabular} \\
\hline
\end{tabular}
```

ALTER TABLE vertices ADD is_contracted BOOLEAN DEFAULT false
ALTER TABLE
ALTER TABLE vertices ADD contracted vertices BIGINT[];
ALTER TABLE
ALTER TABLE edges ADD is_new BOOLEAN DEFAULT false;
ALTER TABLE
ALTER TABLE edges ADD contracted vertices BIGINT[];
ALTER TABLE

```

Store contraction information

Store the contraction results in a table

\section*{SELECT * INTO contraction_results \\ FROM pgr_contraction(}
'SELECT id, source, target, cost, reverse_cost FROM edges',
array[1, 2], directed => false)
SELECT 7

\section*{The vertex table update}

Use is_contracted column to indicate the vertices that are contracted.
```

UPDATE vertices
SET is contracted = true
WHERE id IN (SELECT unnest(contracted_vertices) FROM contraction_results);
UPDATE 10

```

Fill contracted_vertices with the information from the results tha belong to the vertices.
```

UPDATE vertices

```
SET contracted_vertices \(=\) contraction_results.\(c o n t r a c t e d \_v e r t i c e s\)
FROM contraction_results
WHERE type = 'v' AND vertices.id = contraction_results.id;
UPDATE 3

The modified vertices table:
```

SELECT id, contracted_vertices, is_contracted
FROM vertices
ORDER BY id;
id | contracted_vertices | is_contracted

| 11 | t |
| :---: | :---: |
| 21 | t |
| 3\| | \|t |
| $4 \mid\{2\}$ | \|f |
| 51 | \|t |
| 61 | \|t |
| $7 \mid\{1,3\}$ | \|f |
| 8\| | \|t |
| 9\| | \|t |
| 10\| | \|f |
| 11\| | \|f |
| 12 \| | \|f |
| $13 \mid$ | \| t |
| $14 \mid\{13\}$ | \|f |
| 15\| | \| t |
| 16 \| | \|f |
| 17\| | \| t |

(17 rows)

```

\section*{The edge table update}

\section*{Insert the new edges generated by pgr_contraction.}
```

INSERT INTO edges(source, target, cost, reverse_cost, contracted_vertices, is_new
SELECT source, target, cost, -1, contracted_vertices, true
FROM contraction results
WHERE type = 'e';
INSERT 0 4

```

The modified edge_table.

\section*{SELECT id, source, target, cost, reverse_cost, contracted_vertices, is_new}

FROM edges
ORDER BY id
id | source | target | cost | reverse_cost | contracted_vertices | is_new
\begin{tabular}{|c|c|c|c|c|c|}
\hline 1| & 51 & \(6 \mid\) & 1| & 1| & | f \\
\hline \(2 \mid\) & 61 & 10| & -1| & 1 | & | f \\
\hline 31 & 10| & 151 & -1| & \(1 \mid\) & |f \\
\hline 4| & 61 & 7 | & 1| & 1| & | f \\
\hline 51 & 10| & 11| & 1 | & -1| & |f \\
\hline 61 & 1 | & \(3 \mid\) & 1| & 1| & |f \\
\hline 71 & 31 & 7| & 1| & 1| & |f \\
\hline 8। & 71 & 11| & 1| & 1| & | f \\
\hline 91 & 11| & \(16 \mid\) & 1| & \(1 \mid\) & |f \\
\hline 10| & 7| & 8| & 1| & 1| & |f \\
\hline 11| & 11| & 12| & 1| & -1| & |f \\
\hline 12| & 8| & 12| & 1| & -1| & |f \\
\hline 13| & \(12 \mid\) & 17| & 1| & -1| & |f \\
\hline 14| & 8| & 9| & 1| & 1| & f \\
\hline 151 & \(16 \mid\) & \(17 \mid\) & 1| & 1| & |f \\
\hline \(16 \mid\) & 15 | & \(16 \mid\) & 1| & 1 | & |f \\
\hline 17| & \(2 \mid\) & \(4 \mid\) & 1| & 1| & f \\
\hline 18| & 13 | & \(14 \mid\) & 1| & 1| & | f \\
\hline 19| & 7| & 10| & 21 & \(-1 \mid\{5,6\}\) & 1 t \\
\hline 201 & 71 & \(12 \mid\) & 21 & \(-1 \mid\{8,9\}\) & | t \\
\hline 21| & 12 | & \(16 \mid\) & 21 & \(-1 \mid\{17\}\) & | t \\
\hline \(22 \mid\) & 10| & \(16 \mid\) & 21 & \(-1 \mid\{15\}\) & | t \\
\hline \multicolumn{6}{|l|}{(22 rows)} \\
\hline
\end{tabular}

\section*{The contracted graph}

Vertices that belong to the contracted graph.
```

SELECT id
FROM vertices
WHERE is_contracted = false
ORDER BY id;
id
4
10
12
14
(7 rows

```

Edges that belong to the contracted graph.
```

WITH
vertices_in_graph AS (
SELECT id
FROM vertices
WHERE is_contracted = false
)
SELECT id, source, target, cost, reverse_cost, contracted_vertices
FROM edges
WHERE source IN (SELECT * FROM vertices_in_graph)
AND target IN (SELECT * FROM vertices_in_graph)
ORDER BY id;
id | source | target | cost | reverse_cost | contracted_vertices

| 5\| | 10\| | 11\| | 1\| | -1\| |
| :---: | :---: | :---: | :---: | :---: |
| 8\| | 71 | 11\| | 1\| | 1\| |
| 91 | 11\| | $16 \mid$ | 1\| | $1 \mid$ |
| 11\| | 11\| | $12 \mid$ | 1 \| | -1\| |
| 19 \| | 71 | 10\| | 21 | $-1 \mid\{5,6\}$ |
| $20 \mid$ | 7\| | $12 \mid$ | 21 | $-1 \mid\{8,9\}$ |
| 21\| | $12 \mid$ | 16\| | 21 | $-1 \mid\{17\}$ |
| 221 | $10 \mid$ | 16 \| | 21 | $-1 \mid\{15\}$ |
| (8 rows) |  |  |  |  |

```


Using the contracted graph

\section*{Using the contracted graph with pgr_dijkstra}

There are three cases when calculating the shortest path between a given source and target in a contracted graph:
- Case 1: Both source and target belong to the contracted graph.
- Case 2: Source and/or target belong to an edge subgraph.
- Case 3: Source and/or target belong to a vertex.

Case 1: Both source and target belong to the contracted graph.

Using the Edges that belong to the contracted graph. on lines 10 to 19.
```

1 CREATE OR REPLACE FUNCTION my_dijkstra(
2 departure BIGINT, destination BIGINT,
OUT seq INTEGER, OUT path_seq INTEGER,
OUT node BIGINT, OUT edge BIGINT,
OUT cost FLOAT, OUT agg_cost FLOAT)
6 RETURNS SETOF RECORD AS
7 $BODY$
8 SELECT * FROM pgr_dijkstra(

$$
O WITH
vertices_in_graph AS (
    SELECT id
    FROM vertices
    WHERE is_contracted = false
15 )
SELECT id, source, target, cost, reverse_cost
FROM edges
WHERE source IN (SELECT * FROM vertices_in_graph)
AND target IN (SELECT * FROM vertices_in_graph)
$$,

1 departure, destination, false);
22 $BODY$
23 LANGUAGE SQL VOLATILE;
24 CREATE FUNCTION

```

\section*{Case 1}

When both source and target belong to the contracted graph, a path is found.
```

SELECT * FROM my_dijkstra(10, 12);
seq | path_seq | node | edge | cost | agg_cost

```
```

1| 1| 10| 5| 1| 0

```
1| 1| 10| 5| 1| 0
    2| 11| 11| 1| 1
    2| 11| 11| 1| 1
    3| 12| -1| 0| 2
    3| 12| -1| 0| 2
(3 rows)
```

(3 rows)

```

\section*{Case 2}

When source and/or target belong to an edge subgraph then a path is not found.
In this case, the contracted graph do not have an edge connecting with nodel(4\\).
```

SELECT * FROM my_dijkstra(15, 12);
seq | path_seq | node | edge | cost | agg_cost
(0 rows)

```

\section*{Case 3}

When source and/or target belong to a vertex then a path is not found.

In this case, the contracted graph do not have an edge connecting with nodel(7\\) and of node \(\backslash(4 \backslash)\) of the second case.
```

SELECT * FROM my dijkstra(15, 1);
seq | path_seq | node | edge | cost | agg_cost
(0 rows)

```

Case 2: Source and/or target belong to an edge subgraph.

Refining the above function to include nodes that belong to an edge.
- The vertices that need to be expanded are calculated on lines 10 to 16 .
- Adding to the contracted graph that additional section on lines 25 to 27.
```

1 CREATE OR REPLACE FUNCTION my_dijkstra(
2 departure BIGINT, destination BIGINT,
O OUT seq INTEGER, OUT path_seq INTEGER,
OUT node BIGINT, OUT edge BIGINT,
5 OUT cost FLOAT, OUT agg_cost FLOAT)
6 RETURNS SETOF RECORD AS
7 $BODY$
8 SELECT * FROM pgr dijkstra(

9 $$
WITH
edges to expand AS (
SELECT id
FROM edges
WHERE ARRAY[
$$ || departure || \$\$]::BIGINT[] <@ contracted_vertices

OR ARRAY[$$
|| destination ||
$$]::BIGINT[] <@ contracted_vertices
),
vertices_in_graph AS (
SELECT id
FROM vertices
WHERE is_contracted = false
UNION
SELECT unnest(contracted_vertices)
FROM edges
WHERE id IN (SELECT id FROM edges_to_expand)
8)
29
SELECT id, source, target, cost, reverse_cost
FROM edges
WHERE source IN (SELECT * FROM vertices_in_graph)
AND target IN (SELECT * FROM vertices_in_graph)
\$\$,

departure, destination, false);
36 $BODY$
37 LANGUAGE SQL VOLATILE;
38 CREATE FUNCTION

```

\section*{Case 1}

When both source and target belong to the contracted graph, a path is found.
```

SELECT * FROM my_dijkstra(10, 12);
seq | path_seq | node | edge | cost | agg_cost
1| 10| 5| 1| 0
2| 11| 11| 1| 1
3| 3| 12| -1| 0| 2
(3 rows)

```

\section*{Case 2}

When source and/or target belong to an edge subgraph, now, a path is found.
The routing graph now has an edge connecting with node\\(4)).
```

SELECT * FROM my_dijkstra(15, 12);
seq | path_seq | node | edge | cost | agg_cost
1| 15| 16| 1| 0
2| 16| 21| 2| 1
3| 3| 12|-1| 0| 3
(3 rows)

```

\section*{Case 3}

When source and/or target belong to a vertex then a path is not found.
In this case, the contracted graph do not have an edge connecting with nodel(7)).
```

SELECT * FROM my_dijkstra(15, 1)
seq | path_seq | node | edge | cost | agg_cost
(0 rows)

```

Case 3: Source and/or target belong to a vertex.

Refining the above function to include nodes that belong to an edge.
- The vertices that need to be expanded are calculated on lines 18 to 23 .
- Adding to the contracted graph that additional section on lines 38 to 40 .
```

1 CREATE OR REPLACE FUNCTION my dijkstra
2 departure BIGINT, destination BIGINT,
OUT seq INTEGER, OUT path_seq INTEGER,
4 OUT node BIGINT, OUT edge BIGINT,
5 OUT cost FLOAT, OUT agg_cost FLOAT)
6 RETURNS SETOF RECORD AS
7 \$BODY\$
8 SELECT * FROM pgr_dijkstra(
\$\$
WITH
edges_to_expand AS
SELECT id
FROM edges
WHERE ARRAY[\$\$ || departure || \$\$]::BIGINT[] <@ contracted_vertices
OR ARRAY[\$\$ || destination || \$\$]::BIGINT[] <@ contracted_vertices
),
vertices_to_expand AS (
SELECT id
FROM vertices
WHERE ARRAY[\$\$ || departure || \$\$]::BIGINT[] <@ contracted_vertices
OR ARRAY[\$\$ || destination || \$\$]::BIGINT[] <@ contracted_vertices
),
vertices_in_graph AS (
SELECT id
FROM vertices
WHERE is_contracted = false
UNION
SELECT unnest(contracted vertices)
FROM edges
WHERE id IN (SELECT id FROM edges_to_expand)
UNION
SELECT unnest(contracted_vertices
FROM vertices
WHERE id IN (SELECT id FROM vertices_to_expand)
)
SELECT id, source, target, cost, reverse_cost
FROM edges
WHERE source IN (SELECT * FROM vertices_in_graph
AND target IN (SELECT * FROM vertices_in_graph)
\$\$,
departure, destination, false);
9 \$BODY\$
50 LANGUAGE SQL VOLATILE;
51 CREATE FUNCTION

```

\section*{Case 1}

When both source and target belong to the contracted graph, a path is found.
```

SELECT * FROM my dijkstra(10, 12)
seq | path_seq | node | edge | cost | agg_cost
1| 10| 5| 1| 0
2| 11| 11| 1| 1
3| 3| 12|-1| 0| 2
(3 rows)

```

The code change do not affect this case so when source and/or target belong to an edge subgraph, a path is still found.

SELECT * FROM my_dijkstra(15, 12);
seq | path_seq | node | edge | cost | agg_cost
```

1| 15| 16| 1| 0
2| 16| 21| 2| 1
3| 12|-1| 0| 3

```
(3 rows)

\section*{Case 3}

When source and/or target belong to a vertex, now, a path is found.
Now, the routing graph has an edge connecting with node\\(7\\).
```

SELECT * FROM my_dijkstra(15, 1);
seq | path_seq | node | edge | cost | agg_cost
l15| 3| 1| 0
2| 10| 19| 2| 1
3| 7| 7| 1| 3
4| 3| 6| 1| 4
5| 5| 1| -1| 0| 5
(5 rows)

```

\section*{See Also}
```

pgr_contraction
Sample Data
https://www.cs.cmu.edu/afs/cs/academic/class/15210-f12/www/lectures/lecture16.pdf
https://algo2.iti.kit.edu/documents/routeplanning/geisberger_dipl.pdf

```

\section*{Indices and tables}

\section*{Index}
- Search Page

\section*{- Supported versions: Latest (3.3) 3.23 .13 .0}
- Unsupported versions: 2.6 2.5 2.42 .32 .2

\section*{Dijkstra - Family of functions}
- pgr_dijkstra - Dijkstra's algorithm for the shortest paths.
- pgr_dijkstraCost-Get the aggregate cost of the shortest paths.
- pgr_dijkstraCostMatrix - Use pgr_dijkstra to create a costs matrix.
- pgr_drivingDistance - Use pgr_dijkstra to calculate catchament information.
- pgr_KSP - Use Yen algorithm with pgr_dijkstra to get the K shortest paths.

\section*{Proposed}

\section*{Warning}

Proposed functions for next mayor release.
- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.

\footnotetext{
- pgr_dijkstraVia - Proposed - Get a route of a seuence of vertices.
- pgr_dijkstraNear - Proposed - Get the route to the nearest vertex.
- pgr_dijkstraNearCost - Proposed - Get the cost to the nearest vertex.
}

Supported versions: Latest (3.3) 3.23 .13 .0
-
pgr_dijkstra - Shortest path(s) using Dijkstra algorithm.

Boost Graph Inside

\section*{Availability}
- Version 3.1.0
- New Proposed functions: - pgr_dijkstra (Combinations)
- Version 3.0.0
- Official functions
- Version 2.2.0
- New proposed functions:
- pgr_dijkstra (One to Many)
- pgr_dijkstra (Many to One)
- pgr_dijkstra (Many to Many)
- Version 2.1.0
- Signature change on pgr_dijkstra (One to One)
- Version 2.0.0
- Official pgr_dijkstra (One to One)

\section*{Description}

Dijkstra's algorithm, conceived by Dutch computer scientist Edsger Dijkstra in 1956. It is a graph search algorithm that solves the shortest path problem for a graph with non-negative edge path costs, producing a shortest path from a starting vertex to an ending vertex. This implementation can be used with a directed graph and an undirected graph.
- Process is done only on edges with positive costs.
- A negative value on a cost column is interpreted as the edge does not exist.
- Values are returned when there is a path.
- When there is no path:
- When the starting vertex and ending vertex are the same.
- The aggregate cost of the non included values \(\backslash((v, v) \backslash)\) is \(\backslash(0 \backslash)\)
- When the starting vertex and ending vertex are the different and there is no path:
- The aggregate cost the non included values \(\backslash((u, v) \backslash)\) is \(\backslash(\backslash i n f t y \backslash)\)
- For optimization purposes, any duplicated value in the starting vertices or on the ending vertices are ignored.
- Running time: \(\backslash(\mathrm{O}(\mid\) start \(\backslash\) vids \(\mid *(\mathrm{~V} \backslash \log \mathrm{~V}+\mathrm{E})) \backslash)\)

\section*{Signatures}

\section*{Summary}
```

pgr_dijkstra(Edges SQL, start vid, end vid, [directed])
pgr_dijkstra(Edges SQL, start vid, end vids , [directed])
pgr_dijkstra(Edges SQL, start vids, end vid, [directed])
pgr_dijkstra(Edges SQL, start vids, end vids , [directed])
pgr_dijkstra(Edges SQL, Combinations SQL , [directed])
RETURNS SET OF (seq, path_seq, [start_vid], [end_vid], node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{One to One}
```

pgr_dijkstra(Edges SQL, start vid, end vid, [directed])
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertex \(\backslash(6 \backslash)\) to vertex \(\backslash(10 \backslash)\) on a directed graph
```

SELECT * FROM pgr_Dijkstra
'select id, source, target, cost, reverse_cost from edges',
6,10, true);
seq | path_seq | node | edge | cost | agg_cost
1| 6| 4| 1| 0
2| 7| 8| 1| 1
3| 11| 9| 1| 2
4| 16| 16| 1| 3
5| 15| 3| 1| 4
6| 10|-1| 0| 5
(6 rows)

```

\section*{One to Many}
```

pgr_dijkstra(Edges SQL, start vid, end vids , [directed])
RETURNS SET OF (seq, path seq, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertex \(\backslash(6 \backslash)\) to vertices \(\backslash(\backslash\{10,17 \backslash\} \backslash)\) on a directed
```

SELECT * FROM pgr_Dijkstra(
'select id, source, target, cost, reverse_cost from edges',
6, ARRAY[10, 17]);
seq | path_seq | end_vid | node | edge | cost | agg_cost
1| 1| 10| 6| 4| 1| 0
2| 2| 10| 7| 8| 1| 1
3| 3| 10| 11| 9| 1| 2
5| 4| 10| 16| 16| 1| 3
5| 5| 10| 15| 3| 1| 4
6| 6| 10| 10| -1| 0| 5
8| 2|
9| 3| 17| 11| 9| 1| 2
10| 4| 17| 16| 15| 1| 3
11| 5| 17| 17| -1| 0| 4
(11 rows)

```

\section*{Many to One}
```

pgr_dijkstra(Edges SQL, start vids, end vid , [directed])
RETURNS SET OF (seq, path seq, start vid, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{6,1 \backslash\} \backslash)\) to vertex \(\backslash(17 \backslash)\) on a directed graph
```

SELECT * FROM pgr_Dijkstra(
'select id, source, target, cost, reverse_cost from edges',
ARRAY[6, 1], 17);
seq | path_seq | start_vid | node | edge | cost | agg_cost
1| 1| 1| 1| 6| 1| 0
2| 2| 1| 3| 7| 1| 1
3| 1| 7| 8| 1| 2
4| 1| 11| 11| 1| 3
1| 12| 13| 1| 4
1| 17| -1| 0| 5
6| 6| 4| 1| 0
6| 7| 8| 1| 1
6 11 | 11 | 1 |
6| 12| 13| 1| 3
6| 17|-1| 0| 4
11| 5

```

\section*{Many to Many}

RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

\section*{Example:}

From vertices \(\backslash(\backslash\{6,1 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,17 \backslash\} \backslash)\) on an undirected graph
```

SELECT * FROM pgr_Dijkstra(
'select id, source, target, cost, reverse_cost from edges',
ARRAY[6, 1], ARRAY[10, 17],
directed => false);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $1 \mid$ | $10 \mid$ | $1 \mid$ | $6 \mid$ | $1 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $1 \mid$ | $10 \mid$ | $3 \mid$ | $7 \mid$ | $1 \mid$ | 1 |
| $3 \mid$ | $3 \mid$ | $1 \mid$ | $10 \mid$ | $7 \mid$ | $4 \mid$ | $1 \mid$ | 2 |
| $4 \mid$ | $4 \mid$ | $1 \mid$ | $10 \mid$ | $6 \mid$ | $2 \mid$ | $1 \mid$ | 3 |
| $5 \mid$ | $5 \mid$ | $1 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 4 |
| $6 \mid$ | $1 \mid$ | $1 \mid$ | $17 \mid$ | $1 \mid$ | $6 \mid$ | $1 \mid$ | 0 |
| $7 \mid$ | $2 \mid$ | $1 \mid$ | $17 \mid$ | $3 \mid$ | $7 \mid$ | $1 \mid$ | 1 |
| $8 \mid$ | $3 \mid$ | $1 \mid$ | $17 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 2 |
| $9 \mid$ | $4 \mid$ | $1 \mid$ | $17 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 3 |
| $10 \mid$ | $5 \mid$ | $1 \mid$ | $17 \mid$ | $16 \mid$ | $15 \mid$ | $1 \mid$ | 4 |
| $11 \mid$ | $6 \mid$ | $1 \mid$ | $17 \mid$ | $17 \mid$ | $-1 \mid$ | $0 \mid$ | 5 |
| $12 \mid$ | $1 \mid$ | $6 \mid$ | $10 \mid$ | $6 \mid$ | $2 \mid$ | $1 \mid$ | 0 |
| $13 \mid$ | $2 \mid$ | $6 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |
| $14 \mid$ | $1 \mid$ | $6 \mid$ | $17 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| $15 \mid$ | $2 \mid$ | $6 \mid$ | $17 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 1 |
| $16 \mid$ | $3 \mid$ | $6 \mid$ | $17 \mid$ | $11 \mid$ | $11 \mid$ | $1 \mid$ | 2 |
| $17 \mid$ | $4 \mid$ | $6 \mid$ | $17 \mid$ | $12 \mid$ | $13 \mid$ | $1 \mid$ | 3 |
| $18 \mid$ | $5 \mid$ | $6 \mid$ | $17 \mid$ | $17 \mid$ | $-1 \mid$ | $0 \mid$ | 4 |

```

\section*{Combinations}
```

pgr_dijkstra(Edges SQL, Combinations SQL , [directed])
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Using a combinations table on anundirected graph

The combinations table:
```

SELECT source, target FROM combinations;
source | target
5| 6
5| 10
6| 5
6| 15
6| 14
(5 rows)

```

The query
```

SELECT * FROM pgr_Dijkstra(
SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT source, target FROM combinations',
false);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $5 \mid$ | $6 \mid$ | $5 \mid$ | $1 \mid$ | $1 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $5 \mid$ | $6 \mid$ | $6 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |
| $3 \mid$ | $1 \mid$ | $5 \mid$ | $10 \mid$ | $5 \mid$ | $1 \mid$ | $1 \mid$ | 0 |
| $4 \mid$ | $2 \mid$ | $5 \mid$ | $10 \mid$ | $6 \mid$ | $2 \mid$ | $1 \mid$ | 1 |
| $5 \mid$ | $3 \mid$ | $5 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 2 |
| $6 \mid$ | $1 \mid$ | $6 \mid$ | $5 \mid$ | $6 \mid$ | $1 \mid$ | $1 \mid$ | 0 |
| $7 \mid$ | $2 \mid$ | $6 \mid$ | $5 \mid$ | $5 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |
| $8 \mid$ | $1 \mid$ | $6 \mid$ | $15 \mid$ | $6 \mid$ | $2 \mid$ | $1 \mid$ | 0 |
| $9 \mid$ | $2 \mid$ | $6 \mid$ | $15 \mid$ | $10 \mid$ | $3 \mid$ | $1 \mid$ | 1 |
| $10 \mid$ | $3 \mid$ | $6 \mid$ | $15 \mid$ | $15 \mid$ | $-1 \mid$ | $0 \mid$ | 2 |

(10 rows)

```

\section*{Parameters}
Column Type Description
Edges SQL TEXT Edges SQL as described below
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}

Optional parameters


Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL -1 & Weight of the edge (target, source) \\
& & \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& & \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & ANY- & Identifier of the departure vertex. \\
& INTEGER & \\
\hline target & \begin{tabular}{l} 
ANY- \\
\\
INTEGER
\end{tabular} & Identifier of the arrival vertex. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Result Columns}

Returns set of (seq, path_seq [, start_vid] [, end_vid], node, edge, cost, agg_cost)
\begin{tabular}{|c|c|c|}
\hline Column & Type & Description \\
\hline seq & INTEGER & Sequential value starting from 1. \\
\hline path_seq & INTEGER & Relative position in the path. Has value1 for the beginning of a path. \\
\hline \multirow[t]{2}{*}{start_vid} & BIGINT & Identifier of the starting vertex. Returned when multiple starting vetrices are in the query. \\
\hline & & \begin{tabular}{l}
Many to One \\
Many to Many
\end{tabular} \\
\hline end_vid & BIGINT & Identifier of the ending vertex. Returned when multiple ending vertices are in the query. \\
\hline & & \begin{tabular}{l}
One to Many \\
Many to Many
\end{tabular} \\
\hline node & BIGINT & Identifier of the node in the path fromstart_vid to end_vid. \\
\hline
\end{tabular}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline edge & BIGINT & \begin{tabular}{l} 
Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for \\
the last node of the path.
\end{tabular} \\
\hline cost & FLOAT & Cost to traverse from node using edge to the next node in the path sequence. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Demonstration of repeated values are ignored, and result is sorted.
```

SELECT * FROM pgr_Dijkstra
'select id, source, target, cost, reverse_cost from edges',
ARRAY[7, 10, 15, 10, 10, 15], ARRAY[10, 7, 10, 15]);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1\| | 1\| | 7\| | 10\| | 7\| | 8\| | 1\| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | 21 | 7\| | 10\| | 11\| | 91 | 1\| | 1 |
| 3\| | 31 | 71 | 10\| | $16 \mid$ | 16\| | 1\| | 2 |
| 4\| | 4 \| | 71 | 10\| | 15 | 31 | 1\| | 3 |
| 5\| | 51 | 71 | 10\| | 10\| | -1\| | 01 | 4 |
| 61 | 1 \| | 71 | 15\| | 71 | 81 | 1\| | 0 |
| 7\| | 21 | 71 | 15\| | 11\| | 91 | $1 \mid$ |  |
| 8\| | 31 | 71 | 15\| | 16\| | 16\| | $1 \mid$ | 2 |
| 9\| | 4 \| | 71 | 15\| | 15\| | -1\| | 01 | 3 |
| 10\| | 1\| | 10\| | 71 | 10\| | 5। | $1 \mid$ | 0 |
| 11\| | 21 | 10\| | 71 | 11\| | 8। | $1 \mid$ | 1 |
| 12\| | 31 | 10\| | 71 | 71 | -1\| | 01 | 2 |
| 13\| | 1 \| | 10\| | $15 \mid$ | 10\| | 51 | 1\| | 0 |
| 14\| | 21 | 10\| | $15 \mid$ | 11\| | 91 | 1\| | 1 |
| 15\| | 31 | 10\| | $15 \mid$ | 16\| | $16 \mid$ | 1\| | 2 |
| 16\| | 4 \| | 10\| | 15 \| | 15\| | \| 1 | | 01 | 3 |
| 17\| | 1\| | 15\| | 7\| | 15\| | $16 \mid$ | 1\| | 0 |
| 18\| | 21 | 15\| | 71 | $16 \mid$ | 91 | 11 | 1 |
| 19\| | 31 | 15\| | 71 | 11\| | 8। | 1\| | 2 |
| 201 | 4 \| | 15\| | 71 | 71 | -1\| | 01 | 3 |
| 21\| | 1 \| | 15\| | 10\| | 15\| | \| 31 | 1\| | 0 |
| 221 | 21 | 15\| | $10 \mid$ | $10 \mid$ | -1\| | 01 | 1 |

```

\section*{Example 2:}

Making start_vids the same as end_vids
```

SELECT * FROM pgr_Dijkstra(
'select id, source, target, cost, reverse_cost from edges',
ARRAY[7, 10, 15], ARRAY[7, 10, 15]);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1\| | 1\| | 71 | 101 | 71 | 81 | 1\| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2 | $2 \mid$ | 71 | $10 \mid$ | 11\| | 91 | 1\| | 1 |
| 3 | 3 | 71 | 10\| | $16 \mid$ | 16\| | 1\| | 2 |
| 4 | 4 | 71 | $10 \mid$ | 15 | 3\| | 1\| | 3 |
| 5 | 5 \| | 71 | 10\| | 10\| | -1\| | 0 \| | 4 |
| 6 | 1\| | 71 | 15\| | 7\| | 8। | 1\| | 0 |
| 7\| | 2 \| | 71 | 15\| | 11\| | 9\| | 1 \| | 1 |
| 8 | 3\| | 71 | 15\| | 16\| | 16\| | 1 \| | 2 |
| 9 | 4\| | 71 | 15\| | 15\| | -1\| | 01 | 3 |
| $10 \mid$ | $1 \mid$ | $10 \mid$ | 71 | 10\| | 51 | 1 \| | 0 |
| 11\| | 21 | $10 \mid$ | 71 | 11\| | 8\| | 1 \| | 1 |
| $12 \mid$ | 31 | $10 \mid$ | 71 | 71 | -1\| | 01 | 2 |
| 131 | $1 \mid$ | $10 \mid$ | $15 \mid$ | $10 \mid$ | 51 | 1 \| | 0 |
| $14 \mid$ | 21 | $10 \mid$ | $15 \mid$ | 11\| | 91 | 1 \| | 1 |
| 151 | 31 | $10 \mid$ | $15 \mid$ | $16 \mid$ | 16\| | 1 \| | 2 |
| $16 \mid$ | 4 \| | $10 \mid$ | $15 \mid$ | 15 | -1\| | 01 | 3 |
| $17 \mid$ | $1 \mid$ | $15 \mid$ | 71 | 15\| | 16\| | $1 \mid$ | 0 |
| 18 \| | 21 | 15 \| | 71 | 16\| | 9\| | 1\| | 1 |
| 19\| | 31 | 15\| | 71 | 11\| | 8। | 1\| | 2 |
| $20 \mid$ | 4 \| | 15\| | 71 | 7 \| | -1\| | 0\| | 3 |
| 21\| | 1 \| | 15\| | 10\| | 15 | 31 | 1 \| | 0 |
| $22 \mid$ | $2 \mid$ | 15\| | 10\| | $10 \mid$ | -1\| | $0 \mid$ | 1 |
| (22 row |  |  |  |  |  |  |  |

```

\section*{Example:}

Manually assigned vertex combinations.

\section*{SELECT * FROM pgr_Dijkstra(}
'SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT * FROM (VALUES \((6,10),(6,7),(12,10))\) AS combinations (source, target)');
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost


The examples of this section are based on theSample Data network.
- For directed graphs with cost and reverse_cost columns
- 1) Path from \(\backslash(6 \backslash)\) to \(\backslash(10 \backslash)\)
- 2) Path from \(\backslash(6 \backslash)\) to \(\backslash(7 \backslash)\)
- 3) Path from \(\backslash(12 \backslash)\) to \(\backslash(10 \backslash)\)
- 4) Path from \(\backslash(12 \backslash)\) to \(\backslash(7 \backslash)\)
- 5) Using One to Many to get the solution of examples 1 and 2
- 6) Using Many to One to get the solution of examples 2 and 4
- 7) Using Many to Many to get the solution of examples 1 to \(\mathbf{4}\)
- 8) Using Combinations to get the solution of examples \(\mathbf{1}\) to \(\mathbf{3}\)
- For undirected graphs with cost and reverse_cost columns
- 9) Path from \(\backslash(6 \backslash)\) to \(\backslash(10 \backslash)\)
- 10) Path from \(\backslash(6 \backslash)\) to \(\backslash(7 \backslash)\)
- 11) Path from \(\backslash(12 \backslash)\) to \(\backslash(10 \backslash)\)
- 12) Path from \(\backslash(12 \backslash)\) to \(\backslash(7 \backslash)\)
- 13) Using One to Many to get the solution of examples 9 and 10
- 14) Using Many to One to get the solution of examples \(\mathbf{1 0}\) and \(\mathbf{1 2}\)
- 15) Using Many to Many to get the solution of examples 9 to 12
- 16) Using Combinations to get the solution of examples 9 to 11
- For directed graphs only with cost column
- 17) Path from \(\backslash(6 \backslash)\) to \(\backslash(10 \backslash)\)
- 18) Path from \(\backslash(6 \backslash)\) to \(\backslash(7 \backslash)\)
- 19) Path from \(\backslash(12 \backslash)\) to \(\backslash(10 \backslash)\)
- 20) Path from \(\backslash(12 \backslash)\) to \(\backslash(7 \backslash)\)
- 21) Using One to Many to get the solution of examples 17 and 18
- 22) Using Many to One to get the solution of examples 18 and 20
- 23) Using Many to Many to get the solution of examples 17 to 20
- 24) Using Combinations to get the solution of examples 17 to 19
- For undirected graphs only with cost column
- 25) Path from \\(6) to \\(10\\)
- 26) Path from \(\backslash(6 \backslash)\) to \(\backslash(7 \backslash)\)
- 27) Path from \(\backslash(12 \backslash)\) to \(\backslash(10 \backslash)\)
- 28) Path from \(\backslash(12 \backslash)\) to \(\backslash(7 \backslash)\)
- 29) Using One to Many to get the solution of examples 25 and 26
- 30) Using Many to One to get the solution of examples 26 and 28
- 31) Using Many to Many to get the solution of examples 25 to 28
- 32) Using Combinations to get the solution of examples \(\mathbf{2 5}\) to \(\mathbf{2 7}\)
- Equalences between signatures
- 33) Using One to One
- 34) Using One to Many
- 35) Using Many to One
- 36) Using Many to Many
- 37) Using Combinations


Directed graph with cost and reverse cost columns
1) Path from \(\backslash(6 \backslash)\) to \(\backslash(10 \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6,10
);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 6| 4| 1| 0
2| 2| 7| 8| 1| 1
l-3| 11| 9| 1| 2
4| 4| 16| 16| 1| 3
5|
6| 6| 10|-1| 0| 5
(6 rows)

```
2) Path from \(\backslash(6 \backslash)\) to \(\backslash(7 \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6,7
);
seq| path seq| node | edge | cost | agg cost
1| 1| 6| 4| 1| 0
2| 2| 7| -1| 0| 1
(2 rows)

```
3) Path from \(\backslash(12 \backslash)\) to \(\backslash(10 \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
12,10
);
seq | path_seq | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $12 \mid$ | $13 \mid$ | $1 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $17 \mid$ | $15 \mid$ | $1 \mid$ | 1 |
| $3 \mid$ | $3 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 2 |
| $4 \mid$ | $4 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ | 3 |
| $5 \mid$ | $5 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 4 |

(5 rows)

```

\section*{SELECT * FROM pgr_dijkstra(}
'SELECT id, source, target, cost, reverse_cost FROM edges',
12, 7
);
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{rrrrr|r}
\(1 \mid\) & \(1 \mid\) & \(12 \mid\) & \(13 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(17 \mid\) & \(15 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(3 \mid\) & \(16 \mid\) & \(9 \mid\) & \(1 \mid\) & 2 \\
\(4 \mid\) & \(4 \mid\) & \(11 \mid\) & \(8 \mid\) & \(1 \mid\) & 3 \\
\(5 \mid\) & \(5 \mid\) & \(7 \mid\) & \(-1 \mid\) & \(0 \mid\) & 4 \\
(5 rows) & & & & &
\end{tabular}

\section*{5) Using One to Many to get the solution of examples 1 and 2}

Paths \\(\\{6\\}\rightarrow\\{10,7\\}\\)
```

SELECT * FROM pgr_dijkstra
'SELECT id, source, target, cost, reverse_cost FROM edges',
6, ARRAY[10, 7]
seq | path_seq | end_vid | node | edge | cost | agg_cos

| $1 \mid$ | $1 \mid$ | $7 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| ---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $7 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ |
| $3 \mid$ | $1 \mid$ | $10 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| $4 \mid$ | $2 \mid$ | $10 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ |
| $5 \mid$ | $3 \mid$ | $10 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ |
| $6 \mid$ | $4 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ |
| $7 \mid$ | $5 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ |
| $8 \mid$ | $6 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ |
| 8 rows) |  |  |  |  |  |

```

\section*{6) Using Many to One to get the solution of examples 2 and 4}

Paths \(\backslash(\backslash\{6,12 \backslash\} \backslash\) rightarrow \(\backslash\{7 \backslash\} \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[6, 12], }
);
seq | path_seq | start_vid | node | edge | cost | agg_cost
1| 1.-----+--------+-----+-----+-----+-------
3| 1| 12| 12| 13| 1| 0
4| 2| 12| 17| 15| 1| 1
5| 3| 12| 16| 9| 1| 2

```

```

(7 rows)

```

\section*{7) Using Many to Many to get the solution of examples \(\mathbf{1}\) to \(\mathbf{4}\)}

Paths \(\backslash(\backslash\{6,12 \backslash\} \backslash\) rightarrow \(\backslash\{10,7 \backslash\} \backslash)\)
```

SELECT * FROM pgr dijkstra
SELECT id, source, target, cost, reverse cost FROM edges',
ARRAY[6, 12], ARRAY[10,7]
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

```


Paths \(\backslash(\backslash\{6 \backslash\} \backslash\) rightarrow \(\backslash\{10,7 \backslash\} \backslash c u p \backslash\{12 \backslash\} \backslash\) rightarrow \(\backslash\{10 \backslash\} \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)'
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

```


\section*{For undirected graphs with cost and reverse_cost columns}


Undirected graph with cost and reverse cost columns

\section*{9) Path from \(\backslash(6 \backslash)\) to \(\backslash(10 \backslash)\)}
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6,10,
false
);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 6| 2| 1| 0
(2 rows)

```
10) Path from \(\backslash(6 \backslash)\) to \(\backslash(7 \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6,7,
fals
);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 6| 4| 1| 0
2| 2| 7| -1| 0| 1
(2 rows)

```
11) Path from \(\backslash(12 \backslash)\) to \(\backslash(10 \backslash)\)
```

SELECT * FROM pgr_dijkstra(
SELECT id, source, target, cost, reverse_cost FROM edges,
12, 10,
false
);
seq | path_seq | node | edge | cost | agg_cost
1|---------------------------------
2| 2| 11| 5| 1| 1
3| 3| 10| -1| 0| 2
(3 rows)

```
12) Path from \(\backslash(12 \backslash)\) to \(\backslash(7 \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges'
12, 7,
false
);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 12| 12| 1| 0
2| 2| 8| 10| 1| 1
3| 3| 7| -1| 0| 2
(3 rows)

```
13) Using One to Many to get the solution of examples 9 and 10

Paths \(\backslash(\backslash\{6 \backslash\} \backslash\) rightarrow \(\backslash\{10,7 \backslash\} \backslash)\)
```

SELECT * FROM pgr_dijkstra(
SELECT id, source, target, cost, reverse cost FROM edges',
6, ARRAY[10,7]
false
);
seq | path_seq | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $7 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $7 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ |
| $3 \mid$ | $1 \mid$ | $10 \mid$ | $6 \mid$ | $2 \mid$ | $1 \mid$ |
| $4 \mid$ | $2 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ |
| (4 rows) |  |  |  |  |  |

```

\section*{14) Using Many to One to get the solution of examples 10 and 12}

Paths \(\backslash(\backslash\{6,12 \backslash\} \backslash\) rightarrow \(\backslash\{7 \backslash\} \backslash)\)
```

SELECT * FROM pgr_dijkstra(
SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[6,12], 7,
false
);
seq | path_seq | start_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $6 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $6 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ |
| $3 \mid$ | $1 \mid$ | $12 \mid$ | $12 \mid$ | $12 \mid$ | $1 \mid$ |
| $4 \mid$ | $2 \mid$ | $12 \mid$ | $8 \mid$ | $10 \mid$ | $1 \mid$ |
| $5 \mid$ | $3 \mid$ | $12 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ |
| (5 rows) |  |  |  |  |  |

```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[6, 12], ARRAY[10,7],
false
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
\begin{tabular}{cccccccc}
\(1 \mid\) & \(1 \mid\) & \(6 \mid\) & \(7 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(6 \mid\) & \(7 \mid\) & \(7 \mid\) & \(-1 \mid\) & \(0 \mid\) & 1 \\
\(3 \mid\) & \(1 \mid\) & \(6 \mid\) & \(10 \mid\) & \(6 \mid\) & \(2 \mid\) & \(1 \mid\) & 0 \\
\(4 \mid\) & \(2 \mid\) & \(6 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 1 \\
\(5 \mid\) & \(1 \mid\) & \(12 \mid\) & \(7 \mid\) & \(12 \mid\) & \(12 \mid\) & \(1 \mid\) & 0 \\
\(6 \mid\) & \(2 \mid\) & \(12 \mid\) & \(7 \mid\) & \(8 \mid\) & \(10 \mid\) & \(1 \mid\) & 1 \\
\(7 \mid\) & \(3 \mid\) & \(12 \mid\) & \(7 \mid\) & \(7 \mid\) & \(-1 \mid\) & \(0 \mid\) & 2 \\
\(8 \mid\) & \(1 \mid\) & \(12 \mid\) & \(10 \mid\) & \(12 \mid\) & \(11 \mid\) & \(1 \mid\) & 0 \\
\(9 \mid\) & \(2 \mid\) & \(12 \mid\) & \(10 \mid\) & \(11 \mid\) & \(5 \mid\) & \(1 \mid\) & 1 \\
\(10 \mid\) & \(3 \mid\) & \(12 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 2 \\
\((10\) rows \()\)
\end{tabular}
16) Using Combinations to get the solution of examples 9 to 11

Paths \(\backslash(\backslash\{6 \backslash\} \backslash\) rightarrow \(\backslash\{10,7 \backslash\} \backslash c u p \backslash\{12 \backslash\} \backslash\) rightarrow \(\backslash\{10 \backslash\} \backslash)\)
```

SELECT * FROM pgr_dijkstra(
SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)',
false
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
1| 1| 6| 7|------------------------------+-----+------+
2|
3| 1| 6| 10| 6| 2| 1| 0
4| 2| 6| 10| 10| -1| 0| 1
5| 1| 12| 10| 12| 11| 1| 0
6|
7| 3| 12| 10| 10| -1| 0| 2
(7 rows)

```

For directed graphs only with cost column


Directed graph only with cost column
17) Path from \(\backslash(6 \backslash)\) to \(\backslash(10 \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost FROM edges',
6,10
);
seq | path_seq | node | edge | cost | agg_cost
(0 rows)

```
18) Path from \(\backslash(6 \backslash)\) to \(\backslash(7 \backslash)\)
```

SELECT * FROM pgr_dijkstra
SELECT id, source, target, cost FROM edges',
6,7
);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 6| 4| 1| 0
2| 2| 7| -1| 0| 1
(2 rows)

```
19) Path from \(\backslash(12 \backslash)\) to \(\backslash(10 \backslash)\)
```

SELECT * FROM pgr_dijkstra
'SELECT id, source, target, cost FROM edges',
12,10
seq | path_seq | node | edge | cost | agg_cost
(0 rows)

```
20) Path from \(\backslash(12 \backslash)\) to \(\backslash(7 \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost FROM edges',
12,7
);
seq | path_seq | node | edge | cost | agg_cost
(0 rows)

```

\section*{21) Using One to Many to get the solution of examples 17 and 18}

Paths \(\backslash(\backslash\{6 \backslash\} \backslash\) rightarrow \(\backslash\{10,7 \backslash\} \backslash)\)
```

SELECT * FROM pgr dijkstra(
'SELECT id, source, target, cost FROM edges',
6, ARRAY[10,7]
);
seq | path_seq | end_vid | node | edge | cost | agg_cost
1| 1| 7| 6| 4| 1| 0
2| 2| 7| 7| -1| 0|
(2 rows)

```

\section*{22) Using Many to One to get the solution of examples 18 and 20}

Paths \(\backslash(\backslash\{6,12 \backslash\} \backslash\) rightarrow \(\backslash\{7 \backslash\} \backslash)\)
```

SELECT * FROM pgr dijkstra
'SELECT id, source, target, cost FROM edges'
ARRAY[6,12], 7
);
seq| path seq| start vid | node | edge | cost | agg cost
1| 1| 6| 6| 4| 1| 0
2| 2| 6| 7|-1| 0| 1
(2 rows)

```
\begin{tabular}{|c|c|c|c|c|c|c|c|}
\hline 1 & 1 & 6 & 71 & 6 & 4 & 1| & 0 \\
\hline 2 & 2 & 6 & 7 & 71 & -1 & 0 & \\
\hline
\end{tabular}
(2 rows)

\section*{24) Using Combinations to get the solution of examples \(\mathbf{1 7}\) to 19}

Paths \(\backslash(\backslash\{6 \backslash\} \backslash\) rightarrow \(\backslash\{10,7 \backslash\} \backslash c u p \backslash\{12 \backslash\} \backslash\) rightarrow \(\backslash\{10 \backslash\} \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost FROM edges',
SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)'
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1\| | 1 \| | 61 | 71 | 6 | 4 \| |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | 21 | 61 | 71 | 71 | -1\| | 0 |  |

(2 rows)

```

For undirected graphs only with cost column


Undirected graph only with cost column
25) Path from \(\backslash(6 \backslash)\) to \(\backslash(10 \backslash)\)
```

SELECT * FROM pgr_dijkstra(
SELECT id, source, target, cost FROM edges',
6, 10,
false
);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 6| 4| 1| 0
2| 7| 8| 1| 1
3| 11| 5| 1| 2
4| 10| -1| 0| 3
(4 rows)

```
```

SELECT * FROM pgr_dijkstra

```
    'SELECT id, source, target, cost FROM edges',
6, 7,
fals
);
seq | path_seq | node | edge | cost | agg_cost
1| \(1|6| 4|1| 0\)
\(2|\quad 2| 7|-1| 0 \mid \quad 1\)
(2 rows)
27) Path from \(\backslash(12 \backslash)\) to \(\backslash(10 \backslash)\)
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost FROM edges',
12, 10,
false
);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 12| 11| 1|--------------------------
2| 2| 11| 5| 1| 1
3| 3| 10| -1| 0| 2
(3 rows)

```

\section*{28) Path from \(\backslash(12 \backslash)\) to \(\backslash(7 \backslash)\)}
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost FROM edges'
12, 7,
false
);
seq | path_seq | node | edge | cost | agg_cost

```

```

$2|2| 8|10| 1 \mid 1$
$3|\quad 3| 7|-1| 0 \mid \quad 2$
(3 rows)

```
29) Using One to Many to get the solution of examples 25 and 26

Paths \(\backslash(\backslash\{6 \backslash\} \backslash\) rightarrow \(\backslash\{10,7 \backslash\} \backslash)\)
```

SELECT * FROM pgr_dijkstra
SELECT id, source, target, cost FROM edges',
6, ARRAY[10,7]
false
);
seq | path_seq | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $7 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $7 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |
| $3 \mid$ | $1 \mid$ | $10 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| $4 \mid$ | $2 \mid$ | $10 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 1 |
| $5 \mid$ | $3 \mid$ | $10 \mid$ | $11 \mid$ | $5 \mid$ | $1 \mid$ | 2 |
| $6 \mid$ | $4 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 3 |

(6 rows)

```

\section*{30) Using Many to One to get the solution of examples 26 and 28}

Paths \(\backslash(\backslash\{6,12 \backslash\} \backslash\) rightarrow \(\backslash\{7 \backslash\} \backslash)\)
```

SELECT * FROM pgr dijkstra(
'SELECT id, source, target, cost FROM edges',
ARRAY[6,12], 7,
false
);
seq | path_seq | start_vid | node | edge | cost | agg_cost
1| 1| 6| 6| 4| 1| 0
2| 6| 7| -1| 0| 1
1| 12| 12| 12| 1| 0
2| 12| 8| 10| 1| 1
3| 12| 7| -1| 0| 2
(5 rows)

```
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost FROM edges',
ARRAY[6, 12], ARRAY[10,7],
false
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $6 \mid$ | $7 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $6 \mid$ | $7 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |
| $3 \mid$ | $1 \mid$ | $6 \mid$ | $10 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| $4 \mid$ | $2 \mid$ | $6 \mid$ | $10 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 1 |
| $5 \mid$ | $3 \mid$ | $6 \mid$ | $10 \mid$ | $11 \mid$ | $5 \mid$ | $1 \mid$ | 2 |
| $6 \mid$ | $4 \mid$ | $6 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 3 |
| $7 \mid$ | $1 \mid$ | $12 \mid$ | $7 \mid$ | $12 \mid$ | $12 \mid$ | $1 \mid$ | 0 |
| $8 \mid$ | $2 \mid$ | $12 \mid$ | $7 \mid$ | $8 \mid$ | $10 \mid$ | $1 \mid$ | 1 |
| $9 \mid$ | $3 \mid$ | $12 \mid$ | $7 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ | 2 |
| $10 \mid$ | $1 \mid$ | $12 \mid$ | $10 \mid$ | $12 \mid$ | $11 \mid$ | $1 \mid$ | 0 |
| $11 \mid$ | $2 \mid$ | $12 \mid$ | $10 \mid$ | $11 \mid$ | $5 \mid$ | $1 \mid$ | 1 |
| $12 \mid$ | $3 \mid$ | $12 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 2 |
| $(12$ rows) |  |  |  |  |  |  |  |

```

\section*{32) Using Combinations to get the solution of examples 25 to 27}

\section*{Paths \\(\\{6\\}\rightarrow\\{10,7\\}\cup\\{12\\}\rightarrow\\{10\\}\\)}
```

SELECT * FROM pgr_dijkstra
'SELECT id, source, target, cost FROM edges',
'SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)',
false
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1\| | 1\| | 61 | 71 | 61 | 4 \| | 1\| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2 \| | $2 \mid$ | 61 | 71 | 71 | -1\| | 01 | 1 |
| 31 | 1\| | 61 | 10\| | 6\| | 4\| | 1\| | 0 |
| 4 \| | 21 | 61 | 10\| | 7\| | 8\| | 1\| | 1 |
| 51 | 31 | $6 \mid$ | 10\| | 11\| | 51 | 1\| | 2 |
| 61 | 4 \| | $6 \mid$ | 10\| | $10 \mid$ | -1\| | 01 | 3 |
| 71 | 1\| | $12 \mid$ | $10 \mid$ | $12 \mid$ | 11\| | 1 \| | 0 |
| 81 | 21 | $12 \mid$ | $10 \mid$ | 11\| | 5\| | 1\| | 1 |
| 91 | 31 | $12 \mid$ | $10 \mid$ | $10 \mid$ | -1\| | 01 | 2 |

```

\section*{quvalences between signature}

The following examples find the path for \(\backslash(\backslash\{6 \backslash\} \mid\) rightarrow \(\backslash\{10 \backslash\} \backslash)\)

\section*{33) Using One to One}
```

SELECT * FROM pgr_dijkstra(
SELECT id, source, target, cost, reverse cost FROM edges',
6,10
);
seq | path_seq| node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ |
| $3 \mid$ | $3 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ |
| $4 \mid$ | $4 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ |
| $5 \mid$ | $5 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ |
| $6 \mid$ | $6 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ |
| $(6$ rows $)$ |  |  |  |  |

```

\section*{34) Using One to Many}
```

SELECT * FROM pgr_dijkstra
'SELECT id, source, target, cost, reverse_cost FROM edges',
6, ARRAY[10]
seq | path_seq | end_vid | node | edge | cost | agg_cos

| $1 \mid$ | $1 \mid$ | $10 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $10 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 1 |
| $3 \mid$ | $3 \mid$ | $10 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 2 |
| $4 \mid$ | $4 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 3 |
| $5 \mid$ | $5 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ | 4 |
| $6 \mid$ | $6 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 5 |

(6 rows)

```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[6], 10
);
seq | path_seq | start_vid | node | edge | cost | agg_cost
\begin{tabular}{rrrrcc}
\(1 \mid\) & \(1 \mid\) & \(6 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(2 \mid\) & \(6 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(3 \mid\) & \(6 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(4 \mid\) & \(6 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) \\
\(5 \mid\) & \(5 \mid\) & \(6 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(6 \mid\) & \(6 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\hline 6 lows) & & & & &
\end{tabular}
36) Using Many to Many
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[6], ARRAY[10]
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
1| 1| 6| 10| 6| 4| 1| 0
2| 2| 6| 10| 7| 8| 1| 1

```

```

5|
5|
(6 rows)

```

\section*{37) Using Combinations}
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT * FROM (VALUES(6, 10)) AS combinations (source, target)'
);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1\| | 1\| | 61 | $10 \mid$ | 61 | 4 \| | 1\| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2 \| | $2 \mid$ | 61 | $10 \mid$ | 7\| | 8। | 1 \| | 1 |
| 3 | 31 | 61 | 10\| | 11\| | 9\| | 1\| | 2 |
| 41 | 41 | 61 | $10 \mid$ | $16 \mid$ | $16 \mid$ | 1 \| | 3 |
| 51 | 51 | 61 | $10 \mid$ | 15\| | 3\| | 1 | 4 |
| 6 \| | 61 | 61 | $10 \mid$ | $10 \mid$ | -1\| | $0 \mid$ | 5 |

(6 rows)

```

\section*{See Also}
- https://en.wikipedia.org/wiki/Dijkstra\%27s_algorithm
- The queries use the Sample Data network.

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.62 .52 .42 .32 .3
pgr_dijkstraCost
pgr_dijkstraCost - Total cost of the shortest path(s) using Dijkstra algorithm.
boost

Boost Graph Inside

\section*{Availability}
- Version 3.1.0
- New proposed signature:
- pgr_dijkstraCost (Combinations)
- Version 2.2.0
- New Official function

\section*{Description}

The pgr_dijkstraCost function sumarizes of the cost of the shortest path(s) using Dijkstra Algorithm.
Dijkstra's algorithm, conceived by Dutch computer scientist Edsger Dijkstra in 1956. It is a graph search algorithm that solves the shortest path problem for a graph with non-negative edge path costs, producing a shortest path from a starting vertex to an ending vertex. This implementation can be used with a directed graph and an undirected graph.
- Process is done only on edges with positive costs.
- A negative value on a cost column is interpreted as the edge does not exist.
- Values are returned when there is a path.
- When there is no path:
- When the starting vertex and ending vertex are the same.
- The aggregate cost of the non included values \(\backslash((v, v) \backslash)\) is \(\backslash(0 \backslash)\)
- When the starting vertex and ending vertex are the different and there is no path: - The aggregate cost the non included values \(\backslash((u, v) \backslash)\) is \(\backslash(\backslash i n f t y \backslash)\)
- For optimization purposes, any duplicated value in the starting vertices or on the ending vertices are ignored.
- Running time: \(\backslash\left(\mathrm{O}\left(\mid\right.\right.\) start \(\backslash\) vids \(\left.\left.\left.\right|^{*}(\mathrm{~V} \backslash \log \mathrm{~V}+\mathrm{E})\right) \backslash\right)\)
- It does not return a path.
- Returns the sum of the costs of the shortest path of each pair combination of nodes requested.
- Let be the case the values returned are stored in a table, so the unique index would be the pair:(start_vid, end_vid).
- Depending on the function and its parameters, the results can be symmetric.
- The aggregate cost of \(\backslash((u, v) \backslash)\) is the same as for \(\backslash((v, u) \backslash)\).
- Any duplicated value in the start or end vertex identifiers are ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending

\section*{Signatures}

\section*{Summary}
```

pgr_dijkstraCost(Edges SQL, start vid, end vid , [directed])
pgr_dijkstraCost(Edges SQL, start vid, end vids, [directed])
pgr_dijkstraCost(Edges SQL, start vids, end vid, [directed])
pgr_dijkstraCost(Edges SQL, start vids, end vids, [directed])
pgr_dijkstraCost(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{One to One}
```

pgr_dijkstraCost(Edges SQL, start vid, end vid , [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertex \(\backslash(6 \backslash)\) to vertex \(\backslash(10 \backslash)\) on a directed graph
```

SELECT * FROM pgr_dijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6,10, true);
start_vid | end_vid | agg_cost
6| 10| 5
(1 row)

```

RETURNS SET OF (start_vid, end_vid, agg_cost)

\section*{Example:}

From vertex \(\backslash(6 \backslash)\) to vertices \(\backslash(\backslash\{10,17 \backslash\} \backslash)\) on a directed graph
```

SELECT * FROM pgr_dijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6, ARRAY[10, 17]);
start_vid | end_vid | agg_cost
6| 10| 5
-6| 17| 4
(2 rows)

```

\section*{Many to One}
```

pgr_dijkstraCost(Edges SQL, start vids, end vid , [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{6,1 \backslash\} \backslash)\) to vertex \(\backslash(17 \backslash)\) on a directed graph
```

SELECT * FROM pgr_dijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[6, 1], 17);
start_vid | end_vid | agg_cost
1| 17| 5
6| 17| 4
(2 rows)

```

Many to Many
```

pgr_dijkstraCost(Edges SQL, start vids, end vids, [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{6,1 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,17 \backslash\} \backslash)\) on an undirected graph
```

SELECT * FROM pgr_dijkstraCost
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[6, 1], ARRAY[10, 17],
directed => false);
start_vid | end_vid | agg_cost
1| 10| 4
1| 17| 5
6 10| 1
6| 17 | 4
(4 rows)

```

\section*{Combinations}
```

pgr_dijkstraCost(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Using a combinations table on anundirected graph
The combinations table:

The query:
```

SELECT * FROM pgr_dijkstraCost(
SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT source, target FROM combinations',
false);
start_vid | end_vid | agg_cost
5| 6| 1
5| 10| 2
6| 5|
(4 rows)

```

\section*{Parameters}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}

Optional parameters
\begin{tabular}{lllllll} 
Column & Type & Default & Description \\
\hline directed & BOOLEAN true & \(\ddots\) & When true the graph is considered Directed & \\
& & & When false the graph is considered as \\
& & Undirected.
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (target, source) \\
& & \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Combinations SQL
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & ANY- & Identifier of the departure vertex. \\
& INTEGER & \\
\hline target & ANY- & Identifier of the arrival vertex. \\
& INTEGER & \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Result Columns}

Set of (start_vid, end_vid, agg_cost)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline start_vid & BIGINT & Identifier of the starting vertex. \\
\hline end_vid & BIGINT & Identifier of the ending vertex. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to end_vid. \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example 1:}

Demonstration of repeated values are ignored, and result is sorted.
```

SELECT * FROM pgr_dijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[7, 10, 15, 10, 10, 15], ARRAY[10, 7, 10, 15]);
start_vid | end_vid | agg_cost

```


\section*{Example 2:}

Making start_vids the same as end_vids
```

SELECT * FROM pgr_dijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[7, 10, 15], ARRAY[7, 10, 15]);
start_vid | end_vid | agg_cost
710| 4
7| 15| 3
10| 7| 2
10| 15| 3
15| 7| 3
(6 rows)

```

\section*{Example 3:}

Manually assigned vertex combinations.
```

SELECT * FROM pgr_dijkstraCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)');
start_vid | end_vid | agg_cost
6| 7|--------------
6| 10| 5
12| 10| 4
(3 rows)

```

\section*{See Also}

\section*{- Dijkstra - Family of functions}
- Sample Data
- https://en.wikipedia.org/wiki/Dijkstra\%27s_algorithm

\section*{Indices and tables}
- Index
- Search Page

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- Official function
- Version 2.3.0
- New proposed function

\section*{Description}

Using Dijkstra algorithm, calculate and return a cost matrix.
Dijkstra's algorithm, conceived by Dutch computer scientist Edsger Dijkstra in 1956. It is a graph search algorithm that solves the shortest path problem for a graph with non-negative edge path costs, producing a shortest path from a starting vertex to an ending vertex. This implementation can be used with a directed graph and an undirected graph.

The main Characteristics are:
- Can be used as input to pgr_TSP.
- Use directly when the resulting matrix is symmetric and there is nd(\infty\\) value.
- It will be the users responsibility to make the matrix symmetric.
- By using geometric or harmonic average of the non symmetric values.
- By using max or min the non symmetric values.
- By setting the upper triangle to be the mirror image of the lower triangle.
- By setting the lower triangle to be the mirror image of the upper triangle.
- It is also the users responsibility to fix an\\(\infty\\) value.
- Each function works as part of the family it belongs to.
- It does not return a path.
- Returns the sum of the costs of the shortest path for pair combination of nodes in the graph.
- Process is done only on edges with positive costs.
- Values are returned when there is a path.
- When the starting vertex and ending vertex are the same, there is no path.
- The aggregate cost in the non included values \((v, v)\) is 0 .
- When the starting vertex and ending vertex are the different and there is no path.
- The aggregate cost in the non included values \((u, v)\) is \(\backslash(\backslash i n f t y \backslash)\).
- Let be the case the values returned are stored in a table:
- The unique index would be the pair: (start_vid, end_vid).
- Depending on the function and its parameters, the results can be symmetric.
- The aggregate cost of \((u, v)\) is the same as for \((v, u)\).
- Any duplicated value in the start vids are ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending

Signatures

\section*{Summary}
```

pgr_dijkstraCostMatrix(Edges SQL, start vids, [directed])
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Symmetric cost matrix for vertices \(\backslash(\backslash\{5,6,10,15 \backslash\} \backslash)\) on an undirected graph
```

SELECT * FROM pgr dijkstraCostMatrix
'SELECT id, source, target, cost, reverse_cost FROM edges',
(SELECT array_agg(id)
FROM vertices
WHERE id IN (5, 6, 10, 15)),
false);
start_vid | end_vid | agg_cost
5| 6| 1
5| 10| 2
5| 15| 3
6| 5| 1
| 10| 1
| 15| 2
10| 5| 2
| 6| 1
| 15| 1
5| 5| 3
5| 6| 2
15| 10| 1
(12 rows)

```

Parameters
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline start vids & ARRAY[BIGINT] & \begin{tabular}{l} 
Array of identifiers of starting \\
vertices.
\end{tabular}
\end{tabular}

Optional parameters
\begin{tabular}{llllll} 
Column & Type & Default & Description \\
\hline directed & BOOLEAN true & \begin{tabular}{ll} 
& When true the graph is considered Directed \\
& \\
& \\
& \\
& Uhen halse tirected.
\end{tabular}
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& &
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Set of (start_vid, end_vid, agg_cost)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline start_vid & BIGINT & Identifier of the starting vertex. \\
\hline end_vid & BIGINT & Identifier of the ending vertex. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to end_vid. \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Use with pgr_TSP.
```

SELECT * FROM pgr_TSP(

$$
SELECT * FROM pgr_dijkstraCostMatrix(
    'SELECT id, source, target, cost, reverse_cost FROM edges',
    (SELECT array_agg(id)
    FROM vertices
    WHERE id IN (5, 6, 10, 15)),
    false)
$$);

NOTICE: pgr_TSP no longer solving with simulated annaeling
HINT: Ignoring annaeling parameters
seq | node | cost | agg_cost
1| 5| 0| 0
2| 6| 1| 1
3| 10| 1| 2
4| 15| 1| 3
5| 5| 3| 6
(5 rows)

```

\section*{See Also}
- Dijkstra - Family of functions
- Cost Matrix - Category
- Traveling Sales Person - Family of functions
- Sample Data

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1) 3.0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2 2.1 2.0

\section*{pgr_drivingDistance}
pgr_drivingDistance - Returns the driving distance from a start node.

Boost Graph Inside

\section*{Availability}
- Version 2.1.0:
- Signature change pgr_drivingDistance(single vertex)
- New Official pgr_drivingDistance(multiple vertices)
- Version 2.0.0:
- Official pgr_drivingDistance(single vertex)

\section*{Description}

Using the Dijkstra algorithm, extracts all the nodes that have costs less than or equal to the valutdistance. The edges extracted will conform to the corresponding spanning tree.

\section*{Signatures}
```

pgr_drivingDistance(Edges SQL, Root vid, distance, [directed])
pgr_drivingDistance(Edges SQL, Root vids, distance, [options])
options: [directed, equicost]
RETURNS SET OF (seq, [from_v,] node, edge, cost, agg_cost)

```
```

pgr_drivingDistance(Edges SQL, Root vid, distance, [directed])

```
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)

\section*{Example:}

From vertex \\(11\\) for a distance of \\(3.0\\)
```

SELECT * FROM pgr_drivingDistance(
'SELECT id, source, target, cost, reverse_cost FROM edges',
11, 3.0);
seq | node | edge | cost | agg_cost
1| 11| -1| 0| 0
2| 7| 8| 1| 1
3| 12| 11| 1| 1
| 16| 9| 1| 1
l
llllllll
7| 8| 10| 10| 2
8| 15| 16| 1| 2
9| 17| 15| 1| 2
10| 1| 6| 1| 3
11| 5| 1| 1| 3
12| 9| 14| 1| 3
13| 10| 3| 1| 3
(13 rows)

```

\section*{Multiple Vertices}
pgr_drivingDistance(Edges SQL, Root vids, distance, [options])
options: [directed, equicost]
RETURNS SET OF (seq, from_v, node, edge, cost, agg_cost)

\section*{Example:}

From vertices \(\backslash(\backslash\{11,16 \backslash\} \backslash)\) for a distance of \(\backslash(3.0 \backslash)\) with equi-cost on a directed graph
```

SELECT * FROM pgr_drivingDistance(
'SELECT id, source, target, cost, reverse_cost FROM edges',
array[11, 16], 3.0, equicost => true);
seq | from_v | node | edge | cost | agg_cost
11| 11| -1| 0) 0
11| 7| 8| 1| 1
11| 12| 11| 1| 1
11| 3| 7| 1| 2
11| 6| 4| 1| 2
11| 8| 10| 1| 2
11| 1| 6| 1| 3
11| 5| 1| 1| 3
11| 9| 14| 1| 3
16| 16| -1| 0| 0
16| 15| 16| 1| 1
16| 17| 15| 1| 1
| 16| 10| 3| 1| 2
(13 rows)

```

\section*{Parameters}
\begin{tabular}{|c|c|c|}
\hline Parameter & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below. \\
\hline Root vid & BIGINT & Identifier of the root vertex of the tree. \\
\hline Root vids & ARRAY[ANY-INTEGER] & \begin{tabular}{l}
Array of identifiers of the root vertices. \\
- \(\backslash(0 \backslash)\) values are ignored \\
- For optimization purposes, any duplicated value is ignored.
\end{tabular} \\
\hline distance & FLOAT & Upper limit for the inclusion of a node in the result. \\
\hline
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SmALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT


Driving distance optional parameters
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline equicost & BOOLEAN & true & \begin{tabular}{l} 
When true the node will only appear in the closestfrom_v list. \\
\end{tabular} \\
& & \begin{tabular}{l} 
When false which resembles several calls using the single starting point \\
signatures. Tie brakes are arbitrary.
\end{tabular} \\
& &
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & \\
& & Whent of the edge (target, source) \\
& & \\
& & not part of the graph.
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns SET OF (seq, from_v, node, edge, cost, agg_cost)
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline seq & BIGINT & Sequential value starting from \(\backslash(1 \backslash)\). \\
\hline [from_v] & BIGINT & Identifier of the root vertex. \\
\hline node & BIGINT & \begin{tabular}{l} 
Identifier of node within the limits from \\
from_v.
\end{tabular} \\
\hline edge & BIGINT & Identifier of the edge used to arrive to node. \\
& & \(\bullet \quad \backslash(0 \backslash)\) when node \(=\) from_v. \\
\hline cost & FLOAT & Cost to traverse edge. \\
\hline agg_cost & FLOAT & Aggregate cost from from_v to node. \\
\hline
\end{tabular}

Where:

\author{
ANY-INTEGER: \\ SMALLINT, INTEGER, BIGINT \\ ANY-NUMERIC: \\ SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC
}

\section*{Additional Examples}

\section*{Example:}

From vertices \(\backslash(\backslash\{11,16 \backslash\} \backslash)\) for a distance of \(\backslash(3.0 \backslash)\) on an undirected graph
SELECT * FROM pgr_drivingDistance(
'SELECT id, source, target, cost, reverse_cost FROM edges',
array[11, 16], 3.0, directed => false);
seq | from_v | node | edge | cost | agg_cost
1| 11| 11| -1| 0| 0
2| 11| 7| 8| 1| 1
\(3|-11| 10|5| 1 \mid\)
\(4|11| 12|11| 1 \mid 1\)
\(11|16| 9|1| 1\)
11| 3| 7| 1| 2
11| 6| 2| 1| 2
11| \(8|10| 1 \mid 2\)
11| \(15|3| 1 \mid \quad 2\)
| \(11|17| 15|1| \quad 2\)
11| 1| 6| 1| 3
\begin{tabular}{l|l|l|l|}
\hline \(11 \mid\) & \(5 \mid\) & \(1 \mid\) & 3
\end{tabular}
11| \(9|14| 1 \mid \quad 3\)
16| 16| -1| \(0 \mid 0\)
16| 11| 9| 1|

\(16|17| 15|1| 1\)
16| 7| 8| 1| 2

    \(16|12| 13|1| \quad 2\)
    \(16|3| 7|1| \quad 3\)
    16| 6| 4| 1| 3
    \(16|8| 10|1| 3\)
(23 rows)

\section*{See Also}

\section*{- pgr_alphaShape - Alpha shape computation}
- Sample Data network.

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1) 3.0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2 2.1 2.0
pgr_KSP
pgr_KSP - Yen's algorithm for K shortest paths using Dijkstra.

\section*{boost}

Boost Graph Inside

\section*{Availability}
- Version 2.1.0
- Signature change
- Old signature no longer supported
- Version 2.0.0
- Official function

\section*{Description}

The K shortest path routing algorithm based on Yen's algorithm. " \(K\) " is the number of shortest paths desired.

\section*{Signatures}

\section*{Summary}
pgr_KSP(Edges SQL, start vid, end vid, K, [options])
options: [directed, heap_paths]
RETURNS SET OF (seq, path_id, path_seq, node, edge, cost, agg_cost)

\section*{Example:}

Get 2 paths from \(\backslash(6 \backslash)\) to \(\backslash(17 \backslash)\) on a directed graph
```

SELECT * FROM pgr_KSP(
'SELECT id, source, target, cost, reverse cost FROM edges',
6,17, 2);
seq| path_id | path_seq| node | edge | cost | agg_cost

```
\begin{tabular}{rccccc}
\(1 \mid\) & \(1 \mid\) & \(1 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(1 \mid\) & \(2 \mid\) & \(7 \mid\) & \(10 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(1 \mid\) & \(3 \mid\) & \(8 \mid\) & \(12 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(1 \mid\) & \(4 \mid\) & \(12 \mid\) & \(13 \mid\) & \(1 \mid\) \\
\(5 \mid\) & \(1 \mid\) & \(5 \mid\) & \(17 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(6 \mid\) & \(2 \mid\) & \(1 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(7 \mid\) & \(2 \mid\) & \(2 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(8 \mid\) & \(2 \mid\) & \(3 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) \\
\(9 \mid\) & \(2 \mid\) & \(4 \mid\) & \(16 \mid\) & \(15 \mid\) & \(1 \mid\) \\
\(10 \mid\) & \(2 \mid\) & \(5 \mid\) & \(17 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\((10\) rows \()\)
\end{tabular}

\section*{Parameters}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & SQL query as described. \\
\hline start vid & ANY-INTEGER & Identifier of the departure vertex. \\
\hline end vid & ANY-INTEGER & Identifier of the departure vertex. \\
\hline K & ANY-INTEGER & Number of required paths \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

Optional parameters
\begin{tabular}{lllllll} 
Column & Type & Default & Description \\
\hline directed & BOOLEAN true & \(\circ\) & When true the graph is considered Directed \\
& & & When false the graph is considered as \\
& & Undirected.
\end{tabular}

KSP Optional parameters
Column Type Default Description


Inner Queries

Edges SQL
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL -1 & Weight of the edge (target, source) \\
& & & When negative: edge (target, source) does not exist, therefore it's \\
& & & \\
& & & \\
& & &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Returns set of (seq, path_id, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INTEGER & Sequential value starting from \(\mathbf{1 .}\) \\
\hline path_id & INTEGER & Path identifier. \\
& & Has value \(\mathbf{1}\) for the first of a path fromstart vid to end_vid \\
\hline path_seq & INTEGER & Relative position in the path. Has value \(\mathbf{1}\) for the beginning of a path. \\
\hline node & BIGINT & Identifier of the node in the path fromstart vid to end vid \\
\hline edge & BIGINT & \begin{tabular}{l} 
Identifier of the edge used to go fromnode to the next node in the path sequence. \(\mathbf{- 1}\) for the last node of \\
the path.
\end{tabular} \\
\hline cost & FLOAT & \begin{tabular}{ll} 
Cost to traverse from node using edge to the next node in the path sequence. \\
& \\
\hline agg_cost & FLOAT
\end{tabular} \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Get 2 paths from \(\backslash(6 \backslash)\) to \(\backslash(17 \backslash)\) on an undirected graph
Also get the paths in the heap.
```

SELECT * FROM pgr_KSP(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6,17, 2,
directed => false, heap_paths => true
);
seq | path_id | path_seq | node | edge | cost | agg_cost

| 1\| | 1 \| | 1\| | 6 | 4\| | 1 \| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 21 | 1 \| | $2 \mid$ | 71 | 10\| | 1\| | 1 |
| 31 | 1 \| | 31 | 8\| | 12\| | 1\| | 2 |
| 4 | 1 \| | $4 \mid$ | $12 \mid$ | 13\| | 1 \| | 3 |
| 51 | 1 \| | 51 | $17 \mid$ | -1\| | $0 \mid$ | 4 |
| 61 | 21 | 1\| | 61 | 4\| | 1 \| | 0 |
| 71 | 21 | $2 \mid$ | 7\| | 8\| | 1 \| | 1 |
| 81 | 21 | 31 | 11\| | 11\| | 1\| | 2 |
| 91 | 21 | $4 \mid$ | 12\| | 13\| | $1 \mid$ | 3 |
| $10 \mid$ | 21 | 51 | $17 \mid$ | -1\| | 01 | 4 |
| 11\| | 31 | 1 \| | 61 | 4\| | 1 \| | 0 |
| $12 \mid$ | 31 | $2 \mid$ | 71 | 8\| | 1 \| | 1 |
| 131 | 31 | 31 | 11 | 91 | 1 \| | 2 |
| $14 \mid$ | 31 | 4 \| | $16 \mid$ | 15 \| | 1 \| | 3 |
| 151 | 31 | 5 | $17 \mid$ | -1\| | $0 \mid$ | 4 |
| $16 \mid$ | 4 \| | 1 \| | 61 | $2 \mid$ | 1\| | 0 |
| $17 \mid$ | 4 \| | $2 \mid$ | 10 | 51 | 1 \| | 1 |
| 181 | 4 \| | 31 | 11\| | 91 | 1\| | 2 |
| 19 | 4 \| | 4 \| | $16 \mid$ | $15 \mid$ | 1 \| | 3 |
| $20 \mid$ | 4 \| | 51 | $17 \mid$ | -1\| | 01 | 4 |

(20 rows)

```

\section*{See Also}
- K shortest paths - Category
- Sample Data
- https://en.wikipedia.org/wiki/K_shortest_path_routing

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2
pgr_dijkstraVia - Proposed
pgr_dijkstraVia - Route that goes through a list of vertices.

Proposed functions for next mayor release.
- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.

\section*{boost}

Boost Graph Inside

\section*{Availability}
- Version 2.2.0
- New proposed function

\section*{Description}

Given a list of vertices and a graph, this function is equivalent to finding the shortest path between(vertex_i\\) and \(\backslash\) (vertex_\{i+1\}\\) for all \\(i< size\_of(via\;vertices)\\).

\section*{Route:}
is a sequence of paths.

\section*{Path:}
is a section of the route.

\section*{Signatures}

One Via
```

pgr_dijkstraVia(Edges SQL, via vertices, [options])
options: [directed, strict, U_turn_on_edge]
RETURNS SET OF (seq, path_id, path_seq, start_vid, end_vid, node, edge, cost, agg_cost, route_agg_cost)
OR EMPTY SET

```

\section*{Example:}

Find the route that visits the vertices \(\backslash(\backslash\{5,1,8 \backslash\} \backslash)\) in that order on andirected graph.
```

SELECT * FROM pgr_dijkstraVia(
'SELECT id, source, target, cost, reverse_cost FROM edges order by id',
ARRAY[5, 1, 8]);
seq | path_id | path_seq | start_vid | end_vid | node | edge | cost | agg_cost | route_agg_cost

```


\section*{Parameters}
\begin{tabular}{llll} 
Parameter & Type & Default & Description \\
\hline Edges SQL & TEXT & SQL query as described. \\
\hline via vertices & ARRAY [ ANY-INTEGER ] & \begin{tabular}{l} 
Array of ordered vertices identifiers that are going to be \\
visited.
\end{tabular}
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

Optional parameters
\begin{tabular}{lllllll} 
Column & Type & Default & Description \\
\hline directed & BOOLEAN true & When true the graph is considered Directed \\
& & \begin{tabular}{ll} 
& When false the graph is considered as \\
& \\
& \\
& \\
\end{tabular} &
\end{tabular}

Via optional parameters
\begin{tabular}{llll} 
Parameter & Type & Default & Description \\
\hline strict & BOOLEAN & false & - \\
& & & When true if a path is missing stops and returnsEMPTY SET \\
& & When false ignores missing paths returning all paths found \\
\hline U_turn_on_edge & BOOLEAN true & - & When true departing from a visited vertex will not try to avoid \\
\hline
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& &
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INTEGER & Sequential value starting from \(\mathbf{1}\). \\
\hline path_id & INTEGER & Identifier of a path. Has value \(\mathbf{1}\) for the first path. \\
\hline path_seq & INTEGER & Relative position in the path. Has value \(\mathbf{1}\) for the beginning of a path. \\
\hline start_vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline end_vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline node & BIGINT & Identifier of the node in the path from start_vid to end_vid. \\
\hline edge & BIGINT & \begin{tabular}{l} 
Identifier of the edge used to go fromnode to the next node in the path \\
sequence.
\end{tabular} \\
& & \begin{tabular}{l}
-1 for the last node of the path. \\
\\
\end{tabular} \\
\hline FLOAT & Cost to traverse from node using edge to the next node in the path sequence. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline route_agg_cost & FLOAT & Total cost from start_vid of seq=1 to end_vid of the current seq. \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{- The main query}
- Aggregate cost of the third path.
- Route's aggregate cost of the route at the end of the third path.
- Nodes visited in the route.
- The aggregate costs of the route when the visited vertices are reached.
- Status of "passes in front" or "visits" of the nodes.

All this examples are about the route that visits the vertices \(\backslash(\backslash\{5,7,1,8,15 \backslash\} \backslash)\) in that order on a directed graph.

\section*{SELECT * FROM pgr_dijkstraVia(}
'SELECT id, source, target, cost, reverse_cost FROM edges order by id',
ARRAY[5, 7, 1, 8, 15])
seq | path_id | path_seq | start_vid | end_vid | node | edge | cost | agg_cost | route_agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|}
\hline 1| & 1 | & 1| & 5 & 71 & 5 & 1| & 1 | & 01 & 0 \\
\hline 2 | & 1 | & \(2 \mid\) & 51 & 71 & 61 & 4| & 1 | & \(1 \mid\) & 1 \\
\hline 3| & 1 | & 3| & 51 & 71 & 71 & -1| & \(0 \mid\) & 21 & 2 \\
\hline 4| & \(2 \mid\) & 1| & 71 & 1 | & 71 & 71 & 1 | & 01 & 2 \\
\hline 51 & 21 & \(2 \mid\) & 7| & 1 | & 31 & 61 & 1 | & \(1 \mid\) & 3 \\
\hline \(6 \mid\) & 21 & 31 & 7| & 1 | & 1 | & -1| & \(0 \mid\) & \(2 \mid\) & 4 \\
\hline 71 & 31 & 1| & 1| & 8| & 1 | & 61 & 1| & 01 & 4 \\
\hline 8| & 31 & \(2 \mid\) & 1| & 8| & 31 & 71 & 1 | & \(1 \mid\) & 5 \\
\hline 9| & 31 & 31 & 1| & 8। & 7| & 10| & 1| & \(2 \mid\) & 6 \\
\hline 10| & 31 & 4 | & 1| & 8। & 81 & -1| & 01 & 31 & 7 \\
\hline 11| & 4 | & 1 | & 8 | & 15 & 8 & \(12 \mid\) & 1 | & 01 & 7 \\
\hline \(12 \mid\) & 4 | & \(2 \mid\) & 81 & 15 & 12 & 13 & 1 | & \(1 \mid\) & 8 \\
\hline 13 | & 4 | & 31 & 81 & 15 & 17 & 15 & 1 | & 2 | & 9 \\
\hline 14 | & 4 | & 4 | & 8| & 15 | & 16 & 16 & \(1 \mid\) & 31 & 10 \\
\hline 151 & 4 | & 51 & 8| & 15 | & 15 & -2| & \(0 \mid\) & \(4 \mid\) & 11 \\
\hline (15 ro & & & & & & & & & \\
\hline
\end{tabular}

Aggregate cost of the third path.
```

SELECT agg_cost FROM pgr_dijkstraVia
SELECT id, source, target, cost, reverse_cost FROM edges order by id',
ARRAY[5, 7, 1, 8, 15]
WHERE path_id = 3 AND edge <0;
agg_cost
3
(1 row)

```

Route's aggregate cost of the route at the end of the third path.
```

SELECT route_agg_cost FROM pgr_dijkstraVia(
SELECT id, source, target, cost, reverse_cost FROM edges order by id',
ARRAY[5, 7, 1, 8, 15])
WHERE path_id = 3 AND edge <0;
route_agg_cost
7
(1 row)

```

Nodes visited in the route.
```

SELECT row number() over () as node seq, node
FROM pgr_dijkstraVia(
'SELECT id, source, target, cost, reverse_cost FROM edges order by id',
ARRAY[5, 7, 1, 8, 15]
WHERE edge <> -1 ORDER BY seq
node_seq | node
1| 5
2| 6
3|7
4|
5| 1
6| 3
7| 7
8| 8
9| 12
10| 17
11| 16
12| 15
(12 rows)

```

The aggregate costs of the route when the visited vertices are reached.
\begin{tabular}{cc}
\(1 \mid\) & 2 \\
\(2 \mid\) & 4 \\
\(3 \mid\) & 7 \\
\(4 \mid\) & 11 \\
(4 rows) &
\end{tabular}
```

Status of "passes in front" or "visits" of the nodes

```
```

SELECT seq, route_agg_cost, node, agg_cost ,
CASE WHEN edge = -1 THEN 'visits'
ELSE 'passes in front
END as status
FROM pgr dijkstraVia(
'SELECT id, source, target, cost, reverse cost FROM edges order by id',
ARRAY[5, 7, 1, 8, 15])
WHERE agg_cost <> 0 or seq = 1,
seq |route_agg_cost | node |agg_cost | status
1| 0| 5| 0| passes in front
1| 6| 1|passes in front
2| 7| 2 |visits
3| 1 passes in front
1| 2|visits
3| 1 | passes in fron
6 7 2 2 passes in front
7| 8| 3|visits
8| 12| 1|passes in front
9| 17| 2|passes in front
10| 16| 3 |passes in front
11| 15| 4 |passes in front
(12 rows)

```

See Also

\section*{- Via - Category}
- Dijkstra - Family of functions.
- Sample Data network.
- https://en.wikipedia.org/wiki/Dijkstra\%27s_algorithm

\section*{Indices and tables}
```

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```
- Supported versions: Latest (3.3) 3.2

\section*{pgr_dijkstraNear - Proposed}
pgr_dijkstraNear — Using Dijkstra's algorithm, finds the route that leads to the nearest vertex.

\section*{Warning}

Proposed functions for next mayor release.
- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.

\section*{Availability}
- Version 3.3.0
- Promoted to proposed function
- Version 3.2.0
- New experimental function

\section*{Description}

Given a graph, a starting vertex and a set of ending vertices, this function finds the shortest path from the starting vertex to the nearest ending vertex.

\section*{Characteristics}
- Uses Dijkstra algorithm.
- Works for directed and undirected graphs.
- When there are more than one path to the same vertex with same cost:
- The algorithm will return just one path
- Optionally allows to find more than one path.
- When more than one path is to be returned:

> Results are sorted in increasing order of:
- aggregate cost
- Within the same value of aggregate costs:
- results are sorted by (source, target)
- Running time: Dijkstra running time: \(\backslash(\mathrm{drt}=\mathrm{O}((|\mathrm{E}|+|\mathrm{V}|) \log |\mathrm{V}|) \backslash)\)
- One to Many; \\(drt\\)
- Many to One: \(\backslash(\mathrm{drt}\) )
- Many to Many: \(\backslash(\) drt * |Starting vids|\\)
- Combinations: \(\backslash(\mathrm{drt} *|S t a r t i n g ~ v i d s| \backslash)\)

\section*{Signatures}

\section*{Summary}
```

pgr_dijkstraNear(Edges SQL, start vid, end vids, [options A])
pgr_dijkstraNear(Edges SQL, start vids, end vid, [options A])
pgr_dijkstraNear(Edges SQL, start vids, end vids, [options B])
pgr_dijkstraNear(Edges SQL, Combinations SQL, [options B])
options A: [directed, cap]
options B: [directed, cap, global]
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{One to Many}
```

pgr_dijkstraNear(Edges SQL, start vid, end vids, [options])

```
options: [directed, cap]
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

\section*{Example:}

Departing on car from vertex \(\backslash(6 \backslash)\) find the nearest subway station.
- Using a directed graph for car routing.
- The subway stations are on the following vertices \(\backslash(\backslash\{1,10,11 \backslash\} \backslash)\)
- The defaults used:
- directed \(=>\) true
- cap => 1
```

1 SELECT * FROM pgr_dijkstraNear(
2 'SELECT id, source, target, cost, reverse_cost FROM edges',
3,ARRAY[10,11,1]);
4 seq| path_seq| start_vid | end_vid | node | edge | cost| agg_cost
6
7 2| 2| 6| 11| 7| 8| 1| 1
8| 3| 6| 11| 11| -1| 0| 2
9(3 rows)
10

```

The result shows that station at vertex\\(11\\) is the nearest.

\section*{Many to One}
```

pgr_dijkstraNear(Edges SQL, start vids, end vid,[options])
options: [directed, cap]

```
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

\section*{Example:}

Departing on a car from a subway station find the nearesttwo stations to vertex \(\backslash(2 \backslash)\)
- Using a directed graph for car routing.
- The subway stations are on the following vertices \(\backslash(\backslash\{1,10,11 \backslash\} \backslash)\)
- On line 4: using the positional parameter: directed set to true
- In line 5: using named parameter cap => 2
```

1 SELECT * FROM pgr_dijkstraNear
2 'SELECT id, source, target, cost, reverse_cost FROM edges'
ARRAY[10, 11, 1], 6
true,
cap => 2);
seq|path_seq| start_vid | end_vid | node | edge | cost | agg_cost
8 1----+-------------------+--------------+-----+------+-
2| 2| 10| 6| 6| -1| 0| 1
3| 1| 11| 6| 11| 8| 1| 0
14| 2| 11| 6| 7| 4| 1| 1
2 5| 3| 11| 6| 6| -1| 0| 2
3(5 rows)
14

```

The result shows that station at vertex \(\backslash(10)\) is the nearest and the next best is \(\backslash(11)\).

\section*{Many to Many}
```

pgr_dijkstraNear(Edges SQL, start vids, end vids, [options])
options: [directed, cap, global]

```
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

\section*{Example:}

Find the best pedestrian connection between two lines of buses
- Unsing an undirected graph for pedestrian routing
- The first subway line stations are at \(\backslash(\backslash\{15,16 \backslash\} \backslash)\)
- The second subway line stations stops are at\\(\\{1,10,11\\}\\)
- On line 4: using the named parameter: directed => false
- The defaults used:
- cap => 1
- global => true
```

1 SELECT * FROM pgr_dijkstraNear(
2 'SELECT id, source, target, cost, reverse_cost FROM edges',
3 ARRAY[15, 16], ARRAY[10, 11, 1]
4 directed => false);
5 seq \| path_seq\| start_vid \| end_vid \| node \| edge \| cost \| agg_cost
6 -------------------------+---------------------------------
8 2| 2| 15| 10| 10| -1| 0| 1
9(2 rows)
10

```

For a pedestrian the best connection is to get on/off is at vertex \((15 \backslash)\) of the first subway line and at vertex \((10 \backslash)\) of the second subway line.

Only one route is returned because global is true and cap is 1

\section*{Combinations}
```

pgr_dijkstraNear(Edges SQL, Combinations SQL,[options])
options: [directed, cap, global]
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Find the best car connection between all the stations of two subway lines
- Using a directed graph for car routing.
- The first subway line stations stops are at \(\(\backslash\{1,10,11 \backslash\} \backslash)\)
- The second subway line stations are at \(\backslash(\backslash\{15,16 \backslash\} \backslash)\)

The combinations contents:
```

SELECT unnest(ARRAY[10, 11, 1]) as source, target
FROM (SELECT unnest(ARRAY[15, 16]) AS target) a
UNION
SELECT unnest(ARRAY[15, 16]), target
FROM (SELECT unnest(ARRAY[10, 11, 1]) AS target) b ORDER BY source, target;
source | target

| $1 \mid$ | 15 |
| :---: | :---: |
| $1 \mid$ | 16 |
| $10 \mid$ | 15 |
| $10 \mid$ | 16 |
| $11 \mid$ | 15 |
| $11 \mid$ | 16 |
| $15 \mid$ | 1 |
| $15 \mid$ | 10 |
| $15 \mid$ | 11 |
| $16 \mid$ | 1 |
| $16 \mid$ | 10 |
| $16 \mid$ | 11 |

(12 rows)

```

\section*{The query:}
- lines \(3 \sim 4\) sets the start vertices to be from the first subway line and the ending vertices to be from the second subway line
- lines \(6 \sim 7\) sets the start vertices to be from the first subway line and the ending vertices to be from the first subway line
- On line 8: using the named parameter is global => false
- The defaults used:
- directed \(=>\) true
- cap => 1
```

1 SELECT * FROM pgr_dijkstraNear(
2 'SELECT id, source, target, cost, reverse_cost FROM edges',
3 'SELECT unnest(ARRAY[10, 11, 1]) as source, target
4 FROM (SELECT unnest(ARRAY[15, 16]) AS target) a
UNION
SELECT unnest(ARRAY[15, 16]), target
FROM (SELECT unnest(ARRAY[10, 11, 1]) AS target) b'
global => false);
9 seq| path_seq| start_vid | end_vid | node | edge | cost | agg_cost
11| 1| 11| 16| 11| 9| 1| 0
2| 2| 11| 16| 16| -1| 0)
3| 1| 15| 10| 15| 3| 1| 0
4| 2| 15| 10| 10| -1| 0| 1
5 5| 1| 16| 11| 16| 9| 1| 0
6 6|
llllllllllllll
9 9| 3| lllll

```

```

11| 2| 1| 16| 3| 7| 1| 1
12| 3| 1| 16| 7| 8| 1| 2
23}1013|\mp@code{4|
414| 5| 1| 16| 16| -1| 0| 4
26

```

From the results:
- making a connection from the first subway line \(\backslash(\backslash\{1,10,11 \backslash\} \backslash)\) to the second \(\backslash(\backslash\{15,16 \backslash\} \backslash)\) :
- The best connections from all the stations from the first line are:\\(\{(1 \rightarrow 16) (10 \rightarrow 16) (11 \rightarrow 16)\}\\)
- The best one is \(\backslash((11 \backslash\) rightarrow 16\() \backslash)\) with a cost of \(\backslash(1 \backslash)\) (lines: 11 and 12)
- making a connection from the second subway line \(\backslash(\backslash\{15,16 \backslash\} \backslash)\) to the first \(\backslash(\backslash\{1,10,11 \backslash\} \backslash)\) :
- The best connections from all the stations from the second line are:\\(\{(15 \rightarrow 10) (16 \rightarrow 11)\}\\)
- Both are equaly good as they have the same cost. (lines:13 and 14 and lines: 15 and 16)

\section*{Parameters}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}

Dijkstra optional parameters


Near optional parameters
\begin{tabular}{llll} 
Parameter & Type & Default & Description \\
\hline cap & BIGINT & 1 & Find at most cap number of nearest shortest paths \\
\hline global & BOOLEAN & true & \(-\quad\) When true: only cap limit results will be returned \\
& & & - \\
& & & When false: cap limit per Start vid will be returned \\
\hline
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL & -1 & Weight of the edge (target, source) \\
& & & When negative: edge (target, source) does not exist, therefore it's \\
& & & \\
& & & \\
& & & \\
& & & \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the departure vertex. \\
\hline target & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the arrival vertex. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INTEGER & Sequential value starting from \(\mathbf{1}\). \\
\hline path_seq & INTEGER & Relative position in the path. Has value \(\mathbf{1}\) for the beginning of a path. \\
\hline start_vid & BIGINT & Identifier of the starting vertex of the current path. \\
\hline end_vid & BIGINT & Identifier of the ending vertex of the current path. \\
\hline node & BIGINT & Identifier of the node in the path fromstart_vid to end_vid. \\
\hline edge & BIGINT & \begin{tabular}{l} 
Identifier of the edge used to go fromnode to the next node in the path sequence. \(\mathbf{- 1}\) for \\
the last node of the path.
\end{tabular} \\
\hline cost & FLOAT & Cost to traverse from node using edge to the next node in the path sequence. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline
\end{tabular}

\section*{See Also}
- Dijkstra - Family of functions
- pgr_dijkstraNearCost - Proposed
- Sample Data network.
- boost: https://www.boost.org/libs/graph/doc/table_of_contents.html
- Wikipedia: https://en.wikipedia.org/wiki/Dijkstra\%27s_algorithm

\section*{Indices and tables}
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- Supported versions: Latest (3.3) 3.2
pgr_dijkstraNearCost - Proposed
pgr_dijkstraNearCost — Using dijkstra algorithm, finds the route that leads to the nearest vertex.

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- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.

Boost Graph Inside

\section*{Availability}
- Version 3.3.0
- Promoted to proposed function
- Version 3.2.0
- New experimental function

\section*{Description}

Given a graph, a starting vertex and a set of ending vertices, this function finds the shortest path from the starting vertex to the nearest ending vertex.
- Uses Dijkstra algorithm.
- Works for directed and undirected graphs.
- When there are more than one path to the same vertex with same cost:
- The algorithm will return just one path
- Optionally allows to find more than one path.
- When more than one path is to be returned:
- Results are sorted in increasing order of:
- aggregate cost
- Within the same value of aggregate costs:
- results are sorted by (source, target)
- Running time: Dijkstra running time: \(\backslash(\mathrm{drt}=\mathrm{O}((|\mathrm{E}|+|\mathrm{V}|) \log |\mathrm{V}|) \backslash)\)
- One to Many; \\(drt\\)
- Many to One: \(\backslash(\mathrm{drt} \backslash)\)
- Many to Many: \\(drt * |Starting vids|\\)
- Combinations: \\(drt * |Starting vids|\\)

\section*{Signatures}

\section*{Summary}
```

pgr_dijkstraNearCost(Edges SQL, start vid, end vids, [options A])
pgr_dijkstraNearCost(Edges SQL, start vids, end vid, [options A])
pgr_dijkstraNearCost(Edges SQL, start vids, end vids, [options B])
pgr_dijkstraNearCost(Edges SQL, Combinations SQL, [options B])
options A: [directed, cap]
options B: [directed, cap, global]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{One to Many}
```

pgr_dijkstraNearCost(Edges SQL, start vid, end vids, [options])
options: [directed, cap]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Departing on car from vertex \(\backslash(6 \backslash)\) find the nearest subway station.
- Using a directed graph for car routing.
- The subway stations are on the following vertices \(\backslash(\backslash\{1,10,11 \backslash\} \backslash)\)
- The defaults used:
- directed \(=>\) true
- cap => 1
```

1 SELECT * FROM pgr_dijkstraNearCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6, ARRAY[10, 11, 1]);
start_vid | end_vid | agg_cost
----------------------------
(1 row)
8

```

The result shows that station at vertex \((11 \backslash)\) is the nearest.

\section*{Many to One}
```

pgr_dijkstraNearCost(Edges SQL, start vids, end vid, [options])
options: [directed, cap]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Departing on a car from a subway station find the nearesttwo stations to vertex \(\backslash(6 \backslash)\)

Using a directed graph for car routing.
The subway stations are on the following vertices \(\backslash(\backslash\{1,10,11 \backslash\} \backslash)\)
- On line 4: using the positional parameter: directed set to true

In line 5: using named parameter cap \(=>2\)
```

SELECT * FROM pgr_dijkstraNearCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[10, 11, 1], 6,
true,
cap => 2) ORDER BY agg_cost;
start_vid | end_vid | agg_cost
10| 6| 1
11| 6| 2
(2 rows)
1 1

```

The result shows that station at vertex \(\backslash(10 \backslash)\) is the nearest and the next best is \(\backslash(11 \backslash)\).

\section*{Many to Many}
```

pgr_dijkstraNearCost(Edges SQL, start vids, end vids, [options])
options: [directed, cap, global]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Find the best pedestrian connection between two lines of buses
- Unsing an undirected graph for pedestrian routing
- The first subway line stations are at\\(\\{15, 16\\}\\)
- The second subway line stations stops are at \(\backslash(\backslash\{1,10,11 \backslash\} \backslash)\)
- On line 4: using the named parameter: directed => false
- The defaults used:
- cap => 1
- global => true
```

SELECT * FROM pgr_dijkstraNearCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[15, 16], ARRAY[10, 11, 1],
directed => false);
start_vid | end_vid | agg_cost
6 ----------+--------+---------
(1 row)
9

```

For a pedestrian the best connection is to get on/off is at vertex \((15 \backslash)\) of the first subway line and at vertex \((10 \backslash)\) of the second subway line.

Only one route is returned because global is true and cap is 1

\section*{Combinations}
```

pgr_dijkstraNearCost(Edges SQL, Combinations SQL, [options])
options: [directed, cap, global]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Find the best car connection between all the stations of two subway lines
- Using a directed graph for car routing.
- The first subway line stations stops are at \(\backslash(\backslash\{1,10,11 \backslash\} \backslash)\)
- The second subway line stations are at \(\backslash(\backslash\{15,16 \backslash\} \backslash)\)

The combinations contents:
\begin{tabular}{c|c}
\(1 \mid\) & 15 \\
\(1 \mid\) & 16 \\
\(10 \mid\) & 15 \\
\(10 \mid\) & 16 \\
\(11 \mid\) & 15 \\
\(11 \mid\) & 16 \\
15 & 1 \\
15 & 10 \\
15 & 11 \\
16 & 1 \\
16 & 10 \\
16 & 11
\end{tabular}

The query:
- lines 3~4 sets the start vertices to be from the fist subway line and the ending vertices to be from the second subway line
- lines \(6 \sim 7\) sets the start vertices to be from the first subway line and the ending vertices to be from the first subway line
- On line 8: using the named parameter is global => false
- The defaults used:
- directed \(=>\) true
- cap => 1
```

SELECT * FROM pgr_dijkstraNearCost(
'SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT unnest(ARRAY[10, 11, 1]) as source, target
FROM (SELECT unnest(ARRAY[15, 16]) AS target) a
UNION
SELECT unnest(ARRAY[15, 16]), target
FROM (SELECT unnest(ARRAY[10, 11, 1]) AS target) b',
global => false)
start_vid end_vid agg_cost
----------+--------+-----------
15 10| 10
16| 11| 1
40| 16| 2
(5 rows)
17

```

\section*{From the results:}
- making a connection from the first subway line \(\backslash(\backslash\{1,10,11 \backslash\} \backslash)\) to the second \(\backslash(\backslash\{15,16 \backslash\} \backslash)\) :
- The best connections from all the stations from the first line are:\\(\{(1 \rightarrow 16) (10 \rightarrow 16) (11 \rightarrow 16)\}\\)
- The best one is \\((11 \rightarrow 16)\\) with a cost of \\(1\\) (lines: 1)
- making a connection from the second subway line \(\backslash(\backslash\{15,16 \backslash\} \backslash)\) to the first \(\backslash(\backslash\{1,10,11 \backslash\} \backslash)\) :
- The best connections from all the stations from the second line are:\\(\{(15 \rightarrow 10) (16 \rightarrow 11) \}\\)
- Both are equaly good as they have the same cost. (lines:12 and 13)

\section*{Parameters}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}

\section*{Dijkstra optional parameters}
\begin{tabular}{llllll} 
Column & Type & Default & Description \\
\hline directed & BOOLEAN true & When true the graph is considered Directed \\
& & \begin{tabular}{lll}
Wh n false \\
& & Undirected.
\end{tabular}
\end{tabular}
\begin{tabular}{llll} 
Parameter & Type & Default & Description \\
\hline cap & BIGINT & 1 & Find at most cap number of nearest shortest paths \\
\hline global & BOOLEAN & true & - When true: only cap limit results will be returned \\
& & & When false: cap limit per Start vid will be returned
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& & \\
& & \\
& &
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & ANY- & Identifier of the departure vertex. \\
& INTEGER & \\
\hline target & ANY- & Identifier of the arrival vertex. \\
& INTEGER & \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

Result Columns

Set of (start_vid, end_vid, agg_cost)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline start_vid & BIGINT & Identifier of the starting vertex. \\
\hline end_vid & BIGINT & Identifier of the ending vertex. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to end_vid. \\
\hline
\end{tabular}

\section*{See Also}
- Dijkstra - Family of functions
- pgr_dijkstraNear - Proposed
- Sample Data network.
- boost: https://www.boost.org/libs/graph/doc/table_of_contents.html
- Wikipedia: https://en.wikipedia.org/wiki/Dijkstra\%27s_algorithm

\section*{Indices and tables}
. Index
- Search Page

\section*{Introduction}

Dijkstra's algorithm, conceived by Dutch computer scientist Edsger Dijkstra in 1956. It is a graph search algorithm that solves the shortest path problem for a graph with non-negative edge path costs, producing a shortest path from a starting vertex to an ending vertex. This implementation can be used with a directed graph and an undirected graph.

The main characteristics are:
- Process is done only on edges with positive costs.
- A negative value on a cost column is interpreted as the edge does not exist.
- Values are returned when there is a path.
- When there is no path:
- When the starting vertex and ending vertex are the same.
- The aggregate cost of the non included values \(\backslash((v, v) \backslash)\) is \(\backslash(0 \backslash)\)
- When the starting vertex and ending vertex are the different and there is no path:
- The aggregate cost the non included values \(\backslash((u, v) \backslash)\) is \(\backslash(\backslash i n f t y \backslash)\)

For optimization purposes, any duplicated value in the starting vertices or on the ending vertices are ignored.
Running time: \(\backslash\left(\mathrm{O}\left(\mid\right.\right.\) start \(\backslash\) vids \(\left.\left.\left.\right|^{*}(\mathrm{~V} \backslash \log \mathrm{~V}+\mathrm{E})\right) \backslash\right)\)
The Dijkstra family functions are based on the Dijkstra algorithm.

\section*{Parameters}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}

Optional parameters


Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL -1 & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& & \\
& & \\
& & \\
\hline
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the departure vertex. \\
\hline target & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the arrival vertex. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Advanced documentation}

Given the following query:
pgr_dijkstra(\\(sql, start_\{vid\}, end_\{vid\}, directed \(\backslash\) ))
```

where $sq| = \{(id_i, source_i, target_i, cost_i, reverse\_cost_i)\}$

```
and
- \(\quad\) (source \(=\) \bigcup source_i \(\backslash\) ),
- \(\backslash\) (target \(=\) \bigcup target_i \()\),

The graphs are defined as follows:

\section*{Directed graph}

The weighted directed graph, \(\backslash\left(\mathrm{G}_{-} \mathrm{d}(\mathrm{V}, \mathrm{E}) \backslash\right)\), is definied by:
- the set of vertices \(\backslash(\mathrm{V} \backslash)\)
- \(\backslash(V=\) source \(\backslash c u p\) target \(\backslash c u p\{\) start_\{vid \} \} \cup \{end_\{vid\}\}\\)
- the set of edges \(\backslash(\mathrm{E})\)
- \(\backslash(E=\backslash\) begin \(\{\) cases \(\} \backslash\) text \(\} \backslash\{\) (source_i, target_i, cost_i) \text\{ when \} cost \(>=0 \backslash\} \& \backslash q u a d\) ltext \(\{\) if \(\}\) reverse \(\backslash\) cost \(=\) lvarnothing \(\backslash \backslash \backslash \operatorname{text}\} \backslash \operatorname{text}\} \& \backslash q u a \bar{d} \backslash \operatorname{text}\{\overline{\}} \backslash \backslash \backslash \operatorname{tex} t\} \backslash\{\) (source_i, target_i, cost_i) \text\{ when \} cost \(>=0 \backslash\} \&\) \quad \text\{ \} \\ \cup \\{(target_i, source_i, reverse\_cost_i) \text\{ when \} reverse\_cost_i>=0 \\} \& \quad \text\{if \} reverse\_cost \neq \varnothing \\\end\{cases\}\) }

\section*{Undirected graph}

The weighted undirected graph, \(\backslash\left(\mathrm{G}_{-} \mathrm{u}(\mathrm{V}, \mathrm{E}) \backslash\right)\), is definied by:
- the set of vertices \(\backslash(\mathrm{V} \backslash)\)
- \(\backslash\left(\mathrm{V}=\right.\) source \(\backslash c u p\) target \(\backslash c u p ~\left\{s t a r t \_v\{\right.\) vid \(\left.\left.\}\right\} \backslash c u p ~\left\{e n d \_\{v i d\}\right\} \backslash\right)\)
- the set of edges \(\backslash(E \backslash)\)
- \(\backslash(E=\backslash\) begin \(\{\) cases \(\} \backslash\) text \(\} \backslash\{\) (source_i, target_i, cost_i) \(\backslash\) text \(\{\) when \(\}\) cost \(>=0 \backslash\} \& \backslash q u a d \backslash t e x t\} \backslash \backslash \backslash u p \backslash\{(\) target_i, source_i, cost_i) \text\{ when \} cost \(>=0 \backslash\} \& \backslash q u a d ~ \ t e x t\{i f\} r e v e r s e \backslash c o s t=\mid v a r n o t h i n g ~ \ \backslash \backslash t e x t\{ \} \backslash t e x t\{ \} \& \backslash t e x t\{ \}\) \(\ \backslash \backslash\) text \(\} \backslash\{(\) source_i, target_i, cost_i) \text \(\{\) when \} cost \(>=0 \backslash\} \& \backslash\) text \(\} \backslash \backslash \backslash u p \backslash\{(\) target_i, source_i, cost_i) \text \(\{\)
 \text\{ \} \(\backslash \backslash\) \cup \(\backslash\{(\) source_i, target_i, reverse\_cost_i) \text\{ when \} reverse\_cost_i >=0)\\\(\& \quad } \backslash t e x t\{\) if \(\}\) reverse\_cost \neq \varnothing \\\end\{cases\}\) }

\section*{The problem}

\section*{Given:}
- \\(start_\{vid\} \in V) a starting vertex
- \(\backslash\left(e n d \_\{v i d\}\right.\) (in \(\left.\mathrm{V} \backslash\right)\) an ending vertex
 lend\{cases \(\} \backslash\) )

Then:
```

- $\boldsymbol{\pi} = \{(path\_seq_i, node_i, edge_i, cost_i, agg\cost_i)\}$

```
where:
- \(\backslash(\) path \(\backslash\) seq_i \(=i \backslash)\)
- \(\backslash(\) path \(\backslash\) seq_ \(\{\mid \backslash\) pi \(\mid\}=|\backslash p i| \backslash)\)
- \(\backslash(\) node_i \(\backslash i n \mathrm{~V} \backslash)\)
- \(\backslash(\) node_ \(1=\) start_\{vid \(\} \backslash)\)
- \(\backslash(\) node_ \(\{\mid \backslash\) pi \(\mid\}=\) end_\{vid \(\} \backslash)\)
- \\(\forall i \neq | \pi |, \quad (node_i, node_\{i+1\}, cost_i) \in E\\)
- \\(edge_i = \begin\{cases\} id_\{(node_i, node_\{i+1\},cost_i)\} \& } \(=|\backslash \mathrm{pi}| \backslash \backslash\) lend \(\{\) cases \(\} \backslash)\)
- \(\backslash(\operatorname{cost} \mathrm{i}=\operatorname{cost}\{(\) node_i, node_\{i+1\})\} )
- \\(agg\_cost_i = \begin\{cases\} 0 \& } 1 \text { quad } \backslash \operatorname { t e x t } \{ \text { when } \} ~ i = 1 \backslash \text { displaystyle\sum_\{k=1\}^\{i\} cost_\{(node_\{k-1\}, }


In other words: The algorithm returns a the shortest path between \(\backslash(\) start_\{vid \(\} \backslash\) ) and \(\backslash\left(e n d \_\{v i d\} \backslash\right)\), if it exists, in terms of a sequence of nodes and of edges,
- \(\backslash(p a t h \backslash s e q \backslash)\) indicates the relative position in the path of the\\(node\\) or \(\backslash(e d g e \backslash)\).
- \(\backslash\) (cost \(\backslash\) ) is the cost of the edge to be used to go to the next node.
- \(\backslash(\) agg \(\backslash \cos t \backslash)\) is the cost from the \(\backslash(\) start_ \(\{v i d\} \backslash)\) up to the node.

If there is no path, the resulting set is empty.

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.6 2.5 2.42 .3

\section*{Flow - Family of functions}
pgr_maxFlow - Only the Max flow calculation using Push and Relabel algorithm.
- pgr_boykovKolmogorov - Boykov and Kolmogorov with details of flow on edges.
- pgr_edmondsKarp - Edmonds and Karp algorithm with details of flow on edges.
- pgr_pushRelabel - Push and relabel algorithm with details of flow on edges.
- Applications
- pgr_edgeDisjointPaths - Calculates edge disjoint paths between two groups of vertices.
- pgr_maxCardinalityMatch - Calculates a maximum cardinality matching in a graph.

\section*{Experimental}

\section*{Warning}

Possible server crash
- These functions might create a server crash

\section*{Warning}

Experimental functions
They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need \(c / c++\) coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting
- pgr_maxFlowMinCost - Experimental - Details of flow and cost on edges.
- pgr_maxFlowMinCost_Cost - Experimental - Only the Min Cost calculation.

Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: \(2.6 \mathbf{2} \mathbf{2} \mathbf{2} \mathbf{2} 4\)
pgr_maxFlow
pgr_maxFlow - Calculates the maximum flow in a directed graph from the source(s) to the targets(s) using the Push Relabel algorithm.
- Version 3.2.0
- New proposed signature
- pgr_maxFlow (Combinations)
- Version 3.0.0
- Official function
- Version 2.4.0
- New Proposed function

\section*{Description}

\section*{The main characteristics are:}
- The graph is directed.
- Calculates the maximum flow from the source(s) to the target(s).
- When the maximum flow is \(\mathbf{0}\) then there is no flow and \(\mathbf{0}\) is returned.
- There is no flow when a source is the same as a target.
- Any duplicated value in the source(s) or target(s) are ignored.
- Uses the pgr_pushRelabel algorithm.
- Running time: \(\backslash\left(\mathrm{O}\left(\mathrm{V}^{\wedge} 3\right) \backslash\right)\)

Signatures

\section*{Summary}
```

pgr_maxFlow(Edges SQL, start vid, end vid)
pgr_maxFlow(Edges SQL, start vid, end vids)
pgr_maxFlow(Edges SQL, start vids, end vid)
pgr_maxFlow(Edges SQL, start vids, end vids)
pgr_maxFlow(Edges SQL, Combinations SQL)
RETURNS BIGINT

```

\section*{One to One}
```

pgr_maxFlow(Edges SQL, start vid, end vid)
RETURNS BIGINT

```

\section*{Example:}

From vertex \(\backslash(11 \backslash)\) to vertex \(\backslash(12 \backslash)\)
```

SELECT * FROM pgr_maxFlow(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
11, 12);
pgr_maxflow
230
(1 row)

```

\section*{One to Many}
```

pgr_maxFlow(Edges SQL, start vid, end vids)
RETURNS BIGINT

```

\section*{Example:}

From vertex \(\backslash(11 \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr_maxFlow(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
11, ARRAY[5, 10, 12]);
pgr_maxflow
340
(1 row)

```
```

pgr_maxFlow(Edges SQL, start vids, end vid)
RETURNS BIGINT

```

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertex \\(12\\)
```

SELECT * FROM pgr_maxFlow(
SELECT id, source, target, capacity, reverse_capacity
FROM edges',
ARRAY[11, 3, 17], 12);
pgr_maxflow
230
(1 row)

```

Many to Many
```

pgr_maxFlow(Edges SQL, start vids, end vids)
RETURNS BIGINT

```

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr_maxFlow(
SELECT id, source, target, capacity, reverse_capacity
FROM edges',
ARRAY[11, 3, 17], ARRAY[5, 10, 12]);
pgr_maxflow
360
(1 row)

```

Combinations
pgr_maxFlow(Edges SQL, Combinations SQL)
RETURNS BIGINT

\section*{Example:}

Using a combinations table, equivalent to calculating result from vertices \(\backslash(\backslash\{5,6 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,15,14 \backslash\} \backslash)\).
The combinations table:
```

SELECT source, target FROM combinations
WHERE target NOT IN (5, 6);
source | target
5| 10
6| 15
6| 14
(3 rows)

```

The query:
```

SELECT * FROM pgr_maxFlow(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
'SELECT * FROM combinations WHERE target NOT IN (5, 6)');
pgr_maxflow
80
(1 row)

```

Parameters
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices.
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{|c|c|c|c|}
\hline Column & Type & Default & Description \\
\hline id & ANY-INTEGER & & Identifier of the edge. \\
\hline source & ANY-INTEGER & & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & & Identifier of the second end point vertex of the edge. \\
\hline capacity & ANY-INTEGER & & Weight of the edge (source, target) \\
\hline reverse_capacity & ANY-INTEGER & -1 & Weight of the edge (target, source) \\
\hline & & & - When negative: edge (target, source) does not exist, therefore it's not part of the graph. \\
\hline
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & ANY- & Identifier of the departure vertex. \\
& INTEGER & \\
\hline target & ANY- & Identifier of the arrival vertex. \\
& INTEGER & \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

\section*{SMALLINT, INTEGER, BIGINT}

\section*{Result Columns}
\begin{tabular}{ll} 
Type & Description \\
\hline BIGINT & \begin{tabular}{l} 
Maximum flow possible from the source(s) to the \\
target(s)
\end{tabular} \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Manually assigned vertex combinations.
```

SELECT * FROM pgr_maxFlow(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
'SELECT * FROM (VALUES (5, 10), (6, 15), (6, 14)) AS t(source, target)');
pgr_maxflow
80
(1 row)

```

See Also
- Flow - Family of functions
- pgr_pushRelabel
- https://www.boost.org/libs/graph/doc/push_relabel_max_flow.html
- https://en.wikipedia.org/wiki/Push\%E2\%80\%93relabel_maximum_flow_algorithm

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.6 2.5 2.42 .3
pgr_boykovKolmogorov
pgr_boykovKolmogorov - Calculates the flow on the graph edges that maximizes the flow from the sources to the targets using Boykov Kolmogorov algorithm.

Boost Graph Inside

\section*{Availability}
- Version 3.2.0
- New proposed signature - pgr_boykovKolmogorov (Combinations)
- Version 3.0.0
- Official function
- Version 2.5.0
- Renamed from pgr_maxFlowBoykovKolmogorov
- Proposed function
- Version 2.3.0
- New Experimental function

\section*{Description}

\section*{The main characteristics are:}
- The graph is directed.
- Process is done only on edges with positive capacities.
- When the maximum flow is 0 then there is no flow andEMPTY SET is returned.
- There is no flow when a source is the same as a target.
- Any duplicated value in the source(s) or target(s) are ignored.
- Calculates the flow/residual capacity for each edge. In the output
- Edges with zero flow are omitted.
- Creates a super source and edges to all the source(s), and asuper target and the edges from all the targets(s).
- The maximum flow through the graph is guaranteed to be the value returned bypgr_maxFlow when executed with the same parameters and can be calculated:
- By aggregation of the outgoing flow from the sources
- By aggregation of the incoming flow to the targets
- Running time: Polynomial

Signatures

\section*{Summary}
```

pgr_boykovKolmogorov(Edges SQL, start vid, end vid)
pgr_boykovKolmogorov(Edges SQL, start vid, end vids)
pgr_boykovKolmogorov(Edges SQL, start vids, end vid)
pgr_boykovKolmogorov(Edges SQL, start vids, end vids)
pgr_boykovKolmogorov(Edges SQL, Combinations SQL)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)
OR EMPTY SET

```
```

pgr_boykovKolmogorov(Edges SQL, start vid, end vid)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)
OR EMPTY SET

```

\section*{Example:}

From vertex \(\backslash(11 \backslash)\) to vertex \(\backslash(12 \backslash)\)
```

SELECT * FROM pgr_boykovKolmogorov(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
11, 12);
seq | edge | start_vid | end_vid | flow | residual_capacity

| $1\|10\|$ | $7 \mid$ | $8\|100\|$ | 30 |
| ---: | ---: | :---: | ---: |
| $2\|12\|$ | $8 \mid$ | $12\|100\|$ | 0 |
| $3\|8\|$ | $11 \mid$ | $7\|100\|$ | 30 |
| $4\|11\|$ | $11 \mid$ | $12\|130\|$ | 0 |
| (4 rows) |  |  |  |

```

\section*{One to Many}
```

pgr_boykovKolmogorov(Edges SQL, start vid, end vids)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity
OR EMPTY SET

```

\section*{Example:}

From vertex \(\backslash(11 \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr_boykovKolmogorov(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
11, ARRAY[5, 10, 12]);
seq | edge | start_vid | end_vid | flow | residual_capacity

| $1 \mid$ | $1 \mid$ | $6 \mid$ | $5\|50\|$ |
| :---: | :---: | :---: | :---: |
| $2 \mid$ | $4 \mid$ | $7 \mid$ | $6\|50\|$ |
| $3 \mid$ | $10 \mid$ | $7 \mid$ | $8\|80\|$ |
| $4 \mid$ | $12 \mid$ | $8 \mid$ | $12\|80\|$ |
| $5 \mid$ | $8 \mid$ | $11 \mid$ | $7\|130\|$ |
| $6 \mid$ | $11 \mid$ | $11 \mid$ | $12\|130\|$ |
| $7 \mid$ | $9 \mid$ | $11 \mid$ | $16\|80\|$ |
| $8 \mid$ | $3 \mid$ | $15 \mid$ | $10\|80\|$ |
| $9 \mid$ | $16 \mid$ | $16 \mid$ | $15\|80\|$ |
| 90 |  |  |  |

(9 rows)

```

Many to One
```

pgr_boykovKolmogorov(Edges SQL, start vids, end vid)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertex \(\backslash(12 \backslash)\)
```

SELECT * FROM pgr boykovKolmogorov(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
ARRAY[11, 3, 17], 12)
seq | edge | start_vid | end_vid | flow | residual_capacity

| $1 \mid$ | $7 \mid$ | $3 \mid$ | $7\|50\|$ |
| ---: | ---: | :---: | ---: |
| $2\|10\|$ | $7 \mid$ | $8\|100\|$ | 0 |
| $3\|12\|$ | $8 \mid$ | $12\|100\|$ | 0 |
| $4\|8\|$ | $11 \mid$ | $7\|50\|$ | 80 |
| $5\|11\|$ | $11 \mid$ | $12\|130\|$ | 0 |

(5 rows)

```
```

pgr_boykovKolmogorov(Edges SQL, start vids, end vids)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr_boykovKolmogorov(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
ARRAY[11, 3, 17], ARRAY[5, 10, 12]);
seq | edge | start_vid | end_vid | flow | residual_capacity

| 1\| | 71 | 31 | 7\|50| | 0 |
| :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | 1 \| | 61 | 5\|50| | 80 |
| 31 | 4 \| | 71 | 6\| 50 | | 0 |
| 4 \| | 10\| | 71 | 8\| $100 \mid$ | 30 |
| 51 | 12\| | 81 | 12\| $100 \mid$ | 0 |
| 6 | 8\| | 11\| | 7\|100| | 30 |
| 71 | 11\| | 11\| | 12\| 130 | | 0 |
| 81 | 91 | 11\| | 16\| 80| | 50 |
| 91 | 31 | 151 | 10\| 80 | | 50 |
| $10 \mid$ | $16 \mid$ | $16 \mid$ | 15\| 80 | | 0 |

(10 rows)

```

\section*{Combinations}
```

pgr_boykovKolmogorov(Edges SQL, Combinations SQL)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)
OR EMPTY SET

```

\section*{Example:}

Using a combinations table, equivalent to calculating result from vertices \(\backslash(\backslash\{5,6 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,15,14 \backslash\} \backslash)\).
The combinations table:
```

SELECT source, target FROM combinations
WHERE target NOT IN (5, 6);
source | target
5| 10
6| 15
6| 14
(3 rows)

```

The query:
```

SELECT * FROM pgr_boykovKolmogorov(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
'SELECT * FROM combinations WHERE target NOT IN (5, 6)');
seq | edge | start_vid | end_vid | flow | residual_capacity

| $1\|4\|$ | $6 \mid$ | $7\|80\|$ | 20 |
| ---: | ---: | ---: | ---: |
| $2\|8\|$ | $7 \mid$ | $11\|80\|$ | 20 |
| $3\|9\|$ | $11 \mid$ | $16\|80\|$ | 50 |
| $4\|16\|$ | $16 \mid$ | $15\|80\|$ | 0 |
| (4 rows) |  |  |  |

```

\section*{Parameters}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline capacity & ANY-INTEGER & & Weight of the edge (source, target) \\
\hline reverse_capacity & ANY-INTEGER & -1 & Weight of the edge (target, source) \\
& & & \begin{tabular}{l} 
When negative: edge (target, source) does not exist, therefore it's
\end{tabular} \\
& & & not part of the graph.
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the departure vertex. \\
\hline target & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the arrival vertex. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Result Columns}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INT & Sequential value starting from \(\mathbf{1}\). \\
\hline edge & BIGINT & Identifier of the edge in the original query (edges_sql). \\
\hline start_vid & BIGINT & Identifier of the first end point vertex of the edge. \\
\hline end_vid & BIGINT & Identifier of the second end point vertex of the edge. \\
\hline flow & BIGINT & Flow through the edge in the direction (start_vid, end_vid). \\
\hline residual_capacity & BIGINT & \begin{tabular}{l} 
Residual capacity of the edge in the direction (start_vid, \\
end_vid).
\end{tabular} \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Manually assigned vertex combinations.
```

SELECT * FROM pgr_boykovKolmogorov(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
'SELECT * FROM (VALUES (5, 10), (6, 15), (6, 14)) AS t(source, target)');
seq | edge | start_vid | end_vid | flow | residual_capacity

| $1 \mid$ | $4 \mid$ | $6 \mid$ | $7\|80\|$ | 20 |
| :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $8 \mid$ | $7 \mid$ | $11\|80\|$ | 20 |
| $3 \mid$ | $9 \mid$ | $11 \mid$ | $16\|80\|$ | 50 |
| $4\|16\|$ | $16 \mid$ | $15\|80\|$ | 0 |  |

(4 rows)

```

\section*{See Also}
```

- Flow - Family of functions
    - pgr_edmondsKarp
    - pgr_pushRelabel
- https://www.boost.org/libs/graph/doc/boykov_kolmogorov_max_flow.html

```

Supported versions: Latest (3.3) 3.23 .13 .0
Unsupported versions: 2.6 2.5 2.42 .3
pgr_edmondsKarp
pgr_edmondsKarp - Calculates the flow on the graph edges that maximizes the flow from the sources to the targets using Edmonds Karp Algorithm.

Boost Graph Inside

\section*{Availability}
- Version 3.2.0
- New proposed signature - pgr_edmondsKarp (Combinations)
- Version 3.0.0
- Official function
- Version 2.5.0
- Renamed from pgr_maxFlowEdmondsKarp
- Proposed function
- Version 2.3.0
- New Experimental function

\section*{Description}

\section*{The main characteristics are:}
- The graph is directed.
- Process is done only on edges with positive capacities.
- When the maximum flow is 0 then there is no flow andEMPTY SET is returned.
- There is no flow when a source is the same as a target.
- Any duplicated value in the source(s) or target(s) are ignored.
- Calculates the flow/residual capacity for each edge. In the output
- Edges with zero flow are omitted.
- Creates a super source and edges to all the source(s), and asuper target and the edges from all the targets(s).
- The maximum flow through the graph is guaranteed to be the value returned bypgr_maxFlow when executed with the same parameters and can be calculated:
- By aggregation of the outgoing flow from the sources
- By aggregation of the incoming flow to the targets
- Running time: \(\backslash\left(O\left(V * E^{\wedge} 2\right) \backslash\right)\)

\section*{Signatures}

\section*{Summary}
```

pgr_edmondsKarp(Edges SQL, start vid, end vid)
pgr_edmondsKarp(Edges SQL, start vid, end vids)
pgr_edmondsKarp(Edges SQL, start vids, end vid)
pgr_edmondsKarp(Edges SQL, start vids, end vids)
pgr_edmondsKarp(Edges SQL, Combinations SQL)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)
OR EMPTY SET

```

\section*{One to One}

RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity) OR EMPTY SET

\section*{Example:}

From vertex \\(11) to vertex \\(12\\)

SELECT * FROM pgr_edmondsKarp(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
11, 12);
seq | edge | start_vid | end_vid | flow | residual_capacity
\begin{tabular}{rrrr}
\(----+-----+---------------+-----+-------------100 \mid\) & 30 \\
\(1|10|\) & \(7 \mid\) & \(8|100|\) & 0 \\
\(2|12|\) & \(8 \mid\) & \(12|100|\) & 30 \\
\(3|8|\) & \(11 \mid\) & \(7|100|\) & 0 \\
\(4|11|\) & \(11 \mid\) & \(12|130|\) & \\
(4 rows)
\end{tabular}

\section*{One to Many}

\section*{pgr_edmondsKarp(Edges SQL, start vid, end vids) \\ RETURNS SET OF (seq, edge, start vid, end vid, flow, residual capacity) OR EMPTY SET}

\section*{Example:}

From vertex \\(11\\) to vertices \\(\\{5,10,12\\}\\)
```

SELECT * FROM pgr_edmondsKarp(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
11, ARRAY[5, 10, 12]);
seq | edge | start_vid | end_vid | flow | residual_capacity
1| 1| 6| 5| 50| ---------------------------------------
2| 4| 7| 6| 50|
3| 10| 7| 8| 80| 50
4| 12| 8| 12| 80| 20
5| 8|
7| 9| 11| 16| 80| 50
8| 3| 15| 10| 80| 50
9| 16| 16| 15| 80| 0
(9 rows)

```

\section*{Many to One}
```

pgr_edmondsKarp(Edges SQL, start vids, end vid)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertex \\(12\\)
```

SELECT * FROM pgr_edmondsKarp(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
ARRAY[11, 3, 17], 12);
seq | edge | start_vid | end_vid | flow | residual_capacity

1| 7| 3| $7 |$| $7 \mid$ | 0 |
| :--- | :--- | :--- |

$2|10| \quad 7|\quad 8| 100 \mid \quad 30$

```

```

$4|8| \quad 11|\quad 7| 50 \mid \quad 80$
5| 11
(5 rows)

```

Many to Many
```

pgr_edmondsKarp(Edges SQL, start vids, end vids)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)

```

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr_edmondsKarp(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
ARRAY[11, 3, 17], ARRAY[5, 10, 12]);
seq | edge | start_vid | end_vid | flow | residual_capacity
1| 7| 3| 7| 50|
2| 1| 6| 5| 50| 80
3| 4| 7| 6| 50| 0
4| 10| 7| 8| 100| 30
5| 12| 8| 12| 100| 0
6| 8| 11| 7| 100| 30
7| 11| 11| 12| 130| 0
8| 9| 11| 16| 80| 50
9| 3| 15| 10| 80| 50
10| 16| 16| 15| 80| 0
(10 rows)

```

\section*{Combinations}
pgr_edmondsKarp(Edges SQL, Combinations SQL)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity) OR EMPTY SET

\section*{Example:}

Using a combinations table, equivalent to calculating result from vertices \(\backslash(\backslash\{5,6 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,15,14 \backslash\} \backslash)\).
The combinations table:
```

SELECT source, target FROM combinations
WHERE target NOT IN (5, 6);
source | target
5| 10
6| 15
6| 14
(3 rows)

```

The query:
```

SELECT * FROM pgr_edmondsKarp(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
'SELECT * FROM combinations WHERE target NOT IN (5, 6)');
seq | edge | start_vid | end_vid | flow | residual_capacity

| $1 \mid$ | $4 \mid$ | $6 \mid$ | $7\|80\|$ | 20 |
| ---: | ---: | ---: | ---: | ---: |
| $2 \mid$ | $8 \mid$ | $7 \mid$ | $11\|80\|$ | 20 |
| $3 \mid$ | $9 \mid$ | $11 \mid$ | $16\|80\|$ | 50 |
| $4 \mid$ | $16 \mid$ | $16 \mid$ | $15\|80\|$ | 0 |

(4 rows)

```

\section*{Parameter}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}

\section*{Inner Queries}

Edges SQL
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline id & ANY-INTEGER & & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & & Identifier of the second end point vertex of the edge. \\
\hline capacity & ANY-INTEGER & & Weight of the edge (source, target) \\
\hline reverse_capacity & ANY-INTEGER & -1 & Weight of the edge (target, source) \\
& & & \begin{tabular}{l} 
When negative: edge (target, source) does not exist, therefore it's \\
\\
\end{tabular} \\
& & & not part of the graph.
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & ANY- & Identifier of the departure vertex. \\
& INTEGER & \\
\hline target & ANY- & Identifier of the arrival vertex. \\
& INTEGER & \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Result Columns}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INT & Sequential value starting from \(\mathbf{1 .}\) \\
\hline edge & BIGINT & Identifier of the edge in the original query (edges_sql). \\
\hline start_vid & BIGINT & Identifier of the first end point vertex of the edge. \\
\hline end_vid & BIGINT & Identifier of the second end point vertex of the edge. \\
\hline flow & BIGINT & Flow through the edge in the direction istart_vid, end_vid). \\
\hline residual_capacity & BIGINT & \begin{tabular}{l} 
Residual capacity of the edge in the direction (start_vid, \\
end_vid).
\end{tabular} \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Manually assigned vertex combinations.
```

SELECT * FROM pgr_edmondsKarp(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
'SELECT * FROM (VALUES (5, 10), (6, 15), (6, 14)) AS t(source, target)');
seq | edge | start_vid | end_vid | flow | residual_capacity
1| 4| 6| 7| 80| -------------------
2| 8| 7| 11| 80| 20
3| 9| 11| 16| 80| 50
4| 16| 16| 15| 80| 0
(4 rows)

```

\section*{See Also}
- Flow - Family of functions
- pgr_boykovKolmogorov
- pgr_pushRelabel
- https://www.boost.org/libs/graph/doc/edmonds_karp_max_flow.html
- https://en.wikipedia.org/wiki/Edmonds\%E2\%80\%93Karp_algorithm

\section*{Indices and tables}
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Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.6 2.5 2.42 .3
pgr_pushRelabel
pgr_pushRelabel - Calculates the flow on the graph edges that maximizes the flow from the sources to the targets using Push Relabel Algorithm.

\section*{boost}

Boost Graph Inside

\section*{Availability}
- Version 3.2.0
- New proposed signature - pgr_pushRelabel (Combinations)
- Version 3.0.0
- Official function
- Version 2.5.0
- Renamed from pgr_maxFlowPushRelabel
- Proposed function
- Version 2.3.0
- New Experimental function

\section*{Description}

\section*{The main characteristics are:}
- The graph is directed.
- Process is done only on edges with positive capacities.
- When the maximum flow is 0 then there is no flow andEMPTY SET is returned.
- There is no flow when a source is the same as atarget.
- Any duplicated value in the source(s) or target(s) are ignored.
- Calculates the flow/residual capacity for each edge. In the output
- Edges with zero flow are omitted.
- Creates a super source and edges to all the source(s), and asuper target and the edges from all the targets(s).
- The maximum flow through the graph is guaranteed to be the value returned bypgr_maxFlow when executed with the same parameters and can be calculated:
- By aggregation of the outgoing flow from the sources
- By aggregation of the incoming flow to the targets
- Running time: \(\left.\backslash\left(O^{\wedge} \vee^{\wedge} 3\right) \backslash\right)\)

\section*{Signatures}

\section*{Summary}
```

pgr_pushRelabel(Edges SQL, start vid, end vid)
pgr_pushRelabel(Edges SQL, start vid, end vids)
pgr_pushRelabel(Edges SQL, start vids, end vid)
pgr_pushRelabel(Edges SQL, start vids, end vids)
pgr_pushRelabel(Edges SQL, Combinations SQL)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)
OR EMPTY SET

```

\section*{One to One}
```

pgr_pushRelabel(Edges SQL, start vid, end vid)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)

```

\section*{Example:}

From vertex \\(11\\) to vertex \\(12\\)
```

SELECT * FROM pgr_pushRelabel(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
11, 12);
seq | edge | start_vid | end_vid | flow | residual_capacity

| $1 \mid$ | $10 \mid$ | $7 \mid$ | $8\|100\|$ |
| ---: | ---: | ---: | ---: |
| $2 \mid$ | $12 \mid$ | $8 \mid$ | $12\|100\|$ |
| $3 \mid$ | $8 \mid$ | $11 \mid$ | $7\|100\|$ |
| $4 \mid$ | $11 \mid$ | $11 \mid$ | $12\|130\|$ |
|  | 0 |  |  |

(4 rows)

```

\section*{One to Many}
```

pgr_pushRelabel(Edges SQL, start vid, end vids)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)

```
OR EMPTY SET

\section*{Example:}

From vertex \(\backslash(11 \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr_pushRelabel(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
11, ARRAY[5, 10, 12])
seq | edge | start_vid | end_vid | flow | residual_capacity

| $1 \mid$ | $6 \mid$ | $1 \mid$ | $3 \mid$ |
| :---: | :---: | :---: | :---: |
| $2 \mid$ | $60 \mid$ | 0 |  |
| $3 \mid$ | $7 \mid$ | $3 \mid$ | $7 \mid$ |
| $2 \mid$ | $50 \mid$ | 50 |  |
| $4 \mid$ | $1 \mid$ | $6 \mid$ | $5 \mid$ |
| $5 \mid$ | $7 \mid$ | $7 \mid$ | $3 \mid$ |
| $6 \mid$ | $4 \mid$ | $7 \mid$ | $6\|3\|$ |
| $7\|10\|$ | $7 \mid$ | $8\|100\|$ | 100 |
| $8\|12\|$ | $8 \mid$ | $12\|100\|$ | 80 |
| $9 \mid$ | $8 \mid$ | $11 \mid$ | $7\|130\|$ |
| $10 \mid$ | $11 \mid$ | $11 \mid$ | $12\|130\|$ |
| $11 \mid$ | $9 \mid$ | $11 \mid$ | $16\|80\|$ |
| $12 \mid$ | $3 \mid$ | $15 \mid$ | $10\|80\|$ |
| $13 \mid$ | $16 \mid$ | $16 \mid$ | $15\|80\|$ |

(13 rows)

```

\section*{Many to One}
```

pgr_pushRelabel(Edges SQL, start vids, end vid)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertex \(\backslash(12 \backslash)\)
```

SELECT * FROM pgr_pushRelabel(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
ARRAY[11, 3, 17], 12)
seq | edge | start_vid | end_vid | flow | residual_capacity

| $1 \mid$ | $10 \mid$ | $7 \mid$ | $8\|100\|$ |
| ---: | ---: | ---: | ---: |
| $2\|12\|$ | $8 \mid$ | $12\|100\|$ | 30 |
| $3 \mid$ | $12 \mid$ | $11 \mid$ | $7\|100\|$ |
| $4\|11\|$ | $11 \mid$ | $12\|130\|$ | 30 |
| 4 rows) |  |  | 0 |

```

Many to Many

RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity) OR EMPTY SET

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr_pushRelabel(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
ARRAY[11, 3, 17], ARRAY[5, 10, 12]);
seq | edge | start_vid | end_vid | flow | residual_capacity

| 1\| | 71 | 31 | 7\| 20| | 30 |
| :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | 1 \| | 61 | 5\|50| | 80 |
| 3\| | 4 \| | 71 | 6\| 50| | 0 |
| 4\| | 10\| | 71 | 8\|100| | 30 |
| 51 | 12\| | 8 | 12\| $100 \mid$ | 0 |
| 61 | 8 \| | 11\| | 7\|130| | 0 |
| 7\| | 11\| | 11\| | 12\| 130 | | 0 |
| 8\| | 91 | 11\| | 16\| 80| | 50 |
| 91 | 31 | 15 | 10\| $80 \mid$ | 50 |
| $10 \mid$ | $16 \mid$ | 16 \| | 15\| $80 \mid$ | 0 |

(10 rows)

```

\section*{Combinations}
```

pgr_pushRelabel(Edges SQL, Combinations SQL)
RETURNS SET OF (seq, edge, start_vid, end_vid, flow, residual_capacity)
OR EMPTY SET

```

\section*{Example:}

Using a combinations table, equivalent to calculating result from vertices \(\backslash(\backslash\{5,6 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,15,14 \backslash\} \backslash)\).
The combinations table:
```

SELECT source, target FROM combinations
WHERE target NOT IN (5,6);
source | target
5| 10
6| 15
6| 14
(3 rows)

```

The query:
```

SELECT * FROM pgr_pushRelabel(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
'SELECT * FROM combinations WHERE target NOT IN (5, 6)');
seq | edge | start_vid | end_vid | flow | residual_capacity

| 1\| | 4\| | 6\| | 7\| | 80\| | 20 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | 8। | 71 | 11\| | $80 \mid$ | 20 |
| 31 | 11\| | 11\| | 12 \| | 50\| | 80 |
| 41 | 9 | 11\| | $16 \mid$ | 30\| | 100 |
| 51 | 13\| | 12\| | 17\| |  | 50 |
| 61 | 16 | $16 \mid$ | 15 | $80 \mid$ | 0 |
| 71 | 15\| | 17\| | 16 \| |  | 0 |

(7 rows)

```

Parameters
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline capacity & ANY-INTEGER & & Weight of the edge (source, target) \\
\hline reverse_capacity & ANY-INTEGER & -1 & Weight of the edge (target, source) \\
& & & \begin{tabular}{l} 
When negative: edge (target, source) does not exist, therefore it's
\end{tabular} \\
& & & not part of the graph.
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the departure vertex. \\
\hline target & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the arrival vertex. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Result Columns}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INT & Sequential value starting from \(\mathbf{1}\). \\
\hline edge & BIGINT & Identifier of the edge in the original query (edges_sql). \\
\hline start_vid & BIGINT & Identifier of the first end point vertex of the edge. \\
\hline end_vid & BIGINT & Identifier of the second end point vertex of the edge. \\
\hline flow & BIGINT & Flow through the edge in the direction (start_vid, end_vid). \\
\hline residual_capacity & BIGINT & \begin{tabular}{l} 
Residual capacity of the edge in the direction (start_vid, \\
end_vid).
\end{tabular} \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Manually assigned vertex combinations.
```

SELECT * FROM pgr_pushRelabel(
'SELECT id, source, target, capacity, reverse_capacity
FROM edges',
'SELECT * FROM (VALUES (5, 10), (6, 15), (6, 14)) AS t(source, target)');
seq | edge | start_vid | end_vid | flow | residual_capacity

| $1 \mid$ | $4 \mid$ | $6 \mid$ | $7\|80\|$ |
| ---: | ---: | ---: | ---: |
| $2\|8\|$ | $7 \mid$ | $11\|80\|$ | 20 |
| $3 \mid$ | $11 \mid$ | $11 \mid$ | $12\|50\|$ |
| $4 \mid$ | $9 \mid$ | $11 \mid$ | $16\|30\|$ |
| $5\|13\|$ | $12 \mid$ | $17\|50\|$ | 80 |
| $6\|16\|$ | $16 \mid$ | $15\|80\|$ | 50 |
| $7\|15\|$ | $17 \mid$ | $16\|50\|$ | 0 |
| (7 rows) |  |  | 0 |

```

\section*{See Also}

Flow - Family of functions
- pgr_boykovKolmogorov
- pgr_edmondsKarp
https://www.boost.org/libs/graph/doc/push_relabel_max_flow.html
https://en.wikipedia.org/wiki/Push\%E2\%80\%93relabel_maximum_flow_algorithm

\section*{Indices and tables}
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- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.6 2.5 2.42 .3
pgr_edgeDisjointPaths
pgr_edgeDisjointPaths - Calculates edge disjoint paths between two groups of vertices.

\section*{boost}

Boost Graph Inside

\section*{Availability}
```

- Version 3.2.0

```
    - New proposed function:
        - pgr_edgeDisjointPaths(Combinations)
- Version 3.0.0
    - Official function
- Version 2.5.0
    - Proposed function
- Version 2.3.0
    - New Experimental function

\section*{Description}

Calculates the edge disjoint paths between two groups of vertices. Utilizes underlying maximum flow algorithms to calculate the paths.

The main characterics are:
- Calculates the edge disjoint paths between any two groups of vertices.
- Returns EMPTY SET when source and destination are the same, or cannot be reached.
- The graph can be directed or undirected.
- Uses pgr_boykovKolmogorov to calculate the paths.

\section*{Signatures}

\section*{Summary}
pgr_edgeDisjointPaths(Edges SQL, start vid, end vid, [directed])
pgr_edgeDisjointPaths(Edges SQL, start vid, end vids, [directed])
pgr_edgeDisjointPaths(Edges SQL, start vids, end vid, [directed])
pgr_edgeDisjointPaths(Edges SQL, start vids, end vids, [directed])
pgr_edgeDisjointPaths(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (seq, path_id, path_seq, [start_vid,] [end_vid,] node, edge, cost, agg_cost) OR EMPTY SET

\section*{One to One}
```

pgr_edgeDisjointPaths(Edges SQL, start vid, end vid, [directed])
RETURNS SET OF (seq, path_id, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertex \\(11\\) to vertex \\(12\\)

\section*{SELECT * FROM pgr_edgeDisjointPaths(}
'SELECT id, source, target, cost, reverse_cost
FROM edges',
11, 12);
seq | path_id | path_seq | node | edge | cost | agg_cost
\begin{tabular}{rlllll}
\(1 \mid\) & \(1 \mid\) & \(1 \mid\) & \(11 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(1 \mid\) & \(2 \mid\) & \(7 \mid\) & \(10 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(1 \mid\) & \(3 \mid\) & \(8 \mid\) & \(12 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(1 \mid\) & \(4 \mid\) & \(12 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(5 \mid\) & \(2 \mid\) & \(1 \mid\) & \(11 \mid\) & \(11 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(2 \mid\) & \(2 \mid\) & \(12 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
(6 rows)
\end{tabular}

One to Many
```

pgr_edgeDisjointPaths(Edges SQL, start vid, end vids, [directed])
RETURNS SET OF (seq, path_id, path_seq, end_vid, node, edge, cost, agg_cost)

``` OR EMPTY SET

\section*{Example:}

From vertex \(\backslash(11 \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr edgeDisjointPaths(
'SELECT id, source, target, cost, reverse_cost
FROM edges',
11, ARRAY[5, 10, 12])
seq | path_id | path_seq | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $1 \mid$ | $5 \mid$ | $11 \mid$ | $8 \mid$ | $1 \mid$ | 0 |
| ---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $2 \mid$ | $1 \mid$ | $2 \mid$ | $5 \mid$ | $7 \mid$ | $4 \mid$ | $1 \mid$ | 1 |
| $3 \mid$ | $1 \mid$ | $3 \mid$ | $5 \mid$ | $6 \mid$ | $1 \mid$ | $1 \mid$ | 2 |
| $4 \mid$ | $1 \mid$ | $4 \mid$ | $5 \mid$ | $5 \mid$ | $-1 \mid$ | $0 \mid$ | 3 |
| $5 \mid$ | $2 \mid$ | $1 \mid$ | $10 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 0 |
| $6 \mid$ | $2 \mid$ | $2 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 1 |
| $7 \mid$ | $2 \mid$ | $3 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ | 2 |
| $8 \mid$ | $2 \mid$ | $4 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 3 |
| $9 \mid$ | $3 \mid$ | $1 \mid$ | $12 \mid$ | $11 \mid$ | $8 \mid$ | $1 \mid$ | 0 |
| $10 \mid$ | $3 \mid$ | $2 \mid$ | $12 \mid$ | $7 \mid$ | $10 \mid$ | $1 \mid$ | 1 |
| $11 \mid$ | $3 \mid$ | $3 \mid$ | $12 \mid$ | $8 \mid$ | $12 \mid$ | $1 \mid$ | 2 |
| $12 \mid$ | $3 \mid$ | $4 \mid$ | $12 \mid$ | $12 \mid$ | $-1 \mid$ | $0 \mid$ | 3 |
| $13 \mid$ | $4 \mid$ | $1 \mid$ | $12 \mid$ | $11 \mid$ | $11 \mid$ | $1 \mid$ | 0 |
| $14 \mid$ | $4 \mid$ | $2 \mid$ | $12 \mid$ | $12 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |
| $(14$ rows $)$ |  |  |  |  |  |  |  |

```

Many to One
```

pgr_edgeDisjointPaths(Edges SQL, start vids, end vid, [directed])
RETURNS SET OF (seq, path_id, path_seq, start_vid, node, edge, cost, agg_cost)

``` OR EMPTY SET

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertex \(\backslash(12 \backslash)\)

pgr_edgeDisjointPaths(Edges SQL, start vids, end vids, [directed])
RETURNS SET OF (seq, path_id, path_seq, start_vid, end_vid, node, edge, cost, agg_cost) OR EMPTY SET

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr_edgeDisjointPaths(
'SELECT id, source, target, cost, reverse_cost
FROM edges',
ARRAY[11, 3, 17], ARRAY[5, 10, 12]);
seq | path_id | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1\| | 1 \| | 1\| | 31 | 5 | 31 | 7 \| 1 | 1 \| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 21 | 1 \| | $2 \mid$ | 3\| | 51 | 71 | 4\| 1 | 1 \| | 1 |
| 31 | 1 \| | 31 | 3\| | 5 | 61 | 1\| 1 | 1 \| | 2 |
| 4 \| | 1 \| | 4 \| | 3\| | 51 | 51 | -1\| 0 | 01 | 3 |
| 51 | 21 | 1\| | 3\| | $10 \mid$ | 31 | 7\| | 1\| | 0 |
| 61 | 21 | $2 \mid$ | 31 | $10 \mid$ | 71 | 8\| | 1\| | 1 |
| 71 | 2 \| | 31 | 3\| | $10 \mid$ | 11\| | 91 | 1\| | 2 |
| 81 | 21 | 4 \| | 3\| | $10 \mid$ | 16\| | 16\| | 1 \| | 3 |
| 91 | 21 | 51 | 3\| | $10 \mid$ | 15\| | 31 | $1 \mid$ | 4 |
| $10 \mid$ | 21 | 61 | 3\| | 10\| | $10 \mid$ | -1\| | 01 | 5 |
| 11\| | 31 | 1 \| | 31 | $12 \mid$ | 31 | 71 | 1 \| | 0 |
| $12 \mid$ | 31 | $2 \mid$ | 31 | $12 \mid$ | 71 | 8। | 1\| | 1 |
| 131 | 31 | 31 | 3\| | $12 \mid$ | 11\| | 11\| | 1 \| | 2 |
| $14 \mid$ | 31 | $4 \mid$ | 3\| | $12 \mid$ | $12 \mid$ | -1\| | 01 | 3 |
| 151 | $4 \mid$ | 1 \| | 11\| | 5\| | 11\| | 8 \| | 1\| | 0 |
| $16 \mid$ | 4 \| | $2 \mid$ | 11\| | 5\| | 7\| | 4\| | 1 \| | 1 |
| $17 \mid$ | 4 \| | 31 | 11\| | 5\| | 61 | 1\| | 1\| | 2 |
| 18\| | 4 \| | $4 \mid$ | 11\| | 51 | 51 | -1\| | 01 | 3 |
| 19\| | 51 | $1 \mid$ | 11\| | 10\| | 11 | 91 | 1\| | 0 |
| $20 \mid$ | 51 | 21 | 11\| | 10\| | 16 | $16 \mid$ | \| 1 | | 1 |
| 21\| | 51 | 31 | 11\| | 10\| | 15 | 31 | 1\| | 2 |
| $22 \mid$ | 51 | 4 \| | 11\| | 10\| | 10 | \| -1 | | 01 | 3 |
| 231 | 61 | 1 \| | 11\| | $12 \mid$ | 11 | 8\| | 1\| | 0 |
| $24 \mid$ | 61 | $2 \mid$ | 11\| | 12\| | 71 | 10\| | 1\| | 1 |
| 251 | 61 | 31 | 11\| | $12 \mid$ | 8 | 12\| | 1\| | 2 |
| $26 \mid$ | 61 | 41 | 11\| | 12\| | 12 | \| -1| | 01 | 3 |
| $27 \mid$ | 7\| | 1\| | 11\| | $12 \mid$ | 11\| | 11\| | 1\| | 0 |
| 281 | 7\| | 21 | 11\| | $12 \mid$ | 12 | -1\| | $0 \mid$ |  |
| 291 | 8\| | 1\| | 17\| | 5\| | 17\| | 15\| | 1\| | 0 |
| 301 | 8\| | $2 \mid$ | 17\| | 5\| | $16 \mid$ | 16\| | 1\| | 1 |
| 31\| | 8\| | 31 | 17\| | 5\| | 15\| | 31 | 1\| | 2 |
| 321 | 8\| | 41 | 17\| | 51 | 101 | 21 | 1\| | 3 |
| 331 | 8\| | 51 | 17\| | 5\| | $6 \mid$ | 1\| | 1\| | 4 |
| 341 | 8\| | 61 | 17\| | 51 | 5 | -1\| | 01 | 5 |
| 351 | 91 | 1\| | 17\| | $10 \mid$ | $17 \mid$ | \| 15| | 1\| | 0 |
| 361 | 91 | 21 | 17\| | $10 \mid$ | 16 | \| 16| | 1\| | 1 |
| 371 | 91 | 31 | 17\| | $10 \mid$ | 15 | 3\| | 1\| | 2 |
| 381 | 91 | 41 | 17\| | $10 \mid$ | 10 | \| -1 | | 01 | 3 |
| 391 | $10 \mid$ | 1 | $17 \mid$ | 12 \| | 17 | \| 15| | \| 1| | 0 |
| 401 | $10 \mid$ | 2 | 17 \| | 12 \| | 16 | \| 9| | \| 1| | 1 |
| 41\| | $10 \mid$ | 3 | $17 \mid$ | $12 \mid$ | 11 | \| 11| | \| 1| | 2 |
| $42 \mid$ | $10 \mid$ | 4 | 17 \| | 12 | 12 | $\|-1\|$ | \| 0 | | 3 |

```

\section*{Combinations}
```

pgr_edgeDisjointPaths(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (seq, path_id, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Using a combinations table, equivalent to calculating result from vertices \(\backslash(\backslash\{5,6 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,15,14 \backslash\} \backslash)\) on \(a n\) undirected graph.

The combinations table:

The query:
```

SELECT * FROM pgr_edgeDisjointPaths(
'SELECT id, source, target, cost, reverse_cost
FROM edges',
'SELECT * FROM combinations WHERE target NOT IN (5, 6)',
directed => false);
seq | path_id | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $1 \mid$ | $5 \mid$ | $10 \mid$ | $5 \mid$ | $1 \mid$ | $1 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $1 \mid$ | $2 \mid$ | $5 \mid$ | $10 \mid$ | $6 \mid$ | $2 \mid$ | $-1 \mid$ | 1 |
| $3 \mid$ | $1 \mid$ | $3 \mid$ | $5 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 0 |
| $4 \mid$ | $2 \mid$ | $1 \mid$ | $6 \mid$ | $15 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| $5 \mid$ | $2 \mid$ | $2 \mid$ | $6 \mid$ | $15 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 1 |
| $6 \mid$ | $2 \mid$ | $3 \mid$ | $6 \mid$ | $15 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 2 |
| $7 \mid$ | $2 \mid$ | $4 \mid$ | $6 \mid$ | $15 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 3 |
| $8 \mid$ | $2 \mid$ | $5 \mid$ | $6 \mid$ | $15 \mid$ | $15 \mid$ | $-1 \mid$ | $0 \mid$ | 4 |
| $9 \mid$ | $3 \mid$ | $1 \mid$ | $6 \mid$ | $15 \mid$ | $6\|\mid$ | $2 \mid$ | $-1 \mid$ | 0 |
| $10 \mid$ | $3 \mid$ | $2 \mid$ | $6 \mid$ | $15 \mid$ | $10 \mid$ | $3 \mid$ | $-1 \mid$ | -1 |
| $11 \mid$ | $3 \mid$ | $3 \mid$ | $6 \mid$ | $15 \mid$ | $15 \mid$ | $-1 \mid$ | $0 \mid$ | -2 |

```

\section*{Parameters}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}

Optional parameters


Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& &
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Combinations SQL
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the departure vertex. \\
\hline target & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the arrival vertex. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Result Column}

Set of (seq, path_id, path_seq [, start_vid] [, end_vid], node, edge, cost, agg_cost)
\begin{tabular}{|c|c|c|}
\hline Column & Type & Description \\
\hline seq & INTEGER & Sequential value starting from 1. \\
\hline path_id & INTEGER & \begin{tabular}{l}
Path identifier. \\
- Has value \(\mathbf{1}\) for the first of a path fromstart_vid to end_vid.
\end{tabular} \\
\hline path_seq & INTEGER & Relative position in the path. Has value1 for the beginning of a path. \\
\hline start_vid & BIGINT & Identifier of the starting vertex. Returned when multiple starting vetrices are in the query.
Many to One
Many to Many
Combinations \\
\hline end_vid & BIGINT & Identifier of the ending vertex. Returned when multiple ending vertices are in the query.
One to Many
Many to Many
Combinations \\
\hline node & BIGINT & Identifier of the node in the path fromstart_vid to end_vid. \\
\hline edge & BIGINT & Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for the last node of the path. \\
\hline cost & FLOAT & Cost to traverse from node using edge to the next node in the path sequence. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Manually assigned vertex combinations on an undirected graph.
```

SELECT * FROM pgr_edgeDisjointPaths(
'SELECT id, source, target, cost, reverse_cost
FROM edges',
'SELECT * FROM (VALUES (5, 10), (6, 15), (6, 14)) AS t(source, target)',
directed => false);
seq | path_id | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

```


See Also
- Flow - Family of functions

\section*{ndices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: \(2.6 \mathbf{2}\).5 2.42 .3
pgr_maxCardinalityMatch
pgr_maxCardinalityMatch - Calculates a maximum cardinality matching in a graph.

\section*{boost}

Boost Graph Inside

\section*{Availability}
- Version 3.3.3
- directed optional parameter ignored and removed from the documentation as the algorithm works only on undirected graphs
- Version 3.0.0
- Official function
- Version 2.5.0
- Renamed from pgr_maximumCardinalityMatching
- Proposed function
- Version 2.3.0
- New Experimental function

\section*{Description}

The main characteristics are:
- Works for undirected graphs.
- A matching or independent edge set in a graph is a set of edges without common vertices.
- A maximum matching is a matching that contains the largest possible number of edges.
- There may be many maximum matchings.
- Calculates one possible maximum cardinality matching in a graph.
- Running time: \(\backslash\left(O\left(E^{*} V\right.\right.\) * \(\backslash\) alpha \(\left.\left.(E, V)\right) \backslash\right)\)
- \(\backslash(\backslash a l p h a(E, V) \backslash)\) is the inverse of theAckermann function.

\section*{Signatures}
```

pgr_maxCardinalityMatch(Edges SQL)
RETURNS SET OF (seq, edge_id, source, target)
OR EMPTY SET

```

\section*{Example:}

Using all edges.
```

SELECT * FROM pgr_maxCardinalityMatch(
'SELECT id, source, target, cost AS going, reverse_cost AS coming
FROM edges');
seq | edge | source | target

| $1 \mid$ | $6 \mid$ | $1 \mid$ | 3 |
| :---: | :---: | :---: | :---: |
| $2 \mid$ | $17 \mid$ | $2 \mid$ | 4 |
| $3 \mid$ | $1 \mid$ | $5 \mid$ | 6 |
| $4 \mid$ | $14 \mid$ | $8 \mid$ | 9 |
| $5 \mid$ | $3 \mid$ | $10 \mid$ | 15 |
| $6 \mid$ | $9 \mid$ | $11 \mid$ | 16 |
| $7 \mid$ | $13 \mid$ | $12 \mid$ | 17 |
| $8 \mid$ | $18 \mid$ | $13 \mid$ | 14 |
| (8 rows) |  |  |  |

```

Edges SQL

SQL query, which should return a set of rows with the following columns:
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline going & ANY-NUMERICAL & \begin{tabular}{l} 
A positive value represents the existence of the edge source, \\
target).
\end{tabular} \\
\hline coming & ANY-NUMERICAL & -1
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INT & Sequential value starting from \(\mathbf{1 .}\) \\
\hline edge & BIGINT & Identifier of the edge in the original query. \\
\hline source & BIGINT & Identifier of the first end point of the edge. \\
\hline target & BIGINT & \begin{tabular}{l} 
Identifier of the second end point of the \\
edge.
\end{tabular} \\
\hline
\end{tabular}

\section*{See Also}
- Flow - Family of functions
- https://www.boost.org/libs/graph/doc/maximum_matching.html
- https://en.wikipedia.org/wiki/Matching_\%28graph_theory\%29
- https://en.wikipedia.org/wiki/Ackermann_function

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
pgr_maxFlowMinCost - Experimental
pgr_maxFlowMinCost - Calculates the edges that minimizes the total cost of the maximum flow on a graph


Boost Graph Inside

\section*{Warning}

Possible server crash
- These functions might create a server crash

\section*{Warning}

Experimental functions
- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need c/c++ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting

\section*{Availability}
- Version 3.2.0
- New experimental function:
- pgr_maxFlowMinCost (Combinations)
- Version 3.0.0
- New experimental function

\section*{Description}

\section*{The main characteristics are:}
- The graph is directed.
- Process is done only on edges with positive capacities.
- When the maximum flow is 0 then there is no flow andEMPTY SET is returned.
- There is no flow when a source is the same as a target.
- Any duplicated value in the source(s) or target(s) are ignored.
- Calculates the flow/residual capacity for each edge. In the output
- Edges with zero flow are omitted.
- Creates a super source and edges to all the source(s), and asuper target and the edges from all the targets(s).
- The maximum flow through the graph is guaranteed to be the value returned bypgr_maxFlow when executed with the same parameters and can be calculated:
- By aggregation of the outgoing flow from the sources
- By aggregation of the incoming flow to the targets
- TODO check which statement is true:
- The cost value of all input edges must be nonnegative.
- Process is done when the cost value of all input edges is nonnegative.
- Process is done on edges with nonnegative cost.
- Running time: \(\backslash(\mathrm{O}(\mathrm{U} *(\mathrm{E}+\mathrm{V} * \log \mathrm{~V})) \backslash)\)
- where \(\backslash(U \backslash)\) is the value of the max flow.
- \(\backslash(U \backslash)\) is upper bound on number of iterations. In many real world cases number of iterations is much smaller than(U\\).

\section*{Signatures}

\section*{Summary}
```

pgr_maxFlowMinCost(Edges SQL, start vid, end vid)
pgr_maxFlowMinCost(Edges SQL, start vid, end vids)
pgr_maxFlowMinCost(Edges SQL, start vids, end vid)
pgr_maxFlowMinCost(Edges SQL, start vids, end vids)
pgr_maxFlowMinCost(Edges SQL, Combinations SQL)
RETURNS SET OF (seq, edge, source, target, flow, residual_capacity, cost, agg_cost)
OR EMPTY SET

```

\section*{One to One}
```

pgr_maxFlowMinCost(Edges SQL, start vid, end vid)
RETURNS SET OF (seq, edge, source, target, flow, residual_capacity, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

\section*{SELECT * FROM pgr_maxFlowMinCost(}
'SELECT id, source, target, capacity, reverse_capacity, cost, reverse_cost
FROM edges',
11, 12);
seq | edge | source | target | flow | residual_capacity | cost | agg_cost
\begin{tabular}{rrrr|r}
\(1 \mid\) & \(10 \mid\) & \(7 \mid\) & \(8|100|\) & \(30|100|\) \\
\(2 \mid\) & \(12 \mid\) & \(8 \mid\) & \(12|100|\) & \(0|100|\) \\
\(3 \mid\) & \(8 \mid\) & \(11 \mid\) & \(7|100|\) & \(30|100|\) \\
\(4 \mid\) & \(11 \mid\) & \(11 \mid\) & \(12|130|\) & \(0|130|\) \\
\begin{tabular}{rl}
4 \\
(4 rows)
\end{tabular} & & & & 430
\end{tabular}
(4 rows)

\section*{One to Many}

\section*{pgr_maxFlowMinCost(Edges SQL, start vid, end vids) \\ RETURNS SET OF (seq, edge, source, target, flow, residual_capacity, cost, agg_cost) OR EMPTY SET}

\section*{Example:}

From vertex \\(11\\) to vertices \\(\\{5,10,12\\}\\)
```

SELECT * FROM pgr_maxFlowMinCost(
'SELECT id, source, target, capacity, reverse capacity, cost, reverse cost
FROM edges',
11, ARRAY[5, 10, 12]);
seq | edge | source | target | flow | residual_capacity | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $6 \mid$ | $5 \mid$ | $30 \mid$ | $100\|30\|$ | 30 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $2 \mid$ | $4 \mid$ | $7 \mid$ | $6 \mid$ | $30 \mid$ | $20\|30\|$ | 60 |

    3| 10| 7| 8| 100| 30| 100| 160
    4| 12| 8| 12| 100| 0| 100| 260
    5| 8| 11| 7| 130| 0| 130| 390
    6| 11| 11| 12| 130| 0| 130| 520
    7| 9| 11| 16| 80| 50| 80| 600
    8| 3| 15| 10| 80| 50| 80| 680
    9| 16| 16| 15| 80| 0| 80| 760
(9 rows)

```

\section*{Many to One}
```

pgr_maxFlowMinCost(Edges SQL, start vids, end vid)
RETURNS SET OF (seq, edge, source, target, flow, residual_capacity, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertex \(\backslash(12 \backslash)\)
```

SELECT * FROM pgr_maxFlowMinCost(
'SELECT id, source, target, capacity, reverse_capacity, cost, reverse_cost
FROM edges',
ARRAY[11, 3, 17], 12);
seq | edge | source | target | flow | residual_capacity | cost | agg_cost

| 1\| | 71 | 31 | 7\| 50| | 0\|50| | 50 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $10 \mid$ | 7\| | 8\|100| | 30\| 100| | 150 |
| $3 \mid$ | 12\| | 8\| | 12\| $100 \mid$ | 0\| 100| | 250 |
| 4 \| | 8\| | 11\| | 7\| 50| | 80\| $50 \mid$ | 300 |
| 51 | 11\| | 11\| | 12\| 130 | | 0\|130| | 430 |

(5 rows)

```

\section*{Many to Many}
```

pgr_maxFlowMinCost(Edges SQL, start vids, end vids)
RETURNS SET OF (seq, edge, source, target, flow, residual_capacity, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)

\section*{SELECT * FROM pgr_maxFlowMinCost(}
'SELECT id, source, target, capacity, reverse_capacity, cost, reverse_cos
FROM edges',
ARRAY[11, 3, 17], ARRAY[5, 10, 12]);
seq | edge | source | target | flow | residual_capacity | cost | agg_cost
\begin{tabular}{rcccccc}
\(1 \mid\) & \(7 \mid\) & \(3 \mid\) & \(7 \mid\) & \(50 \mid\) & \(0 \mid\) & \(50 \mid\) \\
\(2 \mid\) & \(1 \mid\) & \(6 \mid\) & \(5 \mid\) & \(50 \mid\) & \(80 \mid\) & \(50 \mid\) \\
\(3 \mid\) & \(4 \mid\) & \(7 \mid\) & \(6|50|\) & \(0 \mid\) & \(50 \mid\) & 150 \\
\(4 \mid\) & \(10 \mid\) & \(7 \mid\) & \(8|100|\) & \(30|100|\) & 250 \\
\(5 \mid\) & \(12 \mid\) & \(8 \mid\) & \(12|100|\) & \(0|100|\) & 350 \\
\(6 \mid\) & \(8 \mid\) & \(11 \mid\) & \(7|100|\) & \(30|100|\) & 450 \\
\(7 \mid\) & \(11 \mid\) & \(11 \mid\) & \(12|130|\) & \(0|130|\) & 580 \\
\(8 \mid\) & \(9 \mid\) & \(11 \mid\) & \(16|30|\) & \(100|30|\) & 610 \\
\(9 \mid\) & \(3 \mid\) & \(15 \mid\) & \(10 \mid\) & \(80 \mid\) & \(50|80|\) & 690 \\
\(10 \mid\) & \(16 \mid\) & \(16 \mid\) & \(15 \mid\) & \(80 \mid\) & \(0 \mid\) & \(80 \mid\) \\
\(11 \mid\) & \(15 \mid\) & \(17 \mid\) & \(16 \mid\) & \(50 \mid\) & \(0 \mid\) & \(70 \mid\) \\
10 & 820
\end{tabular}
(11 rows)

\section*{Combinations}
```

pgr maxFlowMinCost(Edges SQL, Combinations SQL)
RETURNS SET OF (seq, edge, source, target, flow, residual_capacity, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Using a combinations table, equivalent to calculating result from vertices \(\backslash(\backslash\{5,6 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,15,14 \backslash\} \backslash)\).
The combinations table:
```

SELECT source, target FROM combinations
WHERE target NOT IN (5, 6);
source | target
5| 10
6| 15
6| 14
(3 rows)

```

The query
```

SELECT * FROM pgr_maxFlowMinCost(
SELECT id, source, target, capacity, reverse_capacity, cost, reverse_cost
FROM edges',
'SELECT * FROM combinations WHERE target NOT IN (5, 6)');
seq | edge | source | target | flow | residual_capacity | cost | agg_cost
1| 4| 6| 7| 80| 20| 80| 80
2| 8| 7| 11| 80| 20| 80| 160
3| 9| 11| 16| 80|
4| 16| 16| 15| 80| 0| 80| 320
(4 rows)

```

\section*{Parameters}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}

\section*{Inner Querie}

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline Column & Type & Default & Description \\
\hline \multirow[t]{2}{*}{capacity} & ANY-INTEGER & & Capacity of the edge (source, target) \\
\hline & & & - When negative: edge (target, source) does not exist, therefore it's not part of the graph. \\
\hline \multirow[t]{2}{*}{reverse_capacity} & ANY-INTEGER & -1 & Capacity of the edge (target, source) \\
\hline & & & - When negative: edge (target, source) does not exist, therefore it's not part of the graph. \\
\hline cost & ANY-NUMERICAL & & Weight of the edge (source, target) if it exist \\
\hline reverse_cost & ANY-NUMERICAL & \\(-1\\) & Weight of the edge (target, source) if it exist \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & ANY- & Identifier of the departure vertex. \\
& INTEGER & \\
\hline target & ANY- & Identifier of the arrival vertex. \\
& INTEGER & \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Result Columns}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INT & Sequential value starting from \(\mathbf{1}\). \\
\hline edge & BIGINT & Identifier of the edge in the original query (edges_sql). \\
\hline source & BIGINT & Identifier of the first end point vertex of the edge. \\
\hline target & BIGINT & Identifier of the second end point vertex of the edge. \\
\hline flow & BIGINT & Flow through the edge in the direction (source, target). \\
\hline residual_capacity & BIGINT & Residual capacity of the edge in the direction (source, target). \\
\hline cost & FLOAT & \begin{tabular}{l} 
The cost of sending this flow through the edge in the direction (source, \\
\\
agg_cost
\end{tabular} \\
\hline & FLOAT & The aggregate cost.
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Manually assigned vertex combinations.
```

SELECT * FROM pgr_maxFlowMinCost(
'SELECT id, source, target, capacity, reverse_capacity, cost, reverse_cost
FROM edges',
'SELECT * FROM (VALUES (5, 10), (6, 15), (6, 14)) AS t(source, target)');
seq | edge | source | target | flow | residual_capacity | cost | agg_cost

| $1 \mid$ | $4 \mid$ | $6 \mid$ | $7\|80\|$ | $20\|80\|$ | 80 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $8 \mid$ | $7 \mid$ | $11 \mid$ | $80 \mid$ | $20 \mid$ |
| $3 \mid$ | $9 \mid$ | $11 \mid$ | $16 \mid$ | $80 \mid$ | $50 \mid$ |
| $4 \mid$ | $16 \mid$ | $16 \mid$ | $15 \mid$ | $80 \mid$ | 160 |
| 240 |  |  |  |  |  |

$4|16| 16|15| 80|\quad 0| 80 \mid 320$
(4 rows)

```

\section*{See Also}

\section*{- Flow - Family of functions}
- https://www.boost.org/libs/graph/doc/successive_shortest_path_nonnegative_weights.html

\section*{Indices and tables}
pgr_maxFlowMinCost_Cost - Calculates the minimum total cost of the maximum flow on a graph

\section*{boost}

Boost Graph Inside

\section*{Warning}

Possible server crash
- These functions might create a server crash

\section*{Warning}

Experimental functions
- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need \(c / c++\) coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting

\section*{Availability}
- Version 3.2.0
- New experimental function:
- pgr_maxFlowMinCost_Cost (Combinations)
- Version 3.0.0
- New experimental function

\section*{Description}

\section*{The main characteristics are:}
- The graph is directed.
- Process is done only on edges with positive capacities.
- When the maximum flow is 0 then there is no flow andEMPTY SET is returned.
- There is no flow when a source is the same as atarget.
- Any duplicated value in the source(s) or target(s) are ignored.
- Calculates the flow/residual capacity for each edge. In the output
- Edges with zero flow are omitted.
- Creates a super source and edges to all the source(s), and asuper target and the edges from all the targets(s).
- The maximum flow through the graph is guaranteed to be the value returned bypgr_maxFlow when executed with the same parameters and can be calculated:
- By aggregation of the outgoing flow from the sources
- By aggregation of the incoming flow to the targets

The main characteristics are:

The graph is directed.
- The cost value of all input edges must be nonnegative.
- When the maximum flow is 0 then there is no flow and0 is returned.
- There is no flow when a source is the same as a target.
- Any duplicated value in the source(s) or target(s) are ignored.
- Uses pgr_maxFlowMinCost - Experimental.
- Running time: \(\backslash(\mathrm{O}(\mathrm{U} *(\mathrm{E}+\mathrm{V}\) * \(\log \mathrm{V})) \backslash)\)
- where \(\backslash(U \backslash)\) is the value of the max flow.
- \(\backslash(U \backslash)\) is upper bound on number of iterations. In many real world cases number of iterations is much smaller than(U\\).

\section*{Signatures}

\section*{Summary}
```

pgr_maxFlowMinCost_Cost(Edges SQL, start vid, end vid)
pgr_maxFlowMinCost_Cost(Edges SQL, start vid, end vids)
pgr_maxFlowMinCost_Cost(Edges SQL, start vids, end vid)
pgr_maxFlowMinCost_Cost(Edges SQL, start vids, end vids)
pgr_maxFlowMinCost_Cost(Edges SQL, Combinations SQL)
RETURNS FLOAT

```

\section*{One to One}
```

pgr_maxFlowMinCost_Cost(Edges SQL, start vid, end vid)
RETURNS FLOAT

```

\section*{Example:}

From vertex \(\backslash(11 \backslash)\) to vertex \(\backslash(12 \backslash)\)
```

SELECT * FROM pgr_maxFlowMinCost_Cost(
'SELECT id, source, target, capacity, reverse_capacity, cost, reverse_cost
FROM edges',
11, 12);
pgr_maxflowmincost_cost
430
(1 row)

```

\section*{One to Many}

\section*{pgr_maxFlowMinCost_Cost(Edges SQL, start vid, end vids)}
```

RETURNS FLOAT

```

\section*{Example:}

From vertex \(\backslash(11 \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr_maxFlowMinCost_Cost(
'SELECT id, source, target, capacity, reverse_capacity, cost, reverse_cost
FROM edges',
ARRAY[11, 3, 17], 12);
pgr_maxflowmincost_cost
4 3 0
(1 row)

```

\section*{Many to One}
\[
\begin{aligned}
& \text { pgr_maxFlowMinCost_Cost(Edges SQL, start vids, end vid) } \\
& \text { RETURNS FLOAT }
\end{aligned}
\]

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertex \(\backslash(12 \backslash)\)
760
(1 row)

\section*{Many to Many}
```

pgr_maxFlowMinCost_Cost(Edges SQL, start vids, end vids)
RETURNS FLOAT

```

\section*{Example:}

From vertices \(\backslash(\backslash\{11,3,17 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{5,10,12 \backslash\} \backslash)\)
```

SELECT * FROM pgr_maxFlowMinCost_Cost(
'SELECT id, source, target, capacity, reverse_capacity, cost, reverse_cost
FROM edges',
ARRAY[11, 3, 17], ARRAY[5, 10, 12]);
pgr_maxflowmincost_cost
820
(1 row)

```

\section*{Combinations}
```

pgr_maxFlowMinCost_Cost(Edges SQL, Combinations SQL)
RETURNS FLOAT

```

\section*{Example:}

Using a combinations table, equivalent to calculating result from vertices \(\backslash(\backslash\{5,6 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,15,14 \backslash\} \backslash)\).
The combinations table:
```

SELECT source, target FROM combinations
WHERE target NOT IN (5, 6);
source | target
5| 10
6| 15
6| 14
(3 rows)

```

The query:
```

SELECT * FROM pgr_maxFlowMinCost_Cost(
SELECT id, source, target, capacity, reverse_capacity, cost, reverse_cost
FROM edges',
SELECT * FROM combinations WHERE target NOT IN (5, 6)');
pgr_maxflowmincost_cost
3 2 0
(1 row)

```

\section*{Parameter}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline end vids & ARRAY[BIGINT] & Array of identifiers of ending vertices. \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline Column & Type & Default & Description \\
\hline id & ANY-INTEGER & & Identifier of the edge. \\
\hline source & ANY-INTEGER & & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & & Identifier of the second end point vertex of the edge. \\
\hline capacity & ANY-INTEGER & & \begin{tabular}{l}
Capacity of the edge (source, target) \\
- When negative: edge (target, source) does not exist, therefore it's not part of the graph.
\end{tabular} \\
\hline reverse_capacity & ANY-INTEGER & -1 & \begin{tabular}{l}
Capacity of the edge (target, source) \\
- When negative: edge (target, source) does not exist, therefore it's not part of the graph.
\end{tabular} \\
\hline cost & ANY-NUMERICAL & & Weight of the edge (source, target) if it exist \\
\hline reverse_cost & ANY-NUMERICAL & \\(-1\\) & Weight of the edge (target, source) if it exist \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & ANY- & Identifier of the departure vertex. \\
& INTEGER & \\
\hline target & ANY- & Identifier of the arrival vertex. \\
& INTEGER & \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Resturn Columns}
\begin{tabular}{ll} 
Type & Description \\
\hline FLOAT & \begin{tabular}{l} 
Minimum Cost Maximum Flow possible from the source(s) to the \\
target(s)
\end{tabular}
\end{tabular}

\section*{Additional Examples}

\section*{Example:}

Manually assigned vertex combinations.
```

SELECT * FROM pgr_maxFlowMinCost_Cost(
'SELECT id, source, target, capacity, reverse_capacity, cost, reverse_cost
FROM edges',
'SELECT * FROM (VALUES (5, 10), (6, 15), (6, 14)) AS t(source, target)');
pgr_maxflowmincost_cost
3 2 0
(1 row)

```

See Also
- Flow - Family of functions
- https://www.boost.org/libs/graph/doc/successive_shortest_path_nonnegative_weights.html

\section*{Indices and tables}
- Index
- Search Page

Flow Functions General Information
- The graph is directed.
- Process is done only on edges with positive capacities.
- When the maximum flow is 0 then there is no flow andEMPTY SET is returned.
- There is no flow when a source is the same as a target.
- Any duplicated value in the source(s) or target(s) are ignored.
- Calculates the flow/residual capacity for each edge. In the output
- Edges with zero flow are omitted.
- Creates a super source and edges to all the source(s), and asuper target and the edges from all the targets(s).
- The maximum flow through the graph is guaranteed to be the value returned bypgr_maxFlow when executed with the same parameters and can be calculated:
- By aggregation of the outgoing flow from the sources
- By aggregation of the incoming flow to the targets
pgr_maxFlow is the maximum Flow and that maximum is guaranteed to be the same on the functionsgr_pushRelabel, pgr_edmondsKarp, pgr_boykovKolmogorov, but the actual flow through each edge may vary.

Inner Queries

Edges SQL

\section*{Capacity edges}
pgr_pushRelabel
- pgr_edmondsKarp
- pgr_boykovKolmogorov
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline id & ANY-INTEGER & & Identifier of the edge. \\
\hline source & ANY-INTEGER & & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & & Identifier of the second end point vertex of the edge. \\
\hline capacity & ANY-INTEGER & & Weight of the edge (source, target) \\
\hline reverse_capacity & ANY-INTEGER & -1 & Weight of the edge (target, source) \\
& & & When negative: edge (target, source) does not exist, therefore it's \\
& & & \\
& & & \\
& & &
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Capacity-Cost edges}
- pgr_maxFlowMinCost - Experimental
- pgr_maxFlowMinCost_Cost - Experimental
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline id & ANY-INTEGER & & Identifier of the edge.
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Cost edges}
- pgr_edgeDisjointPaths
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & \\
& & Weight of the edge (target, source) \\
& & \\
& & not part of the graph.
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the departure vertex. \\
\hline target & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the arrival vertex. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Result Columns}

Used in
- pgr_pushRelabel
- pgr_edmondsKarp
- pgr_boykovKolmogorov
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INT & Sequential value starting from \(\mathbf{1}\). \\
\hline edge & BIGINT & Identifier of the edge in the original query (edges_sql). \\
\hline start_vid & BIGINT & Identifier of the first end point vertex of the edge. \\
\hline end_vid & BIGINT & Identifier of the second end point vertex of the edge. \\
\hline flow & BIGINT & Flow through the edge in the direction istart_vid, end_vid). \\
\hline residual_capacity & BIGINT & \begin{tabular}{l} 
Residual capacity of the edge in the direction (start_vid, \\
end_vid).
\end{tabular} \\
\hline
\end{tabular}

For pgr_maxFlowMinCost - Experimental
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INT & Sequential value starting from \(\mathbf{1 .}\) \\
\hline edge & BIGINT & Identifier of the edge in the original query (edges_sql). \\
\hline source & BIGINT & Identifier of the first end point vertex of the edge. \\
\hline target & BIGINT & Identifier of the second end point vertex of the edge. \\
\hline flow & BIGINT & Flow through the edge in the direction (source, target). \\
\hline residual_capacity & BIGINT & Residual capacity of the edge in the direction (source, target). \\
\hline cost & FLOAT & \begin{tabular}{l} 
The cost of sending this flow through the edge in the direction (source, \\
\\
\hline agg_cost
\end{tabular} \\
\hline
\end{tabular}

\section*{Adcanced Documentation}

A flow network is a directed graph where each edge has a capacity and a flow. The flow through an edge must not exceed the
capacity of the edge. Additionally, the incoming and outgoing flow of a node must be equal except for source which only has outgoing flow, and the destination(sink) which only has incoming flow.

Maximum flow algorithms calculate the maximum flow through the graph and the flow of each edge.
The maximum flow through the graph is guaranteed to be the same with all implementations, but the actual flow through each edge may vary.

Given the following query:
pgr_maxFlow \\((edges\_sql, source\_vertex, sink \_vertex)\\)
where \(\backslash\left(\right.\) edges \(\backslash\) _sql \(=\backslash\left\{\left(i d \_i\right.\right.\), source_i, target_i, capacity_i, reverse \(\left.\left.\left.\backslash c a p a c i t y \_i\right) \backslash\right\} \backslash\right)\)

\section*{Graph definition}

The weighted directed graph, \(\backslash(G(V, E) \backslash)\), is defined as:
- the set of vertices \(\backslash(\mathrm{V} \backslash)\)
- \\(source\_vertex \cup sink\_vertex \bigcup source_i \bigcup target_i\\)
- the set of edges \(\backslash(\mathrm{E})\)
- \\(E = \begin\{cases\} \text\{ \} \\{(source_i, target_i, capacity_i) \text\{ when \} capacity > } 0 \text { \\} \& \quad \text\{ if \} }
 \(0 \backslash\} \& \backslash t e x t\} \backslash \backslash \operatorname{lcup} \backslash\{(\) target_i, source_i, reverse\_capacity_i) \text\{ when \} reverse\_capacity_i > 0) \\\(\& \quad } \backslash t e x t\{\) if \} reverse\_capacity \neq \varnothing \(\backslash \backslash\) \end } \{ \text { cases } \} \backslash \text { ) }

\section*{Maximum flow problem}

Given:
- \(\backslash(G(V, E) \backslash)\)
- \\(source\_vertex \in \(\mathrm{V} \backslash\) ) the source vertex
- \\(sink\_vertex \in \(\mathrm{V} \backslash\) ) the sink vertex

Then:
- \(\backslash\left(\right.\) pgr \(\backslash \_m a x F l o w(e d g e s \backslash s q l\), source, sink) \(=\) \boldsymbol \(\{\backslash\) Phi \(\} \backslash\) )
- \(\backslash\left(\backslash b o l d s y m b o l\{\backslash P h i\}=\left\{\left(i d \_i\right.\right.\right.\), edge\_id_i, source_i, target_i, flow_i, residual\_capacity_i) \(\} \backslash\) )

Where:
\(\backslash(\backslash\) boldsymbol \(\{\backslash \mathrm{Phi}\} \backslash)\) is a subset of the original edges with their residual capacity and flow. The maximum flow through the graph can be obtained by aggregating on the source or sink and summing the flow from/to it. In particular:
- \(\backslash\left(i d \_i=i \backslash\right)\)
- \\(edge\_id = id_i\\) in edges_sql
- \\(residual\_capacity_i = capacity_i - flow_i\\)

\section*{See Also}
- https://en.wikipedia.org/wiki/Maximum_flow_problem

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0

\section*{Kruskal - Family of functions}
- pgr_kruskal
- pgr_kruskalBFS
- pgr_kruskaIDD
- pgr_kruskaIDFS
pgr_kruskal
pgr_kruskal - Minimum spanning tree of a graph using Kruskal's algorithm.

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- New Official function

\section*{Description}

This algorithm finds the minimum spanning forest in a possibly disconnected graph using Kruskal's algorithm.

\section*{The main Characteristics are:}
- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.
- The total weight of all the edges in the tree or forest is minimized.
- Kruskal's running time: \(\backslash(\mathrm{O}(\mathrm{E} * \log \mathrm{E}) \backslash)\)
- EMPTY SET is returned when there are no edges in the graph.

Signatures

\section*{Summary}
```

pgr_kruskal(Edges SQL)
RETURNS SET OF (edge, cost)
OR EMPTY SET

```

\section*{Example:}

Minimum spanning forest
```

SELECT * FROM pgr_kruskal(
'SELECT id, source, target, cost, reverse_cost
FROM edges ORDER BY id'
) ORDER BY edge;
edge | cost
1| 1
2| 1
3:1
7| 1
10| 1
11| 1
12
13|
14 |
15|}
16| 1
17| 1
18| 1
(14 rows)

```

\section*{Parameters}
\begin{tabular}{lllll} 
Parameter & Type & Description \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below.
\end{tabular} & & \\
\hline
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & \\
\hline reverse_cost & ANY-NUMERICAL & -1
\end{tabular}
- When negative: edge (target, source) does not exist, therefore it's not part of the graph.

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns SET OF (edge, cost)
\begin{tabular}{llll} 
Column & Type & Description \\
\hline edge & BIGINT & Identifier of the edge. \\
\hline cost & FLOAT & \begin{tabular}{l} 
Cost to traverse the \\
edge.
\end{tabular} \\
\hline
\end{tabular}

\section*{See Also}
- Spanning Tree - Category
- Kruskal - Family of functions
- The queries use the Sample Data network.
- Boost: Kruskal's algorithm
- Wikipedia: Kruskal's algorithm

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
pgr_kruskalBFs
pgr_kruskalBFS - Kruskal's algorithm for Minimum Spanning Tree with breadth First Search ordering.

\section*{boost}

Boost Graph Inside

\section*{Availability}
```

- Version 3.0.0
- New Official function

```

Visits and extracts the nodes information in Breath First Search ordering of the Minimum Spanning Tree created using Kruskal's algorithm.

\section*{The main Characteristics are:}
- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.
- The total weight of all the edges in the tree or forest is minimized.
- Kruskal's running time: \\(O(E * log E)\\)
- Returned tree nodes from a root vertex are on Breath First Search order
- Breath First Search Running time: \(\backslash(O(E+V) \backslash)\)

\section*{Signatures}
```

pgr_kruskalBFS(Edges SQL, root vid, [max_depth])
pgr_kruskalBFS(Edges SQL, root vids, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Single vertex}
```

pgr_kruskalBFS(Edges SQL, root vid, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Example:}

The Minimum Spanning Tree having as root vertex \(\backslash(6 \backslash)\)
```

SELECT * FROM pgr_kruskalBFS(
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
6);
seq | depth | start_vid | node | edge | cost | agg_cost
1| 0| 6| 6| -1| 0| 0
2| 1| 6| 5| 1| 1| 1

```

```

4| 2| 6| 15| 3| 1| 2
5| 3| 6| 16| 16| 1| 3
6| 4| 6| 17| 15| 1| 4
5| 6| 12| 13| 1| 5
6| 6| 11| 11| 1| 6
9| 6| 6| 8| 12| 1| 6
10| 7| 6| 7| 10| 1| 7
11| 7| 6| 9| 14| 1| 7
12| 8| 6| 3| 7| 1| 8
13| 9| 6| 1| 6| 1| 9
(13 rows)

```

\section*{Multiple vertices}
```

pgr_kruskalBFS(Edges SQL, root vids, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Example:}

The Minimum Spanning Tree starting on vertices \(\backslash \backslash\{\{9,6 \backslash\} \backslash)\) with \(\backslash(\) depth \(\backslash\) leq 3\\)

\section*{SELECT * FROM pgr_kruskalBFS(}
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
ARRAY[9, 6], max_depth => 3);
seq | depth | start vid | node | edge | cost | agg cost
\begin{tabular}{ccccccc}
\(1 \mid\) & \(0 \mid\) & \(6 \mid\) & \(6 \mid\) & \(-1 \mid\) & \(0 \mid\) & 0 \\
\(2 \mid\) & \(1 \mid\) & \(6 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(1 \mid\) & \(6 \mid\) & \(10 \mid\) & \(2 \mid\) & \(1 \mid\) & 1 \\
\(4 \mid\) & \(2 \mid\) & \(6 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) & 2 \\
\(5 \mid\) & \(3 \mid\) & \(6 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) & 3 \\
\(6 \mid\) & \(0 \mid\) & \(9 \mid\) & \(9 \mid\) & \(-1 \mid\) & \(0 \mid\) & 0 \\
\(7 \mid\) & \(1 \mid\) & \(9 \mid\) & \(8 \mid\) & \(14 \mid\) & \(1 \mid\) & 1 \\
\(8 \mid\) & \(2 \mid\) & \(9 \mid\) & \(7 \mid\) & \(10 \mid\) & \(1 \mid\) & 2 \\
\(9 \mid\) & \(2 \mid\) & \(9 \mid\) & \(12 \mid\) & \(12 \mid\) & \(1 \mid\) & 2 \\
\(10 \mid\) & \(3 \mid\) & \(9 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) & 3 \\
\(11 \mid\) & \(3 \mid\) & \(9|11|\) & \(11 \mid\) & \(1 \mid\) & 3 \\
\(12 \mid\) & \(3 \mid\) & \(9 \mid\) & \(17 \mid\) & \(13 \mid\) & \(1 \mid\) & 3 \\
\((12\) rows \()\) & & & & &
\end{tabular}

Parameters
\begin{tabular}{|c|c|c|}
\hline Parameter & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below. \\
\hline \multirow[t]{2}{*}{root vid} & BIGINT & Identifier of the root vertex of the tree. \\
\hline & & When value is \(\backslash(0 \backslash)\) then gets the spanning forest starting in aleatory nodes for each tree in the forest. \\
\hline \multirow[t]{3}{*}{root vids} & ARRAY [ ANY-INTEGER ] & Array of identifiers of the root vertices. \\
\hline & & - \\(0\\) values are ignored \\
\hline & & - For optimization purposes, any duplicated value is ignored. \\
\hline
\end{tabular}

Where:
ANY-INTEGER:
SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

BFS optional parameters
\begin{tabular}{lll} 
Parameter & Type Default & Description \\
\hline max_depth & BIGINT \(\backslash(9223372036854775807 \backslash)\) & Upper limit of the depth of the tree. \\
& & \begin{tabular}{l} 
When negative throws an \\
error.
\end{tabular}
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& & \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

\section*{SMALLINT, INTEGER, BIGINT}

ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)
\begin{tabular}{|c|c|c|}
\hline Parameter & Type & Description \\
\hline seq & BIGINT & Sequential value starting from \\(1). \\
\hline \multirow[t]{2}{*}{depth} & BIGINT & Depth of the node. \\
\hline & & - \(\backslash(0 \backslash)\) when node = start_vid. \\
\hline start_vid & BIGINT & Identifier of the root vertex. \\
\hline node & BIGINT & Identifier of node reached using edge. \\
\hline \multirow[t]{2}{*}{edge} & BIGINT & Identifier of the edge used to arrive to node. \\
\hline & & - \(\backslash(-1 \backslash)\) when node \(=\) start_vid. \\
\hline cost & FLOAT & Cost to traverse edge. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

See Also
- Spanning Tree-Category
- Kruskal - Family of functions
- Sample Data
- Boost: Kruskal's algorithm
- Wikipedia: Kruskal's algorithm

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) \(3.2 \mathbf{3 . 1} 3.0\)
pgr_kruskaldD
pgr_kruskalDD - Catchament nodes using Kruskal's algorithm.

\section*{8boost}

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- New Official function

\section*{Description}

Using Kruskal's algorithm, extracts the nodes that have aggregate costs less than or equal to adistance from a root vertex (or vertices) within the calculated minimum spanning tree.

\section*{The main Characteristics are:}
- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.
- The total weight of all the edges in the tree or forest is minimized.
- Kruskal's running time: \(\backslash(\mathrm{O}(\mathrm{E} * \log \mathrm{E}) \backslash)\)
- Extracts all the nodes that have costs less than or equal to the value distance.
- The edges extracted will conform to the corresponding spanning tree.
- Edge \(\backslash((u, v) \backslash)\) will not be included when:
- The distance from the root to \(\backslash(u \backslash)>\) limit distance.
- The distance from the root to \(\backslash(v \backslash)>\) limit distance.
- No new nodes are created on the graph, so when is within the limit and is not within the limit, the edge is not included.
- Returned tree nodes from a root vertex are on Depth First Search order.
- Depth First Search running time: \(\backslash(O(E+V) \backslash)\)

\section*{Signatures}
```

pgr_kruskaIDD(Edges SQL, root vid, distance)
pgr_kruskalDD(Edges SQL, root vids, distance)
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Single vertex}
```

pgr_kruskalDD(Edges SQL, root vid, distance)
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Example:}

The Minimum Spanning Tree starting on vertex \\(6\\) with \\(distance \leq 3.5\\)
```

SELECT * FROM pgr_kruskaIDD(
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
6,3.5);
seq | depth | start_vid | node | edge | cost | agg_cost
1| 0| 6| 6| -1| 0| 0
2| 1| 6| 5| 1| 1| 1
3| 1| 6| 10| 2| 1| 1
4| 2| 6| 15| 3| 1| 2
5| 3| 6| 16| 16| 1| 3
(5 rows)

```

\section*{Multiple vertices}
```

pgr_kruskalDD(Edges SQL, root vids, distance)
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Example:}

The Minimum Spanning Tree starting on vertices \(\backslash(\backslash\{9,6 \backslash\} \backslash)\) with \(\backslash(\) distance \(\backslash\) leq \(3.5 \backslash)\)
```

SELECT * FROM pgr_kruskaIDD(
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
ARRAY[9, 6], 3.5);
seq | depth | start vid | node | edge | cost | agg cost

| $1 \mid$ | $0 \mid$ | $6 \mid$ | $6 \mid$ | $-1 \mid$ | $0 \mid$ |
| ---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $1 \mid$ | $6 \mid$ | $5 \mid$ | $1 \mid$ | $1 \mid$ |
| $3 \mid$ | $1 \mid$ | $6 \mid$ | $10 \mid$ | $2 \mid$ | $1 \mid$ |
| $4 \mid$ | $2 \mid$ | $6 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ |
| $5 \mid$ | $3 \mid$ | $6 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ |
| $6 \mid$ | $0 \mid$ | $9 \mid$ | $9 \mid$ | $-1 \mid$ | $0 \mid$ |
| $7 \mid$ | $1 \mid$ | $9 \mid$ | $8 \mid$ | $14 \mid$ | $1 \mid$ |
| $8 \mid$ | $2 \mid$ | $9 \mid$ | $7 \mid$ | $10 \mid$ | $1 \mid$ |
| $9 \mid$ | $3 \mid$ | $9 \mid$ | $3 \mid$ | $7 \mid$ | $1 \mid$ |
| $10 \mid$ | $2 \mid$ | $9 \mid$ | $12 \mid$ | $12 \mid$ | $1 \mid$ |
| $11 \mid$ | $3 \mid$ | $9 \mid$ | $11 \mid$ | $11 \mid$ | $1 \mid$ |
| $12 \mid$ | $3 \mid$ | $9 \mid$ | $17 \mid$ | $13 \mid$ | $1 \mid$ |
| $(12$ rows) |  |  |  | 3 |  |

```

\section*{Parameters}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below. \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|}
\hline Parameter & Type & Description \\
\hline Root vid & BIGINT & Identifier of the root vertex of the tree. \\
\hline Root vids & ARRAY[ANY-INTEGER] & \begin{tabular}{l}
Array of identifiers of the root vertices.
\\(O) values are ignored \\
- For optimization purposes, any duplicated value is ignored.
\end{tabular} \\
\hline distance & FLOAT & Upper limit for the inclusion of a node in the result. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Inner Queries}

Edges SQL
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline id & ANY-INTEGER & & Identifier of the edge. \\
\hline source & ANY-INTEGER & & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL & -1 & \begin{tabular}{l} 
Weight of the edge (target, source)
\end{tabular} \\
& & & \begin{tabular}{l} 
When negative: edge (target, source) does not exist, therefore it's \\
not part of the graph.
\end{tabular} \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline seq & BIGINT & Sequential value starting from \(\backslash(1 \backslash)\). \\
\hline depth & BIGINT & \begin{tabular}{l} 
Depth of the node. \\
\\
\end{tabular} \\
\hline Btart_vid & BIGINT & Identifier of the root vertex. \\
\hline node & BIGINT & Identifier of node reached using edge. \\
\hline edge & BIGINT & \begin{tabular}{l} 
Identifier of the edge used to arrive to \\
node.
\end{tabular} \\
\hline cost & & \begin{tabular}{l} 
FLOAT
\end{tabular} \\
\hline agg_cost & FLOAT & Cost to traverse edge. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

See Also
- Spanning Tree - Category
- Kruskal - Family of functions
- Sample Data
- Boost: Kruskal's algorithm
- Wikipedia: Kruskal's algorithm

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
pgr_kruskalDFS
pgr_kruskalDFS - Kruskal's algorithm for Minimum Spanning Tree with Depth First Search ordering.

\section*{boost}

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- New Official function

\section*{Description}

Visits and extracts the nodes information in Depth First Search ordering of the Minimum Spanning Tree created using Kruskal's algorithm.

\section*{The main Characteristics are:}
- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.
- The total weight of all the edges in the tree or forest is minimized.
- Kruskal's running time: \(\backslash(\mathrm{O}(\mathrm{E} * \log \mathrm{E}) \backslash)\)
- Returned tree nodes from a root vertex are on Depth First Search order
- Depth First Search Running time: \\( \(O(E+V) \backslash)\)

\section*{Signatures}
pgr_kruskalDFS(Edges SQL, root vid, [max_depth])
pgr_kruskalDFS(Edges SQL, root vids, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

\section*{Single vertex}
```

pgr_kruskalDFS(Edges SQL, root vid, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Example:}

The Minimum Spanning Tree having as root vertex \(\backslash(6 \backslash)\)
```

SELECT * FROM pgr_kruskalDFS

```
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
6);
seq | depth | start_vid | node | edge | cost | agg_cost
\(1|0| \quad 6|6|-1|0| \quad 0\)
\(2|1| \quad 6|5| 1|1| \quad 1\)
\(3|1|-6|10| 2|1|-1\)
\(4|2| \quad 6|15| 3|1|\)
\(5|3| \quad 6|16| 16|1| \quad 3\)
| \(4|\quad 6| 17|15| 1 \mid \quad 4\)

\(6|-6| 12|13| 1 \mid-5\)

\(\begin{array}{lllll}10 \mid & 6 \mid & 7 \mid & 10 \mid & 1 \mid \\ 7\end{array}\)
\(\left.\begin{array}{lllll}\mid & 8 \mid & 6 \mid & 3 \mid & 7 \mid \\ 1\end{array} \right\rvert\,\)

\(\begin{array}{lllll}13 \mid & 7 \mid & 6 \mid & 9 \mid & 14 \mid \\ 13 \mid & 7\end{array}\)
(13 rows)

\section*{Multiple vertices}
pgr_kruskalDFS(Edges SQL, root vids, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

\section*{Example:}

The Minimum Spanning Tree starting on vertices \(\backslash(\backslash\{9,6 \backslash\} \backslash)\) with \(\backslash(\) depth \(\backslash\) leq \(3 \backslash)\)
```

SELECT * FROM pgr_kruskalDFS
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
ARRAY[9, 6], max_depth => 3);
seq | depth | start_vid | node | edge | cost | agg_cost
1| 0|--------------------------------------
1| 6| 5| 1| 1| 1
1| 6| 10| 2| 1| 1
2| 6| 15| 3| 1| 2
3| 6| 16| 16| 1| 3
0| 9| 9| -1| 0| 0
1-9| 8| 14| 1| 1
2| 9| 7| 10| 1| 2
3| 9| 3| 7| 1| 3
2| 9| 12| 12| 1| 2
3| 9| 11| 11| 1| 3
12| 3| 9| 17| 13| 1| 3
(12 rows)

```

Parameters
\begin{tabular}{|c|c|c|}
\hline Parameter & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below. \\
\hline \multirow[t]{2}{*}{root vid} & BIGINT & Identifier of the root vertex of the tree. \\
\hline & & When value is \(\backslash(O \backslash)\) then gets the spanning forest starting in aleatory nodes for each tree in the forest. \\
\hline \multirow[t]{2}{*}{root vids} & ARRAY [ ANY-INTEGER ] & Array of identifiers of the root vertices. \\
\hline & & \begin{tabular}{l}
- \(\backslash(0 \backslash)\) values are ignored \\
- For optimization purposes, any duplicated value is ignored.
\end{tabular} \\
\hline
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

\section*{DFS optional parameters}
Parameter Type Default Description
max_depth BIGINT \\(9223372036854775807\\) Upper limit of the depth of the tree.
- When negative throws an error.

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)
\begin{tabular}{|c|c|c|}
\hline Parameter & Type & Description \\
\hline seq & BIGINT & Sequential value starting from \\(1). \\
\hline \multirow[t]{2}{*}{depth} & BIGINT & Depth of the node. \\
\hline & & - \(\backslash(0 \backslash)\) when node = start_vid. \\
\hline start_vid & BIGINT & Identifier of the root vertex. \\
\hline node & BIGINT & Identifier of node reached using edge. \\
\hline \multirow[t]{2}{*}{edge} & BIGINT & Identifier of the edge used to arrive to node. \\
\hline & & - \(\backslash(-1 \backslash)\) when node = start_vid. \\
\hline cost & FLOAT & Cost to traverse edge. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERIC:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

See Also
- Spanning Tree - Category
- Kruskal - Family of functions
- Sample Data
- Boost: Kruskal's algorithm
- Wikipedia: Kruskal's algorithm

\section*{Indices and tables}
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- Search Page

\section*{Description}

Kruskal's algorithm is a greedy minimum spanning tree algorithm that in each cycle finds and adds the edge of the least possible weight that connects any two trees in the forest.

\section*{The main Characteristics are:}
- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.
- The total weight of all the edges in the tree or forest is minimized.
- Kruskal's running time: \(\backslash(\mathrm{O}(\mathrm{E} * \log \mathrm{E}) \backslash)\)

Inner Queries
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

See Also
- Spanning Tree - Category
- Boost: Kruskal's algorithm
- Wikipedia: Kruskal's algorithm

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0

Prim - Family of functions
- pgr_prim
- pgr_primBFS
- pgr_primDD
- pgr_primDFS

Boost Graph Inside
- Supported versions: Latest (3.3) 3.2 3.13 .0
pgr_prim
pgr_prim - Minimum spanning forest of a graph using Prim's algorithm.

Boost Graph Inside

\section*{Availability}

\section*{Description}

This algorithm finds the minimum spanning forest in a possibly disconnected graph using Prim's algorithm.

\section*{The main characteristics are:}
- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.
- Prim's running time: \(\backslash(\mathrm{O}(\mathrm{E} * \log \mathrm{~V}) \backslash)\)
- EMPTY SET is returned when there are no edges in the graph.

Signatures

\section*{Summary}
```

pgr_prim(Edges SQL)
RETURNS SET OF (edge, cost)
OR EMPTY SET

```

\section*{Example:}

Minimum spanning forest of a subgraph
```

SELECT edge, cost FROM pgr_prim(
'SELECT id, source, target, cost, reverse_cost
FROM edges WHERE id < 14'
ORDER BY edge
edge | cost
|
|
|
4| 1
| 1
7
8| 1
9| 1
10| }
12| 1
13| 1
(11 rows)

```

Parameters
\begin{tabular}{lllll} 
Parameter & Type & Description & \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below.
\end{tabular} & as
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{llll} 
Column & Type & Default & Description \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL & -1 & Weight of the edge (target, source)
\end{tabular}
- When negative: edge (target, source) does not exist, therefore it's not part of the graph.

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Column}

Returns SET OF (edge, cost)
\begin{tabular}{llll} 
Column & Type & Description \\
\hline edge & BIGINT & Identifier of the edge. \\
\hline cost & FLOAT & \begin{tabular}{l} 
Cost to traverse the \\
edge.
\end{tabular} \\
\hline
\end{tabular}

\section*{See Also}
- Spanning Tree - Category
- Prim - Family of functions
- The queries use the Sample Data network.
- Boost: Prim's algorithm documentation
- Wikipedia: Prim's algorithm

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
pgr_primBFS
pgr_primBFS — Prim's algorithm for Minimum Spanning Tree with Depth First Search ordering.

\section*{boost}

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- New Official function

\section*{Description}

Visits and extracts the nodes information in Breath First Search ordering of the Minimum Spanning Tree created with Prims's algorithm.

\section*{The main Characteristics are:}
- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.
- Prim's running time: \(\backslash(\mathrm{O}(\mathrm{E} * \log \mathrm{~V}) \backslash)\)
- Returned tree nodes from a root vertex are on Breath First Search order
- Breath First Search Running time: \(\backslash(\mathrm{O}(\mathrm{E}+\mathrm{V}) \backslash)\)

\section*{Signatures}

RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

\section*{Single vertex}
```

pgr_primBFS(Edges SQL, root vid, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Example:}

The Minimum Spanning Tree having as root vertex \\(6\\)
```

SELECT * FROM pgr_primBFS(
SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
6);
seq | depth | start_vid | node | edge | cost | agg_cost
1| 0| 6| 6|-1| 0| 0
2| 1| 6| 5| 1| 1| 1
3| 1| 6| 10| 2| 1| 1
4| 1| 6| 7| 4| 1| 1
5| 2| 6| 15| 3| 1| 2
6| 2| 6| 11| 5| 1| 2
7| 2| 6| 3| 7| 1| 2
8| 2| 6| 8| 10| 1| 2
9| 3| 6| 16| 9| 1| 3
10| 3| 6| 12| 11| 1| 3
11| 3| 6| 1| 6| 1| 3
12| 3| 6| 9| 14| 1| 3
13| 4| 6| 17| 13| 1| 4
(13 rows)

```

\section*{Multiple vertices}
```

pgr_primBFS(Edges SQL, root vids, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Example:}

The Minimum Spanning Tree starting on vertices \(\backslash(\backslash\{9,6 \backslash\} \backslash)\) with \(\backslash(\) depth \(\backslash\) leq \(3 \backslash)\)
```

SELECT * FROM pgr_primBFS(
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
ARRAY[9, 6], max depth => 3);
seq | depth | start_vid | node | edge | cost | agg_cost
1| 0| 6| 6|-1| 0| 0
2| 1| 6| 5| 1| 1| 1
3| 1| 6| 10| 2| 1| 1
1| 6| 7| 4| 1| 1
2| 6| 15| 3| 1| 2
2| 6| 11| 5| 1| 2
2| 6| 3| 7| 1| 2
2| 6| 8| 10| 1| 2
3| 6| 16| 9| 1| 3
3| 6| 12| 11| 1| 3
10| 6| 1| 6| 1| 3
3| 6| 9| 14| 1| 3
0| 9| 9| -1| 0| 0
1| 9| 8| 14| 1| 1
15| 2| 9| 7| 10| 1| 2
16| 3| 9| 6| 4| 1| 3
17| 3| 9| 3| 7| 1| 3
(17 rows)

```

Parameters
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below. \\
\hline root vid & BIGINT & Identifier of the root vertex of the tree. \\
& & \begin{tabular}{l} 
When value is \(\backslash(0 \backslash)\) then gets the spanning forest starting in aleatory nodes \\
\\
\\
\\
\end{tabular} \\
&
\end{tabular}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline root vids & ARRAY [ANY-INTEGER ] & Array of identifiers of the root vertices. \\
& & \begin{tabular}{l} 
- \((0 \backslash)\) values are ignored \\
\\
\end{tabular} \\
& & For optimization purposes, any duplicated value is ignored.
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

BFS optional parameters
\begin{tabular}{lll} 
Parameter & Type & Default \\
\hline max_depth & BIGINT \(\backslash(9223372036854775807 \backslash)\) & Upescription limit of the depth of the tree. \\
& & \begin{tabular}{l} 
When negative throws an \\
error.
\end{tabular}
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& & \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)
\begin{tabular}{|c|c|c|}
\hline Parameter & Type & Description \\
\hline seq & BIGINT & Sequential value starting from \(\backslash(1 \backslash)\). \\
\hline \multirow[t]{2}{*}{depth} & BIGINT & Depth of the node. \\
\hline & & - \(\backslash(0 \backslash)\) when node = start_vid. \\
\hline start_vid & BIGINT & Identifier of the root vertex. \\
\hline node & BIGINT & Identifier of node reached using edge. \\
\hline \multirow[t]{2}{*}{edge} & BIGINT & Identifier of the edge used to arrive to node. \\
\hline & & - \\(-1\\) when node = start_vid. \\
\hline cost & FLOAT & Cost to traverse edge. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

Spanning Tree - Category
- Prim - Family of functions
- The queries use the Sample Data network.
- Boost: Prim's algorithm documentation
- Wikipedia: Prim's algorithm

\section*{Indices and tables}

Index
- Search Page
. Supported versions: Latest (3.3) 3.2 3.1 3.0
pgr_primDD
pgr_primDD - Catchament nodes using Prim's algorithm.

\section*{boost}

Boost Graph Inside

\section*{Availability}
- Version 3.0.0
- New Official function

\section*{Description}

Using Prim's algorithm, extracts the nodes that have aggregate costs less than or equal to adistance from a root vertex (or vertices) within the calculated minimum spanning tree.

\section*{The main Characteristics are:}
- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.
- Prim's running time: \(\backslash(\mathrm{O}(\mathrm{E} * \log \mathrm{~V}) \backslash)\)
- Extracts all the nodes that have costs less than or equal to the value distance.
- The edges extracted will conform to the corresponding spanning tree.
- Edge \(\backslash((u, v) \backslash)\) will not be included when:
- The distance from the root to \(\backslash(u \backslash)>\) limit distance.
- The distance from the root to \(\backslash(v \backslash)>\) limit distance.
- No new nodes are created on the graph, so when is within the limit and is not within the limit, the edge is not included.
- Returned tree nodes from a root vertex are on Depth First Search order.
- Depth First Search running time: \(\backslash(O(E+V) \backslash)\)

Signatures
```

pgr_primDD(Edges SQL, root vid, distance)
pgr_primDD(Edges SQL, root vids, distance)
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Single vertex}
```

pgr_primDD(Edges SQL, root vid, distance)
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Example:}

The Minimum Spanning Tree starting on vertex \\(6\\) with \\(distance \leq 3.5\\)

\section*{SELECT * FROM pgr_primDD(}
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
6,3.5);
seq | depth | start_vid | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline 1 | & 0 | & 6 & 6 & -1| & \(0 \mid\) & 0 \\
\hline \(2 \mid\) & 1| & \(6 \mid\) & 5| & 1| & 1 | & 1 \\
\hline 31 & 1| & 61 & 10 & 21 & 1 | & 1 \\
\hline \(4 \mid\) & 2 | & \(6 \mid\) & 15 & 3 & \(1 \mid\) & 2 \\
\hline 51 & 21 & \(6 \mid\) & 11| & 51 & 1 | & 2 \\
\hline 61 & 31 & 61 & 16 | & 91 & 1 | & 3 \\
\hline 71 & 3| & 6 & \(12 \mid\) & 11| & 1 | & 3 \\
\hline 8| & 1| & \(6 \mid\) & 7| & 4| & 1 | & 1 \\
\hline 91 & 21 & 61 & 31 & 7| & 1| & 2 \\
\hline \(10 \mid\) & 3 & 6 & 1 & 61 & 1| & 3 \\
\hline 11| & 2 & 6 & 81 & 10| & \(1 \mid\) & 2 \\
\hline \(12 \mid\) & 3 & 61 & 9 & 14 & 1 | & 3 \\
\hline \multicolumn{7}{|l|}{(12 rows)} \\
\hline
\end{tabular}

\section*{Multiple vertices}
```

pgr_primDD(Edges SQL, root vids, distance)
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Example:}

The Minimum Spanning Tree starting on vertices \\(\\{9, 6\\}\\) with \\(distance \leq 3.5\\)
```

SELECT * FROM pgr_primDD(
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
ARRAY[9, 6], 3.5);
seq | depth | start_vid | node | edge | cost | agg_cost
1| 0| 6| 6| -1| 0| 0
2| 1| 6 6| 5| 1| 1| 1
4| 2| 6| 15| 3| 1| 2
5| 2| 6| 11| 5| 1| 2
6| 3| 6| 16| 9| 1| 3
8| 3| 6| 12| 11| 1| 3
8| 1| 6| 7| 4| 1| 1
9| 2| 6| 3| 7| 1| 2
10| 3| 6| 1| 6| 1| 3
2| 6| 8| 10| 1| 2
12| 3| 6| 9| 14| 1| 3
13| 0| 9| 9|-1| 0| 0
14| 1| 9| 8| 14| 1| 1
15| 2| 9| 7| 10| 1| 2
16| 3| 9| 6| 4| 1| 3
17| 3| 9| 3| 7| 1| 3
(17 rows)

```

\section*{Parameters}
\begin{tabular}{|c|c|c|}
\hline Parameter & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below. \\
\hline Root vid & BIGINT & Identifier of the root vertex of the tree. \\
\hline Root vids & ARRAY[ANY-INTEGER] & Array of identifiers of the root vertices.
\(\backslash(0 \backslash)\) values are ignored
For optimization purposes, any duplicated value is ignored. \\
\hline distance & FLOAT & Upper limit for the inclusion of a node in the result. \\
\hline
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERIC:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& & \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)
\begin{tabular}{|c|c|c|}
\hline Parameter & Type & Description \\
\hline seq & BIGINT & Sequential value starting from \\(1). \\
\hline \multirow[t]{2}{*}{depth} & BIGINT & Depth of the node. \\
\hline & & - \(\backslash(0 \backslash)\) when node = start_vid. \\
\hline start_vid & BIGINT & Identifier of the root vertex. \\
\hline node & BIGINT & Identifier of node reached using edge. \\
\hline \multirow[t]{2}{*}{edge} & BIGINT & Identifier of the edge used to arrive to node. \\
\hline & & - \(\backslash(-1 \backslash\) when node \(=\) start_vid. \\
\hline cost & FLOAT & Cost to traverse edge. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERIC:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

See Also
- Spanning Tree-Category
- Prim - Family of functions
- Sample Data
- Boost: Prim's algorithm documentation
- Wikipedia: Prim's algorithm

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
pgr_primDFS
pgr_primDFS - Prim algorithm for Minimum Spanning Tree with Depth First Search ordering.

\section*{Availability}
- Version 3.0.0
- New Official function

\section*{Description}

Visits and extracts the nodes information in Depth First Search ordering of the Minimum Spanning Tree created using Prims's algorithm.

\section*{The main Characteristics are:}
- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.
- Prim's running time: \(\backslash(\mathrm{O}(\mathrm{E} * \log \mathrm{~V}) \backslash)\)
- Returned tree nodes from a root vertex are on Depth First Search order
- Depth First Search Running time: \(\backslash(\mathrm{O}(\mathrm{E}+\mathrm{V}) \backslash)\)

\section*{Signatures}
```

pgr_primDFS(Edges SQL, root vid, [max_depth])
pgr_primDFS(Edges SQL, root vids, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

```

\section*{Single vertex}
pgr_primDFS(Edges SQL, root vid, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

\section*{Example:}

The Minimum Spanning Tree having as root vertex \\(6\\)
```

SELECT * FROM pgr_primDFS(
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
6);
seq | depth | start_vid | node | edge | cost | agg_cost
0| 6| 6| -1| 0| 0
1| 6| 5| 1| 1| 1
1| 6| 10| 2| 1| 1
2| 6| 15| 3| 1| 2
2| 6| 11| 5| 1| 2
3| 6| 16| 9| 1| 3
3| 6| 12| 11| 1| 3
4| 6| 17| 13| 1| 4
6| 7| 4| 1| 1
6| 3| 7| 1| 2
6| 1| 6| 1| 3
6| 8| 10| 1| 2
6| 9| 14| 1| 3
(3 3

```

\section*{Multiple vertices}
pgr_primDFS(Edges SQL, root vids, [max_depth])
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

\section*{Example:}

The Minimum Spanning Tree starting on vertices \(\backslash(\backslash\{9,6 \backslash\} \backslash)\) with \(\backslash(\) depth \(\backslash\) leq \(3 \backslash)\)

\section*{SELECT * FROM pgr_primDFS}

SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
ARRAY[9, 6], max_depth => 3);
seq | depth | start_vid | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline 1| & 01 & 61 & 61 & -1| & 0| & 0 \\
\hline \(2 \mid\) & 1| & 61 & 51 & 1| & 1| & 1 \\
\hline 31 & 1| & 61 & \(10 \mid\) & 21 & 1| & 1 \\
\hline 4 | & 21 & 61 & 15| & 31 & 1 | & 2 \\
\hline 51 & 21 & 61 & 11| & 51 & 1| & 2 \\
\hline \(6 \mid\) & 31 & 61 & 16| & 91 & 1| & 3 \\
\hline 7| & 31 & 61 & 12| & 11| & 1 | & 3 \\
\hline 8| & 1| & 61 & 71 & \(4 \mid\) & 1 | & 1 \\
\hline 91 & 21 & 61 & 31 & 7| & 1 | & 2 \\
\hline \(10 \mid\) & 31 & 61 & \(1 \mid\) & \(6 \mid\) & 1| & 3 \\
\hline 11| & 21 & 61 & 81 & 10| & 1 | & 2 \\
\hline 12| & 31 & 61 & 91 & 14| & \(1 \mid\) & 3 \\
\hline 13| & 0 & 91 & 91 & -1| & 01 & 0 \\
\hline 14 | & 1 | & 91 & 81 & \(14 \mid\) & \(1 \mid\) & 1 \\
\hline 15 | & 21 & 91 & 71 & 10| & \(1 \mid\) & 2 \\
\hline 16| & 31 & 91 & 61 & \(4 \mid\) & 1 | & 3 \\
\hline 17| & 31 & 91 & 31 & 71 & 1| & 3 \\
\hline
\end{tabular}

\section*{Parameters}
\begin{tabular}{|c|c|c|}
\hline Parameter & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below. \\
\hline root vid & BIGINT & \begin{tabular}{l}
Identifier of the root vertex of the tree. \\
When value is \(\backslash(O \backslash)\) then gets the spanning forest starting in aleatory nodes for each tree in the forest.
\end{tabular} \\
\hline root vids & ARRAY [ ANY-INTEGER ] & \begin{tabular}{l}
Array of identifiers of the root vertices. \\
- \(\backslash(0 \backslash)\) values are ignored \\
- For optimization purposes, any duplicated value is ignored.
\end{tabular} \\
\hline
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

\section*{DFS optional parameters}
\begin{tabular}{llll} 
Parameter & Type & Default & Description \\
\hline max_depth & BIGINT \(\backslash(9223372036854775807 \backslash)\) & Upper limit of the depth of the tree. \\
& & \begin{tabular}{l} 
When negative throws an \\
error.
\end{tabular} \\
\hline
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{|c|c|c|c|}
\hline Column & Type & Default & Description \\
\hline id & ANY-INTEGER & & Identifier of the edge. \\
\hline source & ANY-INTEGER & & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL & -1 & Weight of the edge (target, source) \\
\hline & & & - When negative: edge (target, source) does not exist, therefore it's not part of the graph. \\
\hline
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline seq & BIGINT & Sequential value starting from \(\backslash(1 \backslash)\). \\
\hline depth & BIGINT & Depth of the node. \\
& & \(\backslash(0 \backslash)\) when node \(=\) start_vid. \\
\hline start_vid & BIGINT & Identifier of the root vertex. \\
\hline node & BIGINT & Identifier of node reached using edge. \\
\hline edge & BIGINT & \begin{tabular}{l} 
Identifier of the edge used to arrive to \\
node.
\end{tabular} \\
& & \begin{tabular}{l} 
FLOAT
\end{tabular} \\
\hline cost & Cost to traverse edge. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERIC:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

\section*{See Also}

\section*{- Spanning Tree - Category}
- Prim - Family of functions
- Sample Data
- Boost: Prim's algorithm documentation
- Wikipedia: Prim's algorithm

\section*{Indices and tables}

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- Search Page

\section*{Description}

The prim algorithm was developed in 1930 by Czech mathematician Vojtěch Jarník. It is a greedy algorithm that finds a minimum spanning tree for a weighted undirected graph. This means it finds a subset of the edges that forms a tree that includes every vertex, where the total weight of all the edges in the tree is minimized. The algorithm operates by building this tree one vertex at a time, from an arbitrary starting vertex, at each step adding the cheapest possible connection from the tree to another vertex.

This algorithms find the minimum spanning forest in a possibly disconnected graph; in contrast, the most basic form of Prim's algorithm only finds minimum spanning trees in connected graphs. However, running Prim's algorithm separately for each connected component of the graph, then it is called minimum spanning forest.

\section*{The main characteristics are:}
- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.
- Prim's running time: \(\backslash(\mathrm{O}(\mathrm{E} * \log \mathrm{~V}) \backslash)\)

\section*{Note}

From boost Graph: "The algorithm as implemented in Boost.Graph does not produce correct results on graphs with parallel edges."

\section*{Inner Queries}
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline Column & Type & Default & Description \\
\hline source & ANY-INTEGER & & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & & Weight of the edge (source, target) \\
\hline \multirow[t]{2}{*}{reverse_cost} & ANY-NUMERICAL & -1 & Weight of the edge (target, source) \\
\hline & & & - When negative: edge (target, source) does not exist, therefore it's not part of the graph. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

See Also
- Spanning Tree - Category
- Boost: Prim's algorithm
- Wikipedia: Prim's algorithm

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3)

\section*{Reference}
- pgr_version
- pgr_full_version
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2 \(2.1 \mathbf{2 . 0}\)
pgr_version
pgr_version - Query for pgRouting version information.

\section*{Availability}
- Version 3.0.0
- Breaking change on result columns
- Support for old signature ends
- Version 2.0.0
- Official function

\section*{Description}

Returns pgRouting version information.

\section*{Signature}
```

pgr_version()

```

RETURNS TEXT

\section*{Example:}
pgRouting Version for this documentation
```

SELECT pgr_version();

```
pgr_version
3.3.4
(1 row)

\section*{Result Columns}
\begin{tabular}{ll} 
Type & Description \\
\hline TEXT & \begin{tabular}{l} 
pgRouting \\
\\
\\
version
\end{tabular}
\end{tabular}

\section*{See Also}
- Reference
- pgr_full_version

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
pgr_full_version
pgr_full_version - Get the details of pgRouting version information.

\section*{Availability}
- Version 3.0.0
- New official function

\section*{Description}

Get complete details of pgRouting version information

\section*{Signatures}
\[
\begin{aligned}
& \text { pgr_full_version() } \\
& \text { RETURNS (version, build_type, compile_date, library, system, PostgreSQL, compiler, boost, hash) }
\end{aligned}
\]

\section*{Example:}

Information about when this documentation was built
```

SELECT version, library FROM pgr_full_version();
version | library
3.3.4 | pgrouting-3.3.4
(1 row)

```

\section*{Return columns}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline version & TEXT & pgRouting version \\
\hline build_type & TEXT & The Build type \\
\hline compile_date & TEXT & Compilation date \\
\hline library & TEXT & Library name and version \\
\hline system & TEXT & Operative system \\
\hline postgreSQL & TEXT & pgsql used \\
\hline compiler & TEXT & Compiler and version \\
\hline boost & TEXT & Boost version \\
\hline
\end{tabular}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline hash & TEXT & \begin{tabular}{l} 
Git hash of pgRouting \\
build
\end{tabular}
\end{tabular}

\section*{See Also}
- Reference
- pgr_version

\section*{Indices and tables}
- Index
- Search Page

See Also

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2 2.12 .0

\section*{Topology - Family of Functions}

The pgRouting's topology of a network represented with a graph in form of two tables: and edge table and a vertex table.
Attributes associated to the tables help to indicate if the graph is directed or undirected, if an edge is one way on a directed graph, and depending on the final application needs, suitable topology(s) need to be created.
pgRouting suplies some functions to create a routing topology and to analyze the topology.
Additional functions to create a graph:
- Contraction - Family of functions

Additional functions to analyze a graph:
- Components - Family of functions

The following functions modify the database directly therefore the user must have special permissions given by the administrators to use them.
pgr_createTopology - create a topology based on the geometry.
- pgr_createVerticesTable - reconstruct the vertices table based on the source and target information.
- pgr_analyzeGraph - to analyze the edges and vertices of the edge table.
- pgr_analyzeOneWay - to analyze directionality of the edges.
- pgr_nodeNetwork -to create nodes to a not noded edge table.

\section*{Proposed}

\section*{Warning}

Proposed functions for next mayor release.
- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.

These proposed functions do not modify the database.
- pgr_extractVertices - Proposed - Extracts vertex information based on the edge table information.

Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: \(2.6 \mathbf{2 . 5} \mathbf{2 . 4} \mathbf{2 . 3} \mathbf{2 . 2} 2.1 \mathbf{2 . 0}\)
pgr_createTopology
pgr_createTopology - Builds a network topology based on the geometry information.

\section*{Availability}
- Version 2.0.0
- Renamed from version 1.x
- Official function

\section*{Description}

The function returns:
- OK after the network topology has been built and the vertices table created.
- FAIL when the network topology was not built due to an error.

\section*{Signatures}
```

pgr_createTopology(edge_table, tolerance, [options])
options: [the_geom, id, source, target, rows_where, clean]

```
RETURNS VARCHAR

\section*{Parameters}

The topology creation function accepts the following parameters:
```

edge_table:
text Network table name. (may contain the schema name AS well)

```

\section*{tolerance:}
float8 Snapping tolerance of disconnected edges. (in projection unit)
the_geom:
text Geometry column name of the network table. Default value is the_geom.
id:
text Primary key column name of the network table. Default value is id.
source:
text Source column name of the network table. Default value issource.
target:
text Target column name of the network table. Default value is target.

\section*{rows_where:}
text Condition to SELECT a subset or rows. Default value is true to indicate all rows that wheresource or target have a null value, otherwise the condition is used.

\section*{clean:}
boolean Clean any previous topology. Default value is false.

\section*{Warning}

The edge_table will be affected
- The source column values will change.
- The target column values will change.
- An index will be created, if it doesn't exists, to speed up the process to the following columns:
- id
- the_geom
- source
- target

The function returns:
- OK after the network topology has been built.
- Creates a vertices table: <edge_table>_vertices_pgr.
- Fills id and the_geom columns of the vertices table.
- Fills the source and target columns of the edge table referencing theid of the vertices table.
- FAIL when the network topology was not built due to an error:
- A required column of the Network table is not found or is not of the appropriate type.
- The condition is not well formed.
- The names of source , target or id are the same.
- The SRID of the geometry could not be determined.

\section*{The Vertices Table}

The vertices table is a requirement of thepgr_analyzeGraph and the pgr_analyzeOneWay functions.
The structure of the vertices table is:

\section*{id:}
bigint Identifier of the vertex.
cnt:
integer Number of vertices in the edge_table that reference this vertex. Seepgr_analyzeGraph.
chk:
integer Indicator that the vertex might have a problem. See pgr_analyzeGraph.
ein:
integer Number of vertices in the edge_table that reference this vertex AS incoming. Seepgr_analyzeOneWay.
eout:
integer Number of vertices in the edge_table that reference this vertex AS outgoing. Seepgr_analyzeOneWay.
the_geom:
geometry Point geometry of the vertex.

Usage when the edge table's columns MATCH the default values:

The simplest way to use pgr_createTopology is:
```

SELECT pgr_createTopology('edges', 0.001, 'geom');
NOTICE: PROCESSING:
NOTICE: pgr_createTopology('edges', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'true', clean := f)
NOTICE: Performing checks, please wait.
NOTICE: Creating Topology, Please wait..
NOTICE: ----------->> TOPOLOGY CREATED FOR 18 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)

```

\section*{When the arguments are given in the order described:}

We get the same result AS the simplest way to use the function.
```

SELECT pgr_createTopology('edges', 0.001,
'geom', 'id', 'source', 'target');
NOTICE: PROCESSING:
NOTICE: pgr_createTopology('edges', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'true', clean := f)
NOTICE: Performing checks, please wait.
NOTICE: Creating Topology, Please wait..
NOTICE: --------------> TOPOLOGY CREATED FOR 18 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)

```

\section*{Warning}

An error would occur when the arguments are not given in the appropriate order: In this example, the columnid of the table ege_table is passed to the function as the geometry column, and the geometry columnthe_geom is passed to the function as the id column.
NOTICE: Performing checks, please wait ..
NOTICE: ----> PGR ERROR in pgr_createTopology: Wrong type of Column id:geom
HINT: ----> Expected type of geom is integer,smallint or bigint but USER-DEFINED was found
NOTICE: Unexpected error raise_exception
pgr_createtopology
FAIL
(1 row)

\section*{When using the named notation}

Parameters defined with a default value can be omitted, as long as the value matches the default And The order of the parameters would not matter.
```

SELECT pgr_createTopology('edges', 0.001,
the_geom:='geom', id:='id', source:='source', target:='target');
pgr_Createtopology
OK
(1 row)

```
SELECT pgr_createTopology('edges', 0.001 ,
    source:='source', id:='id', target:='target', the_geom:='geom');
pgr_createtopology
OK
(1 row)
```

SELECT pgr_createTopology('edges', 0.001, 'geom', source:='source');
pgr_createtopology
OK
(1 row)

```

\section*{Selecting rows using rows_where parameter}

Selecting rows based on the id.
```

SELECT pgr_createTopology('edges', 0.001, 'geom', rows_where:='id < 10');
pgr_createtopology
OK
(1 row)

```

Selecting the rows where the geometry is near the geometry of row withid= 5 .
```

SELECT pgr_createTopology('edges', 0.001, 'geom'
rows_where:='geom \&\& (SELECT st_buffer(geom, 0.05) FROM edges WHERE id=5)');
pgr_createtopology
OK
(1 row)

```

Selecting the rows where the geometry is near the geometry of the row withgid \(=100\) of the tableothertable.
```

CREATE TABLE otherTable AS (SELECT 100 AS gid, st_point(2.5, 2.5) AS other_geom);
SELECT }
SELECT pgr_createTopology('edges', 0.001, 'geom',
rows_where:='geom \&\& (SELECT st_buffer(other_geom, 1) FROM otherTable WHERE gid=100)');
pgr_createtopology
OK
(1 row)

```

\section*{Using positional notation:}

The arguments need to be given in the order described in the parameters.
Note that this example uses clean flag. So it recreates the whole vertices table.
```

SELECT pgr_createTopology('mytable', 0.001, 'mygeom', 'gid', 'src', 'tgt', clean := TRUE);
pgr_createtopology
OK
(1 row)

```

\section*{Warning}

An error would occur when the arguments are not given in the appropiriate order: In this example, the columngid of the table mytable is passed to the function AS the geometry column, and the geometry columnmygeom is passed to the function AS the id column.
```

SELECT pgr_createTopology('mytable', 0.001, 'gid', 'mygeom', 'src', 'tgt');
NOTICE: PROCESSING:
NOTICE: pgr_createTopology('mytable', 0.001, 'gid', 'mygeom', 'src', 'tgt', rows_where := 'true', clean := f)
NOTICE: Performing checks, please wait.
NOTICE: ----> PGR ERROR in pgr_createTopology: Wrong type of Column id:mygeom
HINT: ---> Expected type of mygeom is integer,smallint or bigint but USER-DEFINED was found
NOTICE: Unexpected error raise_exception
pgr_createtopology
FAIL
(1 row)

```

\section*{When using the named notation}

In this scenario omitting a parameter would create an error because the default values for the column names do not match the column names of the table. The order of the parameters do not matter:
```

SELECT pgr_createTopology('mytable', 0.001, the_geom:='mygeom', id:='gid', source:='src', target:='tgt');
pgr_createtopology
OK
(1 row)

```
SELECT pgr_createTopology('mytable', 0.001, source:='src', id:='gid', target:='tgt', the_geom:='mygeom');
pgr_createtopology
OK
(1 row)

\section*{Selecting rows using rows_where parameter}

Based on id:
```

SELECT pgr_createTopology('mytable', 0.001, 'mygeom', 'gid', 'src', 'tgt', rows_where:='gid < 10');

```
pgr_createtopology

\section*{OK}
(1 row)
SELECT pgr_createTopology('mytable', 0.001, source:='src', id:='gid', target:='tgt', the_geom:='mygeom', rows_where:='gid < 10'); pgr_createtopology

OK
(1 row)
SELECT pgr_createTopology('mytable', 0.001, 'mygeom', 'gid', 'src', 'tgt',
rows_where:='mygeom \&\& (SELECT st_buffer(mygeom, 1) FROM mytable WHERE gid=5)');
pgr_createtopology
OK
(1 row)
SELECT pgr_createTopology('mytable', 0.001, source:='src', id:='gid', target:='tgt', the_geom:='mygeom', rows_where:='mygeom \&\& (SELECT st_buffer(mygeom, 1) FROM mytable WHERE gid=5)'); pgr_createtopology

OK
(1 row)

Selecting the rows where the geometry is near the geometry of the row withgid \(=100\) of the tableothertable.
```

SELECT pgr_createTopology('mytable', 0.001, 'mygeom', 'gid', 'src', 'tgt',
rows_where:='mygeom \&\& (SELECT st_buffer(other_geom, 1) FROM otherTable WHERE gid=100)');
pgr_createtopology
OK
(1 row)
SELECT pgr_createTopology('mytable', 0.001, source:='src', id:='gid', target:='tgt', the_geom:='mygeom',
rows_where:='mygeom \&\& (SELECT st_buffer(other_geom, 1) FROM otherTable WHERE gid=100)');
pgr_createtopology
OK
(1 row)

```

\section*{Additional Examples}
```

- Create a routing topology

```
- Make sure the database does not have the vertices_table
- Clean up the columns of the routing topology to be created
- Create the vertices table
- Inspect the vertices table
- Create the routing topology on the edge table
- Inspect the routing topology
- With full output

\section*{Create a routing topology}

An alternate method to create a routing topology usepgr_extractVertices - Proposed

Make sure the database does not have the vertices_table
```

DROP TABLE IF EXISTS vertices table;
NOTICE: table "vertices_table" does not exist, skipping
DROP TABLE

```

Clean up the columns of the routing topology to be created
```

UPDATE edges
SET source = NULL, target = NULL,
x1 = NULL, y1 = NULL,
x2 = NULL, y2 = NULL;
UPDATE }1

```

\section*{Create the vertices table}
- When the LINESTRING has a SRID then use geom::geometry(POINT, <SRID>)
- For big edge tables that are been prepared,
- Create it as UNLOGGED and
```

SELECT * INTO vertices_table
FROM pgr_extractVertices('SELECT id, geom FROM edges ORDER BY id');
SELECT 17

```

Inspect the vertices table


\section*{Create the routing topology on the edge table}

\section*{Updating the source information}
```

WITH
out_going AS(
SELECT id AS vid, unnest(out_edges) AS eid, x, y
FROM vertices_table
)
UPDATE edges
SET source = vid, x1 = x, y1 = y
FROM out_going WHERE id = eid;
UPDATE 18

```

\section*{Updating the target information}
```

WITH
in_coming AS (
SELECT id AS vid, unnest(in_edges) AS eid, x, y
FROM vertices_table
)
UPDATE edges
SET target = vid, x2 = x, y2 = y
FROM in_coming WHERE id = eid;
UPDATE }1

```

\section*{Inspect the routing topology}
\begin{tabular}{|c|c|c|c|c|}
\hline \multicolumn{5}{|l|}{SELECT id, source, target, \(\mathrm{x} 1, \mathrm{y} 1, \mathrm{x} 2, \mathrm{y} 2\) FROM edges ORDER BY id; id | source | target | x1 | y1 | x2} \\
\hline \(1 \mid\) & 5। & \(6|2| 0 \mid\) & & | 1 \\
\hline 21 & & 10| \(2|1|\) & & | 1 \\
\hline 31 & 10| & 15| 3| 1| & & | 1 \\
\hline \(4 \mid\) & 61 & 7| 2| 1| & & 12 \\
\hline 51 & 10| & 11| 3| 1| & & | 2 \\
\hline 61 & 1| & 3| 0 | \(2 \mid\) & & 12 \\
\hline 71 & 31 & 7| 1| \(2 \mid\) & & 12 \\
\hline 81 & 71 & 11| 2| \(2 \mid\) & & | 2 \\
\hline 91 & 11| & 16| 3| \(2 \mid\) & & | 2 \\
\hline 10 & 71 & 8| \(2|2|\) & & | 3 \\
\hline 11 & 11| & 12| 3| \(2 \mid\) & & 3| 3 \\
\hline 12 & 8| & 12| \(2|3|\) & & | 3 \\
\hline 13 & 12| & 17| 3| 3| & & \(4 \mid 3\) \\
\hline 14 & 8| & 9| \(2|3|\) & & | 4 \\
\hline 15 & \(16 \mid\) & 17| \(4|2|\) & & \(4 \mid 3\) \\
\hline 16 & \(15 \mid\) & 16| \(4|1|\) & & | 2 \\
\hline 17 & 2| & \(4|0.5| 3.5 \mid\) & 1.9999999 & 9999999 \\
\hline 18 & 13| & 14|3.5|2.3 & & \(3.5 \mid 4\) \\
\hline \multicolumn{5}{|l|}{(18 rows)} \\
\hline
\end{tabular}


\section*{Generated topology}

\section*{With full output}

This example start a clean topology, with 5 edges, and then its incremented to the rest of the edges.
```

SELECT pgr_createTopology('edges', 0.001, 'geom', rows_where:='id < 6', clean := true);
NOTICE: PROCESSING:
NOTICE: pgr_createTopology('edges', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'id < 6', clean := t)
NOTICE: Performing checks, please wait
NOTICE: Creating Topology, Please wait.
NOTICE: ------------> TOPOLOGY CREATED FOR 5 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE:Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)
SELECT pgr_createTopology('edges', 0.001, 'geom');
NOTICE: PROCESSING
NOTICE: pgr_createTopology('edges', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'true', clean := f)
NOTICE: Performing checks, please wait
NOTICE: Creating Topology, Please wait..
NOTICE: -----------> TOPOLOGY CREATED FOR 13 edges
NOTICE:Rows with NULL geometry or NULL id: 0
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)

```

The example uses the Sample Data network.

See Also
- Topology - Family of Functions
- pgr_createVerticesTable
- pgr_analyzeGraph

\section*{Indices and tables}
```

- Index
- Search Page

```
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2 2 .1 2.0
pgr_createVerticesTable - Reconstructs the vertices table based on the source and target information.

\section*{Availability}
- Version 2.0.0
- Renamed from version 1.x
- Official function

\section*{Description}

The function returns:
- OK after the vertices table has been reconstructed.
- FAIL when the vertices table was not reconstructed due to an error.

\section*{Signatures}
```

pgr_createVerticesTable(edge_table, [the_geom, source, target, rows_where])
RETURNS VARCHAR

```

\section*{Parameters}

The reconstruction of the vertices table function accepts the following parameters:
```

edge_table:
text Network table name. (may contain the schema name as well)
the_geom:
text Geometry column name of the network table. Default value is the_geom.
source:
text Source column name of the network table. Default value issource.
target:
text Target column name of the network table. Default value is target.
rows_where:
text Condition to SELECT a subset or rows. Default value is true to indicate all rows.

```

\section*{Warning}

The edge_table will be affected
- An index will be created, if it doesn't exists, to speed up the process to the following columns:
- the_geom
- source
- target

The function returns:
- OK after the vertices table has been reconstructed
- Creates a vertices table: <edge_table>_vertices_pgr.
- Fills id and the_geom columns of the vertices table based on the source and target columns of the edge table.
- FAIL when the vertices table was not reconstructed due to an error.
- A required column of the Network table is not found or is not of the appropriate type.
- The condition is not well formed.
- The names of source, target are the same.
- The SRID of the geometry could not be determined.

\section*{The Vertices Table}

The vertices table is a requierment of thepgr_analyzeGraph and the pgr_analyzeOneWay functions.
The structure of the vertices table is:
id:
bigint Identifier of the vertex.
cnt:
integer Number of vertices in the edge_table that reference this vertex. Seepgr_analyzeGraph.
chk:
integer Indicator that the vertex might have a problem. See pgr_analyzeGraph.
ein:
integer Number of vertices in the edge_table that reference this vertex as incoming. Seepgr_analyzeOneWay

\section*{eout:}
integer Number of vertices in the edge_table that reference this vertex as outgoing. Seepgr_analyzeOneWay.

\section*{the_geom:}
geometry Point geometry of the vertex.

\section*{Example 1:}

The simplest way to use pgr_createVerticesTable
```

SELECT pgr_createVerticesTable('edges', 'geom');
NOTICE: PROCESSING
NOTICE: pgr_createVerticesTable('edges','geom','source','target','true')
NOTICE: Performing checks, please wait.
NOTICE: Populating public.edges_vertices_pgr, please wait.
NOTICE: -----> VERTICES TABLE CREATED WITH 17 VERTICES
NOTICE: FOR 18 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed: 18
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

```

\section*{Additional Examples}

\section*{Example 2:}

When the arguments are given in the order described in the parameters:
```

SELECT pgr_createVerticesTable('edges', 'geom', 'source', 'target')
NOTICE: PROCCESSING
NOTICE: pgr_createVerticesTable('edges','geom','source','target','true')
NOTICE: Performing checks, please wait
NOTICE: Populating public.edges_vertices_pgr, please wait.
NOTICE: -----> VERTICES TABLE CREATED WITH 17 VERTICES
NOTICE: FOR 18 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed: }1
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

```

We get the same result as the simplest way to use the function.

\section*{Warning}

An error would occur when the arguments are not given in the appropriate order: In this example, the column source column source of the tablemytable is passed to the function as the geometry column, and the geometry column the_geom is passed to the function as the source column.
```

SELECT pgr_createVerticesTable('edges', 'source', 'geom',''target');
NOTICE: PROCESSING
NOTICE: pgr_createVerticesTable('edges','source','geom','target','true')
NOTICE: Performing checks, please wait
NOTICE: ----> PGR ERROR in pgr_createVerticesTable: Wrong type of Column source: geom
HINT: ---> Expected type of geom is integer, smallint or bigint but USER-DEFINED was found
NOTICE: Unexpected error raise_exception
pgr_createverticestable

```
FAIL
(1 row)

\section*{When using the named notation}

\section*{Example 3:}

The order of the parameters do not matter:
```

SELECT pgr_createVerticesTable('edges', the_geom:='geom', source:='source', target:='target');

```
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('edges','geom','source','target','true')
NOTICE: Performing checks, please wait
NOTICE: Populating public.edges_vertices_pgr, please wait.
NOTICE: -----> VERTICES TABLE CREATED WITH 17 VERTICES
NOTICE: FOR 18 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed: 18
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

\section*{Example 4:}

Using a different ordering
```

SELECT pgr_createVerticesTable('edges', source:='source', target:='target', the_geom:='geom');
NOTICE: PROCESSING
NOTICE: pgr_createVerticesTable('edges','geom','source','target','true')
NOTICE: Performing checks, please wait.
NOTICE: Populating public.edges_vertices_pgr, please wait.
NOTICE: -----> VERTICES TABLE CREATED WITH 17 VERTICES
NOTICE: FOR 18 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed: 18
NOTICE:Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

```

\section*{Example 5:}

Parameters defined with a default value can be omitted, as long as the value matches the default:
```

SELECT pgr_createVerticesTable('edges', 'geom', source:='source');
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('edges','geom','source','target','true')
NOTICE: Performing checks, please wait.
NOTICE: Populating public.edges_vertices_pgr, please wait..
NOTICE: ----> VERTICES TABLE CREATED WITH 17 VERTICES
NOTICE: FOR 18 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed: }1
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

```

\section*{Selecting rows using rows_where parameter}

\section*{Example 6:}

Selecting rows based on the id.
```

SELECT pgr_createVerticesTable('edges', 'geom', rows_where:='id < 10');
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('edges','geom','source','target','id < 10')
NOTICE: Performing checks, please wait.
NOTICE: Populating public.edges_vertices_pgr, please wait..
NOTICE: ----> VERTICES TABLE CREATED WITH 9 VERTICES
NOTICE: FOR 10 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed: 10
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

```

Example 7:
Selecting the rows where the geometry is near the geometry of row withid \(=5\).

SELECT pgr_createVerticesTable('edges', 'geom',
rows_where:='geom \&\& (select st_buffer(geom,0.5) FROM edges WHERE id=5)');
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('edges','geom','source','target','geom \& \& (select st_buffer(geom,0.5) FROM edges WHERE id=5)')
NOTICE: Performing checks, please wait .
NOTICE: Populating public.edges_vertices_pgr, please wait..
NOTICE: -----> VERTICES TABLE CREATED WITH 9 VERTICES
NOTICE: FOR 9 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed: 9
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

\section*{Example 8:}

Selecting the rows where the geometry is near the geometry of the row withgid \(=100\) of the tableothertable.
```

DROP TABLE IF EXISTS otherTable;
NOTICE: table "othertable" does not exist, skipping
DROP TABLE
CREATE TABLE otherTable AS (SELECT 100 AS gid, st_point(2.5,2.5) AS other_geom) ;
SELECT }
SELECT pgr_createVerticesTable('edges', 'geom',
rows_where:='geom \&\& (select st_buffer(other_geom,0.5) FROM otherTable WHERE gid=100)');
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('edges','geom','source','target','geom \&\& (select st_buffer(other_geom,0.5) FROM otherTable WHERE gid=100)')
NOTICE: Performing checks, please wait.
NOTICE: Populating public.edges_vertices_pgr, please wait.
NOTICE: ----> VERTICES TABLE CREATED WITH 10 VERTICES
NOTICE: FOR 12 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed:12
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

```

\section*{Usage when the edge table's columns DO NOT MATCH the default values:}

Using the following table
```

DROP TABLE IF EXISTS mytable;
NOTICE: table "mytable" does not exist, skipping
DROP TABLE
CREATE TABLE mytable AS (SELECT id AS gid, geom AS mygeom, source AS src ,target AS tgt FROM edges) ;
SELECT }1

```

\section*{Using positional notation:}

\section*{Example 9:}

The arguments need to be given in the order described in the parameters:
```

SELECT pgr_createVerticesTable('mytable', 'mygeom', 'src', 'tgt')
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('mytable','mygeom','src','tgt','true')
NOTICE: Performing checks, please wait ..
NOTICE: Populating public.mytable_vertices_pgr, please wait...
NOTICE: -----> VERTICES TABLE CREATED WITH 17 VERTICES
NOTICE: FOR }18\mathrm{ EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed:18
NOTICE: Vertices table for table public.mytable is: public.mytable_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

```

\section*{Warning}

An error would occur when the arguments are not given in the appropriate order: In this example, the columnsrc of the table mytable is passed to the function as the geometry column, and the geometry columnmygeom is passed to the function as the source column.
```

SELECT pgr_createVerticesTable('mytable', 'src', 'mygeom', 'tgt');

```
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('mytable','src','mygeom','tgt','true')
NOTICE: Performing checks, please wait .
NOTICE: ----> PGR ERROR in pgr_createVerticesTable: Wrong type of Column source: mygeom
HINT: ----> Expected type of mygeom is integer, smallint or bigint but USER-DEFINED was found
NOTICE: Unexpected error raise_exception
pgr_createverticestable
FAIL
(1 row)

\section*{When using the named notation}

\section*{Example 10:}

The order of the parameters do not matter:
```

SELECT pgr_createVerticesTable('mytable',the_geom:='mygeom',source:='src',target:='tgt');
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('mytable','mygeom','src','tgt','true')
NOTICE: Performing checks, please wait
NOTICE: Populating public.mytable_vertices_pgr, please wait..
NOTICE: -----> VERTICES TABLE CREATED WITH 17 VERTICES
NOTICE: FOR 18 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed:18
NOTICE: Vertices table for table public.mytable is: public.mytable_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

```

\section*{Example 11:}

Using a different ordering
In this scenario omitting a parameter would create an error because the default values for the column names do not match the column names of the table.
```

SELECT pgr_createVerticesTable(
mytable', source:='src', target:='tgt',
the geom:='mygeom')
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('mytable','mygeom','src','tgt','true')
NOTICE: Performing checks, please wait
NOTICE: Populating public.mytable_vertices_pgr, please wait..
NOTICE: -----> VERTICES TABLE CREATED WITH 17 VERTICES
NOTICE: FOR 18 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed: 18
NOTICE:Vertices table for table public.mytable is: public.mytable_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

```

\section*{Selecting rows using rows_where parameter}

\section*{Example 12:}

Selecting rows based on the gid. (positional notation)

\section*{SELECT pgr_createVerticesTable(}
'mytable', 'mygeom', 'src', 'tgt',
rows_where: :='gid < 10');
NOTICE: PROCESSING
NOTICE: pgr_createVerticesTable('mytable','mygeom','src','tgt','gid < 10')
NOTICE: Performing checks, please wait.
NOTICE: Populating public.mytable_vertices_pgr, please wait..
NOTICE: -----> VERTICES TABLE CREATED WITH 9 VERTICES
NOTICE: FOR 10 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed: 10
NOTICE: Vertices table for table public.mytable is: public.mytable_vertices pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

\section*{Example 13:}

Selecting rows based on the gid. (named notation)
```

SELECT pgr_createVerticesTable(
'mytable', source:='src', target:='tgt', the_geom:='mygeom',
rows_where:='gid < 10');
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('mytable','mygeom','src','tgt','gid < 10')
NOTICE: Performing checks, please wait ..
NOTICE: Populating public.mytable_vertices_pgr, please wait.
NOTICE: -----> VERTICES TABLE CREATED WITH 9 VERTICES
NOTICE: FOR 10 EDGES
NOTICE: Edges with NULL geometry,source or target: 0
NOTICE: Edges processed: 10
NOTICE: Vertices table for table public.mytable is: public.mytable_vertices_pgr
NOTICE:
pgr_createverticestable
OK
(1 row)

```

\section*{Example 14:}

Selecting the rows where the geometry is near the geometry of row withgid \(=5\).
```

SELECT pgr_createVerticesTable(
'mytable', 'mygeom', 'src', 'tgt',
rows_where := 'the_geom \&\& (SELECT st_buffer(mygeom,0.5) FROM mytable WHERE gid=5)');
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('mytable','mygeom','src','tgt','the_geom \&\& (SELECT st_buffer(mygeom,0.5) FROM mytable WHERE gid=5)')
NOTICE: Performing checks, please wait .
NOTICE: Got column "the_geom" does not exist
NOTICE: ERROR: Condition is not correct, please execute the following query to test your condition
NOTICE: select * from public.mytable WHERE true AND (the_geom \&\& (SELECT st_buffer(mygeom,0.5) FROM mytable WHERE gid=5)) limit 1
pgr_createverticestable
FAIL
(1 row)

```

\section*{Example 15:}

TBD

\section*{SELECT pgr_createVerticesTable(}
'mytable', source:='src', target:='tgt', the_geom:='mygeom',
rows_where:='mygeom \&\& (SELECT st_buffer(mygeom,0.5) FROM mytable WHERE id=5)');
NOTICE: PROCESSING:
NOTICE: pgr_createVerticesTable('mytable','mygeom','src','tgt','mygeom \&\& (SELECT st_buffer(mygeom,0.5) FROM mytable WHERE id=5)')
NOTICE: Performing checks, please wait .....
NOTICE: Got column "id" does not exist
NOTICE: ERROR: Condition is not correct, please execute the following query to test your condition
NOTICE: select * from public.mytable WHERE true AND (mygeom \& \& (SELECT st_buffer(mygeom,0.5) FROM mytable WHERE id=5)) limit 1
pgr_createverticestable
FAIL
(1 row)

\section*{Example 16:}

Selecting the rows where the geometry is near the geometry of the row withgid \(=100\) of the tableothertable.
```

DROP TABLE IF EXISTS otherTable
DROP TABLE
CREATE TABLE otherTable AS (SELECT 100 AS gid, st_point(2.5,2.5) AS other_geom) ;
SELECT 1

```

\section*{SELECT pgr_createVerticesTable \\ 'mytable', 'mygeom', 'src', 'tgt',}
rows_where:='the_geom \&\& (SELECT st_buffer(othergeom,0.5) FROM otherTable WHERE gid=100)');
NOTICE: PROCESSING
NOTICE: pgr_createVerticesTable('mytable','mygeom','src','tgt','the_geom \&\& (SELECT st_buffer(othergeom,0.5) FROM otherTable WHERE gid=100)')
NOTICE: Performing checks, please wait .
NOTICE: Got column "the_geom" does not exist
NOTICE: ERROR: Condition is not correct, please execute the following query to test your condition
NOTICE: select * from public.mytable WHERE true AND (the_geom \&\& (SELECT st_buffer(othergeom,0.5) FROM otherTable WHERE gid=100)) limit 1 pgr_createverticestable

\section*{FAIL}
(1 row)

\section*{Example 17:}

TBD

\section*{SELECT pgr_createVerticesTable(}
'mytable',source:='src',target:='tgt',the_geom:='mygeom',
rows_where:='the_geom \&\& (SELECT st_buffer(othergeom,0.5) FROM otherTable WHERE gid=100)');
NOTICE: PROCESSING
NOTICE: pgr_createVerticesTable('mytable','mygeom','src','tgt','the_geom \&\& (SELECT st_buffer(othergeom,0.5) FROM otherTable WHERE gid=100)')
NOTICE: Performing checks, please wait
NOTICE: Got column "the_geom" does not exist
NOTICE: ERROR: Condition is not correct, please execute the following query to test your condition
NOTICE: select * from public.mytable WHERE true AND (the_geom \&\& (SELECT st_buffer(othergeom,0.5) FROM otherTable WHERE gid=100)) limit 1 pgr_createverticestable

\section*{FAIL}
(1 row)

The example uses the Sample Data network.

\section*{See Also}
- Topology - Family of Functions for an overview of a topology for routing algorithms.
- pgr_createTopology <pgr_create_topology>' to create a topology based on the geometry.
- pgr_analyzeGraph to analyze the edges and vertices of the edge table.
- pgr_analyzeOneWay to analyze directionality of the edges.

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.6 2.5 2.42 .3 2.2 2.12 .0
pgr_analyzeGraph
pgr_analyzeGraph - Analyzes the network topology.

\section*{Availability}
- Version 2.0.0
- Official function

\section*{Description}

The function returns:
- OK after the analysis has finished.
- FAIL when the analysis was not completed due to an error.
```

pgr_analyzeGraph(edge_table, tolerance, [options])

```
options: [the_geom, id, source, target, rows_where]
RETURNS VARCHAR

\section*{Prerequisites}

The edge table to be analyzed must contain a source column and a target column filled with id's of the vertices of the
segments and the corresponding vertices table <edge_table>_vertices_pgr that stores the vertices information.
- Use pgr_createVerticesTable to create the vertices table.
- Use pgr_createTopology to create the topology and the vertices table.

\section*{Parameters}

The analyze graph function accepts the following parameters:

\section*{edge_table:}
text Network table name. (may contain the schema name as well)

\section*{tolerance:}
float8 Snapping tolerance of disconnected edges. (in projection unit)

\section*{the_geom:}
text Geometry column name of the network table. Default value is the_geom.
id:
text Primary key column name of the network table. Default value is id.

\section*{source:}
text Source column name of the network table. Default value issource.
target:
text Target column name of the network table. Default value is target.
rows_where:
text Condition to select a subset or rows. Default value is true to indicate all rows.

The function returns:
- OK after the analysis has finished.
- Uses the vertices table: <edge_table>_vertices_pgr.
- Fills completely the cnt and chk columns of the vertices table.
- Returns the analysis of the section of the network defined by rows_where
- FAIL when the analysis was not completed due to an error.
- The vertices table is not found.
- A required column of the Network table is not found or is not of the appropriate type.
- The condition is not well formed.
- The names of source , target or id are the same.
- The SRID of the geometry could not be determined.

\section*{The Vertices Table}

The vertices table can be created with pgr_createVerticesTable or pgr_createTopology
The structure of the vertices table is:
id:
bigint Identifier of the vertex.
cnt:
integer Number of vertices in the edge_table that reference this vertex.
chk:
integer Indicator that the vertex might have a problem.
ein:
integer Number of vertices in the edge_table that reference this vertex as incoming. Seepgr_analyzeOneWay. eout:
integer Number of vertices in the edge_table that reference this vertex as outgoing. Seepgr_analyzeOneWay. the_geom:
geometry Point geometry of the vertex.

Usage when the edge table's columns MATCH the default values:

The simplest way to use pgr_analyzeGraph is:
```

SELECT pgr_createTopology('edges',0.001, 'geom', clean := true);
NOTICE: PROCCESSING:

```
NOTICE: pgr_createTopology('edges', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'true', clean := t)
NOTICE: Performing checks, please wait
NOTICE: Creating Topology, Please wait
NOTICE: ------------> TOPOLOGY CREATED FOR 18 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)
SELECT pgr_analyzeGraph('edges',0.001,'geom');
NOTICE: PROCESSING
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','true')
NOTICE: Performing checks, please wait ..
NOTICE: Analyzing for dead ends. Please wait...
NOTICE: Analyzing for gaps. Please wait..
NOTICE: Analyzing for isolated edges. Please wait...
NOTICE: Analyzing for ring geometries. Please wait...
NOTICE: Analyzing for intersections. Please wait...
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 2
NOTICE: Dead ends: 7
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected: 1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

\section*{Arguments are given in the order described in the parameters:}
```

SELECT pgr_analyzeGraph('edges',0.001,'geom','id','source','target');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','true')
NOTICE: Performing checks, please wait ..
NOTICE:Analyzing for dead ends. Please wait..
NOTICE:Analyzing for gaps. Please wait..
NOTICE:Analyzing for isolated edges. Please wait..
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE:Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:2
NOTICE: Dead ends:7
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected:1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

```

We get the same result as the simplest way to use the function.

\section*{Warning}

An error would occur when
the arguments are not given in the appropriate order:
In this example, the column id of the tablemytable is passed to the function as the geometry column, and the geometry column the_geom is passed to the function as the id column.
```

SELECT pgr_analyzeGraph('edges',0.001,'id','geom','source','target');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'id','geom','source','target','rue')
NOTICE: Performing checks, please wait .
NOTICE: Got function st_srid(bigint) does not exist
NOTICE: ERROR: something went wrong when checking for SRID of id in table public.edges
pgr_analyzegraph

```
FAIL
(1 row)

\section*{When using the named notation}

The order of the parameters do not matter:
```

SELECT pgr_analyzeGraph('edges',0.001,the_geom:='geom',id:='id',source:='source',target:='target');
NOTICE: PROZCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','true')
NOTICE: Performing checks, please wait .
NOTICE:Analyzing for dead ends. Please wait.
NOTICE:Analyzing for gaps. Please wait...
NOTICE:Analyzing for isolated edges. Please wait..
NOTICE:Analyzing for ring geometries. Please wait.
NOTICE:Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:2
NOTICE: Dead ends:7
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected:1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

```
```

SELECT pgr_analyzeGraph('edges',0.001,source:='source',id:=''id',target:='target',the_geom:='geom');
NOTICE: PROCESSING
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','rue')
NOTICE: Performing checks, please wait ..
NOTICE:Analyzing for dead ends. Please wait..
NOTICE:Analyzing for gaps. Please wait...
NOTICE:Analyzing for isolated edges. Please wait.
NOTICE:Analyzing for ring geometries. Please wait.
NOTICE:Analyzing for intersections. Please wait...
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:2
NOTICE: Dead ends:7
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected:1
NOTICE: Ring geometries:0
pgr_analyzegraph
OK
(1 row)

```

Parameters defined with a default value can be omitted, as long as the value matches the default:
```

SELECT pgr_analyzeGraph('edges',0.001, 'geom', source:='source');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','rue')
NOTICE: Performing checks, please wait .
NOTICE:Analyzing for dead ends. Please wait.
NOTICE:Analyzing for gaps. Please wait.,
NOTICE: Analyzing for isolated edges. Please wait.
NOTICE:Analyzing for ring geometries. Please wait.
NOTICE:Analyzing for intersections. Please wait.
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:2
NOTICE: Dead ends:7
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected:1
NOTICE: Ring geometries:0
pgr_analyzegraph
OK
(1 row)

```

\section*{Selecting rows using rows_where parameter}

Selecting rows based on the id. Displays the analysis a the section of the network.
```

SELECT pgr_analyzeGraph('edges',0.001, 'geom', rows_where:='id < 10');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','id < 10')
NOTICE: Performing checks, please wait .
NOTICE:Analyzing for dead ends. Please wait.
NOTICE:Analyzing for gaps. Please wait..
NOTICE:Analyzing for isolated edges. Please wait..
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE:Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 0
NOTICE: Dead ends:4
NOTICE: Potential gaps found near dead ends:0
NOTICE: Intersections detected: 0
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

```

Selecting the rows where the geometry is near the geometry of row withid \(=5\)
```

SELECT pgr_analyzeGraph('edges',0.001, 'geom', rows_where:='geom \&\& (SELECT st buffer(geom,0.05) FROM edge table WHERE id=5)')
NOTICE: PROZCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','geom \&\& (SELECT st_buffer(geom,0.05) FROM edge_table WHERE id=5)')
NOTICE: Performing checks, please wait
NOTICE:Got relation "edge table" does not exist
NOTICE: ERROR: Condition is not correct. Please execute the following query to test your condition
NOTICE: select count(*) from public.edges WHERE true AND (geom \&\& (SELECT st_buffer(geom,0.05) FROM edge_table WHERE id=5))
pgr_analyzegraph
FAIL
(1 row)

```

Selecting the rows where the geometry is near the geometry of the row withgid \(=100\) of the tableothertable.
```

CREATE TABLE otherTable AS (SELECT 100 AS gid, st_point(2.5,2.5) AS other_geom) ;
SELECT }
SELECT pgr_analyzeGraph('edges',0.001, 'geom', rows_where:='geom \&\& (SELECT st_buffer(geom,1) FROM otherTable WHERE gid=100)');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','geom \&\& (SELECT st_buffer(geom,1) FROM otherTable WHERE gid=100)')
NOTICE: Performing checks, please wait ..
NOTICE: Analyzing for dead ends. Please wait..
NOTICE:Analyzing for gaps. Please wait..
NOTICE:Analyzing for isolated edges. Please wait..
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE: Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES
NOTICE: Isolated segments:2
NOTICE: Dead ends:7
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected:1
NOTICE: Ring geometries:0
pgr_analyzegraph
OK
(1 row)

```

\section*{Usage when the edge table's columns DO NOT MATCH the default values:}

For the following table
```

CREATE TABLE mytable AS (SELECT id AS gid, source AS src ,target AS tgt , geom AS mygeom FROM edges);
SELECT 18
SELECT pgr_createTopology('mytable',0.001,'mygeom','gid','src','tgt', clean := true);
NOTICE: PROCESSING:
NOTICE: pgr_createTopology('mytable', 0.001, 'mygeom', 'gid', 'src', 'tgt', rows_where := 'true', clean := t)
NOTICE: Performing checks, please wait ..
NOTICE:Creating Topology, Please wait...
NOTICE: ------------> TOPOLOGY CREATED FOR 18 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE: Vertices table for table public.mytable is: public.mytable_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)

```

\section*{Using positional notation:}

The arguments need to be given in the order described in the parameters:
```

SELECT pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt');

```
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt','true')
NOTICE: Performing checks, please wait .
NOTICE: Analyzing for dead ends. Please wait..
NOTICE: Analyzing for gaps. Please wait...
NOTICE: Analyzing for isolated edges. Please wait...
NOTICE: Analyzing for ring geometries. Please wait.
NOTICE: Analyzing for intersections. Please wait...
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 2
NOTICE: Dead ends: 7
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected: 1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

\section*{Warning}

An error would occur when the arguments are not given in the appropriate order: In this example, the columngid of the table mytable is passed to the function as the geometry column, and the geometry columnmygeom is passed to the function as the id column.

SELECT pgr_analyzeGraph('mytable',0.0001,'gid','mygeom','src','tgt');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('mytable',0.0001,'gid','mygeom','src','tgt','true')
NOTICE: Performing checks, please wait ..
NOTICE: Got function st_srid(bigint) does not exist
NOTICE: ERROR: something went wrong when checking for SRID of gid in table public.mytable
pgr_analyzegraph
FAIL
(1 row)

\section*{When using the named notation}

The order of the parameters do not matter:
```

SELECT pgr_analyzeGraph('mytable',0.001,the_geom:='mygeom',id:='gid',source:='src',target:='tgt');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt','true')
NOTICE: Performing checks, please wait ..
NOTICE: Analyzing for dead ends. Please wait..
NOTICE:Analyzing for gaps. Please wait..
NOTICE:Analyzing for isolated edges. Please wait...
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE: Analyzing for intersections. Please wait...
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:2
NOTICE: Dead ends:7
NOTICE: Potential gaps found near dead ends: }
NOTICE: Intersections detected:1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

```
SELECT pgr_analyzeGraph('mytable',0.001,source:='src',id:='gid',target:='tgt',the_geom:='mygeom');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt','true')
NOTICE: Performing checks, please wait ..
NOTICE: Analyzing for dead ends. Please wait...
NOTICE: Analyzing for gaps. Please wait...
NOTICE: Analyzing for isolated edges. Please wait...
NOTICE: Analyzing for ring geometries. Please wait..
NOTICE: Analyzing for intersections. Please wait...
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 2
NOTICE: Dead ends:7
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected: 1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

In this scenario omitting a parameter would create an error because the default values for the column names do not match the column names of the table.

\section*{Selecting rows using rows_where parameter}

Selecting rows based on the id.
```

SELECT pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt',rows_where:='gid < 10');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('mytable',0.001,'mygeom','gid',''src','tgt','gid < 10')
NOTICE: Performing checks, please wait ..
NOTICE: Analyzing for dead ends. Please wait.
NOTICE:Analyzing for gaps. Please wait.,
NOTICE:Analyzing for isolated edges. Please wait...
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE:Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 0
NOTICE: Dead ends:4
NOTICE: Potential gaps found near dead ends: 0
NOTICE: Intersections detected:0
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

```
SELECT pgr_analyzeGraph('mytable',0.001,source:='src','id:='gid',target:='tgt',the_geom:='mygeom',rows_where:='gid < 10');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt','gid < 10')
NOTICE: Performing checks, please wait ...
NOTICE: Analyzing for dead ends. Please wait.
NOTICE: Analyzing for gaps. Please wait.
NOTICE: Analyzing for isolated edges. Please wait...
NOTICE: Analyzing for ring geometries. Please wait...
NOTICE: Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 0
NOTICE: Dead ends:4
NOTICE: Potential gaps found near dead ends: 0
NOTICE: Intersections detected: 0
NOTICE: \(\quad\) Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

Selecting the rows WHERE the geometry is near the geometry of row withid \(=5\).
```

SELECT pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt',
rows_where:='mygeom \&\& (SELECT st_buffer(mygeom,1) FROM mytable WHERE gid=5)');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt','mygeom \&\& (SELECT st_buffer(mygeom,1) FROM mytable WHERE gid=5)')
NOTICE: Performing checks, please wait ..
NOTICE:Analyzing for dead ends. Please wait..
NOTICE: Analyzing for gaps. Please wait..
NOTICE: Analyzing for isolated edges. Please wait...
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE:Analyzing for intersections. Please wait...
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:1
NOTICE: Dead ends:5
NOTICE: Potential gaps found near dead ends: 0
NOTICE: Intersections detected: 1
NOTICE: Ring geometries:0
pgr_analyzegraph
OK
(1 row)

```

SELECT pgr_analyzeGraph('mytable',0.001,source:='src',id:='gid',target:='tgt',the_geom:='mygeom',
rows_where:='mygeom \&\& (SELECT st_buffer(mygeom,1) FROM mytable WHERE gid=5)');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt','mygeom \&\& (SELECT st_buffer(mygeom,1) FROM mytable WHERE gid=5)')
NOTICE: Performing checks, please wait ..
NOTICE: Analyzing for dead ends. Please wait...
NOTICE: Analyzing for gaps. Please wait...
NOTICE: Analyzing for isolated edges. Please wait..
NOTICE: Analyzing for ring geometries. Please wait...
NOTICE: Analyzing for intersections. Please wait...
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 1
NOTICE: Dead ends: 5
NOTICE: Potential gaps found near dead ends: 0
NOTICE: Intersections detected: 1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

Selecting the rows WHERE the geometry is near the place='myhouse' of the tableothertable. (note the use of quote literal)
```

DROP TABLE IF EXISTS otherTable;
DROP TABLE
CREATE TABLE otherTable AS (SELECT 'myhouse'::text AS place, st_point(2.5,2.5) AS other_geom) ;
SELECT 1
SELECT pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt',
rows_where:='mygeom \&\& (SELECT st_buffer(other_geom,1) FROM otherTable WHERE place='||quote_literal('myhouse')||')');
NOTICE: PROCESSING
NOTICE: pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt','mygeom \&\& (SELECT st_buffer(other_geom,1) FROM otherTable WHERE place='myhouse')')
NOTICE: Performing checks, please wait ..
NOTICE: Analyzing for dead ends. Please wait..
NOTICE:Analyzing for gaps. Please wait..
NOTICE:Analyzing for isolated edges. Please wait...
NOTICE: Analyzing for ring geometries. Please wait...
NOTICE:Analyzing for intersections. Please wait...
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:2
NOTICE: Dead ends: 10
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected: 1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

```
```

SELECT pgr_analyzeGraph('mytable',0.001,source:='src',id:='gid',target:='tgt',the_geom:='mygeom',
rows_where:='mygeom \&\& (SELECT st_buffer(other_geom,1) FROM otherTable WHERE place='|quote_literal('myhouse')||')');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('mytable',0.001,'mygeom','gid','src','tgt','mygeom \&\& (SELECT st_buffer(other_geom,1) FROM otherTable WHERE place='myhouse')')
NOTICE: Performing checks, please wait ..
NOTICE:Analyzing for dead ends. Please wait..
NOTICE:Analyzing for gaps. Please wait.,
NOTICE: Analyzing for isolated edges. Please wait...
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE:Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:2
NOTICE: Dead ends: 10
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected:1
NOTICE: Ring geometries:0
pgr_analyzegraph
OK
(1 row)

```

\section*{SELECT pgr_createTopology('edges',0.001, 'geom', clean := true); \\ NOTICE: PROCESSING:}

NOTICE: pgr_createTopology('edges', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'true', clean := t)
NOTICE: Performing checks, please wait
NOTICE: Creating Topology, Please wait..
NOTICE: -------------> TOPOLOGY CREATED FOR 18 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)
SELECT pgr_analyzeGraph('edges', 0.001, 'geom');
NOTICE: PROCESSING
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','true')
NOTICE: Performing checks, please wait ..
NOTICE: Analyzing for dead ends. Please wait..
NOTICE: Analyzing for gaps. Please wait.
NOTICE: Analyzing for isolated edges. Please wait...
NOTICE: Analyzing for ring geometries. Please wait..
NOTICE: Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 2
NOTICE: Dead ends: 7
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected: 1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

SELECT pgr_analyzeGraph('edges',0.001,'geom', rows_where:='id < 10 ');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','id < 10')
NOTICE: Performing checks, please wait .
NOTICE: Analyzing for dead ends. Please wait...
NOTICE: Analyzing for gaps. Please wait..
NOTICE: Analyzing for isolated edges. Please wait...
NOTICE: Analyzing for ring geometries. Please wait...
NOTICE: Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 0
NOTICE: Dead ends: 4
NOTICE: Potential gaps found near dead ends: 0
NOTICE: Intersections detected: 0
NOTICE: \(\quad\) Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

SELECT pgr_analyzeGraph('edges',0.001,'geom', rows_where:='id >= 10 ');
NOTICE: PROCESSING
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','id >= 10')
NOTICE: Performing checks, please wait ..
NOTICE: Analyzing for dead ends. Please wait...
NOTICE: Analyzing for gaps. Please wait..
NOTICE: Analyzing for isolated edges. Please wait...
NOTICE: Analyzing for ring geometries. Please wait...
NOTICE: Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 2
NOTICE: Dead ends: 8
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected: 1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)
```

SELECT pgr_analyzeGraph('edges',0.001,'geom', rows_where:='id < 17');

```
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','id < 17')
NOTICE: Performing checks, please wait
NOTICE: Analyzing for dead ends. Please wait.
NOTICE: Analyzing for gaps. Please wait...
NOTICE: Analyzing for isolated edges. Please wait..
NOTICE: Analyzing for ring geometries. Please wait.
NOTICE: Analyzing for intersections. Please wait.
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 0
NOTICE: Dead ends:3
NOTICE: Potential gaps found near dead ends: 0
NOTICE: Intersections detected: 0
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)
```

SELECT pgr_createTopology('edges', 0.001,'geom', rows_where:='id <17', clean := true);
NOTICE: PROCESSING:
NOTICE: pgr_createTopology('edges', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'id <17', clean := t)
NOTICE: Performing checks, please wait
NOTICE: Creating Topology, Please wait.
NOTICE: ------------> TOPOLOGY CREATED FOR 16 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE:Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)

```
```

SELECT pgr_analyzeGraph('edges', 0.001, 'geom');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','true')
NOTICE: Performing checks, please wait.
NOTICE: Analyzing for dead ends. Please wait.
NOTICE:Analyzing for gaps. Please wait..
NOTICE:Analyzing for isolated edges. Please wait.
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE: Analyzing for intersections. Please wait.
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:0
NOTICE: Dead ends:3
NOTICE: Potential gaps found near dead ends: 0
NOTICE: Intersections detected: 0
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

```

The examples use the Sample Data network.

See Also

Topology - Family of Functions
- pgr_analyzeOneWay
- pgr_createVerticesTable
- pgr_nodeNetwork to create nodes to a not noded edge table.

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.6 2.5 \(2.42 .3 \mathbf{2}\).2 \(2.1 \mathbf{2 . 0}\)

\section*{pgr analyzeOneWay}
pgr_analyzeOneWay - Analyzes oneway Sstreets and identifies flipped segments.
This function analyzes oneway streets in a graph and identifies any flipped segments.

\section*{Availability}
- Version 2.0.0
- Official function

Description

The analyses of one way segments is pretty simple but can be a powerful tools to identifying some the potential problems created by setting the direction of a segment the wrong way. A node is a source if it has edges the exit from that node and no edges enter that node. Conversely, a node is a sink if all edges enter the node but none exit that node. For aource type node it is logically impossible to exist because no vehicle can exit the node if no vehicle and enter the node. Likewise, if you had a sink node you would have an infinite number of vehicle piling up on this node because you can enter it but not leave it.

So why do we care if the are not feasible? Well if the direction of an edge was reversed by mistake we could generate exactly these conditions. Think about a divided highway and on the north bound lane one segment got entered wrong or maybe a sequence of multiple segments got entered wrong or maybe this happened on a round-about. The result would be potentially a source and/or a sink node.

So by counting the number of edges entering and exiting each node we can identify bothsource and sink nodes so that you can look at those areas of your network to make repairs and/or report the problem back to your data vendor.

\section*{Prerequisites}

The edge table to be analyzed must contain a source column and a target column filled with id's of the vertices of the segments and the corresponding vertices table <edge_table>_vertices_pgr that stores the vertices information.
- Use pgr_createVerticesTable to create the vertices table.
- Use pgr_createTopology to create the topology and the vertices table.

\section*{Signatures}
```

pgr_analyzeOneWay(geom_table, s_in_rules, s_out_rules, t_in_rules, t_out_rules, [options])

```
options: [oneway, source, target, two_way_if_null]
```

RETURNS TEXT

```

\section*{Parameters}

\section*{edge_table:}
text Network table name. (may contain the schema name as well)

\section*{s_in_rules:}
text[] source node in rules

\section*{s_out_rules:}
text[] source node out rules
t_in_rules:
text[] target node in rules
t_out_rules:
text[] target node out rules

\section*{oneway:}
text oneway column name name of the network table. Default value isoneway.

\section*{source:}
text Source column name of the network table. Default value issource.

\section*{target:}
text Target column name of the network table. Default value is target.
two_way_if_null:
boolean flag to treat oneway NULL values as bi-directional. Default value is true.

\section*{Note}

It is strongly recommended to use the named notation. Seepgr_createVerticesTable or pgr_createTopology for examples.

The function returns:
- OK after the analysis has finished.
- Uses the vertices table: <edge_table>_vertices_pgr.
- Fills completely the ein and eout columns of the vertices table.
- FAIL when the analysis was not completed due to an error.
- The vertices table is not found.
- A required column of the Network table is not found or is not of the appropriate type.
- The names of source, target or oneway are the same.

The rules are defined as an array of text strings that if match theoneway value would be counted astrue for the source or target in or out condition.

\section*{The Vertices Table}

The vertices table can be created with pgr_createVerticesTable or pgr_createTopology
The structure of the vertices table is:

\section*{id:}
bigint Identifier of the vertex.
cnt:
integer Number of vertices in the edge_table that reference this vertex. Seepgr_analyzeGgraph. chk:
integer Indicator that the vertex might have a problem. See pgr_analyzeGraph.
ein:
integer Number of vertices in the edge_table that reference this vertex as incoming.

\section*{eout:}
integer Number of vertices in the edge_table that reference this vertex as outgoing.

\section*{the_geom:}
geometry Point geometry of the vertex.

\section*{Additional Examples}
```

ALTER TABLE edges ADD COLUMN dir TEXT;
ALTER TABLE
SELECT pgr_createTopology('edges', 0.001, 'geom');
NOTICE: PROCESSING:
NOTICE: pgr_createTopology('edges', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'true', clean := f)
NOTICE: Performing checks, please wait .
NOTICE: Creating Topology, Please wait..
NOTICE: ------------> TOPOLOGY CREATED FOR 0 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)
UPDATE edges SET
dir = CASE WHEN (cost>0 AND reverse_cost>0) THEN 'B'
WHEN (cost>0 AND reverse_cost<0) THEN 'FT'
WHEN (cost<0 AND reverse_cost>0) THEN 'TF'
ELSE " END;
UPDATE }1
SELECT pgr_analyzeOneWay('edges',
ARRAY[", 'B', 'TF'],
ARRAY[", 'B', 'FT'],
ARRAY[", 'B', 'FT'],
ARRAY[", 'B', 'TF'],
oneway:='dir');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeOneway('edges','{"",B,TF}','{"",B,FT}','{"",B,FT}','{"",B,TF}','dir','source','target',t)
NOTICE:Analyzing graph for one way street errors.
NOTICE:Analysis 25% complete
NOTICE: Analysis 50% complete
NOTICE: Analysis 75% complete
NOTICE: Analysis 100% complete
NOTICE: Found 0 potential problems in directionality
pgr_analyzeoneway
OK
(1 row)

```

\section*{See Also}
```

Topology - Family of Functions
pgr_analyzeGraph
pgr_createVerticesTable
Sample Data

```

\section*{Indices and tables}

Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: \(2.6 \mathbf{2}\).5 2.4 2.3 2.2 2.12 .0
pgr_nodeNetwork
pgr_nodeNetwork - Nodes an network edge table.

\section*{Author:}

Nicolas Ribot
Copyright:
Nicolas Ribot, The source code is released under the MIT-X license.

The function reads edges from a not "noded" network table and writes the "noded" edges into a new table.
```

| pgr_nodenetwork(edge_table, tolerance, [options])
| options: [id, text the_geom, table_ending, rows_where, outall]

```
| RETURNS TEXT

\section*{Availability}
- Version 2.0.0
- Official function

\section*{Description}

\section*{The main characteristics are:}

A common problem associated with bringing GIS data into pgRouting is the fact that the data is often not "noded" correctly. This will create invalid topologies, which will result in routes that are incorrect.

What we mean by "noded" is that at every intersection in the road network all the edges will be broken into separate road segments. There are cases like an over-pass and under-pass intersection where you can not traverse from the over-pass to the under-pass, but this function does not have the ability to detect and accommodate those situations.

This function reads the edge_table table, that has a primary key columnid and geometry column named the_geom and intersect all the segments in it against all the other segments and then creates a table edge_table_noded. It uses the tolerance for deciding that multiple nodes within the tolerance are considered the same node.

\section*{Parameters}

\section*{edge_table:}
text Network table name. (may contain the schema name as well)

\section*{tolerance:}
float8 tolerance for coincident points (in projection unit)dd
id:
text Primary key column name of the network table. Default value isid.
the_geom:
text Geometry column name of the network table. Default value isthe_geom.
table_ending:
text Suffix for the new table's. Default value isnoded.

The output table will have for edge_table_noded
id:
bigint Unique identifier for the table
old_id:
bigint Identifier of the edge in original table
sub_id:
integer Segment number of the original edge
source:
integer Empty source column to be used withpgr_createTopology function
target:
integer Empty target column to be used withpgr_createTopology function
the geom:
geometry Geometry column of the noded network

Let's create the topology for the data in Sample Data
```

SELECT pgr_createTopology('edges', 0.001, 'geom', clean := TRUE);
NOTICE: PROCESSING:
NOTICE: pgr_createTopology('edges', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'true', clean := t)
NOTICE: Performing checks, please wait
NOTICE: Creating Topology, Please wait..
NOTICE.

```
\(\qquad\)
```

> TOPOLOGY CREATED FOR 18 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE:Vertices table for table public.edges is: public.edges vertices pgr
NOTICE:
pgr_createtopology
OK
(1 row)

```

Now we can analyze the network.
```

SELECT pgr_analyzegraph('edges', 0.001, 'geom');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','true')
NOTICE: Performing checks, please wait .
NOTICE: Analyzing for dead ends. Please wait.
NOTICE: Analyzing for gaps. Please wait..
NOTICE:Analyzing for isolated edges. Please wait.
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE:Analyzing for intersections. Please wait.
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:2
NOTICE: Dead ends:7
NOTICE: Potential gaps found near dead ends: 1
NOTICE: Intersections detected: 1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

```

The analysis tell us that the network has a gap and an intersection. We try to fix the problem using:
```

SELECT pgr_nodeNetwork('edges', 0.001, the_geom => 'geom');
NOTICE: PROCESSING:
NOTICE: id: id
NOTICE: the geom: geom
NOTICE: table_ending: noded
NOTICE: rows_where:
NOTICE: outall:f
NOTICE: pgr_nodeNetwork('edges', 0.001, 'id', 'geom', 'noded', ", f)
NOTICE: Performing checks, please wait.
NOTICE: Processing, please wait.
NOTICE: Split Edges:3
NOTICE:Untouched Edges: 15
NOTICE: Total original Edges: 18
NOTICE: Edges generated:6
NOTICE:Untouched Edges:15
NOTICE: Total New segments:21
NOTICE: New Table: public.edges_noded
NOTICE:
pgr_nodenetwork
OK
(1 row)

```

Inspecting the generated table, we can see that edges 13,14 and 18 has been segmented

SELECT old_id, sub_id FROM edges_noded ORDER BY old_id, sub_id;
old_id | sub_id
\begin{tabular}{c:c}
1 & 1 \\
\(2 \mid\) & 1 \\
\(3:\) & 1 \\
4 & 1 \\
5 & 1 \\
6 & 1 \\
7 & 1 \\
8 & 1 \\
9 & 1 \\
10 & 1 \\
11 & 1 \\
12 & 1 \\
13 & 1 \\
13 & 2 \\
14 & 1 \\
14 & 2 \\
15 & 1 \\
16 & 1 \\
17 & 1 \\
18 & 1 \\
18 & 2
\end{tabular}
(21 rows)

We can create the topology of the new network
SELECT pgr_createTopology ('edges_noded', 0.001, 'geom');
NOTICE: PROCESSING:
NOTICE: pgr_createTopology('edges_noded', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'true', clean := f)
NOTICE: Performing checks, please wait .
NOTICE: Creating Topology, Please wait..
NOTICE: -------------> TOPOLOGY CREATED FOR 21 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE: Vertices table for table public.edges_noded is: public.edges_noded_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)

Now let's analyze the new topology
```

SELECT pgr_analyzegraph('edges_noded', 0.001, 'geom');
NOTICE: PROCESSING
NOTICE: pgr_analyzeGraph('edges_noded',0.001,'geom','id','source','target','true')
NOTICE: Performing checks, please wait .
NOTICE:Analyzing for dead ends. Please wait.
NOTICE:Analyzing for gaps. Please wait.,
NOTICE:Analyzing for isolated edges. Please wait...
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE:Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments: 0
NOTICE: Dead ends:6
NOTICE: Potential gaps found near dead ends: 0
NOTICE: Intersections detected: 0
NOTICE: Ring geometries:0
pgr_analyzegraph
OK
(1 row)

```

\section*{Images}

\section*{Before Image}


After Image

Comparing the results

Comparing with the Analysis in the original edge_table, we see that.
\begin{tabular}{|c|c|c|}
\hline & Before & After \\
\hline Table name & edge_table & edge_table_noded \\
\hline Fields & All original fields & Has only basic fields to do a topology analysis \\
\hline Dead ends & Edges with 1
dead end:
\(1,6,24\)
Edges with 2
dead ends
17,18 & \[
\begin{aligned}
& \text { Edges with } 1 \text { dead end: } \\
& 1-1,6-1,14-2,18-117-1 \\
& 18-2
\end{aligned}
\] \\
\hline Isolated segments & two isolated segments: 17 and 18 both they have 2 dead ends & \begin{tabular}{l}
No Isolated segments \\
- Edge 17 now shares a node with edges 141 and 14-2 \\
Edges 18-1 and 18-2 share a node with edges 13-1 and 13-2
\end{tabular} \\
\hline Gaps & There is a gap between edge 17 and 14 because edge 14 is near to the right node of edge 17 & Edge 14 was segmented Now edges: 14-1 14-2 17 share the same node The tolerance value was taken in account \\
\hline Intersections & Edges 13 and 18 were intersecting & Edges were segmented, So, now in the interection's point there is a node and the following edges share it: 13-1 13-2 18-1 18-2 \\
\hline
\end{tabular}

Now, we are going to include the segments 13-1, 13-2 14-1, 14-2 ,18-1 and 18-2 into our edge-table, copying the data for dir,cost, and reverse cost with tho following steps:
- Add a column old_id into edge_table, this column is going to keep track the id of the original edge
- Insert only the segmented edges, that is, the ones whose max(sub_id) >1
```

alter table edges drop column if exists old_id;
NOTICE: column "old_id" of relation "edges" does not exist, skipping
ALTER TABLE
alter table edges add column old_id integer;
ALTER TABLE
insert into edges (old_id, cost, reverse_cost, geom)
(with
segmented as (select old_id,count(*) as i from edges_noded group by old_id)
select segments.old_id, cost, reverse_cost, segments.geom
from edges as edges join edges_noded as segments on (edges.id = segments.old_id)
where edges.id in (select old_id from segmented where i>1) );
INSERT 0 6

```

We recreate the topology:

\section*{SELECT pgr_createTopology('edges', 0.001, 'geom'); \\ NOTICE: PROCESSING}

NOTICE: pgr_createTopology ('edges', 0.001, 'geom', 'id', 'source', 'target', rows_where := 'true', clean := f)
NOTICE: Performing checks, please wait
NOTICE: Creating Topology, Please wait..
NOTICE: -------------> TOPOLOGY CREATED FOR 6 edges
NOTICE: Rows with NULL geometry or NULL id: 0
NOTICE: Vertices table for table public.edges is: public.edges_vertices_pgr
NOTICE:
pgr_createtopology
OK
(1 row)

To get the same analysis results as the topology of edge_table_noded, we do the following query:
```

SELECT pgr_analyzegraph('edges', 0.001, 'geom', rows_where:='id not in (select old_id from edges where old_id is not null)');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target',id not in (select old_id from edges where old_id is not null)')
NOTICE: Performing checks, please wait .
NOTICE: Analyzing for dead ends. Please wait..
NOTICE:Analyzing for gaps. Please wait..
NOTICE:Analyzing for isolated edges. Please wait.
NOTICE:Analyzing for ring geometries. Please wait.,
NOTICE:Analyzing for intersections. Please wait...
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:0
NOTICE: Dead ends:6
NOTICE: Potential gaps found near dead ends: 0
NOTICE: Intersections detected:0
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

```

To get the same analysis results as the original edge_table, we do the following query:
```

SELECT pgr_analyzegraph('edges', 0.001, 'geom', rows_where:='old_id is null');
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','old_id is null')
NOTICE: Performing checks, please wait ..
NOTICE:Analyzing for dead ends. Please wait..
NOTICE:Analyzing for gaps. Please wait...
NOTICE:Analyzing for isolated edges. Please wait...
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE: Analyzing for intersections. Please wait...
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:2
NOTICE: Dead ends:7
NOTICE: Potential gaps found near dead ends: }
NOTICE: Intersections detected:1
NOTICE: Ring geometries: 0
pgr_analyzegraph
OK
(1 row)

```

Or we can analyze everything because, maybe edge 18 is an overpass, edge 14 is an under pass and there is also a street level juction, and the same happens with edges 17 and 13.
```

SELECT pgr_analyzegraph('edges', 0.001, 'geom')
NOTICE: PROCESSING:
NOTICE: pgr_analyzeGraph('edges',0.001,'geom','id','source','target','true')
NOTICE: Performing checks, please wait ..
NOTICE:Analyzing for dead ends. Please wait.
NOTICE:Analyzing for gaps. Please wait.,
NOTICE: Analyzing for isolated edges. Please wait..
NOTICE:Analyzing for ring geometries. Please wait..
NOTICE:Analyzing for intersections. Please wait..
NOTICE: ANALYSIS RESULTS FOR SELECTED EDGES:
NOTICE: Isolated segments:0
NOTICE: Dead ends:3
NOTICE: Potential gaps found near dead ends: 0
NOTICE: Intersections detected:5
NOTICE: Ring geometries:0
pgr_analyzegraph
OK
(1 row)

```

Topology - Family of Functions for an overview of a topology for routing algorithms.pgr_analyzeOneWay to analyze directionality of the edges. pgr_createTopology to create a topology based on the geometry. pgr_analyzeGraph to analyze the edges and vertices of the edge table.

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
pgr_extractVertices - Proposed
pgr_extractVertices - Extracts the vertices information

\section*{Warning}

Proposed functions for next mayor release.
- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.

\section*{Availability}
```

- Version 3.3.0

```
- Classified as proposed function
- Version 3.0.0
- New experimental function

\section*{Description}

This is an auxiliary function for extracting the vertex information of the set of edges of a graph.
- When the edge identifier is given, then it will also calculate the in and out edges

\section*{Signatures}
```

pgr_extractVertices(Edges SQL, [dryrun])
RETURNS SETOF (id, in_edges,out_edges, x, y, geom)
OR EMTPY SET

```

\section*{Example:}

Extracting the vertex information
```

SELECT * FROM pgr_extractVertices
'SELECT id, geom F\overline{ROM edges');}
id|in_edges |out_edges | x | y | geom
1| ll{6} | | 0| 2 |0101000000000000000000000000000000000000040
3|{6} |{7} | 1| 2|0101000000000000000000F03F0000000000000040
4|{17} | | 1.999999999999|3.5|010100000068EEFFFFFFFFFF3F0000000000000C40
5| |{1} | 2| 0|010100000000000000000000400000000000000000
6|{1} |{2,4} | 2 1 1|01010000000000000000000040000000000000F03F
7|{4,7} |{8,10} | 2| 2|010100000000000000000000400000000000000040
8|{10} |{12,14} | 2| 3|010100000000000000000000400000000000000840
9|{14} | | 2 4|010100000000000000000000400000000000001040
10|{2} |{3,5} | 3| 1|01010000000000000000000840000000000000F03F
11|{5,8} |{9,11} | 3| 2|010100000000000000000008400000000000000040
2|{11,12} |{13} | 3| 3|010100000000000000000008400000000000000840
13| |{18} | 3.5|2.3|01010000000000000000000C406666666666660240
14|{18} | | 3.5| 4|01010000000000000000000C400000000000001040
15|{3} |{16} | 4| 1|01010000000000000000001040000000000000F03F
16|{9,16} |{15} | 4 2|010100000000000000000010400000000000000040
17|{13,15} | | 4| 3|010100000000000000000010400000000000000840
(17 rows)

```

\section*{Parameter}
\begin{tabular}{lllll} 
Parameter & Type & Description \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below
\end{tabular}
\end{tabular}

\section*{Optional parameters}
\begin{tabular}{llll} 
Parameter & Type & Default & Description \\
\hline dryrun & BOOLEAN & false & \begin{tabular}{l} 
When true do not process and get in a NOTICE the resulting \\
query.
\end{tabular}
\end{tabular}

\section*{Inner Queries}

\section*{- Edges SQL}
- When line geometry is known
- When vertex geometry is known
- When identifiers of vertices are known

Edges SQL

When line geometry is known
\begin{tabular}{lllll} 
Column & Type & Description & & \\
\hline id & BIGINT & \begin{tabular}{l} 
(Optional) identifier of the \\
edge.
\end{tabular} \\
\hline geom & LINESTRING & Geometry of the edge. & \\
\hline
\end{tabular}

This inner query takes precedence over the next two inner query, therefore other columns are ignored whergeom column appears.
- Ignored columns:
- startpoint
- endpoint
- source
- target

When vertex geometry is known

To use this inner query the columngeom should not be part of the set of columns.
\begin{tabular}{lll} 
Column & Type & Description \\
\hline id & BIGINT & (Optional) identifier of the edge. \\
\hline startpoint & POINT & \begin{tabular}{l} 
POINT geometry of the starting \\
vertex.
\end{tabular} \\
\hline endpoint & POINT & POINT geometry of the ending vertex. \\
\hline
\end{tabular}

This inner query takes precedence over the next inner query, therefore other columns are ignored wherstartpoint and endpoint columns appears.
- Ignored columns:
- source
- target

When identifiers of vertices are known
To use this inner query the columns geom, startpoint and endpoint should not be part of the set of columns.
\begin{tabular}{lll} 
Column & Type & Description \\
\hline id & BIGINT & (Optional) identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & \begin{tabular}{l} 
Identifier of the second end point vertex of the \\
edge.
\end{tabular} \\
\hline
\end{tabular}

Result Columns
\begin{tabular}{|c|c|c|}
\hline Column & Type & Description \\
\hline id & BIGINT & Vertex identifier \\
\hline in_edges & BIGINT[] & \begin{tabular}{l}
Array of identifiers of the edges that have the vertexid as first end point. \\
- NULL When the id is not part of the inner query
\end{tabular} \\
\hline out_edges & BIGINT[] & \begin{tabular}{l}
Array of identifiers of the edges that have the vertexid as second end point. \\
NULL When the id is not part of the inner query
\end{tabular} \\
\hline \(x\) & FLOAT & \begin{tabular}{l}
\(X\) value of the point geometry \\
- NULL When no geometry is provided
\end{tabular} \\
\hline y & FLOAT & \begin{tabular}{l}
\(X\) value of the point geometry \\
- NULL When no geometry is provided
\end{tabular} \\
\hline geom & POINT & \begin{tabular}{l}
Geometry of the point \\
- NULL When no geometry is provided
\end{tabular} \\
\hline
\end{tabular}

\section*{Additional Examples}

\section*{- Dryrun execution}
- Create a routing topology
- Make sure the database does not have the vertices_table
- Clean up the columns of the routing topology to be created
- Create the vertices table
- Inspect the vertices table
- Create the routing topology on the edge table
- Inspect the routing topology
- Crossing edges
- Adding split edges
- Adding new vertices
- Updating edges topology
- Removing the surplus edges
- Updating vertices topology
- Checking for crossing edges
- Graphs without geometries
- Insert the data
- Find the shortest path
- Vertex information

Dryrun execution

To get the query generated used to get the vertex information, usedryrun :=true.
The results can be used as base code to make a refinement based on the backend development needs.
    main_sql AS
        SELECT id, geom FROM edges
    ),
    the_out AS
        SELECT id::BIGINT AS out_edge, ST_StartPoint(geom) AS geom
        FROM main_sql
    ),
    agg_out AS (
        SELECT array_agg(out_edge ORDER BY out_edge) AS out_edges, ST_x(geom) AS x, ST_Y(geom) AS y, geom
        FROM the_out
        GROUP BY geom
    ),
    the_in AS (
        SELECT id::BIGINT AS in_edge, ST_EndPoint(geom) AS geom
        FROM main_sql
    ),
    agg_in AS (
        SELECT array_agg(in_edge ORDER BY in_edge) AS in_edges, ST_x(geom) AS x, ST_Y(geom) AS y, geom
        FROM the in
        GROUP BY geom
    ),
    the_points AS (
        SELECT in_edges, out_edges, coalesce(agg_out.geom, agg_in.geom) AS geom
        FROM agg_out
        FULL OUTER JOIN agg_in USING ( \(x\), y)
    )
    SELECT row_number() over(ORDER BY ST_X(geom), ST_Y(geom)) AS id, in_edges, out_edges, ST_X(geom), ST_Y(geom), geom
    FROM the points;
id | in_edges | out_edges \(|x| y \mid\) geom
(0 rows)

\section*{Create a routing topology}

\section*{Make sure the database does not have the vertices_table}
```

DROP TABLE IF EXISTS vertices table;
NOTICE: table "vertices_table" does not exist, skipping
DROP TABLE

```

Clean up the columns of the routing topology to be created
```

UPDATE edges
SET source = NULL, target = NULL,
x1 = NULL, y1 = NULL,
x2 = NULL, y2 = NULL;
UPDATE }1

```

\section*{Create the vertices table}
- When the LINESTRING has a SRID then use geom::geometry(POINT, <SRID>)
- For big edge tables that are been prepared,
- Create it as UNLOGGED and
- After the table is created ALTER TABLE .. SET LOGGED
```

SELECT * INTO vertices_table
FROM pgr_extractVertices('SELECT id, geom FROM edges ORDER BY id');
SELECT 17

```

\section*{SELECT *}

FROM vertices_table;
id | in_edges | out_edges | x | y | geom
\begin{tabular}{|c|c|}
\hline \(1|\quad|\{6\}\) & 0| \(2 \mid 010100000000000000000000000000000000000040\) \\
\hline \(2|\quad|\{17\}\) & 0.5 | 3.5 | 0101000000000000000000 E 03 F 0000000000000 C 40 \\
\hline \(3|\{6\} \quad|\{7\}\) & 1| 2 |0101000000000000000000F03F0000000000000040 \\
\hline \(4|\{17\} \quad| \quad \mid 1\) & 1.999999999999 | 3.5 | 010100000068 EEFFFFFFFFFF3F0000000000000C40 \\
\hline \(5|\quad|\{1\}\) & \(2|0| 010100000000000000000000400000000000000000\) \\
\hline \(6|\{1\} \quad|\{2,4\}\) & \(2|1| 01010000000000000000000040000000000000 \mathrm{~F} 03 \mathrm{~F}\) \\
\hline \(7|\{4,7\} \quad|\{8,10\}\) & \(2|2| 010100000000000000000000400000000000000040\) \\
\hline \(8|\{10\} \quad|\{12,14\}\) & \(2|3| 010100000000000000000000400000000000000840\) \\
\hline \(9 \mid\{14\}\) & \(2|4| 010100000000000000000000400000000000001040\) \\
\hline \(10|\{2\} \quad|\{3,5\}\) & \(3|1| 01010000000000000000000840000000000000 F 03 F\) \\
\hline \(11|\{5,8\} \quad|\{9,11\}\) & \(3|2| 010100000000000000000008400000000000000040\) \\
\hline \(12|\{11,12\}|\{13\}\) & \(3|3| 010100000000000000000008400000000000000840\) \\
\hline 13| | 18\(\}\) & 3.5 | 2.3 | \(01010000000000000000000 C 406666666666660240\) \\
\hline \(14 \mid\{18\}\) & \(3.5|4| 01010000000000000000000 C 400000000000001040\) \\
\hline 15|\{3\} | \(\{16\}\) & \(4|1| 01010000000000000000001040000000000000 F 03 F\) \\
\hline \(16|\{9,16\}|\{15\}\) & \(4|2| 010100000000000000000010400000000000000040\) \\
\hline \(17|\{13,15\}|\) & \(4|3| 010100000000000000000010400000000000000840\) \\
\hline (17 rows) & \\
\hline
\end{tabular}

\section*{Create the routing topology on the edge table}

\section*{Updating the source information}
```

WITH
out_going AS
SELECT id AS vid, unnest(out_edges) AS eid, x, y
FROM vertices_table
)
UPDATE edges
SET source = vid, x1 = x, y1 = y
FROM out_going WHERE id = eid;
UPDATE }1

```

\section*{Updating the target information}
```

WITH
in_coming AS(
SELECT id AS vid, unnest(in_edges) AS eid, x, y
FROM vertices_table
)
UPDATE edges
SET target = vid, x2 = x, y2 = y
FROM in_coming WHERE id = eid;
UPDATE }1

```

\section*{Inspect the routing topology}



\section*{Generated topology}

\section*{Crossing edges}

To get the crossing edges:
```

SELECT a.id, b.id
FROM edges AS a, edges AS b
WHERE a.id < b.id AND st_crosses(a.geom, b.geom);
id | id
13|18
(1 row)

```


That information is correct, for example, when in terms of vehicles, is it a tunnel or bride crossing over another road.
It might be incorrect, for example:
1. When it is actually an intersection of roads, where vehicles can make turns.
2. When in terms of electrical lines, the electrical line is able to switch roads even on a tunnel or bridge.

When it is incorrect, it needs fixing:
1. For vehicles and pedestrians
- If the data comes from OSM and was imported to the database usingosm2pgrouting, the fix needs to be done in theOSM portal and the data imported again.
- In general when the data comes from a supplier that has the data prepared for routing vehicles, and there is a problem, the data is to be fixed from the supplier
2. For very specific applications
- The data is correct when from the point of view of routing vehicles or pedestrians.
- The data needs a local fix for the specific application.

Once analyzed one by one the crossings, for the ones that need a local fix, the edges need to besplit.
```

SELECT ST_AsText((ST_Dump(ST_Split(a.geom, b.geom))).geom)
FROM edges AS a, edges AS b
WHERE a.id = 13 AND b.id = 18
UNION
SELECT ST_AsText((ST_Dump(ST_Split(b.geom, a.geom))).geom)
FROM edges AS a, edges AS b
WHERE a.id = 13 AND b.id = 18
st_astext
LINESTRING(3.5 2.3,3.5 3)
LINESTRING(3 3,3.5 3)
LINESTRING(3.5 3,4 3)
LINESTRING(3.5 3,3.5 4)
(4 rows)

```

The new edges need to be added to the edges table, the rest of the attributes need to be updated in the new edges, the old edges need to be removed and the routing topology needs to be updated.

\section*{Adding split edges}

For each pair of crossing edges a process similar to this one must be performed.
The columns inserted and the way are calculated are based on the application. For example, if the edges have a trainame, then that column is to be copied.

\section*{For pgRouting calculations}
- factor based on the position of the intersection of the edges can be used to adjust thecost and reverse_cost columns.
- Capacity information, used on the Flow - Family of functions functions does not need to change when splitting edges.
```

WITH
first_edge AS(
SELECT (ST_Dump(ST_Split(a.geom, b.geom))).path[1],
(ST_Dump(ST_Split(a.geom, b.geom))).geom,
ST_LineLocatePoint(a.geom,ST_Intersection(a.geom,b.geom)) AS factor
FROM edges AS a, edges AS b
WHERE a.id = 13 AND b.id = 18),
first_segments AS
SELECT path, first_edge.geom,
capacity, reverse_capacity,
CASE WHEN path=1 THEN factor * cost
ELSE (1 - factor) * cost END AS cost,
CASE WHEN path=1 THEN factor * reverse_cost
ELSE (1 - factor) * reverse_cost END AS reverse_cost
FROM first_edge , edges WHERE id = 13),
second_edge AS
SELECT (ST_Dump(ST_Split(b.geom, a.geom))).path[1],
(ST_Dump(ST_Split(b.geom, a.geom))).geom,
ST_LineLocatePoint(b.geom,ST_Intersection(a.geom,b.geom)) AS factor
FROM edges AS a, edges AS b
WHERE a.id = 13 AND b.id = 18)
second_segments AS
SELECT path, second_edge.geom,
capacity, reverse_capacity,
CASE WHEN path=1 THEN factor * cost
ELSE (1 - factor) * cost END AS cost,
CASE WHEN path=1 THEN factor * reverse_cost
ELSE (1 - factor) * reverse cost END AS reverse cost
FROM second_edge, edges WHERE id = 18),
all_segments AS (
SELECT * FROM first_segments
UNION
SELECT * FROM second_segments)
INSERT INTO edges
(capacity, reverse_capacity,
cost, reverse cost,
x1, y1, x2, y2,
geom)
(SELECT capacity, reverse_capacity, cost, reverse_cost,
ST X(ST StartPoint(geom)), ST Y(ST StartPoint(geom)),
ST_X(ST_EndPoint(geom)), ST_Y(ST_EndPoint(geom)),
geom
FROM all_segments)
INSERT 0 4

```

\section*{Adding new vertices}

After adding all the split edges required by the application, the newly created vertices need to be added to the vertices table.

\section*{Updating edges topology}
```

/* -- set the source information */
UPDATE edges AS e
SET source = v.id
FROM vertices AS v
WHERE source IS NULL AND ST_StartPoint(e.geom) = v.geom
UPDATE 4
/* -- set the target information */
UPDATE edges AS e
SET target = v.id
FROM vertices AS v
WHERE target IS NULL AND ST_EndPoint(e.geom) = v.geom;
UPDATE 4

```

\section*{Removing the surplus edges}

Once all significant information needed by the application has been transported to the new edges, then the crossing edges can be deleted.
```

DELETE FROM edges WHERE id IN (13, 18);
DELETE 2

```

There are other options to do this task, like creating a view, or a materialized view.

\section*{Updating vertices topology}

To keep the graph consistent, the vertices topology needs to be updated
```

UPDATE vertices AS v SET
in_edges = nv.in_edges, out_edges = nv.out_edges
FROM (SELECT * FROM pgr_extractVertices('SELECT id, geom FROM edges')) AS nv
WHERE v.geom = nv.geom
UPDATE }1

```

\section*{Checking for crossing edges}

There are no crossing edges on the graph.
```

SELECT a.id, b.id
FROM edges AS a, edges AS b
WHERE a.id < b.id AND st_crosses(a.geom, b.geom);
id | id
(0 rows)

```

\section*{Graphs without geometries}

Using this table design for this example:
```

CREATE TABLE wiki (
id SERIAL,
source INTEGER,
target INTEGER,
cost INTEGER);
CREATE TABLE

```

Insert the data
```

INSERT INTO wiki (source, target, cost) VALUES
(1,2, 7), (1, 3, 9), (1, 6, 14),
(2, 3, 10), (2, 4, 15),
(3,6,2), (3,4, 11),
(4, 5, 6),
(5, 6, 9);
(,6, );

```

To solve this example pgr_dijkstra is used:
```

SELECT * FROM pgr_dijkstra(
'SELECT id, source, target, cost FROM wiki',
1,5, false);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 1| 2| 9| 0
2| 2| 3| 6| 2| 9
3| 3| 6| 9| 9| 11
4|

```

To go from \(\backslash(1 \backslash)\) to \(\backslash(5 \backslash)\) the path goes thru the following vertices: \(\backslash(1\) rightarrow 3 \rightarrow 6 \rightarrow \(5 \backslash)\)


Vertex information

To obtain the vertices information, use pgr_extractVertices - Proposed
```

SELECT id, in_edges, out_edges
FROM pgr_extractVertices('SELECT id, source, target FROM wiki');
id | in_edges | out_edges
3|{2,4} |{6,7}
5|{8} |{9}
4|{5,7} |{8}
2|{1} |{4,5}
1| |{1,2,3}
6 |{3,6,9} |
(6 rows)

```

\section*{See Also}
- Topology - Family of Functions
- pgr_createVerticesTable

\section*{Indices and tables}
- Index
- Search Page

See Also

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: \(2.6 \mathbf{2}\).5 2.42 .3

\section*{Traveling Sales Person - Family of functions}
- pgr_TSP - When input is given as matrix cell information.
- pgr_TSPeuclidean - When input are coordinates.

\footnotetext{
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: \(2.6 \mathbf{2}\).5 2.4 2.3 2.2 2.12 .0
}

\section*{Availability:}
- Version 3.2.1
- Metric Algorithm from Boost library
- Simulated Annealing Algorithm no longer supported
- The Simulated Annealing Algorithm related parameters are ignored: max_processing_time, tries_per_temperature, max_changes_per_temperature, max_consecutive_non_changes, initial_temperature, final_temperature, cooling_factor, randomize
- Version 2.3.0
- Signature change
- Old signature no longer supported
- Version 2.0.0
- Official function

Description

Problem Definition

The travelling salesperson problem (TSP) asks the following question:
Given a list of cities and the distances between each pair of cities, which is the shortest possible route that visits each city exactly once and returns to the origin city?

\section*{Characteristics}
- This problem is an NP-hard optimization problem.
- Metric Algorithm is used
- Implementation generates solutions that are twice as long as the optimal tour in the worst casewhen:
- Graph is undirected
- Graph is fully connected
- Graph where traveling costs on edges obey the triangle inequality.
- On an undirected graph:
- The traveling costs are symmetric:
- Traveling costs fromu to v are just as much as traveling fromv to u
- Can be Used with Cost Matrix - Category functions preferably with directed => false.
- With directed => false
- Will generate a graph that:
- is undirected
- is fully connected (As long as the graph has one component)
- all traveling costs on edges obey the triangle inequality.
- When start_vid \(=0\) OR end_vid \(=0\)
- The solutions generated is garanteed to betwice as long as the optimal tour in the worst case
- When start_vid !=0 AND end_vid != 0 AND start_vid != end_vid
- It is not garanteed that the solution will be, in the worse case, twice as long as the optimal tour, due to the fact that end_vid is forced to be in a fixed position.
- With directed => true
- It is not garanteed that the solution will be, in the worse case, twice as long as the optimal tour
- Will generate a graph that:
- is directed
- is fully connected (As long as the graph has one component)
- some (or all) traveling costs on edges might not obey the triangle inequality.
- As an undirected graph is required, the directed graph is transformed as follows:
- edges \((u, v)\) and \((v, u)\) is considered to be the same edge (denoted \((u, v)\)
- if agg_cost differs between one or more instances of edge \((u, v)\)
- The minimum value of theagg_cost all instances of edge \((u, v)\) is going to be considered as theagg_cost of edge (u, v)
- Some (or all) traveling costs on edges will still might not obey the triangle inequality.
- When the data is incomplete, but it is a connected graph:
- the missing values will be calculated with dijkstra algorithm.

\section*{Summary}
```

pgr_TSP(Matrix SQL, [start_id, end_id])
RETURNS SET OF (seq, node, cost, agg_cost)
OR EMTPY SET

```

\section*{Example:}

Using pgr_dijkstraCostMatrix to generate the matrix information
- Line 4 Vertices \(\backslash(\backslash\{2,4,13,14 \backslash\} \backslash)\) are not included because they are not connected.
```

1 SELECT * FROM pgr_TSP(

$$
SELECT * FROM pgr_dijkstraCostMatrix(
        'SELECT id, source, target, cost, reverse_cost FROM edges',
        (SELECT array_agg(id) FROM vertices WHERE id NOT IN (2, 4, 13, 14)),
        directed => false)
$$);

seq | node | cost | agg_cost
1| 1|-----+-----
1| 1| 0| 0
2| 3| 1| 1
3| 7| 1| 2
4| 6| 1| 3
5| 5| 1| 4
6| 10| 2| 6
7| 11| 1| 7
8| 12| 1| 8
9| 16| 2| 10
10| 15| 1| 11
11| 17| 2| 13
12| 9| 3| 16
13| 8| 1| 17
14| 1| 3| 20
22 (14 rows)
23

```

\section*{Parameters}
\begin{tabular}{lllll} 
Parameter & Type & Description \\
\hline Matrix SQL & TEXT & \begin{tabular}{l} 
Matrix SQL as described \\
below
\end{tabular} & & \\
\hline
\end{tabular}

TSP optional parameters
\begin{tabular}{|c|c|c|c|}
\hline Column & Type & Default & Description \\
\hline \multirow[t]{2}{*}{start_id} & ANY-INTEGER & 0 & The first visiting vertex \\
\hline & & & - When 0 any vertex can become the first visiting vertex. \\
\hline \multirow[t]{3}{*}{end_id} & ANY-INTEGER & 0 & Last visiting vertex before returning tostart_vid. \\
\hline & & & When 0 any vertex can become the last visiting vertex before returning to start_id. \\
\hline & & & - When NOT0 and start_id = 0 then it is the first and last vertex \\
\hline
\end{tabular}

\section*{Inner Queries}

\section*{Matrix SQL}

Matrix SQL: an SQL query, which should return a set of rows with the following columns:
\begin{tabular}{lllll} 
Column & Type & Description \\
\hline start_vid & ANY-INTEGER & Identifier of the starting vertex. & \\
\hline end_vid & ANY-INTEGER & Identifier of the ending vertex. & \\
\hline agg_cost & ANY-NUMERICAL & \begin{tabular}{l} 
Cost for going from start_vid \\
end_vid
\end{tabular} & & \\
\hline
\end{tabular}

\section*{Result Columns}

Returns SET OF (seq, node, cost, agg_cost)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INTEGER & Row sequence
\end{tabular}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline node & BIGINT & Identifier of the node/coordinate/point. \\
\hline cost & FLOAT & \begin{tabular}{l} 
Cost to traverse from the current node to the next node in the path \\
sequence.
\end{tabular} \\
\hline agg_cost & FLOAT & \begin{tabular}{l} 
Aggregate cost from the node at seq \(=1\) to the current node. \\
\\
\end{tabular} \\
& \(0 \quad 0\) for the first row in the tour sequence.
\end{tabular}

\section*{Additional Examples}
```

- Start from vertex $1$
- Using points of interest to generate an asymetric matrix.
- Connected incomplete data

```
```

Start from vertex $1$

```
- Line \(\mathbf{6}\) start_vid => 1
```

1 SELECT * FROM pgr_TSP(

$$
SELECT * FROM pgr_dijkstraCostMatrix(
    'SELECT id, source, target, cost, reverse_cost FROM edges',
    (SELECT array_agg(id) FROM vertices WHERE id NOT IN (2, 4, 13, 14)),
    directed => false)
$$,

    start_id => 1);
    seq | node | cost | agg_cost
    1| 1| 0| 0
    2| 3| 1| 1
    3| 7| 1| 2
    4| 6| 1| 3
    5| 5| 1| 4
    6| 10| 2| 6
    7| 11| 1| 7
    8| 12| 1| 8
    9| 16| 2| 10
    10| 15| 1| 11
    11| 17| 2| 13
    20 12| 9| 3| 16
    21 13| 8| 1| 17
    14| 1| 3| 20
23 (14 rows)
24

```

Using points of interest to generate an asymetric matrix.

To generate an asymmetric matrix:
- Line 4 The side information of pointsOflnterset is ignored by not including it in the query
- Line \(\mathbf{6}\) Generating an asymetric matrix with directed => true
- \(\backslash(\min (a g g \backslash c o s t(u, v), a g g \backslash \operatorname{cost}(v, u)) \backslash)\) is going to be considered as theagg_cost
- The solution that can be larger thantwice as long as the optimal tour because:
- Triangle inequality might not be satisfied.
- start_id ! \(=0\) AND end_id ! \(=0\)
```

1 SELECT * FROM pgr_TSP(

$$
SELECT * FROM pgr_withPointsCostMatrix(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction from pointsOfInterest',
    array[-1, 10, 7, 11, -6],
    directed => true)
$$,

    start_id => 7,
    end_id => 11);
    seq | node | cost | agg_cost
    --------------------------
11 1| 7| 0| 0
12 2| -6| 0.3| 0.3
3 3| -1| 1.3| 1.6
14 4| 10| 1.6| 3.2
5 5| 11| 1| 4.2
16 6| 7| 1| 5.2
17 (6 rows)
18

```

\section*{Connected incomplete data}
```

1 SELECT * FROM pgr_dijkstraCostMatrix(
$q1$SELECT id, source, target, cost, reverse_cost FROM edges WHERE id IN (2, 4, 5, 8, 9, 15)$q1$,
(SELECT ARRAY[6, 7, 10, 11, 16, 17]),
directed => true);
start_vid | end_vid | agg_cost
6 ------------------+-------
6| 11| 2
6| 16| 3
6| 17| 4
7 6|
11| 1
7) 17 2
10| 6| 1
0| 7| 2
0| 11| 1
0| 16| 2
0| 17| 3
11| 6| 2
7|
1| 16| 1
1| 17| 2
6| 6| 3
6| 7| 2
6| 11| 1
6| 17 |
17 6| 4
17|
17| 16| 1
(25 rows)
33

```

Cost value for \\(17 \rightarrow 10\\) do not exist on the matrix, but the value used is taken froml(10 \rightarrow 17\\).
```

1 SELECT * FROM pgr_TSP(

$$
SELECT * FROM pgr_dijkstraCostMatrix(
$q1$SELECT id, source, target, cost, reverse_cost FROM edges WHERE id IN (2, 4, 5, 8, 9, 15)$q1$,
(SELECT ARRAY[6, 7, 10, 11, 16, 17]),
directed => true)
$$);

seq | node | cost | agg_cost
------------------------
8 1| 6| 0| 0
9 2| 7| 1| 1
10 3| 11| 1| 2
11 4| 16| 1| 3
12 5| 17| 1| 4
13 6| 10| 3| 7
14 7| 6| 1| 8
5 (7 rows)
16

```

See Also
- Traveling Sales Person - Family of functions
- Sample Data
- Boost's metric appro's metric approximation
- Wikipedia: Traveling Salesman Problem

\section*{Indices and tables}
- Index
- Search Page

Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.5 2.4 2.3
pgr_TSPeuclidean
- pgr_TSPeuclidean - Aproximation using metric algorithm.

Boost Graph Inside
- Version 3.2.1
- Metric Algorithm from Boost library
- Simulated Annealing Algorithm no longer supported
- The Simulated Annealing Algorithm related parameters are ignored:max_processing_time, tries_per_temperature, max_changes_per_temperature, max_consecutive_non_changes, initial_temperature, final_temperature, cooling_factor, randomize
Version 3.0.0
- Name change from pgr_eucledianTSP
- Version 2.3.0
- New Official function

\section*{Description}

\section*{Problem Definition}

The travelling salesperson problem (TSP) asks the following question:
Given a list of cities and the distances between each pair of cities, which is the shortest possible route that visits each city exactly once and returns to the origin city?

\section*{Characteristics}
- This problem is an NP-hard optimization problem.
- Metric Algorithm is used
- Implementation generates solutions that are twice as long as the optimal tour in the worst casewhen:
- Graph is undirected
- Graph is fully connected
- Graph where traveling costs on edges obey the triangle inequality.
- On an undirected graph:
- The traveling costs are symmetric:
- Traveling costs from \(u\) to \(v\) are just as much as traveling fromv to u
- Any duplicated identifier will be ignored. The coordinates that will be kept is arbitrarly.
- The coordinates are quite similar for the same identifier, for example

1, 3.5, 1
1,3.499999999999 0.9999999
- The coordinates are quite different for the same identifier, for example
```

2, 3.5,1.0

```
2, 3.6, 1.1

\section*{Signatures}

\section*{Summary}
pgr_TSPeuclidean(Coordinates SQL, [start_id, end_id])
RETURNS SET OF (seq, node, cost, agg_cost)
OR EMTPY SET

\section*{Example:}

With default values
```

SELECT * FROM pgr_TSPeuclidean
SEL
SELECT id, st_X(geom) AS x, st_Y(geom)AS y FROM vertices

\$\$);
seq | node | cost | agg_cost
1| 1| 0| 0
2| 6| 2.2360679775| 2.2360679775
3| 5| 1| 3.2360679775
10| 1.41421356237 |.65028153987
7| 1.41421356237|6.06449510225
| 2.12132034356 8.1858154458
9| 1.58113883008|9.76695427589
4| 0.5|10.2669542759
9| 14| 1.58113883009 | 11.848093106
10 | 17 | 1.11803398875 | 12.9661270947
11| 16| 1|13.9661270947
12| 15| 1|14.9661270947
13| 11| 1.41421356237 | 16.3803406571
14| 13 |.0.583095189485 | 16.9634358466
15| 12|0.860232526704 | 17.8236683733
16| 8| 1| 18.8236683733
17| 3| 1.41421356237 | 20.2378819357
18| 1| 1|21.2378819357
(18 rows)
```

Parameters

| Parameter | Type | Description |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| Coordinates SQL | TEXT | Coordinates <br> below | SQL as | described |  |

TSP optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| start_id | ANY-INTEGER | 0 | The first visiting vertex |
|  |  |  | - When 0 any vertex can become the first visiting vertex. |

Inner Queries

Coordinates SQL

Coordinates SQL: an SQL query, which should return a set of rows with the following columns:

| Column | Type | Description |
| :--- | :--- | :--- | :--- |
| $\mathbf{i d}$ | ANY-INTEGER | Identifier of the starting <br> vertex. |
| $\mathbf{x}$ | ANY-NUMERICAL | X value of the coordinate. |
| $\mathbf{y}$ | ANY-NUMERICAL | $Y$ value of the coordinate. |

## Result Columns

Returns SET OF (seq, node, cost, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Row sequence. |
| node | BIGINT | Identifier of the node/coordinate/point. |
| cost | FLOAT | Cost to traverse from the current node to the next node in the path <br> sequence. |
|  |  | $0 \quad 0$ for the last row in the tour sequence. |

## Additional Examples

- Creating a table for the data and storing the data
- Adding a geometry (for visual purposes)
- Total tour cost
- Getting a geometry of the tour
- Visual results

Test 29 cities of Western Sahara

This example shows how to make performance tests using University of Waterloo'sexample data using the 29 cities of Western Sahara dataset

Creating a table for the data and storing the data

```
CREATE TABLE wi29 (id BIGINT, x FLOAT, y FLOAT, geom geometry);
INSERT INTO wi29 (id, x, y) VALUES
```

(1,20833.3333,17100.0000),
(2,20900.0000,17066.6667),
$(3,21300.0000,13016.6667)$,
(4,21600.0000,14150.0000),
(5,21600.0000,14966.6667), (6,21600.0000,16500.0000), (7,22183.3333,13133.3333), (8,22583.3333,14300.0000), (9,22683.3333,12716.6667), (10,23616.6667,15866.6667), (11,23700.0000,15933.3333), (12,23883.3333,14533.3333), (13,24166.6667,13250.0000), (14,25149.1667,12365.8333), (15,26133.3333,14500.0000), (16,26150.0000,10550.0000), (17,26283.3333,12766.6667), (18,26433.3333,13433.3333), (19,26550.0000,13850.0000), (20,26733.3333,11683.3333), (21,27026.1111,13051.9444), (22,27096.1111,13415.8333), (23,27153.6111,13203.3333), (24,27166.6667,9833.3333), (25,27233.3333,10450.0000), (26,27233.3333,11783.3333), (27,27266.6667,10383.3333), ( $28,27433.3333,12400.0000$ ), (29,27462.5000,12992.2222);

## Adding a geometry (for visual purposes)

```
UPDATE wi29 SET geom = ST_makePoint(x,y);
```


## Total tour cost

Getting a total cost of the tour, compare the value with the length of an optimal tour is 27603 , given on the dataset

```
SELECT
FROM pgr_TSPeuclidean($$SELECT * FROM wi29$$)
WHERE seq = 30;
seq | node | cost | agg_cost
30| 1 | 2266.91173136 | 28777.4854127
(1 row)
```

Getting a geometry of the tour

WITH
tsp_results AS (SELECT seq, geom FROM pgr_TSPeuclidean(\$\$SELECT * FROM wi29\$\$) JOIN wi29 ON (node = id))
SELECT ST_MakeLine(ARRAY(SELECT geom FROM tsp_results ORDER BY seq));

01020000001E000000F085C9545558D4400000000000B3D040000000000069D440107A36ABAAAAD040000000000018D54000000000001DD040107A36AB2A10D (1 row)

## Visual result

Visualy, The first image is the optimal solution and the second image is the solution obtained with pgr_TSPeuclidean.


## See Also

- Traveling Sales Person - Family of functions
- Sample Data network
- Boost's metric appro's metric approximation
- University of Waterloo TSP
- Wikipedia: Traveling Salesman Problem


## Indices and tables

```
- Index
- Search Page
```


## Table of Contents

- General Information
- Problem Definition
- Origin
- Characteristics
- TSP optional parameters
- See Also

General Information

Problem Definition

The travelling salesperson problem (TSP) asks the following question:
Given a list of cities and the distances between each pair of cities, which is the shortest possible route that visits each city exactly once and returns to the origin city?

Origin

The traveling sales person problem was studied in the 18th century by mathematiciansSir William Rowam Hamilton and Thomas Penyngton Kirkman.

A discussion about the work of Hamilton \& Kirkman can be found in the bookgraph Theory (Biggs et al. 1976).

- ISBN-13: 978-0198539162
- ISBN-10: 0198539169

It is believed that the general form of the TSP have been first studied by Kalr Menger in Vienna and Harvard. The problem was later promoted by Hassler, Whitney \& Merrill at Princeton. A detailed description about the connection between Menger \& Whitney, and the development of the TSP can be found in On the history of combinatorial optimization (till 1960)

To calculate the number of different tours through $\backslash(\mathrm{n} \backslash)$ cities:

- Given a starting city,
- There are $\backslash(n-1 \backslash)$ choices for the second city,
- And $\backslash(n-2 \backslash)$ choices for the third city, etc.
- Multiplying these together we get $\backslash((n-1)!=(n-1)(n-2) \ldots 1 \backslash)$.
- Now since the travel costs do not depend on the direction taken around the tour:
- this number by 2
- $\backslash((n-1)!/ 2 \backslash)$.


## Characteristics

- This problem is an NP-hard optimization problem.
- Metric Algorithm is used
- Implementation generates solutions that are twice as long as the optimal tour in the worst casewhen:
- Graph is undirected
- Graph is fully connected
- Graph where traveling costs on edges obey the triangle inequality.
- On an undirected graph:
- The traveling costs are symmetric:
- Traveling costs from $u$ to $v$ are just as much as traveling fromv to u

TSP optional parameters

| Column | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| start_id | ANY-INTEGER | 0 | The first visiting vertex |
|  |  |  | - When 0 any vertex can become the first visiting vertex. |
| end_id | ANY-INTEGER | 0 | Last visiting vertex before returning tostart_vid. |
|  |  |  | - When 0 any vertex can become the last visiting vertex before returning to start_id. |
|  |  |  | - When NOT0 and start_id = 0 then it is the first and last vertex |

See Also

## References

- Boost's metric appro's metric approximation
- University of Waterloo TSP
- Wikipedia: Traveling Salesman Problem


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0

Spanning Tree - Category

- Kruskal - Family of functions
- Prim - Family of functions

A spanning tree of an undirected graph is a tree that includes all the vertices of $G$ with the minimum possible number of edges.
For a disconnected graph, there there is no single tree, but a spanning forest, consisting of a spanning tree of each connected component.

## Characteristics:

- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.


## See Also

Boost: Prim's algorithm

- Boost: Kruskal's algorithm
- Wikipedia: Prim's algorithm
- Wikipedia: Kruskal's algorithm


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1) 3.0
- Unsupported versions: 2.5 2.4 2.6

K shortest paths - Category

- pgr_KSP - Yen's algorithm based on pgr_dijkstra


## Proposed

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.
- pgr_withPointsKSP - Proposed - Yen's algorithm based on pgr_withPoints


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: $2.6 \mathbf{2}$.5 2.4 2.3 2.2 2.12 .0


## pgr_trsp - Turn Restriction Shortest Path (TRSP)

pgr_trsp - Returns the shortest path with support for turn restrictions.

## Availability

- Version 2.1.0
- New Via prototypes
- pgr_trspViaVertices
- pgr_trspViaEdges
- Version 2.0.0
- Official function


## Description

The turn restricted shorthest path (TRSP) is a shortest path algorithm that can optionally take into account complicated turn restrictions like those found in real world navigable road networks. Performamnce wise it is nearly as fast as the $A^{*}$ search but has many additional features like it works with edges rather than the nodes of the network. Returns a set of (seq, id1, id2, cost) or (seq, id1, id2, id3, cost) rows, that make up a path.

```
pgr_trsp(sql text, source integer, target integer,
    directed boolean, has_rcost boolean [,restrict_sql text]);
RETURNS SETOF (seq, id1, id2, cost)
```

```
pgr_trsp(sql text, source_edge integer, source_pos float8,
    target_edge integer, target pos float8,
    directed boolean, has_rcost boolean [,restrict_sql text]);
RETURNS SETOF (seq, id1, id2, cost)
```

pgr_trspViaVertices(sql text, vids integer[],
directed boolean, has_rcost boolean
[, turn_restrict_sql text]);
RETURNS SETOF (seq, id1, id2, id3, cost)

```
pgr_trspViaEdges(sql text, eids integer[], pcts float8[],
    directed boolean, has_rcost boolean
    [, turn_restrict_sql text]);
RETURNS SETOF (seq, id1, id2, id3, cost)
```


## The main characteristics are:

The Turn Restricted Shortest Path algorithm (TRSP) is similar to the shooting star in that you can specify turn restrictions.
The TRSP setup is mostly the same asDijkstra shortest path with the addition of an optional turn restriction table. This provides an easy way of adding turn restrictions to a road network by placing them in a separate table.
sql:
a SQL query, which should return a set of rows with the following columns:

## SELECT id, source, target, cost, [, reverse_cost] FROM edge_table

## id:

int4 identifier of the edge
source:
int4 identifier of the source vertex

## target:

int4 identifier of the target vertex
cost:
float8 value, of the edge traversal cost. A negative cost will prevent the edge from being inserted in the graph.
reverse_cost:
(optional) the cost for the reverse traversal of the edge. This is only used when thedirected and has_rcost parameters are true (see the above remark about negative costs).

## source:

int4 NODE id of the start point
target:
int4 NODE id of the end point
directed:
true if the graph is directed
has_rcost:
if true, the reverse_cost column of the SQL generated set of rows will be used for the cost of the traversal of the edge in the opposite direction.

## restrict_sql:

(optional) a SQL query, which should return a set of rows with the following columns:

```
SELECT to_cost, target_id, via_path FROM restrictions
```


## to_cost:

float8 turn restriction cost
target_id:
int4 target id

## via_path:

text comma separated list of edges in the reverse order ofrule

Another variant of TRSP allows to specify EDGE id of source and target together with a fraction to interpolate the position:

```
source_edge:
```

int4 EDGE id of the start edge
source_pos:
float8 fraction of 1 defines the position on the start edge
target_edge:
int4 EDGE id of the end edge

## target_pos:

float8 fraction of 1 defines the position on the end edge
Returns set of:

## seq:

row sequence
id1:
node ID
id2:
edge ID (-1 for the last row)
cost:

Support for Vias

## Warning

The Support for Vias functions are prototypes. Not all corner cases are being considered.

We also have support for vias where you can say generate a from A to B to C, etc. We support both methods above only you pass an array of vertices or and array of edges and percentage position along the edge in two arrays.
sql:
a SQL query, which should return a set of rows with the following columns:

## SELECT id, source, target, cost, [,reverse cost] FROM edge table

id:
int4 identifier of the edge

## source:

int4 identifier of the source vertex
target:
int4 identifier of the target vertex
cost:
float8 value, of the edge traversal cost. A negative cost will prevent the edge from being inserted in the graph

## reverse_cost:

(optional) the cost for the reverse traversal of the edge. This is only used when thedirected and has_rost parameters are true (see the above remark about negative costs).

## vids:

int4[] An ordered array of NODE id the path will go through from start to end.
directed:
true if the graph is directed
has_rcost:
if true, the reverse_cost column of the SQL generated set of rows will be used for the cost of the traversal of the edge in the opposite direction.
restrict_sql:
(optional) a SQL query, which should return a set of rows with the following columns:

```
SELECT to_cost, target_id, via_path FROM restrictions
```


## to_cost:

float8 turn restriction cost

## target_id:

int4 target id
via_path:
text commar separated list of edges in the reverse order ofrule

Another variant of TRSP allows to specify EDGE id together with a fraction to interpolate the position:
eids:
int4 An ordered array of EDGE id that the path has to traverse
pcts:
float8 An array of fractional positions along the respective edges ineids, where 0.0 is the start of the edge and 1.0 is the end of the eadge.

Returns set of:

## seq:

row sequence
id1:
route ID
id2:
node ID
id3:
edge ID (-1 for the last row)
cost:
cost to traverse from id2 using id3

## Example:

Without turn restrictions

```
SELECT * FROM pgr_trsp(
    'SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost FROM edges',
    1,17, false, false
);
seq id1 | id2 | cos
    0| 1| 6| 1
    1| 3| 7| 1
    2| 7| 8|
    3| 11| 9| 1
    4|16| 15| 1
    5| 17 -1 0
(6 rows)
```


## Example:

With turn restrictions
Then a query with turn restrictions is created as:

```
SELECT * FROM pgr_trsp(
    'SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost FROM edges',
    6, 1, false, false,
    SELECT to cost, target id::int4,
    from_edge || coalesce("," || via_path, "") AS via_path
    FROM restrictions'
);
seq|id1 | id2 | cost
    0| 6| 4| 1
    1| 7| 10| 1
    2| 8| 12| 
    3| 12| 11| 1
    4| 11| 8| 1
    5| 7| 7| 1
    6| 3| 6| 1
    7| 1|-1| 0
(8 rows)
SELECT * FROM pgr_trsp(
    'SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost FROM edges',
    1, 12, false, false,
    SELECT to_cost, target_id::int4,
    from_edge || coalesce("," || via_path, "") AS via_path
    FROM restrictions
);
seq | id1 | id2 | cos
    0| 1| 6| 1
    1| 3| 7| 1
    2| 7| 8| 1
    3| 11| 9| 1
    4| 16| 15| 1
    5| 17| 13| 1
    6| 12|-1| 0
(7 rows)
```

An example query using vertex ids and via points:

SELECT * FROM pgr_trspViaVertices(
'SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost FROM edges',
ARRAY[6,1,12]::INTEGER[],
false, false,
'SELECT to_cost, target_id::int4, from_edge ||
coalesce(","||via_path,"'") AS via_path FROM restrictions');
seq | id1 | id2 | id3 | cost
--------+-----+----+--
2| $1|7| 10 \mid 1$
3|1| 8|12| 1
$4|1| 12|11| 1$
$5|1| 11|8| 1$
$6|1| 7|7| 1$
$7|1| 3|6| 1$

9| $2|3| 7 \mid 1$
$10|2| 7|8| 1$
11| $2|11| 9 \mid 1$
$12|2| 16|15| 1$
13| $2|17| 13 \mid 1$
$14|2| 12|-1| 0$
(14 rows)

An example query using edge ids and vias:

```
SELECT * FROM pgr_trspViaEdges(
    'SELECT id::INTEGER, source::INTEGER, target:INTEGER, cost,
    reverse_cost FROM edges',
    ARRAY[2,7,11]::INTEGER[]
    ARRAY[0.5, 0.5, 0.5]::FLOAT[],
    true,
    true,
    'SELECT to_cost, target_id::int4, from_edge ||
    coalesce(","||via_path,"") AS via_path FROM restrictions');
seq | id1 | id2 | id3 | cost
    1| 1|-1| 2| 0.5
    | 1| 6| 4| 1
    | 1| 7| 8| 1
    | 1| 11| 9| 1
    1| 16| 16| 1
    1| 15| 3| 1
    1| 10| 5| 1
    1| 11| 8| 1
    1| 7| 7| 1
    2| 7| 8| 1
    | 2| 11| 9| 1
    | 2| 16| 16| 1
    2| 15| 3| 1
    2| 10| 5| 1
    15| 2| 11| 11|0.5
(15 rows)
```

The queries use the Sample Data network.

## Known Issues

## Introduction

pgr_trsp code has issues that are not being fixed yet, but as time passes and new functionality is added to pgRouting with wrappers to hide the issues, not to fix them.

For clarity on the queries:

- _pgr_trsp (internal_function) is the original code
- pgr_trsp (lower case) represents the wrapper calling the original code
- pgr_TRSP (upper case) represents the wrapper calling the replacement function, depending on the function, it can be:
- pgr_dijkstra
- pgr_dijkstraVia
- pgr_withPoints
- _pgr_withPointsVia (internal function)


## The restrictions

The restriction used in the examples does not have to do anything with the graph:

- No vertex has id: 25,32 or 33
- No edge has id: 25,32 or 33

A restriction is assigned as:
SELECT 100:::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path; to_cost | target id | via_path

```
--------------------+----------
(1 row)
```

The back end code has that same restriction as follows

```
SELECT 1 AS id, 100::float AS cost, 25::INTEGER AS target_id, ARRAY[33, 32, 25] AS path;
id|cost|target_id| path
1| 100 | 25|{33,32,25}
(1 row)
```

therefore the shortest path expected are as if there was no restriction involved

## 1 Different ways to represent 'no path found

- Sometimes represents with EMPTY SET a no path found
- Sometimes represents with Error a no path found


## Returning EMPTY SET to represent no path found

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    7,4, true, true
);
seq | id1 | id2 | cost
(0 rows)
```

pgr_trsp calls pgr_dijkstra when there are no restrictions which returnsEMPTY SET when a path is not found

```
SELECT * FROM pgr_dijkstra(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    7,4
);
seq | path_seq | node | edge | cost | agg_cost
(0 rows)
```


## Throwing EXCEPTION to represent no path found

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    7, 4, true, true,
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
ERROR: Error computing path: Path Not Found
CONTEXT: PL/pgSQL function pgr_trsp(text,integer,integer,boolean,boolean,text) line 53 at RAISE
```

pgr_trsp use the original code when there are restrictions, even if they have nothing to do with the graph, which will throw an EXCEPTION to represent no path found.

## 1 Routing from/to same location

When routing from location $\backslash(1 \backslash)$ to the same location $\backslash(1)$, no path is needed to reach the destination, its already there. Therefore is expected to return an EMPTY SET or an EXCEPTION depending on the parameters

- Sometimes represents with EMPTY SET no path found (expected)
- Sometimes represents with EXCEPTION no path found (expected)
- Sometimes finds a path (not expected)


## Returning expected EMPTY SET to represent no path found

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    7,7, true, true
);
seq | id1 | id2 | cost
----+----
```

pgr_trsp calls pgr_dijkstra when there are no restrictions which returns the expected to returnEMPTY SET to represent no path found.

## Returning expected EXCEPTION to represent no path found

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    2, 2, true, true,
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
ERROR: Error computing path: Path Not Found
CONTEXT: PL/pgSQL function pgr_trsp(text,integer,integer,boolean,boolean,text) line 53 at RAISE
```

In this case pgr_trsp calls the original code when there are restrictions, even if they have nothing to do with the graph, in this case that code throws the expected EXCEPTION

## Returning unexpected path

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    5, 5, true, true,
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
seq | id1 | id2 | cost
    0| 5| 1| 1
    1| 6| 4| 1
    2| 7| 8| 1
    3| 11| 9| 1
    4| 16| 16| 1
    5| 15| 3| 1
    6| 10| 2| 1
    7| 6| 1| 1
    8| 5|-1| 0
(9 rows)
```

In this case pgr_trsp calls the original code when there are restrictions, even if they have nothing to do with the graph, in this case that code finds an unexpected path.

## 1 User contradictions

pgr_trsp unlike other pgRouting functions does not autodectect the existence ofreverse_cost column. Therefor it has has_rcost parameter to check the existence of reverse_cost column. Contradictions happen:

- When the reverse_cost is missing, and the flaghas_rcost is set to true
- When the reverse_cost exists, and the flaghas_rcost is set to false


## When the reverse_cost is missing, and the flag has_rcost is set to true.

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost FROM edges$$,
    6, 10, false, true
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
ERROR: Error, reverse_cost is used, but query did't return 'reverse_cost' column
CONTEXT: PL/pgSQL function pgr_trsp(text,integer,integer,boolean,boolean,text) line 24 at RAISE
```

An EXCEPTION is thrown.

## When the reverse_cost exists, and the flag has_rcost is set to false

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    6, 10, false, false,
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
seq | id1 | id2 | cost
0| 6| 4| 1
    1| 7| 8| 1
    2| 11| 5| 1
    3| 10|-1| 0
(4 rows)
```

The reverse_cost column will be effectively removed and will cost execution time

```
The "Edges" signature version
```

pgr_trsp(sql text, source_edge integer, source_pos float8,
target_edge integer, target_pos float8,
directed boolean, has_rcost boolean [,restrict_sql text]);

## 2 Different ways to represent 'no path found

- Sometimes represents with EMPTY SET a no path found
- Sometimes represents with EXCEPTION a no path found

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    1, 0.5, 17, 0.5, true, true
);
seq | id1 | id2 | cost
(0 rows)
```

pgr_trsp calls pgr_withPoints - Proposed when there are no restrictions which returnsEMPTY SET when a path is not found

## Throwing EXCEPTION to represent no path found

```
SELECT * FROM _pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    1, 0.5, 17, 0.5, true, true,
    $$SELECT 100::float AS to cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
ERROR: Error computing path: Path Not Found
```

pgr_trsp use the original code when there are restrictions, even if they have nothing to do with the graph, which will throw an EXCEPTION to represent no path found.

## Paths with equal number of vertices and edges

A path is made of $N$ vertices and $N-1$ edges.

- Sometimes returns $N$ vertices and N-1 edges.
- Sometimes returns N-1 vertices and N-1 edges.


## Returning $\boldsymbol{N}$ vertices and $\boldsymbol{N - 1}$ edges.

```
SELECT * FROM pgr_TRSP(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    1, 0.5, 1, 0.8, true, true
);
seq | id1 | id2 | cost
    0|-1| 1| 0.3
    1|-2|-1| 0
(2 rows)
```

pgr_trsp calls pgr_withPoints - Proposed when there are no restrictions which returns the correct number of rows that will include all the vertices. The last row will have $\mathrm{a}-1$ on the edge column to indicate the edge number is invalidu for that row.

## Returning $\boldsymbol{N}-1$ vertices and $\boldsymbol{N}-1$ edges.

```
SELECT * FROM pgr_TRSP(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    1, 0.5, 1, 0.8, true, true,
    $$SELECT 100::float AS to cost, 25::INTEGER AS target id, '32, 33'::TEXT AS via path$$
);
seq | id1 | id2 | cost
----+----------+-----
(1 row)
```

pgr_trsp use the original code when there are restrictions, even if they have nothing to do with the graph, and will not return the last vertex of the path.

## 2 Routing from/to same location

When routing from the same edge and position to the same edge and position, no path is needed to reach the destination, its already there. Therefore is expected to return an EMPTY SET or an EXCEPTION depending on the parameters, non of which is happening.

## A path with 2 vertices and edge cost 0

```
SELECT * FROM pgr_TRSP(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    1,0.5, 1, 0.5, true, true
);
seq | id1 | id2 | cost
    0| -1| 1| 0
    1| -2| -1 | 0
(2 rows)
```

pgr_trsp calls pgr_withPoints - Proposed setting the first<br>((edge, position)<br>) with a differenct point id from the second $\backslash$ ((edge, position)<br>) making them different points. But the cost using the edge, is<br>(0)).

## A path with 1 vertices and edge cost 0

```
SELECT * FROM pgr_TRSP(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    1,0.5, 1, 0.5, true, true
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
seq | id1 | id2 | cost
---+------------
(1 row)
```

pgr_trsp use the original code when there are restrictions, even if they have nothing to do with the graph, and will not have the row for the vertex $\backslash(-2 \backslash)$.

## 2 User contradictions

pgr_trsp unlike other pgRouting functions does not autodectect the existence ofreverse_cost column. Therefor it has has_rcost parameter to check the existence of reverse_cost column. Contradictions happen:

- When the reverse_cost is missing, and the flaghas_rcost is set to true
- When the reverse_cost exists, and the flaghas_rcost is set to false


## When the reverse_cost is missing, and the flag has_rcost is set to true.

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost FROM edges$$,
    1, 0.5, 1, 0.8, false, true,
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
ERROR: Error, reverse_cost is used, but query did't return 'reverse_cost' column
CONTEXT: PL/pgSQL function pgr_trsp(text,integer,double precision,integer,double precision,boolean,boolean,text) line 36 at RAISE
```

An EXCEPTION is thrown.

## When the reverse_cost exists, and the flag has_rcost is set to false

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse cost FROM edges$$
    1,0.5,1,0.8, false, false
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
seq | id1 | id2 | cost
--------+----------
    0|-1| 1| 0.3
(1 row)
```

The reverse_cost column will be effectively removed and will cost execution time

## Using a points of interest table

Given a set of points of interest:

```
SELECT * FROM pointsOfInterest;
pid | x y | edge_id | side | fraction | geom | newpoint
    1|1.8|0.4| 1|| | 0.4|0101000000CDCCCCCCCCCCFC3F9A9999999999D93F | 010100000000000000000000409A9999999999D93F
    2|4.2|2.4| 15|r | 0.4|0101000000CDCCCCCCCCCC10403333333333330340|010100000000000000000010403333333333330340
    3|2.6|3.2| 12|| | 0.6|0101000000CDCCCCCCCCCC04409A99999999990940|0101000000CDCCCCCCCCCC04400000000000000840
    4|0.3|1.8| 6|r | 0.3|0101000000333333333333D33FCDCCCCCCCCCCFC3F|0101000000333333333333D33F0000000000000040
    5|2.9|1.8| 5|| | 0.8|01010000003333333333330740CDCCCCCCCCCCFC3F|01010000000000000000000840CDCCCCCCCCCCFC3F
    6|2.2|1.7 | 4 | b | 0.7|01010000009A99999999990140333333333333FB3F | 01010000000000000000000040333333333333FB3F
(6 rows)
```


## Using pgr_trsp

```
SELECT * FROM pgr_TRSP(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    (SELECT edge_id::INTEGER FROM pointsOflnterest WHERE pid = 1),
    (SELECT fraction FROM pointsOflnterest WHERE pid = 1)
    (SELECT edge_id::INTEGER FROM pointsOfInterest WHERE pid = 6),
    (SELECT fraction FROM pointsOflnterest WHERE pid =6),
    true, true,
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
seq | id1 | id2 | cost
    0| -1 | 1 1 0.6
    1| 6| 4| 0.7
(2 rows)
```

On pgr_trsp, to be able to use the table information:

- Each parameter has to be extracted explicitly from the table
- Regardles of the point pid original value
- will always be -1 for the first point
- will always be -2 for the second point
- the row reaching point -2 will not be shown


## Using pgr_withPoints - Proposed

```
SELECT * FROM pgr withPoints(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    $$SELECT pid, edge_id, fraction FROM pointsOflnterest$$,
    -1, -6
);
seq | path_seq | node | edge | cost | agg_cost
    1| -1| 1| 0.6| 0
    2| 2| 6| 4| 0.7| 0.6
3| 3| -6| -1| 0| 1.3
(3 rows)
```

Suggestion: use pgr_withPoints - Proposed when there are no turn restrictions:

- Results are more complete
- Column names are meaningful


## Routing from a vertex to a point

Solving a shortest path from vertex $\backslash(6 \backslash)$ to pid 1 using a points of interest table

## Using pgr_trsp

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    8, 1,
    (SELECT edge_id::INTEGER FROM pointsOfInterest WHERE pid = 1),
    (SELECT fraction FROM pointsOflnterest WHERE pid = 1),
    true, true,
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
seq | id1 | id2 | cost
----+---+----+----
    1| 7 4 4 1
    2| 6| 1| 0.6
(3 rows)
```

```
SELECT * FROM pgr_trsp(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse cost FROM edges$$,
    11,0,
    (SELECT edge_id::INTEGER FROM pointsOfInterest WHERE pid = 1),
    (SELECT fraction FROM pointsOfInterest WHERE pid = 1),
    true, true,
    $$SELECT 100::float AS to_cost, 25::INTEGER AS target_id, '32, 33'::TEXT AS via_path$$
);
seq | id1 | id2 | cost
--------------------
    0| 11| 8| 1
    1| 7| 4| 1
    2| 6| 1| 0.6
(3 rows)
```

- Vertex 6 is also edge 11 at 0 fraction


## Using pgr_withPoints - Proposed

```
SELECT * FROM pgr_withPoints(
    $$SELECT id::INTEGER, source::INTEGER, target::INTEGER, cost, reverse_cost FROM edges$$,
    $$SELECT pid, edge id, fraction FROM pointsOflnterest$$,
    11, -
);
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{rrrrrr}
\(1 \mid\) & \(1 \mid\) & \(11 \mid\) & \(8 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(7 \mid\) & \(4 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(3 \mid\) & \(6 \mid\) & \(1 \mid\) & \(0.6 \mid\) & 2 \\
\(4 \mid\) & \(4 \mid\) & \(-1 \mid\) & \(-1 \mid\) & \(0 \mid\) & 2.6
\end{tabular}
(4 rows)
```

Suggestion: use pgr_withPoints - Proposed when there are no turn restrictions:

- No need to choose where the vertex is located.
- Results are more complete
- Column names are meaningful


## prototypes

pgr_trspViaVertices and pgr_trspViaEdges were added to pgRouting as prototypes

These functions use the pgr_trsp functions inheriting all the problems mentioned above. When there are no restrictions and have a routing "via" problem with vertices:

- pgr_dijkstraVia - Proposed

See Also

## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: 2.5 2.4 2.6


## Cost - Category

- pgr_aStarCost
- pgr_bdAstarCost
- pgr_dijkstraCost
- pgr_bdDijkstraCost
- pgr_dijkstraNearCost - Proposed


## Proposed

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.
- pgr_withPointsCost - Proposed

General Information

## Characteristics

Each function works as part of the family it belongs to.
The main Characteristics are:

- It does not return a path.
- Returns the sum of the costs of the shortest path of each pair combination of nodes requested.
- Let be the case the values returned are stored in a table, so the unique index would be the pair((start_vid, end_vid).
- Depending on the function and its parameters, the results can be symmetric.
- The aggregate cost of $\backslash((u, v) \backslash)$ is the same as for $\backslash((v, u) \backslash)$.
- Any duplicated value in the start or end vertex identifiers are ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending

See Also

## Indices and tables

## Index

- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: $2.6 \mathbf{2 . 5} 2.4$


## Cost Matrix - Category

```
pgr_aStarCostMatrix
```

- pgr_bdAstarCostMatrix
- pgr_bdDijkstraCostMatrix
- pgr_dijkstraCostMatrix
- pgr_bdDijkstraCostMatrix


## proposed

## Warning

Proposed functions for next mayor release.
They are not officially in the current release.

- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.
- pgr_withPointsCostMatrix - proposed


## General Information

## Synopsis

Traveling Sales Person - Family of functions needs as input a symmetric cost matrix and no edge( $u$, $v$ ) must value $\backslash$ ( Iinfty<br>).

This collection of functions will return a cost matrix in form of a table.

## Characteristics

The main Characteristics are:

- Can be used as input to pgr_TSP.
- Use directly when the resulting matrix is symmetric and there is nd(\infty<br>) value.
- It will be the users responsibility to make the matrix symmetric.
- By using geometric or harmonic average of the non symmetric values.
- By using max or min the non symmetric values.
- By setting the upper triangle to be the mirror image of the lower triangle.
- By setting the lower triangle to be the mirror image of the upper triangle.
- It is also the users responsibility to fix an<br>(\infty<br>) value.
- Each function works as part of the family it belongs to.
- It does not return a path.
- Returns the sum of the costs of the shortest path for pair combination of nodes in the graph.
- Process is done only on edges with positive costs.
- Values are returned when there is a path.
- When the starting vertex and ending vertex are the same, there is no path.
- The aggregate cost in the non included values( $v, v$ ) is 0 .
- When the starting vertex and ending vertex are the different and there is no path.
- The aggregate cost in the non included values $(u, v)$ is $\backslash(\backslash i n f t y \backslash)$.
- Let be the case the values returned are stored in a table:
- The unique index would be the pair: (start_vid, end_vid).
- Depending on the function and its parameters, the results can be symmetric.
- The aggregate cost of $(u, v)$ is the same as for $(v, u)$.
- Any duplicated value in the start vids are ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending


## Parameters

## Used in:

- pgr_aStarCostMatrix
- pgr_dijkstraCostMatrix

| Column | Type | Description |
| :--- | :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting <br> vertices. |

## Used in:

- pgr_withPointsCostMatrix - proposed

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Points SQL | TEXT | Points SQL as described below |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting <br> vertices. |

Optional parameters

| Column | Type | Default | Description |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| directed | BOOLEAN | true |  | When true the graph is considered Directed |  |
|  |  |  |  | Whenfalse the graph is considered Undirected. | as |

Inner Queries

Edges SQL

## Used in:

pgr_dijkstraCostMatrix

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |  |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |  |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |  |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |  |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |
|  |  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  | not part of the graph. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Points SQL

| Parameter | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| pid | ANY-INTEGER | value | Identifier of the point. |

- Use with positive value, as internally will be converted to negative value
- If column is present, it can not be NULL.
- If column is not present, a sequential negativevalue will be given automatically.

| edge_id | ANY-INTEGER |  | Identifier of the "closest" edge to the point. |
| :---: | :---: | :---: | :---: |
| fraction | ANY-NUMERICAL |  | Value in $<0,1\rangle$ that indicates the relative postition from the first end point of the edge. |
| side | CHAR | b | Value in [b, r, l, NULL] indicating if the point is: |
|  |  |  | - In the right $r$, <br> - In the left I, <br> - In both sides b, NULL |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Result Columns

Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

See Also

- Traveling Sales Person - Family of functions


## Indices and tables

- Index
- Search Page

[^0]pgr_drivingDistance - Driving Distance based on Dijkstra's algorithm
pgr_primDD - Driving Distance based on Prim's algorithm
pgr_kruskalDD - Driving Distance based on Kruskal's algorithm
Post pocessing

- pgr_alphaShape - Alpha shape computation


## Proposed

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.
- pgr_withPointsDD - Proposed - Driving Distance based on pgr_withPoints

Supported versions: Latest (3.3) 3.23 .13 .0

- Unsupported versions: 2.62 .52 .42 .3 2.2 2.1 2.0
pgr_alphaShape
pgr_alphaShape - Polygon part of an alpha shape.


## Availability

- Version 3.0.0
- Breaking change on signature
- Old signature no longer supported
- Boost 1.54 \& Boost 1.55 are supported
- Boost 1.56+ is preferable - Boost Geometry is stable on Boost 1.56
- Version 2.1.0
- Added alpha argument with default 0 (use optimal value)
- Support to return multiple outer/inner ring
- Version 2.0.0
- Official function
- Renamed from version 1.x


## Support

## Description

Returns the polygon part of an alpha shape.

## Characteristics

- Input is a geometry and returns a geometry
- Uses PostGis ST_DelaunyTriangles
- Instead of using CGAL's definition of alpha it use the spoon_radius
- $\backslash($ spoon $\$ radius $=$ \sqrt alpha<br>)
- A Triangle area is considered part of the alpha shape when $\backslash($ circumcenter $\backslash$ radius $<$ spoon $\$ radius $\backslash$ )
- The alpha parameter is the spoon radius
- When the total number of points is less than 3, returns an EMPTY geometry


## Signatures

## Summary

## Example:

passing a geometry collection with spoon radius $\backslash(1.5 \backslash)$ using the return variable geom
SELECT ST_Area(pgr_alphaShape((SELECT ST_Collect(geom) FROM vertices), 1.5));
st_area
9.75
(1 row)

## Parameters

| Parameter | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- |
| geometry | geometry | Geometry with at least $\backslash(3 \backslash)$ <br> points |  |
| alpha | FLOAT | 0 | The radius of the spoon. |

## Return Value

| Kind of geometry | Description |  |  |
| :--- | :--- | :--- | :--- |
| GEOMETRY | A Geometry collection of |  |  |
| COLLECTION | Polygons |  |  |

## See Also

- pgr_drivingDistance
- Sample Data network.
- ST_ConcaveHull


## Indices and tables

- Index
- Search Page

Calculate nodes that are within a distance.

- Extracts all the nodes that have costs less than or equal to the value distance.
- The edges extracted will conform to the corresponding spanning tree.
- Edge $\backslash((u, v) \backslash)$ will not be included when:
- The distance from the root to $\backslash(u \backslash)>$ limit distance.
- The distance from the root to $\backslash(v \backslash)>$ limit distance.
- No new nodes are created on the graph, so when is within the limit and is not within the limit, the edge is not included.


## Parameters

| Parameter | Type | Description |
| :---: | :---: | :---: |
| Edges SQL | TEXT | Edges SQL as described below. |
| Root vid | BIGINT | Identifier of the root vertex of the tree. |
| Root vids | ARRAY[ANY-INTEGER] | Array of identifiers of the root vertices. $\backslash(0 \backslash)$ values are ignored For optimization purposes, any duplicated value is ignored. |
| distance | FLOAT | Upper limit for the inclusion of a node in the result. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Inner Queries

Edges SQL

| Column | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER |  | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |
|  |  |  | - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

See Also

## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3)

BFS - Category

- pgr_kruskalBFS
- pgr_primBFS

Traversal using breadth first search.

- It's implementation is only on undirected graph.
- Process is done only on edges with positive costs.
- When the graph is connected
- The resulting edges make up a tree
- When the graph is not connected,
- Finds a minimum spanning tree for each connected component.
- The resulting edges make up a forest.


## Parameters

| Parameter | Type | Description |
| :---: | :---: | :---: |
| Edges SQL | TEXT | Edges SQL as described below. |
| root vid | BIGINT | Identifier of the root vertex of the tree. <br> - When value is $\backslash(0 \backslash)$ then gets the spanning forest starting in aleatory nodes for each tree in the forest. |
| root vids | ARRAY [ ANY-INTEGER ] | Array of identifiers of the root vertices. <br> - $\backslash(0 \backslash)$ values are ignored <br> - For optimization purposes, any duplicated value is ignored. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

BFS optional parameters
Parameter Type Default Description
max_depth BIGINT $\backslash(9223372036854775807 \backslash)$ Upper limit of the depth of the tree.

- When negative throws an error.

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL -1 | Weight of the edge (target, source) |
|  |  |  |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Result Columns

Returns SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

| Parameter | Type | Description |
| :---: | :---: | :---: |
| seq | BIGINT | Sequential value starting from $\backslash(1 \backslash)$. |
| depth | BIGINT | Depth of the node. |
|  |  | - $\backslash(0 \backslash)$ when node = start_vid. |
| start_vid | BIGINT | Identifier of the root vertex. |
| node | BIGINT | Identifier of node reached using edge. |
| edge | BIGINT | Identifier of the edge used to arrive to node. |
|  |  | - $\backslash(-1 \backslash)$ when node $=$ start_vid. |
| cost | FLOAT | Cost to traverse edge. |
| agg_cost | FLOAT | Aggregate cost from start_vid to node. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

See Also

- Boost: Prim's algorithm
- Boost: Kruskal's algorithm
- Wikipedia: Prim's algorithm
- Wikipedia: Kruskal's algorithm


## Indices and tables

- Index
- Search Page


## See Also

## Indices and tables

- Index
- Search Page

All Pairs - Family of Functions

- pgr_floydWarshall - Floyd-Warshall's algorithm.
- pgrjohnson - Johnson's algorithm


## A* - Family of functions

- pgr_aStar - A* algorithm for the shortest path.
- pgr_aStarCost - Get the aggregate cost of the shortest paths.
- pgr_aStarCostMatrix - Get the cost matrix of the shortest paths.


## Bidirectional A* - Family of functions

- pgr_bdAstar - Bidirectional A* algorithm for obtaining paths.
- pgr_bdAstarCost - Bidirectional A* algorithm to calculate the cost of the paths.
- pgr_bdAstarCostMatrix - Bidirectional A* algorithm to calculate a cost matrix of paths.


## Bidirectional Dijkstra - Family of functions

- pgr_bdDijkstra - Bidirectional Dijkstra algorithm for the shortest paths.
- pgr_bdDijkstraCost - Bidirectional Dijkstra to calculate the cost of the shortest paths
- pgr_bdDijkstraCostMatrix - Bidirectional Dijkstra algorithm to create a matrix of costs of the shortest paths.


## Components - Family of functions

- pgr_connectedComponents - Connected components of an undirected graph.
- pgr_strongComponents - Strongly connected components of a directed graph.
- pgr_biconnectedComponents - Biconnected components of an undirected graph.
- pgr_articulationPoints - Articulation points of an undirected graph.
- pgr_bridges - Bridges of an undirected graph.


## Contraction - Family of functions

- pgr_contraction


## Dijkstra - Family of functions

- pgr_dijkstra - Dijkstra's algorithm for the shortest paths.
- pgr_dijkstraCost - Get the aggregate cost of the shortest paths.
- pgr_dijkstraCostMatrix - Use pgr_dijkstra to create a costs matrix.
- pgr_drivingDistance - Use pgr_dijkstra to calculate catchament information.
- pgr_KSP - Use Yen algorithm with pgr_dijkstra to get the K shortest paths.


## Flow - Family of functions

- pgr_maxFlow - Only the Max flow calculation using Push and Relabel algorithm.
pgr_boykovKolmogorov - Boykov and Kolmogorov with details of flow on edges.
- pgr_edmondsKarp - Edmonds and Karp algorithm with details of flow on edges.
- pgr_pushRelabel - Push and relabel algorithm with details of flow on edges.
- Applications
- pgr_edgeDisjointPaths - Calculates edge disjoint paths between two groups of vertices.
- pgr_maxCardinalityMatch - Calculates a maximum cardinality matching in a graph.


## Kruskal - Family of functions

- pgr_kruskal
pgr_kruskaIBFS
- pgr_kruskaIDD
- pgr_kruskaIDFS


## Prim - Family of functions

```
pgr_prim
```

pgr_primBFS

- pgr_primDD
- pgr_primDFS


## Reference

pgr_version

- pgr_full_version


## Topology - Family of Functions

The following functions modify the database directly therefore the user must have special permissions given by the administrators to use them.
pgr_createTopology - create a topology based on the geometry.
pgr_createVerticesTable - reconstruct the vertices table based on the source and target information.

- pgr_analyzeGraph - to analyze the edges and vertices of the edge table.
- pgr_analyzeOneWay - to analyze directionality of the edges.
- pgr_nodeNetwork -to create nodes to a not noded edge table.

Traveling Sales Person - Family of functions

- pgr_TSP - When input is given as matrix cell information.
- pgr_TSPeuclidean - When input are coordinates.
pgr_trsp - Turn Restriction Shortest Path (TRSP) - Turn Restriction Shortest Path (TRSP)


## Functions by categories

## Cost - Category

- pgr_aStarCost
- pgr_bdAstarCost
- pgr_dijkstraCost
- pgr_bdDijkstraCost
- pgr_dijkstraNearCost - Proposed


## Cost Matrix - Category

- pgr_aStarCostMatrix
- pgr_bdAstarCostMatrix
- pgr_bdDijkstraCostMatrix
- pgr_dijkstraCostMatrix
- pgr_bdDijkstraCostMatrix

Driving Distance - Category

- pgr_drivingDistance - Driving Distance based on Dijkstra's algorithm
- pgr_primDD - Driving Distance based on Prim's algorithm
- pgr_kruskalDD - Driving Distance based on Kruskal's algorithm
- Post pocessing
- pgr_alphaShape - Alpha shape computation

K shortest paths - Category

- pgr_KSP - Yen's algorithm based on pgr_dijkstra

Spanning Tree - Category

- Kruskal - Family of functions
- Prim - Family of functions

```
BFS - Category
```

- pgr_kruskalBFS
- pgr_primBFS

DFS - Category

- pgr_kruskaIDFS
- pgr_primDFS

Available Functions but not official pgRouting functions

- Proposed Functions
- Experimental Functions
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: $2.6 \mathbf{2}$.5 2.42 .32 .2


## Proposed Functions

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)


## Families

## Dijkstra - Family of functions

- pgr_dijkstraVia - Proposed - Get a route of a seuence of vertices.
- pgr_dijkstraNear - Proposed - Get the route to the nearest vertex.
- pgr_dijkstraNearCost - Proposed - Get the cost to the nearest vertex.


## withPoints - Family of functions

- pgr_withPoints - Proposed - Route from/to points anywhere on the graph.
- pgr_withPointsCost - Proposed - Costs of the shortest paths.
- pgr_withPointsCostMatrix - proposed - Costs of the shortest paths.
- pgr_withPointsKSP - Proposed - K shortest paths.
- pgr_withPointsDD - Proposed - Driving distance.


## Topology - Family of Functions

These proposed functions do not modify the database.

- pgr_extractVertices - Proposed - Extracts vertex information based on the edge table information.


## Coloring - Family of functions

- pgr_sequentialVertexColoring - Proposed - Vertex coloring algorithm using greedy approach.

Traversal - Family of functions

- pgr_depthFirstSearch - Proposed - Depth first search traversal of the graph.


## - Supported versions: Latest (3.3) 3.2

Traversal - Family of functions

## Proposed

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.
- pgr_depthFirstSearch - Proposed - Depth first search traversal of the graph.


## Experimental

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting
pgr_breadthFirstSearch - Experimental - Breath first search traversal of the graph.
- pgr_binaryBreadthFirstSearch - Experimental - Breath first search traversal of the graph.

Aditionaly there are 2 categories under this family

- BFS - Category
- DFS - Category
- Supported versions: Latest (3.3) 3.2
pgr_depthFirstSearch - Proposed
pgr_depthFirstSearch - Returns a depth first search traversal of the graph. The graph can be directed or undirected.


## boost

Boost Graph Inside

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.


## Availability

- Version 3.3.0
- Promoted to proposed function
- Version 3.2.0
- New experimental signatures:
- pgr_depthFirstSearch (Single Vertex)
- pgr_depthFirstSearch (Multiple Vertices)


## Description

Depth First Search algorithm is a traversal algorithm which starts from a root vertex, goes as deep as possible, and backtracks once a vertex is reached with no adjacent vertices or with all visited adjacent vertices. The traversal continues until all the vertices reachable from the root vertex are visited.

## The main Characteristics are:

- The implementation works for both directed and undirected graphs.
- Provides the Depth First Search traversal order from a root vertex or from a set of root vertices.
- An optional non-negative maximum depth parameter to limit the results up to a particular depth.
- For optimization purposes, any duplicated values in the Root vids are ignored.
- It does not produce the shortest path from a root vertex to a target vertex.
- The aggregate cost of traversal is not guaranteed to be minimal.
- The returned values are ordered in ascending order ofstart_vid.
- Depth First Search Running time: $\backslash(O(E+V) \backslash)$

Signatures

## Summary

```
pgr_depthFirstSearch(Edges SQL, root vid, [options])
pgr_depthFirstSearch(Edges SQL, root vids, [options])
options: [directed, max_depth]
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)
```


## Single vertex

```
pgr_depthFirstSearch(Edges SQL, root vid, [options])
options: [directed, max_depth]
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)
```


## Example:

From root vertex $\backslash(6 \backslash)$ on a directed graph with edges in ascending order of id

```
SELECT * FROM pgr_depthFirstSearch(
    SELECT id, source, target, cost, reverse_cost FROM edges
    ORDER BY id',
6);
seq | depth | start_vid | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|}
\hline 01 & 61 & 61 & -1| & 01 & 0 \\
\hline 1| & 61 & 51 & 1| & 1| & 1 \\
\hline 1| & 61 & 71 & 4 | & 1 | & 1 \\
\hline \(2 \mid\) & 61 & 31 & 7| & 1 | & 2 \\
\hline 31 & 61 & 1| & 61 & 1 | & 3 \\
\hline \(2 \mid\) & 61 & 11| & 81 & 1| & 2 \\
\hline 31 & 61 & 16| & 9| & 1| & 3 \\
\hline 4| & 61 & 17| & 15| & 1 | & 4 \\
\hline 4 | & 61 & 15| & 16| & 1 | & 4 \\
\hline 51 & 6 & 10 & 31 & 1 | & 5 \\
\hline 31 & 6 & 12 & 11| & 1 | & 3 \\
\hline 21 & 6 & 8 & 10| & 1 | & 2 \\
\hline 31 & 6 & 91 & 14| & 1 | & 3 \\
\hline s) & & & & & \\
\hline
\end{tabular}
```

Multiple vertices
pgr_depthFirstSearch(Edges SQL, root vids, [options])
options: [directed, max_depth]
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

## Example:

From root vertices $\backslash(\backslash\{12,6 \backslash\} \backslash)$ on an undirected graph with depth $\backslash(<=2 \backslash)$ and edges in ascending order ofid

```
SELECT * FROM pgr_depthFirstSearch(
    'SELECT id, source, target, cost, reverse_cost FROM edges
    ORDER BY id',
    ARRAY[12, 6], directed => false, max depth => 2);
seq | depth | start_vid | node | edge | cost | agg_cost
\begin{tabular}{cccccc}
\(1 \mid\) & \(0 \mid\) & \(6 \mid\) & \(6 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(2 \mid\) & \(1 \mid\) & \(6 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(1 \mid\) & \(6 \mid\) & \(10 \mid\) & \(2 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(2 \mid\) & \(6 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) \\
\(5 \mid\) & \(2 \mid\) & \(6 \mid\) & \(11 \mid\) & \(5 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(1 \mid\) & \(6 \mid\) & \(7 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(7 \mid\) & \(2 \mid\) & \(6 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) \\
\hline \(8 \mid\) & \(2 \mid\) & \(6 \mid\) & \(8 \mid\) & \(10 \mid\) & \(1 \mid\) \\
\(9 \mid\) & \(0 \mid\) & \(12 \mid\) & \(12 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(10 \mid\) & \(1 \mid\) & \(12 \mid\) & \(11 \mid\) & \(11 \mid\) & \(1 \mid\) \\
\(11 \mid\) & \(2 \mid\) & \(12 \mid\) & \(10 \mid\) & \(5 \mid\) & \(1 \mid\) \\
\(12 \mid\) & \(2 \mid\) & \(12 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(13 \mid\) & \(2 \mid\) & \(12 \mid\) & \(16 \mid\) & \(9 \mid\) & \(1 \mid\) \\
\(14 \mid\) & \(1 \mid\) & \(12 \mid\) & \(8 \mid\) & \(12 \mid\) & \(1 \mid\) \\
\(15 \mid\) & \(2 \mid\) & \(12 \mid\) & \(9 \mid\) & \(14 \mid\) & \(1 \mid\) \\
\(16 \mid\) & \(1 \mid\) & \(12 \mid\) & \(17 \mid\) & \(13 \mid\) & \(1 \mid\) \\
12
\end{tabular}
(16 rows)
```

Parameters

| Parameter | Type | Description |
| :---: | :---: | :---: |
| Edges SQL | TEXT | Edges SQL as described below. |
| root vid | BIGINT | Identifier of the root vertex of the tree. <br> - When value is $\backslash(0 \backslash)$ then gets the spanning forest starting in aleatory nodes for each tree in the forest. |
| root vids | ARRAY [ ANY-INTEGER ] | Array of identifiers of the root vertices. <br> - $\backslash(0 \backslash)$ values are ignored <br> - For optimization purposes, any duplicated value is ignored. |

## Where:

ANY-INTEGER:
SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

## Optional parameters



DFS optional parameters

| Parameter | Type Default | Description |
| :--- | :--- | :--- |
| max_depth | BIGINT $\backslash(9223372036854775807 \backslash)$ | Upper limit of the depth of the tree. |
|  |  | When negative throws an <br> error. |

## Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL -1 | Weight of the edge (target, source) |
|  |  |  |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |
|  |  |  |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Return column

Returns SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

| Parameter | Type | Description |
| :--- | :--- | :--- |
| seq | BIGINT | Sequential value starting from $\backslash(1 \backslash)$. |
| depth | BIGINT | Depth of the node. |
|  |  | (1) |
| start_vid | BIGINT | Identifier of the root vertex. |
| node | BIGINT | Identifier of node reached using edge. |
| edge | BIGINT | Identifier of the edge used to arrive to <br> node. |
|  |  | FLOAT |
| cost | Cost to traverse edge. |  |
| agg_cost | FLOAT | Aggregate cost from start_vid to node. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

## Additional Examples

## Example:

Same as Single vertex but with edges in descending order ofid

```
SELECT * FROM pgr_depthFirstSearch(
    'SELECT id, source, target, cost, reverse_cost FROM edges
    ORDER BY id DESC',
6);
seq | depth | start_vid | node | edge | cost | agg_cost
\begin{tabular}{cccccc}
\(1 \mid\) & \(0 \mid\) & \(6 \mid\) & \(6 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(2 \mid\) & \(1 \mid\) & \(6 \mid\) & \(7 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(2 \mid\) & \(6 \mid\) & \(8 \mid\) & \(10 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(3 \mid\) & \(6 \mid\) & \(9 \mid\) & \(14 \mid\) & \(1 \mid\) \\
\(5 \mid\) & \(3 \mid\) & \(6 \mid\) & \(12 \mid\) & \(12 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(4 \mid\) & \(6 \mid\) & \(17 \mid\) & \(13 \mid\) & \(1 \mid\) \\
\(7 \mid\) & \(5 \mid\) & \(6 \mid\) & \(16 \mid\) & \(15 \mid\) & \(1 \mid\) \\
\(8 \mid\) & 5 \\
\(8 \mid\) & \(6 \mid\) & \(15 \mid\) & \(16 \mid\) & \(1 \mid\) & 6 \\
\(9 \mid\) & \(7 \mid\) & \(6 \mid\) & \(10 \mid\) & \(3 \mid\) & \(1 \mid\) \\
\(10 \mid\) & \(8 \mid\) & \(6 \mid\) & \(11 \mid\) & \(5 \mid\) & \(1 \mid\) \\
\(11 \mid\) & \(2 \mid\) & \(6 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) \\
\(12 \mid\) & \(3 \mid\) & \(6 \mid\) & \(1 \mid\) & \(6 \mid\) & \(1 \mid\) \\
\(13 \mid\) & \(1 \mid\) & \(6 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) \\
13 \\
13
\end{tabular}
(13 rows)
```

The resulting traversal is different.
The left image shows the result with ascending order of ids and the right image shows with descending order of the edge identifiers.


See Also

- DFS - Category
- Sample Data
- Boost: Depth First Search algorithm documentation
- Boost: Undirected DFS algorithm documentation
- Wikipedia: Depth First Search algorithm


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
pgr_breadthFirstSearch - Experimental
pgr_breadthFirstSearch — Returns the traversal order(s) using Breadth First Search algorithm.


## Warning

Possible server crash

- These functions might create a server crash


## Warning

## Experimental functions

- They are not officially of the current release
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need c/c++ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.0.0
- New experimental signature:
- pgr_breadthFirstSearch (Single Vertex)
- pgr_breadthFirstSearch (Multiple Vertices)


## Description

Provides the Breadth First Search traversal order from a root vertex to a particular depth.

## The main Characteristics are:

- The implementation will work on any type of graph.
- Provides the Breadth First Search traversal order from a source node to a target depth level.
- Running time: $\backslash(O(E+V) \backslash)$


## Signatures

## Summary

```
pgr_breadthFirstSearch(Edges SQL, root vid, [options])
pgr_breadthFirstSearch(Edges SQL, root vids, [options])
options: [max_depth, directed]
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)
```


## Single vertex

```
pgr_breadthFirstSearch(Edges SQL, root vid, [options])
options: [max_depth, directed]
RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)
```


## Example:

From root vertex $\backslash(6 \backslash)$ on a directed graph with edges in ascending order of id

## SELECT * FROM pgr_breadthFirstSearch(

'SELECT id, source, target, cost, reverse_cost
FROM edges ORDER BY id',
6);
seq | depth | start_vid | node | edge | cost | agg_cost

| $1 \mid$ | $0 \mid$ | $6 \mid$ | $6 \mid$ | $-1 \mid$ | $0 \mid$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $1 \mid$ | $6 \mid$ | $5 \mid$ | $1 \mid$ | $1 \mid$ |
| $3 \mid$ | $1 \mid$ | $6 \mid$ | $7 \mid$ | $4 \mid$ | $1 \mid$ |
| $4 \mid$ | $2 \mid$ | $6 \mid$ | $3 \mid$ | $7 \mid$ | $1 \mid$ |
| $5 \mid$ | $2 \mid$ | $6 \mid$ | $11 \mid$ | $8 \mid$ | $1 \mid$ |
| $6 \mid$ | $2 \mid$ | $6 \mid$ | $8 \mid$ | $10 \mid$ | $1 \mid$ |
| $7 \mid$ | $3 \mid$ | $6 \mid$ | $1 \mid$ | $6 \mid$ | $1 \mid$ |
| $8 \mid$ | $3 \mid$ | $6 \mid$ | $16 \mid$ | $9 \mid$ | $1 \mid$ |
| $9 \mid$ | $3 \mid$ | $6 \mid$ | $12 \mid$ | $11 \mid$ | $1 \mid$ |
| $10 \mid$ | $3 \mid$ | $6 \mid$ | $9 \mid$ | $14 \mid$ | $1 \mid$ |
| 10 | 3 |  |  |  |  |
| $11 \mid$ | $4 \mid$ | $6 \mid$ | $17 \mid$ | $15 \mid$ | $1 \mid$ |
| $12 \mid$ | $4 \mid$ | $6 \mid$ | $15 \mid$ | $16 \mid$ | $1 \mid$ |
| $13 \mid$ | $5 \mid$ | $6 \mid$ | $10 \mid$ | $3 \mid$ | $1 \mid$ |
| 13 | 4 |  |  |  |  |

## Multiple vertices

pgr_breadthFirstSearch(Edges SQL, root vids, [options])
options: [max_depth, directed]

RETURNS SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

## Example:

From root vertices $\backslash(\backslash\{12,6 \backslash\} \backslash)$ on an undirected graph with depth $\backslash(<=2 \backslash)$ and edges in ascending order of id

```
SELECT * FROM pgr_breadthFirstSearch
    'SELECT id, source, target, cost, reverse_cost
    FROM edges ORDER BY id',
    ARRAY[12, 6], directed => false, max_depth => 2);
seq | depth | start_vid | node | edge | cost | agg_cost
\begin{tabular}{cccccc}
\(1 \mid\) & \(0 \mid\) & \(6 \mid\) & \(6 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(2 \mid\) & \(1 \mid\) & \(6 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(1 \mid\) & \(6 \mid\) & \(10 \mid\) & \(2 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(1 \mid\) & \(6 \mid\) & \(7 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(5 \mid\) & \(2 \mid\) & \(6 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(2 \mid\) & \(6 \mid\) & \(11 \mid\) & \(5 \mid\) & \(1 \mid\) \\
\(7 \mid\) & \(2 \mid\) & \(6 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) \\
\(8 \mid\) & \(2 \mid\) & \(6 \mid\) & \(8 \mid\) & \(10 \mid\) & \(1 \mid\) \\
\(9 \mid\) & \(0 \mid\) & \(12 \mid\) & \(12 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(10 \mid\) & \(1 \mid\) & \(12 \mid\) & \(11 \mid\) & \(11 \mid\) & \(1 \mid\) \\
\(11 \mid\) & \(1 \mid\) & \(12 \mid\) & \(8 \mid\) & \(12 \mid\) & \(1 \mid\) \\
\(12 \mid\) & \(1 \mid\) & \(12 \mid\) & \(17 \mid\) & \(13 \mid\) & \(1 \mid\) \\
\(13 \mid\) & \(2 \mid\) & \(12 \mid\) & \(10 \mid\) & \(5 \mid\) & \(1 \mid\) \\
\(14 \mid\) & \(2 \mid\) & \(12 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(15 \mid\) & \(2 \mid\) & \(12 \mid\) & \(16 \mid\) & \(9 \mid\) & \(1 \mid\) \\
\(16 \mid\) & \(2 \mid\) & \(12 \mid\) & \(9 \mid\) & \(14 \mid\) & \(1 \mid\) \\
12
\end{tabular}
```


## Parameter

| Parameter | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below. |
| root vid | BIGINT | Identifier of the root vertex of the tree. |

- When value is $\backslash(0 \backslash)$ then gets the spanning forest starting in aleatory nodes for each tree in the forest
root vids ARRAY [ ANY-INTEGER ] Array of identifiers of the root vertices.
- <br>(0<br>) values are ignored
- For optimization purposes, any duplicated value is ignored.


## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERIC:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true | When true the graph is considered Directed |  |
|  |  | When false the graph is considered as  <br>   <br>   <br>   <br>   |  |

DFS optional parameters

| Parameter | Type | Default |
| :--- | :--- | :--- |
| max_depth | BIGINT $\backslash(9223372036854775807 \backslash)$ | Description |
|  |  | Upper limit of the depth of the tree. <br>  <br>  |
|  |  |  |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL -1 | Weight of the edge (target, source) |
|  |  |  |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |
|  |  |  |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Return columns

Returns SET OF (seq, depth, start_vid, node, edge, cost, agg_cost)

| Parameter | Type | Description |
| :--- | :--- | :--- |
| seq | BIGINT | Sequential value starting from $\backslash(1 \backslash)$. |
| depth | BIGINT | Depth of the node. |
|  |  | $\bullet$ |
| start_vid | BIGINT | Identifier of the root vertex. |
| node | BIGINT | Identifier of node reached using edge. |
| edge | BIGINT | Identifier of the edge used to arrive to <br> node. |
|  |  | 0 |
|  | FLOAT | Cost to traverse edge. |
| cost | FLOAT | Aggregate cost from start_vid to node. |
| agg_cost |  |  |

Where:
ANY-INTEGER:
SMALLINT, INTEGER, BIGINT

## ANY-NUMERIC:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

## Additional Examples

## Example:

Same as Single vertex with edges in ascending order of id.

## SELECT * FROM pgr_breadthFirstSearch(

'SELECT id, source, target, cost, reverse_cost
FROM edges ORDER BY id'
6);
seq | depth | start_vid | node | edge | cost | agg_cost

| $1 \mid$ | $0 \mid$ | $6 \mid$ | $6 \mid$ | $-1 \mid$ | $0 \mid$ |
| ---: | ---: | ---: | ---: | ---: | ---: |
| $2 \mid$ | $1 \mid$ | $6 \mid$ | $5 \mid$ | $1 \mid$ | $1 \mid$ |
| $3 \mid$ | $1 \mid$ | $6 \mid$ | $7 \mid$ | $4 \mid$ | $1 \mid$ |
| $4 \mid$ | $2 \mid$ | $6 \mid$ | $3 \mid$ | $7 \mid$ | $1 \mid$ |
| $5 \mid$ | $2 \mid$ | $6 \mid$ | $11 \mid$ | $8 \mid$ | $1 \mid$ |
| $6 \mid$ | $2 \mid$ | $6 \mid$ | $8 \mid$ | $10 \mid$ | $1 \mid$ |
| $7 \mid$ | $3 \mid$ | $6 \mid$ | $1 \mid$ | $6 \mid$ | $1 \mid$ |
| $8 \mid$ | $3 \mid$ | $6 \mid$ | $16 \mid$ | $9 \mid$ | $1 \mid$ |
| $9 \mid$ | $3 \mid$ | $6 \mid$ | $12 \mid$ | $11 \mid$ | $1 \mid$ |
| $10 \mid$ | $3 \mid$ | $6 \mid$ | $9 \mid$ | $14 \mid$ | $1 \mid$ |
| 10 | 3 |  |  |  |  |
| $11 \mid$ | $4 \mid$ | $6 \mid$ | $17 \mid$ | $15 \mid$ | $1 \mid$ |
| $12 \mid$ | $4 \mid$ | $6 \mid$ | $15 \mid$ | $16 \mid$ | $1 \mid$ |
| $13 \mid$ | $5 \mid$ | $6 \mid$ | $10 \mid$ | $3 \mid$ | $1 \mid$ |
| $(13$ rows) |  |  |  | 5 |  |

## Example:

Same as Single vertex with edges in descending order ofid.

```
SELECT * FROM pgr_breadthFirstSearch(
    'SELECT id, source, target, cost, reverse_cost
    FROM edges ORDER BY id DESC',
6);
seq | depth | start_vid | node | edge | cost | agg_cost
    1| 0| 6| 6| -1| 0| 0
    1| 6| 7| 4| 1|
    1| 6| 5| 1| 1| 1
    2|
    2| 6| 3| 7| 1| 
    3| 6| 9| 14| 1| 3
    3| 6| 12| 12| 1| 3
    3| 6| 16| 9| 1| 3
    3| 6| 1| 6| 1| 3
    11| 4| 6| 17| 13| 1| 4
    12| 4| 6| 15| 16| 1| 4
    13| 5| 6| 10| 3| 1| 5
(13 rows)
```

The resulting traversal is different.
The left image shows the result with ascending order of ids and the right image shows with descending order of the edge identifiers.


## See Also

- BFS - Category
- Sample Data
- Boost: Breadth First Search algorithm documentation
- Wikipedia: Breadth First Search algorithm


## Indices and tables

```
- Index
- Search Page
```

Any graph whose edge-weights belongs to the set $\{0, X\}$, where ' $X$ ' is any non-negative integer, is termed as a 'binary graph'.

## boost

Boost Graph Inside


#### Abstract

Warning Possible server crash - These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need c/c++ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.2.0
- New experimental signature:
- pgr_binaryBreadthFirstSearch(Combinations)
- Version 3.0.0
- New experimental signatures:
- pgr_binaryBreadthFirstSearch(One to One)
- pgr_binaryBreadthFirstSearch(One to Many)
- pgr_binaryBreadthFirstSearch(Many to One)
- pgr_binaryBreadthFirstSearch(Many to Many)


## Description

It is well-known that the shortest paths between a single source and all other vertices can be found using Breadth First Search in $\backslash(O(|E|) \backslash)$ in an unweighted graph, i.e. the distance is the minimal number of edges that you need to traverse from the source to another vertex. We can interpret such a graph also as a weighted graph, where every edge has the weight $\backslash(1)$ ). If not alledges in graph have the same weight, that we need a more general algorithm, like Dijkstra's Algorithm which runs in \} ( $\mathrm{O}(|\mathrm{E}| \mathrm{log}|\mathrm{V}|)$ ) time.

However if the weights are more constrained, we can use a faster algorithm. This algorithm, termed as 'Binary Breadth First Search' as well as '0-1 BFS', is a variation of the standard Breadth First Search problem to solve the SSSP (single-source shortest path) problem in $\backslash(O(|E|) \backslash)$, if the weights of each edge belongs to the set $\{0, X\}$, where ' $X$ ' is any non-negative real integer.

## The main Characteristics are:

- Process is done only on 'binary graphs'. ('Binary Graph': Any graph whose edge-weights belongs to the set $\{0, \mathrm{X}\}$, where ' $X$ ' is any non-negative real integer.)
- For optimization purposes, any duplicated value in thestart_vids or end_vids are ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending


## Summary

pgr_binaryBreadthFirstSearch(Edges SQL, start vid, end vid, [directed])
pgr_binaryBreadthFirstSearch(Edges SQL, start vid, end vids, [directed])
pgr_binaryBreadthFirstSearch(Edges SQL, start vids, end vid, [directed])
pgr_binaryBreadthFirstSearch(Edges SQL, start vids, end vids, [directed])
pgr_binaryBreadthFirstSearch(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (seq, path_seq, [start_vid], [end_vid], node, edge, cost, agg_cost)
OR EMPTY SET

Note: Using the Sample Data Network as all weights are same (i.e<br>(1`))

## One to One

```
pgr_binaryBreadthFirstSearch(Edges SQL, start vid, end vid, [directed])
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(6 \backslash)$ to vertex $\backslash(10 \backslash)$ on a directed graph

```
SELECT * FROM pgr_binaryBreadthFirstSearch
'SELECT id, source, target, cost, reverse_cost from edges',
6,10, true);
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{rcccc}
\(1 \mid\) & \(1 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(2 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(3 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(4 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) \\
\(5 \mid\) & \(5 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(6 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\((6\) rows \()\) & & & &
\end{tabular}
```


## One to Many

```
pgr_binaryBreadthFirstSearch(Edges SQL, start vid, end vids, [directed]
RETURNS SET OF (seq, path_seq, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(6 \backslash)$ to vertices $\backslash(\backslash\{10,17 \backslash\} \backslash)$ on a directed graph

```
SELECT * FROM pgr_binaryBreadthFirstSearch(
    'SELECT id, source, target, cost, reverse_cost from edges',
    6, ARRAY[10, 17]);
seq | path_seq | end_vid | node | edge | cost | agg_cost
    10| 6| 4| 1| 0
    2| 10| 7| 8| 1| 1
    < 3|-10| 11| 9| 1| 2
    | 10| 16| 16| 1| 3
    10| 15| 3| 1| 4
    10| 10| -1| 0| 5
    17| 6| 4| 1| 0
    17| 71 8| 11| 11 1 |
    | 4| 17| 12| 13| 1| < 3
    11| 5| 17| 17| -1| 0| 4
(11 rows)
```


## Many to One

RETURNS SET OF (seq, path_seq, start_vid, node, edge, cost, agg_cost) OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,1 \backslash\} \backslash)$ to vertex $\backslash(17 \backslash)$ on a directed graph

SELECT * FROM pgr_binaryBreadthFirstSearch(
'SELECT id, source, target, cost, reverse_cost from edges',
ARRAY[6, 1], 17),
seq | path_seq | start_vid | node | edge | cost | agg_cost


Many to Many

> pgr_binaryBreadthFirstSearch(Edges SQL, start vids, end vids, [directed])
> RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)

OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,1 \backslash\} \backslash)$ to vertices $\backslash(\backslash\{10,17 \backslash\} \backslash)$ on an undirected graph


## Combinations

```
pgr_binaryBreadthFirstSearch(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

Using a combinations table on anundirected graph

The combinations table:

The query:

```
SELECT * FROM pgr_binaryBreadthFirstSearch(
    'SELECT id, source, target, cost, reverse_cost FROM edges',
    'SELECT source, target FROM combinations',
    false);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
    1| 1| 5| 6| 5| 1| 1| 0
2| 2| 5| 6| 6| -1| 0| 1
3| 1| 5| 10| 5| 1| 1| 0
4| 2| 5| 10| 6| 2| 1| 1
5| 3| 5| 10| 10| -1| 0| 2
6| 1| 6| 5| 6| 1| 1| 0
7| 2| 6| 5| 5| -1| 0| 1
9| 2| 6| 15| 10| 3| 1| 1
cllllll
(10 rows)
```

Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | BIGINT | Identifier of the starting vertex of the path. |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting vertices. |
| end vid | BIGINT | Identifier of the ending vertex of the path. |
| end vids | ARRAY[BIGINT] | Array of identifiers of ending vertices. |

Optional Parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true |  When true the graph is considered Directed <br>   <br>   <br>   <br>   <br>  Undirected. n false the graph is considered as |  |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Combinations SQL

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- | Identifier of the departure vertex. |
|  | INTEGER |  |
| target | ANY- <br>  <br> INTEGER | Identifier of the arrival vertex. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Result Columns

Set of (seq, path_id, path_seq [, start_vid] [, end_vid], node, edge, cost, agg_cost)

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. |
| path_id | INTEGER | Path identifier. <br> - Has value $\mathbf{1}$ for the first of a path fromstart_vid to end_vid. |
| path_seq | INTEGER | Relative position in the path. Has value 1 for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex. Returned when multiple starting vetrices are in the query. Many to One Many to Many Combinations |
| end_vid | BIGINT | Identifier of the ending vertex. Returned when multiple ending vertices are in the query. <br> - One to Many <br> - Many to Many <br> - Combinations |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for the last node of the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | FLOAT | Aggregate cost from start_vid to node. |

## Additional Examples

## Example:

Manually assigned vertex combinations.

```
SELECT * FROM pgr_binaryBreadthFirstSearch(
    'SELECT id, source, target, cost, reverse_cost FROM edges',
    'SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)');
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
```

| $1 \mid$ | $1 \mid$ | $6 \mid$ | $7 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| ---: | ---: | ---: | ---: | ---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $6 \mid$ | $7 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |
| $3 \mid$ | $1 \mid$ | $6 \mid$ | $10 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0 |
| $4 \mid$ | $2 \mid$ | $6 \mid$ | $10 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 1 |
| $5 \mid$ | $3 \mid$ | $6 \mid$ | $10 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 2 |
| $6 \mid$ | $4 \mid$ | $6 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 3 |
| $7 \mid$ | $5 \mid$ | $6 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ | 4 |
| $8 \mid$ | $6 \mid$ | $6 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 5 |
| $9 \mid$ | $1 \mid$ | $12 \mid$ | $10 \mid$ | $12 \mid$ | $13 \mid$ | $1 \mid$ | 0 |
| $10 \mid$ | $2 \mid$ | $12 \mid$ | $10 \mid$ | $17 \mid$ | $15 \mid$ | $1 \mid$ | 1 |
| $11 \mid$ | $3 \mid$ | $12 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ | 2 |
| $12 \mid$ | $4 \mid$ | $12 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ | 3 |
| $13 \mid$ | $5 \mid$ | $12 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 4 |
| $(13$ rows) |  |  |  |  |  |  |  |

## See Also

## - Sample Data

- https://cp-algorithms.com/graph/01_bfs.html
- https://en.wikipedia.org/wiki/Dijkstra\'s_algorithm\#Specialized_variants
- Index
- Search Page


## See Also

## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2


## Coloring - Family of functions

## Proposed

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.
- pgr_sequentialVertexColoring - Proposed - Vertex coloring algorithm using greedy approach.


## Experimental

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions
They are not officially of the current release.

- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting
- pgr_bipartite -Experimental - Bipartite graph algorithm using a DFS-based coloring approach.
- pgr_edgeColoring - Experimental - Edge Coloring algorithm using Vizing's theorem.
- Supported versions: Latest (3.3) 3.2


## boost

Boost Graph Inside

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.


## Availability

- Version 3.3.0
- Promoted to proposed signature
- Version 3.2.0
- New experimental signature


## Description

Sequential vertex coloring algorithm is a graph coloring algorithm in which color identifiers are assigned to the vertices of a graph in a sequential manner, such that no edge connects two identically colored vertices.

## The main Characteristics are:

- The implementation is applicable only for undirected graphs.
- Provides the color to be assigned to all the vertices present in the graph.
- Color identifiers values are in the Range $\backslash([1,|\mathrm{~V}|] \backslash)$
- The algorithm tries to assign the least possible color to every vertex.
- Efficient graph coloring is an NP-Hard problem, and therefore, this algorithm does not always produce optimal coloring. It follows a greedy strategy by iterating through all the vertices sequentially, and assigning the smallest possible color that is not used by its neighbors, to each vertex.
- The returned rows are ordered in ascending order of the vertex value.
- Sequential Vertex Coloring Running Time: $\backslash\left(\mathrm{O}\left(|\mathrm{V}|^{*}(\mathrm{~d}+\mathrm{k})\right) \backslash\right)$
- where $\backslash(|\mathrm{V}| \backslash)$ is the number of vertices,
- $\backslash(\mathrm{d} \backslash)$ is the maximum degree of the vertices in the graph,
- $\backslash(\mathrm{k} \backslash)$ is the number of colors used.


## Signatures

pgr_sequentialVertexColoring(Edges SQL)
RETURNS SET OF (vertex_id, color_id)
OR EMPTY SET

## Example:

Graph coloring of pgRouting Sample Data

```
SELECT * FROM pgr_sequentialVertexColoring
    'SELECT id, source, target, cost, reverse_cost FROM edges
    ORDER BY id
);
vertex_id | color_id
    1| 1
    2| 1
    3|}
    4| 2
    5| 1
    6| 2
    7| 1
    8| 2
    9| 1
    1| 2
    12| 1
    3| 1
    14| 2
    5| 2
    16| 1
    17| 2
(17 rows)
```


## Parameters

| Parameter | Type | Description |
| :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL | -1 |

- When negative: edge (target, source) does not exist, therefore it's not part of the graph.


## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Result Columns

Returns SET OF (vertex_id, color_id)

| Column | Type | Description |
| :--- | :---: | :--- |
| vertex_id | BIGINT | Identifier of the vertex. |
| color_id | BIGINT | Identifier of the color of the vertex. |
|  |  | The minimum value of color is |
|  |  | 1. |

## See Also

- The queries use the Sample Data network.
- Boost: Sequential Vertex Coloring algorithm documentation
- Wikipedia: Graph coloring


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2


## pgr_bipartite -Experimental

pgr_bipartite - Disjoint sets of vertices such that no two vertices within the same set are adjacent.

## boost

Boost Graph Inside

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.2.0
- New experimental signature


## Description

A bipartite graph is a graph with two sets of vertices which are connected to each other, but not within themselves. A bipartite graph is possible if the graph coloring is possible using two colors such that vertices in a set are colored with the same color.

## The main Characteristics are:

- The algorithm works in undirected graph only.
- The returned values are not ordered.
- The algorithm checks graph is bipartite or not. If it is bipartite then it returns the node along with two colors 0 and 1 which represents two different sets.
- If graph is not bipartite then algorithm returns empty set.
- Running time: $\backslash(O(V+E) \backslash)$


## Signatures

```
pgr_bipartite(Edges SQL)
RETURNS SET OF (vertex_id, color_id)
OR EMPTY SET
```


## Example:

When the graph is bipartite

| 21 | 0 |
| :---: | :---: |
| 31 | 1 |
| $4 \mid$ | 1 |
| 51 | 0 |
| 61 | 1 |
| 71 | 0 |
| 8\| | 1 |
| 91 | 0 |
| 10\| | 0 |
| 11\| | 1 |
| $12 \mid$ | 0 |
| 13\| | 0 |
| 14\| | 1 |
| 151 | 1 |
| 16\| | 0 |
| 17\| | 1 |
| rows) |  |

## Parameters

| Parameter | Type | Description |
| :--- | :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |
|  |  |  |
|  |  |  |
|  |  |  |

## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

Returns SET OF (vertex_id, color_id)

| Column | Type | Description |
| :--- | :--- | :--- |
| vertex_id | BIGINT | Identifier of the vertex. |
| color_id | BIGINT | Identifier of the color of the vertex. |
|  |  | 0 |
|  |  | The minimum value of color is |
|  |  |  |

## Additional Example

## Example:

The odd length cyclic graph can not be bipartite.
The edge $\backslash(5$ rightarrow $1 \backslash)$ will make subgraph with vertices $\backslash(\backslash\{1,3,7,6,5 \backslash\} \backslash)$ an odd length cyclic graph, as the cycle has 5 vertices.

```
INSERT INTO edges (source, target, cost, reverse_cost) VALUES
(5, 1, 1, 1);
INSERT 0 }
```

Edges in blue represent odd length cycle subgraph.


```
SELECT * FROM pgr_bipartite(
    $$SELECT id, source, target, cost, reverse_cost FROM edges$$
);
vertex_id | color_id
(0 rows)
```


## See Also

- Boost: is_bipartite
- Wikipedia: bipartite graph
- Sample Data network.


## Indices and tables

- Index
- Search Page


## - Supported versions: Latest (3.3)

## pgr_edgeColoring - Experimental

pgr_edgeColoring - Returns the edge coloring of undirected and loop-free graphs

## 8 boost

Boost Graph Inside

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.3.0
- New experimental signature


## Description

Edge Coloring is an algorithm used for coloring of the edges for the vertices in the graph. It is an assignment of colors to the edges of the graph so that no two adjacent edges have the same color.

## The main Characteristics are:

- The implementation is for undirected and loop-free graphs


## - loop free:

no self-loops and no parallel edges.

- Provides the color to be assigned to all the edges present in the graph.
- At most <br>(\Delta $+1 \backslash$ ) colors are used, where $\backslash(\backslash$ Delta $\backslash)$ is the degree of the graph.
- This is optimal for some graphs, and by Vizing's theorem it uses at most one color more than the optimal for all others.
- When the graph is bipartite
- the chromatic number $\backslash\left(x^{\prime}(G) \backslash\right)$ (minimum number of colors needed for proper edge coloring of graph) is equal to the degree $\backslash(\backslash$ Delta $+1 \backslash)$ of the graph, $\Lambda\left(x^{\prime}(G)=\backslash\right.$ Delta $\left.\left.\backslash\right)\right)$
- The algorithm tries to assign the least possible color to every edge.
- Does not always produce optimal coloring.
- The returned rows are ordered in ascending order of the edge identifier.
- Efficient graph coloring is an NP-Hard problem, and therefore:
- In this implelentation the running time: $\backslash\left(\mathrm{O}\left(|E|^{*}|\mathrm{~V}|\right) \backslash\right)$ - where $\backslash(|E| \backslash)$ is the number of edges in the graph, - $\backslash(|\mathrm{V}| \backslash)$ is the number of vertices in the graph.


## Signatures

```
pgr_edgeColoring(Edges SQL)
RETURNS SET OF (edge_id, color_id)
OR EMPTY SET
```


## Example:

Graph coloring of pgRouting Sample Data

```
SELECT * FROM pgr_edgeColoring(
    'SELECT id, source, target, cost, reverse_cost FROM edges
    ORDER BY id'
);
edge_id | color_id
1| 3
    2| 2
    3|}
    4| 4
    6| 1
    7| 2
    8| 1
    |
    |
    5
    3
        2
        1
        3
        1
        |
    18| 1
(18 rows)
```

| Parameter | Type | Description |  |
| :--- | :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |  |
|  |  |  |  |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL -1 | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |
|  |  |  |
|  |  |  |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

Returns SET OF (edge_id, color_id)

| Column | Type | Description |
| :--- | :--- | :--- |
| edge_id | BIGINT | Identifier of the edge. |
| color_id | BIGINT | Identifier of the color of the edge. |

- The minimum value of color is 1.


## See Also

- The queries use the Sample Data network.
- Boost: Edge Coloring Algorithm documentation
- Wikipedia: Graph Coloring


## Indices and tables

- Index
- Search Page


## Result Columns

Returns SET OF (vertex_id, color_id)

| Column | Type | Description |
| :--- | :--- | :--- |
| vertex_id | BIGINT | Identifier of the vertex. |
| color_id | BIGINT | Identifier of the color of the vertex. |
|  |  | The minimum value of color is <br>  |
|  |  |  |

Returns SET OF (edge_id, color_id)

| Column | Type | Description |
| :--- | :--- | :--- |
| edge_id | BIGINT | Identifier of the edge. |
| color_id | BIGINT | Identifier of the color of the edge. |
|  |  | The minimum value of color is |
|  |  | 1. |

- Boost: Sequential Vertex Coloring algorithm documentation
- Wikipedia: Graph coloring
- Boost: is_bipartite
- Wikipedia: bipartite graph
- Boost: Edge Coloring Algorithm documentation
- Wikipedia: Graph Coloring


## Indices and tables

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## categories

Cost - Category

- pgr_withPointsCost - Proposed


## Cost Matrix - Category

- pgr_withPointsCostMatrix - proposed

Driving Distance - Category

- pgr_withPointsDD - Proposed - Driving Distance based on pgr_withPoints


## K shortest paths - Category

- pgr_withPointsKSP - Proposed - Yen's algorithm based on pgr_withPoints


## Via - Category

- pgr_dijkstraVia - Proposed
withPoints - Category
- withPoints - Family of functions - Functions based on Dijkstra algorithm.
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: $2.6 \mathbf{2}$.5 $2.4 \mathbf{2 . 3} \mathbf{2 . 2}$
withPoints - Family of functions
When points are also given as input:


## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.
pgr_withPoints - Proposed - Route from/to points anywhere on the graph.
- pgr_withPointsCost - Proposed - Costs of the shortest paths.
pgr_withPointsCostMatrix - proposed - Costs of the shortest paths.
pgr_withPointsKSP - Proposed - K shortest paths.
pgr_withPointsDD - Proposed - Driving distance.
pgr_withPoints - Returns the shortest path in a graph with additional temporary vertices.


## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.


## boost

Boost Graph Inside

## Availability

- Version 3.2.0
- New proposed function:
- pgr_withPoints(Combinations)
- Version 2.2.0
- New proposed function


## Description

Modify the graph to include points defined by points_sql. Using Dijkstra algorithm, find the shortest path(s)

## The main characteristics are:

- Process is done only on edges with positive costs.
- Vertices of the graph are:
- positive when it belongs to the edges_sql
- negative when it belongs to the points_sql
- Values are returned when there is a path.
- When the starting vertex and ending vertex are the same, there is no path. - The agg_cost the non included values (v, v) is 0
- When the starting vertex and ending vertex are the different and there is no path: - The agg_cost the non included values ( $u, v$ ) is $\infty$
- For optimization purposes, any duplicated value in the start_vids or end_vids are ignored.
- The returned values are ordered: - start_vid ascending - end_vid ascending
- Running time: $\backslash(\mathrm{O}(\mid$ start $\backslash \mathrm{vids} \mid \backslash \operatorname{times}(\mathrm{V} \backslash \log \mathrm{V}+\mathrm{E})) \backslash)$


## Signatures

## Summary

```
pgr_withPoints(Edges SQL, Points SQL, start vid, end vid, [options])
pgr_withPoints(Edges SQL, Points SQL, start vid, end vids, [options])
pgr_withPoints(Edges SQL, Points SQL, start vids, end vid, [options])
pgr_withPoints(Edges SQL, Points SQL, start vids, end vids, [options])
pgr_withPoints(Edges SQL, Points SQL, Combinations SQL, [options])
options: [directed, driving_side, details])
RETURNS SET OF (seq, path_seq, [start_pid], [end_pid], node, edge, cost, agg_cost)
OR EMTPY SET
```

pgr_withPoints(Edges SQL, Points SQL, start vid, end vid, [options])
options: [directed, driving_side, details])
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMTPY SET

## Example:

From point $\backslash(1 \backslash)$ to vertex $\backslash(10 \backslash)$ with details

```
SELECT * FROM pgr_withPoints(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    SELECT pid, edge id, fraction, side from pointsOfInterest'
    -1, 10,
    details => true);
seq | path_seq | node | edge | cost | agg_cost
    1| -1| 1| 0.6| 0
    2| 6| 4| 0.7| 0.6
    3| -6| 4| 0.3 1. 1.3
    4| 7| 8| 1| 1.6
    5| 11| 9| 1| 2.6
    6| 16| 16| 1| 3.6
    7| 15| 3| 1| 4.6
    8| 10| -1| 0| 5.6
8 rows
```

One to Many
pgr_withPoints(Edges SQL, Points SQL, start vid, end vids, [options])
options: [directed, driving_side, details])
RETURNS SET OF (seq, path_seq, end_pid, node, edge, cost, agg_cost)
OR EMTPY SET

## Example:

From point $\backslash(1 \backslash)$ to point $\backslash(3 \backslash)$ and vertex $\backslash(7 \backslash)$ on an undirected graph

```
SELECT * FROM pgr_withPoints(
    SELECT id, source, target, cost, reverse cost FROM edges ORDER BY id',
    SELECT pid, edge_id, fraction, side from pointsOfInterest',
    -1, ARRAY[-3, 7],
eq| path seq|end pid | node|edge | cost | agg_cos
1| -3| -1| 1| 0.6 | 0
2|-3| 6| 4| 1| 0.6
3| -3| 7| 10| 10 1.6
4| -3| 8| 12| 0.6| 2.6
5| -3| -3| -1| 0| 3.2
1| 7| -1| 1| 0.6| 0
2| 7| 6| 4| 1| 0.6
8| 3| 7| 7| -1| 0| 1.6
(8 rows)
```


## Many to One

pgr_withPoints(Edges SQL, Points SQL, start vids, end vid, [options])
options: [directed, driving_side, details])
RETURNS SET OF (seq, path_seq, start_pid, node, edge, cost, agg_cost) OR EMTPY SET

## Example:

From point $\backslash(1 \backslash)$ and vertex $\backslash(6 \backslash)$ to point $\backslash(3 \backslash)$

```
SELECT * FROM pgr_withPoints(
```

    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOfInterest',
    ARRAY[-1, 6], -3);
seq | path_seq | start_pid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $-1 \mid$ | $-1 \mid$ | $1 \mid$ | $0.6 \mid$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $-1 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| $3 \mid$ | $3 \mid$ | $-1 \mid$ | $7 \mid$ | $10 \mid$ | $1 \mid$ |
| $4 \mid$ | $4 \mid$ | $-1 \mid$ | $8 \mid$ | $12 \mid$ | $0.6 \mid$ |
| $4 \mid$ | 2.6 |  |  |  |  |
| $5 \mid$ | $5 \mid$ | $-1 \mid$ | $-3 \mid$ | $-1 \mid$ | $0 \mid$ |
| $6 \mid$ | $1 \mid$ | $6 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| $7 \mid$ | $2 \mid$ | $6 \mid$ | $7 \mid$ | $10 \mid$ | $1 \mid$ |
| $8 \mid$ | $3 \mid$ | $6 \mid$ | $8 \mid$ | $12 \mid$ | $0.6 \mid$ |
| $9 \mid$ | $4 \mid$ | $6 \mid$ | $-3 \mid$ | $-1 \mid$ | $0 \mid$ |
| 9 rows) |  |  |  | 2.6 |  |

## Many to Many

pgr_withPoints(Edges SQL, Points SQL, start vids, end vids, [options])
options: [directed, driving_side, details])

RETURNS SET OF (seq, path_seq, start_pid, end_pid, node, edge, cost, agg_cost) OR EMTPY SET

## Example:

From point $\backslash(1 \backslash)$ and vertex $\backslash(6 \backslash)$ to point $\backslash(3 \backslash)$ and vertex $\backslash(1 \backslash)$

```
SELECT * FROM pgr_withPoints(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id'
    'SELECT pid, edge id, fraction, side from pointsOfInterest',
ARRAY[-1, 6], ARRAY[-3, 1]);
seq | path_seq | start_pid | end_pid | node | edge | cost | agg_cost
    1| 1| -1| -3| -1| 1| 0.6| 0
2| 2| -1| -3| 6| 4| 1| 0.6
4| 4| -1| -3| 8| 12| 0.6 | 2.6
5| 5| -1| -3| -3| -1| 0| 3. 
1| -1| 1| -1| 1| 0.6| 0
2| -1| 1| 6| 4| 1| 0.6
1| 7| 7| 1| 1.6
3| 6| 1| 2.6
| 1| 1| -1| 0| 3.6
6| -3| 6| 4| 1| 0
|
| -3| -3| -1| 0| | 2.6
| 1| 6| 4| 1| 0
| 1| 7| 7| 1| 1
|llllll
|lllllll
(18 rows)
```


## Combinations

```
pgr_withPoints(Edges SQL, Points SQL, Combinations SQL, [options])
options: [directed, driving_side, details])
RETURNS SET OF (seq, path_seq, start_pid, end_pid, node, edge, cost, agg_cost)
OR EMTPY SET
```


## Example:

## Two combinations

From point $\backslash(1 \backslash)$ to vertex $\backslash(10 \backslash)$, and from vertex $\backslash(6 \backslash)$ to point $\backslash(3 \backslash)$ with right side driving.

```
SELECT * FROM pgr_withPoints(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOfInterest',
    SELECT * FROM (VALUES (-1, 10), (6, -3)) AS combinations(source, target)'
    driving_side => 'r', details => true);
seq | path_seq | start_pid | end_pid | node | edge | cost | agg_cost
\begin{tabular}{rccccccc}
\(1 \mid\) & \(1 \mid\) & \(-1 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(1 \mid\) & \(0.4 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(-1 \mid\) & \(10 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) & 0.4 \\
\(3 \mid\) & \(3 \mid\) & \(-1 \mid\) & \(10 \mid\) & \(6 \mid\) & \(4 \mid\) & \(0.7 \mid\) & 1.4 \\
\(4 \mid\) & \(4 \mid\) & \(-1 \mid\) & \(10 \mid\) & \(-6 \mid\) & \(4 \mid\) & \(0.3 \mid\) & 2.1 \\
\(5 \mid\) & \(5 \mid\) & \(-1 \mid\) & \(10 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 2.4 \\
\(6 \mid\) & \(6 \mid\) & \(-1 \mid\) & \(10 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) & 3.4 \\
\(7 \mid\) & \(7 \mid\) & \(-1 \mid\) & \(10 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) & 4.4 \\
\(8 \mid\) & \(8 \mid\) & \(-1 \mid\) & \(10 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) & 5.4 \\
\(9 \mid\) & \(9 \mid\) & \(-1 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 6.4 \\
\(10 \mid\) & \(1 \mid\) & \(6 \mid\) & \(-3 \mid\) & \(6 \mid\) & \(4|0.7|\) & 0 \\
\(11 \mid\) & \(2 \mid\) & \(6 \mid\) & \(-3 \mid\) & \(-6 \mid\) & \(4 \mid\) & \(0.3 \mid\) & 0.7 \\
\(12 \mid\) & \(3 \mid\) & \(6 \mid\) & \(-3 \mid\) & \(7 \mid\) & \(10 \mid\) & \(1 \mid\) & 1 \\
\(13 \mid\) & \(4 \mid\) & \(6 \mid\) & \(-3 \mid\) & \(8 \mid\) & \(12 \mid\) & \(0.6 \mid\) & 2 \\
\(14 \mid\) & \(5 \mid\) & \(6 \mid\) & \(-3 \mid\) & \(-3 \mid\) & \(-1 \mid\) & \(0 \mid\) & 2.6 \\
\((14\) rows) & & & & &
\end{tabular}
```

Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Points SQL | TEXT | Points SQL as described below |
| Combinations SQL | TEXT | Combinations SQL as described below <br> start vid <br> start vids <br> end vid |
| ARRAY[BIGINT] | Identifier of the starting vertex of the path. Negative value is for point's <br> identifier. |  |
| Array of identifiers of starting vertices. Negative values are for point's |  |  |
| identifiers. |  |  |

## Optional parameters

Column Type Default Description

| directed BOOLEAN true | When true the graph is considered Directed |
| :--- | :--- |
|  | When false the graph is considered as |
| Undirected. |  |

With points optional parameters

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| driving_side | CHAR | b | Value in [r, l, b] indicating if the driving side is: |
|  |  |  | - r for right driving side. |
|  |  |  | - I for left driving side. |
|  |  |  | - b for both. |
| details | BOOLEAN | false | - When true the results will include the points that are in the path. |
|  |  |  | - When false the results will not include the points that are in the path. |

nner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |

Where:

ANY-INTEGER:
SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Points SQL

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| pid | ANY-INTEGER | value | Identifier of the point. <br> - Use with positive value, as internally will be converted to negative value <br> - If column is present, it can not be NULL. <br> - If column is not present, a sequential negativevalue will be given automatically. |
| edge_id | ANY-INTEGER |  | Identifier of the "closest" edge to the point. |
| fraction | ANY-NUMERICAL |  | Value in $\langle 0,1\rangle$ that indicates the relative postition from the first end point of the edge. |
| side | CHAR | b | Value in $[b, r, l$, NULL] indicating if the point is: <br> - In the right $r$, <br> - In the left I, <br> - In both sides b, NULL |

## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Combinations SQL

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- <br> INTEGER | Identifier of the departure vertex. |
| target | ANY- <br> INTEGER | Identifier of the arrival vertex. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Result Columns

Returns set of (seq, path_seq[, start_pid] [, end_pid], node, edge, cost, agg_cost)

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. |
| path_seq | INTEGER | Relative position in the path. <br> 1 For the first row of the path. |
| start_pid | BIGINT | Identifier of a starting vertex/point of the path. <br> - When positive is the identifier of the starting vertex. <br> - When negative is the identifier of the starting point. <br> - Returned on Many to One and Many to Many |
| end_pid | BIGINT | Identifier of an ending vertex/point of the path. <br> - When positive is the identifier of the ending vertex. <br> - When negative is the identifier of the ending point. <br> - Returned on One to Many and Many to Many |
| node | BIGINT | Identifier of the node in the path fromstart_pid to end_pid. <br> - When positive is the identifier of the a vertex. <br> - When negative is the identifier of the a point. |


| Column | Type | Description |
| :--- | :--- | :--- |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path <br> sequence. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
|  |  | Aggregate cost from start_vid to node. |
| agg_cost | FLOAT | $\mathbf{0}$ For the first row of the path. |

## Additional Examples

## - Usage variations

- Passes in front or visits with right side driving.
- Passes in front or visits with left side driving.

Usage variations

All the examples are about traveling from point $\backslash(1 \backslash)$ and vertex $\backslash(5 \backslash)$ to points $\backslash(\backslash\{2,3,6 \backslash\} \backslash)$ and vertices $\backslash(\backslash\{10,11 \backslash\} \backslash)$

SELECT *
FROM pgr_withPoints(
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
'SELECT pid, edge_id, fraction, side from pointsOfinterest',
ARRAY[5, -1], ARRAY[-2, $-3,-6,10,11]$,
driving_side => 'r', details => true);
seq | path_seq | start_pid | end_pid | node | edge | cost | agg_cost

| 1 \| | 1\| | -1 \| | -6\| | -1\| | 1\| 0.4 | | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | -1\| | -6\| | 51 | 1\| 1 | | 0.4 |
| 31 | 31 | -1\| | -6\| | 61 | 4\|0.7| | 1.4 |
| 4\| | 4 \| | -1\| | -6\| | -6\| | -1\| 0 | | 2.1 |
| 51 | 1 \| | -1\| | -3\| | -1\| | 1\| 0.4 | | 0 |
| 61 | 21 | -1\| | -3\| | 51 | 1\| 1 | | 0.4 |
| 71 | 31 | -1\| | -3\| | 61 | $4 \mid 0.7$ \| | 1.4 |
| 8\| | 4\| | -1\| | -3\| | -6\| | 4\|0.3| | 2.1 |
| 9\| | 51 | -1\| | -3\| | 71 | 10\| 1| | 2.4 |
| 10\| | 61 | -1\| | -3\| | 8। | 12\| $0.6 \mid$ | 3.4 |
| 11\| | 71 | -1\| | -3\| | -3\| | -1\| 0 | | 4 |
| 12\| | 1 \| | -1\| | -2\| | -1\| | 1\| 0.4 | | 0 |
| 13\| | 21 | -1\| | -2\| | 5। | 1\| 1| | 0.4 |
| 14\| | 31 | -1\| | -2\| | 61 | 4\|0.7| | 1.4 |
| 15\| | 4 | -1\| | -2\| | -6\| | 4\|0.3| | 2.1 |
| 16\| | 5 | -1\| | -2\| | 71 | 8\| 1| | 2.4 |
| 17\| | 6 | -1\| | -2\| | 11\| | 9\| 1| | 3.4 |
| 18\| | 7 | -1\| | -2\| | $16 \mid$ | 15\|0.4| | 4.4 |
| 19\| | 8 | -1\| | -21 | -2\| | -1\| $0 \mid$ | 4.8 |
| 201 | 1 | -1\| | 10\| | -1\| | 1\| $0.4 \mid$ | 0 |
| 21\| | 2 | -1\| | $10 \mid$ | 51 | 1\| 1| | 0.4 |
| 221 | 3 | -1\| | $10 \mid$ | 61 | 4\|0.7| | 1.4 |
| 23\| | 4 | -1\| | $10 \mid$ | -6\| | 4\|0.3| | 2.1 |
| 24\| | 5 | -1\| | $10 \mid$ | 71 | 8\| 1| | 2.4 |
| 251 | 6 | -1\| | $10 \mid$ | 11\| | 9\|1| | 3.4 |
| $26 \mid$ | 7 | -1\| | $10 \mid$ | $16 \mid$ | 16\| 1| | 4.4 |
| $27 \mid$ | 81 | -1\| | $10 \mid$ | 15\| | 3\| 1 | | 5.4 |
| 28\| | 91 | -1\| | $10 \mid$ | $10 \mid$ | -1\| 0 | | 6.4 |
| 291 | $1 \mid$ | -1\| | 11\| | -1\| | 1\| 0.4 | | 0 |
| 301 | 2 | -1\| | $11 \mid$ | 51 | 1\| 1 | | 0.4 |
| 31\| | 31 | -1\| | 11\| | 61 | $4 \mid 0.7$ \| | 1.4 |
| 321 | 4 | -1\| | 11\| | -6\| | 4\|0.3| | 2.1 |
| 331 | 51 | -1\| | 11\| | 71 | 8\| 1| | 2.4 |
| $34 \mid$ | 6 | -1\| | 11\| | 11\| | -1\| 0 | | 3.4 |
| 351 | $1 \mid$ | 51 | -6\| | 51 | 1\| 1 | | 0 |
| $36 \mid$ | 21 | 51 | -6\| | 61 | 4\|0.7 | | 1 |
| 371 | 31 | 51 | -6\| | -6\| | -1\| 0 | | 1.7 |
| 38\| | 1 \| | 51 | -3\| | 51 | 1\| 1| | 0 |
| 391 | 21 | 51 | -3\| | 61 | 4\|0.7 | | 1 |
| 401 | 31 | 51 | -3\| | -6\| | 4\|0.3| | 1.7 |
| 41\| | 41 | 51 | -3\| | 71 | 10\| 1 | | 2 |
| $42 \mid$ | 51 | 51 | -3\| | 8। | 12\| 0.6 | | 3 |
| 431 | 61 | 51 | -31 | -3\| | -1\| $0 \mid$ | 3.6 |
| $44 \mid$ | 1 \| | 51 | -2। | 51 | 1\| 1| | 0 |
| 451 | 21 | 51 | -2\| | 61 | 4\| 0.7 | | 1 |
| $46 \mid$ | 31 | 51 | -2\| | -6\| | 4\|0.3| | 1.7 |
| 471 | 4 \| | 51 | -2\| | 71 | 8\| 1| | 2 |
| 48\| | 51 | 51 | -2\| | 11\| | 9\| 1| | 3 |
| 491 | 61 | 51 | -2\| | 16\| | 15\| $0.4 \mid$ | 4 |
| 501 | 71 | 51 | -2\| | -2\| | -1\| 0 | | 4.4 |
| 51\| | 1 \| | 51 | $10 \mid$ | 51 | 1\| 1| | 0 |
| 521 | 21 | 51 | $10 \mid$ | 61 | 4\|0.7 | | 1 |
| 531 | 31 | 51 | $10 \mid$ | -6\| | 4\|0.3| | 1.7 |
| 541 | 41 | 51 | $10 \mid$ | 71 | 8\| 1| | 2 |
| 551 | 51 | 51 | $10 \mid$ | 11\| | 9\| 1| | 3 |
| 56\| | 61 | 51 | 10\| | $16 \mid$ | 16\| 1| | 4 |
| 571 | 71 | 51 | $10 \mid$ | 15\| | 3\| 1| | 5 |
| 58\| | 81 | 51 | 10\| | $10 \mid$ | -1\| 0 | | 6 |
| 591 | $1 \mid$ | 51 | 11\| | 51 | 1\| 1| | 0 |
| 601 | 21 | 51 | 11\| | 61 | 4\|0.7 | | 1 |
| 61\| | 31 | 51 | 11\| | -6\| | 4\|0.3| | 1.7 |
| 621 | 41 | 51 | 11\| | 71 | 8\| 1| | 2 |
| 631 | 51 | 51 | 11\| | 11\| | -1\| 0 | | 3 |

## Passes in front or visits with right side driving

For point $\backslash(6 \backslash)$ and vertex $\backslash(11 \backslash)$.

SELECT (start_pid || ' -> ' || end_pid ||' at ' || path_seq || 'th step')::TEXT AS path_at,
CASE WHEN edge $=-1$ THEN ' visits'
ELSE ' passes in front of'
END as status,
CASE WHEN node < 0 THEN 'Point'
ELSE 'Vertex'
END as is_a,
abs(node) as id
FROM pgr_withPoints(
'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
'SELECT pid, edge_id, fraction, side from pointsOfInterest',
ARRAY[5, -1], ARRAY[-2, -3, -6, 10, 11],
driving_side => 'r', details => true)
WHERE node IN (-6, 11);

| path_at | status \| is_a |id |
| :---: | :---: |
| -1 -> -6 at 4th step \| | visits \| Point | 6 |
| $-1->-3$ at 4th step \| | passes in front of \| Point | 6 |
| $-1->-2$ at 4th step \| | passes in front of \| Point | 6 |
| $-1->-2$ at 6th step \| | passes in front of \| Vertex | 11 |
| $-1->10$ at 4th step | passes in front of \| Point | 6 |
| $-1->10$ at 6th step | passes in front of \| Vertex | 11 |
| $-1->11$ at 4th step | passes in front of \| Point | 6 |
| $-1->11$ at 6th step | visits \| Vertex|11 |
| $5->-6$ at 3th step | visits \| Point | 6 |
| $5 \rightarrow-3$ at 3th step | passes in front of \| Point | 6 |
| $5->-2$ at 3th step | passes in front of \| Point | 6 |
| $5->-2$ at 5th step | passes in front of \| Vertex | 11 |
| $5->10$ at 3th step | passes in front of Point \| 6 |
| $5->10$ at 5th step | passes in front of \| Vertex | 11 |
| $5->11$ at 3th step | passes in front of Point \| 6 |
| $5->11$ at 5th step | visits \| Vertex|11 |
| (16 rows) |  |

Passes in front or visits with left side driving.

For point $\backslash(6 \backslash)$ and vertex $\backslash(11 \backslash)$

```
SELECT (start_pid || ' => ' | end_pid ||' at ' || path_seq || 'th step')::TEXT AS path_at,
CASE WHEN edge =-1 THEN ' visits'
    ELSE ' passes in front of'
END as status,
CASE WHEN node < 0 THEN 'Point'
        ELSE 'Vertex'
    END as is_a,
    abs(node) as id
FROM pgr_withPoints(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOflnterest',
    ARRAY[5, -1], ARRAY[-2, -3, -6, 10, 11],
driving_side => ''', details => true)
WHERE node IN (-6, 11);
\begin{tabular}{|c|c|}
\hline path_at & status | is_a | id \\
\hline \(-1=>-6\) at 3th step & visits | Point | 6 \\
\hline \(-1=>-3\) at 3th step & passes in front of | Point \\
\hline \(-1=>-2\) at 3th step & passes in front of | Point \\
\hline \(-1=>-2\) at 5th step & passes in front of | Vertex | 11 \\
\hline \(-1=>10\) at 3th step & passes in front of | Point \\
\hline \(-1=>10\) at 5th step & passes in front of | Vertex | 11 \\
\hline -1 \(=>11\) at 3th step & passes in front of | Point | 6 \\
\hline \(-1=>11\) at 5th step & visits | Vertex|11 \\
\hline \(5 \Rightarrow>-6\) at 4th step & visits | Point | 6 \\
\hline \(5=>-3\) at 4th step & passes in front of | Point \\
\hline \(5=>-2\) at 4th step & passes in front of | Point \\
\hline \(5=>-2\) at 6th step & passes in front of | Vertex | 11 \\
\hline \(5 \Rightarrow>10\) at 4th step & passes in front of | Point \\
\hline \(5 \Rightarrow>10\) at 6th step & passes in front of | Vertex | 11 \\
\hline \(5=>11\) at 4th step & passes in front of | Point | \\
\hline \(5=>11\) at 6 th step (16 rows) & visits | Vertex|11 \\
\hline
\end{tabular}
```

See Also
withPoints - Family of functions
withPoints - Category

- Sample Data


## Indices and tables

- Index
- Search Page

Supported versions: Latest (3.3) 3.23 .13 .0

- Unsupported versions: 2.6 2.5 2.4 $2.3 \mathbf{2 . 2}$


## pgr_withPointsCost - Proposed

pgr_withPointsCost - Calculates the shortest path and returns only the aggregate cost of the shortest path(s) found, for the combination of points given.

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.


## boost

Boost Graph Inside

## Availability

- Version 3.2.0
- New proposed function:

```
        - pgr_withPointsCost(Combinations)
```

- Version 2.2.0
- New proposed function


## Description

Modify the graph to include points defined by points_sql. Using Dijkstra algorithm, return only the aggregate cost of the shortest path(s) found.

The main characteristics are:

- It does not return a path.
- Returns the sum of the costs of the shortest path for pair combination of vertices in the modified graph.
- Vertices of the graph are:
- positive when it belongs to the edges_sql
- negative when it belongs to the points_sql
- Process is done only on edges with positive costs.
- Values are returned when there is a path.
- The returned values are in the form of a set of(start_vid, end_vid, agg_cost).
- When the starting vertex and ending vertex are the same, there is no path.
- The agg_cost in the non included values $(v, v)$ is 0
- When the starting vertex and ending vertex are the different and there is no path.
- The agg_cost in the non included values (u,v) is <br>(\infty<br>)
- If the values returned are stored in a table, the unique index would be the pair:(start_vid, end_vid).
- For undirected graphs, the results are symmetric.
- The agg_cost of $(u, v)$ is the same as for $(v, u)$.
- For optimization purposes, any duplicated value in thestart_vids or end_vids is ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending
- Running time: $\backslash(\mathrm{O}(\mid$ start $\backslash$ vids $\mid *(\mathrm{~V} \backslash \log \mathrm{~V}+\mathrm{E})) \backslash)$


## Signatures

## Summary

```
pgr_withPointsCost(Edges SQL, 'Points SQL`_, start vid, end vid, [options])
pgr_withPointsCost(Edges SQL, 'Points SQL`,, start vid, end vids, [options])
pgr_withPointsCost(Edges SQL, 'Points SQL`, start vids, end vid, [options])
pgr_withPointsCost(Edges SQL, 'Points SQL`_, start vids, end vids, [options])
pgr_withPointsCost(Edges SQL, 'Points SQL`_, Combinations SQL, [options])
options: [directed, driving_side]
RETURNS SET OF (start_pid, end_pid, agg_cost)
OR EMPTY SET
```


## Note

There is no details flag, unlike the other members of the withPoints family of functions.

## One to One

```
pgr_withPointsCost(Edges SQL, 'Points SQL`_, start vid, end vid, [options])
options: [directed, driving_side]
RETURNS SET OF (start_pid, end_pid, agg_cost)
OR EMPTY SET
```


## Example:

From point $\backslash(1 \backslash)$ to vertex $\backslash(10 \backslash)$ with defaults

```
SELECT * FROM pgr_withPointsCost(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge id, fraction, side from pointsOfInterest',
    -1, 10);
start_pid | end_pid | agg_cost
-1| 10| 5.6
(1 row)
```


## One to Many

```
pgr_withPointsCost(Edges SQL, Points SQL, start vid, end vids, [options])
options: [directed, driving_side]
RETURNS SET OF (start_pid, end_pid, agg_cost)
OR EMPTY SET
```


## Example:

From point $\backslash(1 \backslash)$ to point $\backslash(3 \backslash)$ and vertex $\backslash(7 \backslash)$ on an undirected graph

```
SELECT * FROM pgr_withPointsCost(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOfInterest',
    -1, ARRAY[-3, 7],
    directed => false)
start_pid | end_pid | agg_cost
    -1|
    -1| 7| 1.6
(2 rows)
```


## Many to One

```
pgr_withPointsCost(Edges SQL, Points SQL, start vids, end vid, [options])
options: [directed, driving_side]
RETURNS SET OF (start_pid, end_pid, agg_cost)
OR EMPTY SET
```


## Example:

From point $\backslash(1 \backslash)$ and vertex $\backslash(6 \backslash)$ to point $\backslash(3 \backslash)$

## Many to Many

```
pgr_withPointsCost(Edges SQL, Points SQL, start vids, end vids, [options])
options: [directed, driving_side]
RETURNS SET OF (start_pid, end_pid, agg_cost)
OR EMPTY SET
```


## Example:

From point $\backslash(15 \backslash)$ and vertex $\backslash(6 \backslash)$ to point $\backslash(3 \backslash)$ and vertex $\backslash(1 \backslash)$

```
SELECT * FROM pgr_withPointsCost(
    SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOfInterest',
    ARRAY[-1, 6], ARRAY[-3, 1])
start_pid | end_pid |agg_cost
    -1| -3| 3.2
    -1| 1| 3.6
    6| -3| 2.6
    6| 1| 3
(4 rows)
```


## Combinations

```
pgr_withPointsCost(Edges SQL, Points SQL, Combinations SQL, [options])
options: [directed, driving_side]
RETURNS SET OF (start_pid, end_pid, agg_cost)
OR EMPTY SET
```


## Example:

Two combinations
From point $\backslash(1 \backslash)$ to vertex $\backslash(10 \backslash)$, and from vertex $\backslash(6 \backslash)$ to point $\backslash(3 \backslash)$ with right side driving.

```
SELECT * FROM pgr_withPointsCost(
    SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOfInterest',
    'SELECT * FROM (VALUES (-1, 10), (6, -3)) AS combinations(source, target)',
    driving_side => 'r');
start_pid | end_pid | agg_cost
    -1|}10|\quad6.
    6| -3| 2.6
(2 rows)
```


## Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Points SQL | TEXT | Points SQL as described below |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | ARRAY[BIGINT] | Identifier of the starting vertex of the path. Negative value is for point's <br> identifier. |
| start vids | Array of identifiers of starting vertices. Negative values are for point's <br> identifiers. |  |
| end vid | Identifier of the ending vertex of the path. Negative value is for point's <br> identifier. |  |


| Column | Type | Description |
| :--- | :--- | :--- |
| end vids | ARRAY[BIGINT] | Array of identifiers of ending vertices. Negative values are for point's <br> identifiers. |

Optional parameters


With points optional parameters

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| driving_side | CHAR | b | Value in [r, l, b] indicating if the driving side is: |
|  |  |  | - r for right driving side. |
|  |  |  | - I for left driving side. |
|  |  |  | - b for both. |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Points SQL

| Parameter | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| pid | ANY-INTEGER | value | Identifier of the point. |

- Use with positive value, as internally will be converted to negative value
- If column is present, it can not be NULL.
- If column is not present, a sequential negativevalue will be given automatically.

| edge_id | ANY-INTEGER |  | Identifier of the "closest" edge to the point. |
| :---: | :---: | :---: | :---: |
| fraction | ANY-NUMERICAL |  | Value in $<0,1\rangle$ that indicates the relative postition from the first end point of the edge. |
| side | CHAR | b | Value in [b, r, l, NULL] indicating if the point is: |
|  |  |  | - In the right r , <br> - In the left I, <br> - In both sides b, NULL |

## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- | Identifier of the departure vertex. |
|  | INTEGER |  |
| target | ANY- | Identifier of the arrival vertex. |
|  | INTEGER |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Result Columns

| Column | Type | Description |
| :---: | :---: | :---: |
| start_pid | BIGINT | Identifier of the starting vertex or point. <br> - When positive: is a vertex's identifier. <br> - When negative: is a point's identifier. |
| end_pid | BIGINT | Identifier of the ending vertex or point. <br> - When positive: is a vertex's identifier. <br> - When negative: is a point's identifier. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

## Additional Examples

```
- Right side driving topology
- Left side driving topology
- Does not matter driving side driving topology
```

Right side driving topology

Traveling from point $\backslash(1 \backslash)$ and vertex $\backslash(5 \backslash)$ to points $\backslash(\backslash\{2,3,6 \backslash\} \backslash)$ and vertices $\backslash(\backslash\{10,11 \backslash\} \backslash)$

```
SELECT * FROM pgr_withPointsCost(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge id, fraction, side from pointsOfInterest',
    ARRAY[5, -1], ARRAY[-2, -3, -6, 10, 11],
    driving_side => 'r');
start_pid | end_pid | agg_cost
    -1| -6| 2.1
    | -3| 4
    | -2| 4.8
    10| 6.4
    11| }3.
    -6| 1.7
    -3| 3.6
    -2| 4.4
    10| 6
    5| 11| 3
(10 rows)
```

Left side driving topology

Traveling from point $\backslash(1 \backslash)$ and vertex $\backslash(5 \backslash)$ to points $\backslash(\backslash\{2,3,6 \backslash\} \backslash)$ and vertices $\backslash(\backslash\{10,11 \backslash\} \backslash)$

## SELECT * FROM pgr_withPointsCost(

'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id', 'SELECT pid, edge_id, fraction, side from pointsOfInterest',
ARRAY[5, -1], ARRAY[-2, -3, -6, 10, 11],
driving_side => 'l');
start_pid | end_pid | agg_cos

| $-1 \mid$ | $-6 \mid$ | 1.3 |
| ---: | ---: | ---: |
| $-1 \mid$ | $-3 \mid$ | 3.2 |
| $-1 \mid$ | $-2 \mid$ | 5.2 |
| $-1 \mid$ | $10 \mid$ | 5.6 |
| $-1 \mid$ | $11 \mid$ | 2.6 |
| $5 \mid$ | $-6 \mid$ | 1.7 |
| 5 | $-3 \mid$ | 3.6 |
| $5 \mid$ | $-2 \mid$ | 5.6 |
| $5 \mid$ | $10 \mid$ | 6 |
| 5 | $11 \mid$ | 3 |
| (10 rows) |  |  |

Does not matter driving side driving topology

Traveling from point $\backslash(1 \backslash)$ and vertex $\backslash(5 \backslash)$ to points $\backslash(\backslash\{2,3,6 \backslash\} \backslash)$ and vertices $\backslash(\backslash\{10,11 \backslash\} \backslash)$

```
SELECT * FROM pgr_withPointsCost(
    SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    SELECT pid, edge id, fraction, side from pointsOfInterest',
ARRAY[5, -1], ARRAY[-2, -3, -6, 10, 11]);
start_pid | end_pid | agg_cost
    -1| -6| 1.3
    -3| 3.2
    -2| 4
    10| 5.6
    11| 2.6
    -6| 1.7
    -3| 3.6
    -2| 4.4
    10| 6
    (10 rows)
```

The queries use the Sample Data network.

See Also

- withPoints - Family of functions


## Indices and tables

- Index
- Search Page


## - Supported versions: Latest (3.3) 3.23 .13 .0

- Unsupported versions: $2.6 \mathbf{2 . 5} \mathbf{2 . 4} \mathbf{2 . 3}$


## pgr_withPointsCostMatrix - proposed

pgr_withPointsCostMatrix - Calculates a cost matrix using pgr_withPoints - Proposed.

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.


## Availability

- Version 2.2.0
- New proposed function


## Description

Using Dijkstra algorithm, calculate and return a cost matrix.
Dijkstra's algorithm, conceived by Dutch computer scientist Edsger Dijkstra in 1956. It is a graph search algorithm that solves the shortest path problem for a graph with non-negative edge path costs, producing a shortest path from a starting vertex to an ending vertex. This implementation can be used with a directed graph and an undirected graph.

The main Characteristics are:

- Can be used as input topgr_TSP.
- Use directly when the resulting matrix is symmetric and there is nd(\infty<br>) value.
- It will be the users responsibility to make the matrix symmetric.
- By using geometric or harmonic average of the non symmetric values.
- By using max or min the non symmetric values.
- By setting the upper triangle to be the mirror image of the lower triangle.
- By setting the lower triangle to be the mirror image of the upper triangle.
- It is also the users responsibility to fix an<br>(\infty<br>) value.
- Each function works as part of the family it belongs to.
- It does not return a path.
- Returns the sum of the costs of the shortest path for pair combination of nodes in the graph.
- Process is done only on edges with positive costs.
- Values are returned when there is a path.
- When the starting vertex and ending vertex are the same, there is no path.
- The aggregate cost in the non included values $(v, v)$ is 0 .
- When the starting vertex and ending vertex are the different and there is no path. - The aggregate cost in the non included values(u,v) is <br>(\infty<br>).
- Let be the case the values returned are stored in a table:
- The unique index would be the pair: (start_vid, end_vid).
- Depending on the function and its parameters, the results can be symmetric.
- The aggregate cost of $(u, v)$ is the same as for $(v, u)$.
- Any duplicated value in the start vids are ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending


## Signatures

## Summary

```
pgr_withPointsCostMatrix(Edges SQL, Points SQL, start vids, [options])
options: [directed, driving_side]
RETURNS SET OF (start_vid, end_vid, agg_cost)
OR EMPTY SET
```


## Note

There is no details flag, unlike the other members of the withPoints family of functions.

## Example:

Cost matrix for points $\backslash(\backslash\{1,6 \backslash\} \backslash)$ and vertices $\backslash(\backslash\{10,11 \backslash\} \backslash)$ on an undirected graph

- Returning a symmetrical cost matrix
- Using the default side value on the points_sql query
- Using the default driving_side value

```
SELECT * FROM pgr_withPointsCostMatrix(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id'
    SELECT pid, edge_id, fraction from pointsOfInterest',
array[-1, 10, 11, -6], directed := false);
start vid | end vid | agg cost
|| -1 | 1.3
    |
    6| 11| 1.3
    1| -6 1.3
    10| 1.6
    11| 2.6
    -6| 1.7
    -1| 1.6
    -6| 13
    -6 1.3
    1| 10| 1
(12 rows)
```


## Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Points SQL | TEXT | Points SQL as described below |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting <br> vertices. |

Optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true | When true the graph is considered Directed <br>  | W hen false the graph is considered as  <br>  Undirected. |

With points optional parameters

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| driving_side | CHAR | b | Value in [r, l, b] indicating if the driving side is: |
|  |  |  | - r for right driving side. |
|  |  |  | - I for left driving side. |
|  |  |  | - b for both. |

Inner Queries

Edges SQL

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |  |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |

- When negative: edge (target, source) does not exist, therefore it's not part of the graph.


## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Points SQL
Parameter Type Default Description

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| pid | ANY-INTEGER | value | Identifier of the point. <br> - Use with positive value, as internally will be converted to negative value <br> - If column is present, it can not be NULL. <br> - If column is not present, a sequential negativevalue will be given automatically. |
| edge_id | ANY-INTEGER |  | Identifier of the "closest" edge to the point. |
| fraction | ANY-NUMERICAL |  | Value in $\langle 0,1\rangle$ that indicates the relative postition from the first end point of the edge. |
| side | CHAR | b | Value in $[b, r, l$, NULL $]$ indicating if the point is: <br> - In the right $r$, <br> - In the left I, <br> - In both sides b, NULL |

## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

## Note

When start_vid or end_vid columns have negative values, the identifier is for a Point.

## Additional Examples

```
- Use with pgr_TSP.
```

Use with pgr_TSP.

```
SELECT * FROM pgr_TSP(
    $$
SELECT * FROM pgr_withPointsCostMatrix(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction from pointsOflnterest',
    array[-1, 10, 11, -6], directed := false);
$$
);
NOTICE: pgr_TSP no longer solving with simulated annaeling
HINT: Ignoring annaeling parameters
seq | node | cost | agg_cost
    1| -6| 0| 0
    2| -1| 1.3| 1.3
    3| 10| 1.6| 2.9
    4| 11| 1| 3.9
    5| -6| 1.3| 5.2
(5 rows)
```

See Also

- withPoints - Family of functions
- Cost Matrix - Category
- Traveling Sales Person - Family of functions
- Sample Data


## Indices and tables

Index

- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: $2.6 \mathbf{2 . 5} 2.4 \mathbf{2 . 3} \mathbf{2 . 2}$
pgr_withPointsKSP - Proposed
pgr_withPointsKSP - Yen's algorithm for K shortest paths using Dijkstra.


## Warning

Proposed functions for next mayor release.
They are not officially in the current release.

- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.

Boost Graph Inside

## Availability

- Version 2.2.0
- New proposed function


## Description

Modifies the graph to include the points defined in the Points SQL and using Yen algorithm, finds the $\backslash(\mathrm{K} \backslash)$ shortest paths.

## Signatures

```
pgr_withPointsKSP(Edges SQL, Points SQL start vid, end vid, K, [options])
options: [directed, heap_paths, driving_side, details]
RETURNS SET OF (seq, path_id, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

Get 2 paths from Point $\backslash(1 \backslash)$ to point $\backslash(2 \backslash)$ on a directed graph.

- For a directed graph.
- The driving side is set asb both. So arriving/departing to/from the point(s) can be in any direction.
- No details are given about distance of other points of the query.
- No heap paths are returned.


## SELECT * FROM pgr_withPointsKSP(

'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
'SELECT pid, edge_id, fraction, side from pointsOflnterest',
-1, -2, 2);
seq | path_id | path_seq | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $1 \mid$ | $-1 \mid$ | $1 \mid$ | $0.6 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $1 \mid$ | $2 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0.6 |
| $3 \mid$ | $1 \mid$ | $3 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 1.6 |
| $4 \mid$ | $1 \mid$ | $4 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ | 2.6 |
| $5 \mid$ | $1 \mid$ | $5 \mid$ | $16 \mid$ | $15 \mid$ | $0.4 \mid$ | 3.6 |
| $6 \mid$ | $1 \mid$ | $6 \mid$ | $-2 \mid$ | $-1 \mid$ | $0 \mid$ | 4 |
| $7 \mid$ | $2 \mid$ | $1 \mid$ | $-1 \mid$ | $1 \mid$ | $0.6 \mid$ | 0 |
| $8 \mid$ | $2 \mid$ | $2 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ | 0.6 |
| $9 \mid$ | $2 \mid$ | $3 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ | 1.6 |
| $10 \mid$ | $2 \mid$ | $4 \mid$ | $11 \mid$ | $11 \mid$ | $1 \mid$ | 2.6 |
| $11 \mid$ | $2 \mid$ | $5 \mid$ | $12 \mid$ | $13 \mid$ | $1 \mid$ | 3.6 |
| $12 \mid$ | $2 \mid$ | $6 \mid$ | $17 \mid$ | $15 \mid$ | $0.6 \mid$ | 4.6 |
| $13 \mid$ | $2 \mid$ | $7 \mid$ | $-2 \mid$ | $-1 \mid$ | $0 \mid$ | 5.2 |

## Parameter

| Column | Type | Description |  |
| :--- | :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL query as described. |  |
| Points SQL | TEXT | Points SQL query as described. |  |
| start vid | ANY-INTEGER | Identifier of the departure vertex. |  |
| end vid | ANY-INTEGER | Identifier of the destination vertex. <br> point |  |
| K |  | ANY-INTEGER | Number of required paths |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

Optional parameters


KSP Optional parameters


With points optional parameters


Inner Queries

Edges SQL

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |  |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |  |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |  |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |
|  |  |  | When negative: edge (target, source) does not exist, therefore it's <br> not part of the graph. |
|  |  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Points SQL

| Parameter | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| pid | ANY-INTEGER | value | Identifier of the point. |

- Use with positive value, as internally will be converted to negative value
- If column is present, it can not be NULL.
- If column is not present, a sequential negativevalue will be given automatically.

| edge_id | ANY-INTEGER | Identifier of the "closest" edge to the point. |
| :--- | :--- | :--- |
| fraction | ANY-NUMERICAL | Value in $<0,1>$ that indicates the relative postition from the first end <br> point of the edge. |
| side | CHAR | Value in $[b, r, l$, NULL $]$ indicating if the point is: |
|  |  |  |
|  |  | In the right $r$, |
|  |  | In the left, |
|  |  | In both sides $b$, NULL |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Result Columns

Returns set of (seq, path_id, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Sequential value starting from $\mathbf{1 .}$ |
| path_id | INTEGER | Path identifier. |
|  |  | - Has value $\mathbf{1}$ for the first of a path fromstart vid to end_vid |
| path_seq | INTEGER | Relative position in the path. Has value $\mathbf{1}$ for the beginning of a path. |
| node | BIGINT | Identifier of the node in the path fromstart vid to end vid |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence. $\mathbf{- 1}$ for the last node of <br> the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence.  <br>   <br> agg_cost FLOAT |

## Additional Examples

```
- Left driving side
- Right driving side
```

Get $\backslash(2 \backslash)$ paths using left side driving topology, from point $\backslash(1 \backslash)$ to point $\backslash(2 \backslash)$ with details.

```
SELECT * FROM pgr_withPointsKSP(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOflnterest',
    -1,-2, 2,
    driving side := '\', details := true);
seq | path_id | path_seq | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|}
\hline 1| & 1| & 1| & -1| & 1| 0.6 | & 0 \\
\hline 21 & 1| & 21 & \(6 \mid\) & 4|0.7| & 0.6 \\
\hline 31 & 1| & 31 & -6| & 4|0.3| & 1.3 \\
\hline 4| & \(1 \mid\) & 4 | & 7| & 8| 1| & 1.6 \\
\hline 51 & 1 | & 51 & 11| & 11| 1| & 2.6 \\
\hline 61 & 1 | & 61 & 12| & 13| 1| & 3.6 \\
\hline 71 & 1| & 71 & 17| & 15|0.6| & 4.6 \\
\hline 81 & 1 | & 8। & -2| & -1| 0 | & 5.2 \\
\hline 91 & 21 & 1| & -1| & 1| 0.6 | & 0 \\
\hline 10| & 21 & 21 & 61 & 4|0.7| & 0.6 \\
\hline 11| & 21 & 31 & -6| & 4|0.3| & 1.3 \\
\hline \(12 \mid\) & 21 & 4 | & 7| & 8| 1| & 1.6 \\
\hline 13 | & 21 & 51 & 11| & 9| 1| & 2.6 \\
\hline 14| & 21 & 61 & 16| & 15| 1| & 3.6 \\
\hline 15 | & 21 & 71 & 17| & 15|0.6| & 4.6 \\
\hline \(16 \mid\) & 2 | & 81 & \(-2 \mid\) & -1| 0 | & 5.2 \\
\hline
\end{tabular}
```


## ight driving side

Get $\backslash(2 \backslash)$ paths using right side driving topology from, point $\backslash(1 \backslash)$ to point $\backslash(2 \backslash)$ with heap paths and details.

```
SELECT * FROM pgr_withPointsKSP(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOfInterest',
    -1, -2, 2,
    heap paths := true, driving side := 'r', details := true);
seq | path_id | path_seq | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline 1| & 1 | & 1| & -1| & 1| & 0.4 & 0 \\
\hline 21 & 1 | & 21 & 51 & 1| & 1| & 0.4 \\
\hline 31 & 1 | & 31 & 6| & & 0.71 & 1.4 \\
\hline 4 | & 1| & 4| & -6| & \(4 \mid\) & 0.31 & 2.1 \\
\hline 51 & \(1 \mid\) & 51 & 71 & 8| & 1| & 2.4 \\
\hline 61 & \(1 \mid\) & 61 & 11| & 91 & 1| & 3.4 \\
\hline 71 & 1 | & 7| & 16| & 15| & 0.4 & 4.4 \\
\hline 8| & 1| & 8| & -2| & -1| & 01 & 4.8 \\
\hline 91 & 21 & 1| & -1| & 1| & 0.4 & 0 \\
\hline \(10 \mid\) & 21 & 21 & 5| & 1 | & 1| & 0.4 \\
\hline 11| & 21 & 31 & 61 & 4| & 0.7| & 1.4 \\
\hline \(12 \mid\) & 21 & 41 & -6| & 4 | & 0.31 & 2.1 \\
\hline 131 & 21 & 51 & 71 & 8। & 1| & 2.4 \\
\hline \(14 \mid\) & 21 & 61 & 11| & 11 & | 1| & 3.4 \\
\hline 151 & 21 & 71 & 12| & 13 & | 1| & 4.4 \\
\hline \(16 \mid\) & 21 & 81 & 17| & 15 & | 1 | & 5.4 \\
\hline \(17 \mid\) & 21 & 91 & \(16 \mid\) & 15 & | 0.4 | & 6.4 \\
\hline 18| & 21 & \(10 \mid\) & -2| & -1| & | 0| & 6.8 \\
\hline 19| & 31 & 1 | & -1| & \(1 \mid\) & 0.4| & 0 \\
\hline \(20 \mid\) & 31 & 21 & 5। & 1| & 1| & 0.4 \\
\hline 21| & 31 & 31 & 6| & 4| & 0.7| & 1.4 \\
\hline \(22 \mid\) & 31 & 4 & -6| & \(4 \mid\) & 0.31 & 2.1 \\
\hline 231 & 31 & 51 & 7| & 10| & 1| & 2.4 \\
\hline \(24 \mid\) & 31 & 61 & 8। & 12| & 0.6 & 3.4 \\
\hline 251 & 31 & 71 & -3| & 12| & | 0.4 | & 4 \\
\hline \(26 \mid\) & 31 & 81 & 12| & 13 & | 1 | & 4.4 \\
\hline 271 & 31 & 91 & 17| & 15 & | 1| & 5.4 \\
\hline \(28 \mid\) & 31 & \(10 \mid\) & \(16 \mid\) & 15 & | 0.4 & 6.4 \\
\hline 291 & 31 & \(11 \mid\) & \(-2 \mid\) & -1| & | 0| & 6.8 \\
\hline
\end{tabular}
(29 rows)
```

The queries use the Sample Data network.

## See Also

- withPoints - Family of functions
- K shortest paths - Category
- Sample Data


## Indices and tables

- Index
- Search Page

Supported versions: Latest (3.3) 3.23 .13 .0

- Unsupported versions: 2.6 2.5 2.4 $2.3 \mathbf{2 . 2}$


## pgr_withPointsDD - Proposed

pgr_withPointsDD - Returns the driving distance from a starting point.

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.

Boost Graph Inside

## Availability

- Version 2.2.0
- New proposed function


## Description

Modify the graph to include points and using Dijkstra algorithm, extracts all the nodes and points that have costs less than or equal to the value **distance** from the starting point. The edges extracted will conform the corresponding spanning tree.

## Signatures

```
pgr_withPointsDD(Edges SQL, Points SQL, root vid, distance, [options A])
pgr_withPointsDD(Edges SQL, Points SQL, root vids, distance, [options B])
options A: [directed, driving_side, details]
options B: [directed, driving_side, details, equicost]
RETURNS SET OF (seq, [start_vid], node, edge, cost, agg_cost)
OR EMPTY SET
```


## Single vertex

```
pgr_withPointsDD(Edges SQL, Points SQL, root vid, distance, [options])
options: [directed, driving_side, details]
RETURNS SET OF (seq, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

Right side driving topology, from point $\backslash(1 \backslash)$ within a distance of $\backslash(3.3 \backslash)$ with details.

```
SELECT * FROM pgr_withPointsDD
    SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOflnterest',
    -1,3.3,
    driving_side => 'r',
    details => true);
seq | node | edge | cost | agg_cost
    1| -1 | -1| 0| 0
    2| 5| 1| 0.4| 0.4
    3| 6| 1| 1| 1.4
    4 - 6- 4 | 0.7 | 2.1
    5| 7| 4| 0.3| 2.4
(5 rows)
```

Multiple vertices

```
pgr_withPointsDD(Edges SQL, Points SQL, root vids, distance, [options])
options: [directed, driving_side, details, equicost]
RETURNS SET OF (seq, start_vid, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From point $\backslash(1 \backslash)$ and vertex $\backslash(16 \backslash)$ within a distance of $\backslash(3.3 \backslash)$ with equicost on a directed graph

```
SELECT * FROM pgr withPointsDD(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOfInterest',
    ARRAY[-1, 16], 3.3
    driving side => ' '',
    equicost => true);
seq | start_vid | node | edge | cost | agg_cost
    -1|-1| -1| 0| 0
    1| 6| 1| 0.6| 0.6
    1| 7| 4| 1| 1.6
    l|
    | 3| 7| 1| 2.6
    1| 8| 10| 1| 2.6
        16| 16| -1| 0| 0
        6| 11 9| 1| 1
        6| 15| 16| 1| 1
        16| 17| 15| 1| 1
        16| 10| 3| 1| 2
        16| 12| 11| 1| 2
(12 rows)
```


## Parameter

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Points SQL | TEXT | Points SQL as described below |
| Root vid | BIGINT | Identifier of the root vertex of the tree. |
|  |  | Negative values represent a point |

## Optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true | When true the graph is considered Directed <br>  |  Whe n false the graph is considered as <br>  Undirected. |

## With points optional parameters



Driving distance optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| equicost | BOOLEAN | true | When true the node will only appear in the closestfrom_v list. |
|  |  |  | When false which resembles several calls using the single starting point <br> signatures. Tie brakes are arbitrary. |
|  |  |  |  |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Points SQL

| Parameter | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| pid | ANY-INTEGER | value | Identifier of the point. |

- Use with positive value, as internally will be converted to negative value
- If column is present, it can not be NULL.
- If column is not present, a sequential negativevalue will be given automatically.

| edge_id | ANY-INTEGER |  | Identifier of the "closest" edge to the point. |
| :---: | :---: | :---: | :---: |
| fraction | ANY-NUMERICAL |  | Value in $<0,1\rangle$ that indicates the relative postition from the first end point of the edge. |
| side | CHAR | b | Value in [b, r, l, NULL] indicating if the point is: |
|  |  |  | - In the right $r$, <br> - In the left I, <br> - In both sides b, NULL |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

RETURNS SET OF (seq, [start_vid], node, edge, cost, agg_cost)

| Parameter | Type | Description |
| :--- | :--- | :--- |
| seq | BIGINT | Sequential value starting from |
| (1 |  |  |
| ). |  |  |
| [start_vid] | BIGINT | Identifier of the root vertex. <br> node |
| BIGINT | Identifier of node within the limits from <br> from_v. |  |
| edge | BIGINT | Identifier of the edge used to arrive to node. |
| Fost | FLOAT | Cost to traverse edge. |
| agg_cost | FLOAT | Aggregate cost from from_v to node. |

## Where:

ANY-INTEGER:<br>SMALLINT, INTEGER, BIGINT<br>ANY-NUMERIC:<br>SMALLINT, INTEGER, BIGINT, REAL, FLOAT, NUMERIC

Additional Examples

## - Driving side does not matter

## Driving side does not matter

From point $\backslash(1 \backslash)$ within a distance of $\backslash(3.3 \backslash)$, does not matter driving side, with details.

```
SELECT * FROM pgr_withPointsDD(
    'SELECT id, source, target, cost, reverse_cost FROM edges ORDER BY id',
    'SELECT pid, edge_id, fraction, side from pointsOfInterest',
    -1,3.3,
    driving_side => 'b',
    details => true);
seq | node | edge | cost | agg_cost
    1| -1| -1| 0| 0
    2| 5| 1| 0.4| 0.4
    3| 6| 1| 0.6| 0.6
    4| -6| 4| 0.7)}1.
    5| 7| 4| 0.3| 1.6
    6| 3| 7| 1| 2.6
    7| 8| 10| 1| 2.6
    8| 11| 8| 1| 2.6
    9| -3| 12| 0.6| 3.2
    10| -4| 6| 0.7 | 3.3
(10 rows)
```


## See Also

## - pgr_drivingDistance

- pgr_alphaShape
- Sample Data


## Indices and tables

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- Search Page


## Introduction

This family of functions belongs to thewithPoints - Category and the functions that compose them are based one way or another on dijkstra algorithm.

## Depending on the name:

- pgr_withPoints is pgr_dijkstra with points
- pgr_withPointsCost is pgr_dijkstraCost with points
- pgr_withPointsCostMatrix is pgr_dijkstraCostMatrix with points
- pgr_withPointsKSP is pgr_ksp with points
- pgr_withPointsDD is pgr_drivingDistance with points

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Points SQL | TEXT | Points SQL as described below |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | ARRAY[BIGINT] | Identifier of the starting vertex of the path. Negative value is for point's <br> identifier. |
| start vids | Array of identifiers of starting vertices. Negative values are for point's <br> identifiers. |  |
| end vid | ARRAY[BIGINT] | Identifier of the ending vertex of the path. Negative value is for point's <br> identifier. |
| end vids | Array of identifiers of ending vertices. Negative values are for point's <br> identifiers. |  |

Optional parameters

| Column | Type | Default | Description |  |
| :---: | :---: | :---: | :---: | :---: |
| directed | BOOLEAN | true |  | When true the graph is considered Directed |
|  |  |  |  | Whenfalse the graph is considered as Undirected. |

With points optional parameters

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| driving_side | CHAR | b | Value in [r, l, b] indicating if the driving side is: |
|  |  |  | - $\quad$ for right driving side. <br> - I for left driving side. <br> - b for both. |
| details | BOOLEAN | false | - When true the results will include the points that are in the path. <br> - When false the results will not include the points that are in the path. |

## Inner Queries

Edges SQL

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |  |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |  |
| target | ANY-INTEGER |  | Identifier of the second end point vertex of the edge. |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |

- When negative: edge (target, source) does not exist, therefore it's not part of the graph.

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Points SQL

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| pid | ANY-INTEGER | value | Identifier of the point. |
|  |  |  | - Use with positive value, as internally will be converted to negative value <br> - If column is present, it can not be NULL. <br> - If column is not present, a sequential negativevalue will be given automatically. |


| edge_id | ANY-INTEGER | Identifier of the "closest" edge to the point. |
| :--- | :--- | :--- |
| fraction | ANY-NUMERICAL | Value in $<0,1>$ that indicates the relative postition from the first end <br> point of the edge. |


| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| side | CHAR | b | Value in [b, r, l, NULL] indicating if the point is: |
|  |  |  | - In the right $r$, <br> - In the left I, <br> - In both sides b, NULL |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Combinations SQL

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- | Identifier of the departure vertex. |
|  | INTEGER |  |
| target | ANY- | Identifier of the arrival vertex. |
|  | INTEGER |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Advanced Documentation

## Contents

- About points
- Driving side
- Right driving side
- Left driving side
- Driving side does not matter
- Creating temporary vertices
- On a right hand side driving network
- On a left hand side driving network
- When driving side does not matter


## About points

For this section the following city (seeSample Data) some interesing points such as restaurant, supermarket, post office, etc. will be used as example.


- The graph is directed
- Red arrows show the (source, target) of the edge on the edge table
- Blue arrows show the (target, source) of the edge on the edge table
- Each point location shows where it is located with relation of the edge(source, target)
- On the right for points $\mathbf{2}$ and $\mathbf{4}$.
- On the left for points 1, $\mathbf{3}$ and $\mathbf{5}$.
- On both sides for point 6 .

The representation on the data base follows thePoints SQL description, and for this example:

```
SELECT pid, edge_id, fraction, side FROM pointsOfInterest
pid | edge_id | fraction | side
\begin{tabular}{cc}
\(1 \mid\) & \(0.4|\mid\) \\
\(15 \mid\) & \(0.4 \mid r\) \\
\(12 \mid\) & \(0.6|\mid\) \\
\(6 \mid\) & \(0.3 \mid r\) \\
\(5 \mid\) & \(0.8|\mid\) \\
\(4 \mid\) & \(0.7 \mid b\)
\end{tabular}
(6 rows)
```


## Driving side

## In the the folowwing images:

- The squared vertices are the temporary vertices,
- The temporary vertices are added according to the driving side,
- visually showing the differences on how depending on the driving side the data is interpreted.

Right driving side


- Point 1 located on edge $(6,5)$
- Point $\mathbf{2}$ located on edge $(16,17)$
- Point $\mathbf{3}$ located on edge $(8,12)$
- Point $\mathbf{4}$ located on edge $(1,3)$
- Point 5 located on edge $(10,11)$
- Point 6 located on edges $(6,7)$ and $(7,6)$

Left driving side


- Point $\mathbf{1}$ located on edge $(5,6)$
- Point $\mathbf{2}$ located on edge $(17,16)$
- Point $\mathbf{3}$ located on edge $(8,12)$
- Point $\mathbf{4}$ located on edge $(3,1)$
- Point 5 located on edge $(10,11)$
- Point 6 located on edges $(6,7)$ and $(7,6)$


## Driving side does not matter

- Like having all points to be considered in both sidesb
- Prefered usage on undirected graphs

- Point 1 located on edge $(5,6)$ and $(6,5)$
- Point 2 located on edge $(17,16)$ " and " 16,17
- Point $\mathbf{3}$ located on edge $(8,12)$
- Point $\mathbf{4}$ located on edge $(3,1)$ and $(1,3)$
- Point 5 located on edge $(10,11)$
- Point 6 located on edges $(6,7)$ and $(7,6)$

Creating temporary vertices
This section will demonstrate how a temporary vertex is created internally on the graph.

## Problem

## insert point:

```
SELECT pid, edge_id, fraction, side
FROM pointsOfInterest WHERE pid = 2;
pid | edge_id | fraction | side
    2| 15| 0.4|r
(1 row)
```

On a right hand side driving network

## Right driving side



- Arrival to point-2 can be achived only via vertex $\mathbf{1 6}$.
- Does not affects edge $(17,16)$, therefore the edge is kept.
- It only affects the edge $(16,17)$, therefore the edge is removed.
- Create two new edges:
- Edge (16, -2) with cost 0.4 (original cost * fraction $==\backslash(1 * 0.4 \backslash))$
- Edge $(-2,17)$ with cost 0.6 (the remaing cost)
- The total cost of the additional edges is equal to the original cost.
- If more points are on the same edge, the process is repeated recursevly.

On a left hand side driving network

## Left driving side

- Arrival to point-2 can be achived only via vertex 17.
- Does not affects edge $(16,17)$, therefore the edge is kept.
- It only affects the edge $(17,16)$, therefore the edge is removed.
- Create two new edges:
- Work with the original edge $(16,17)$ as the fraction is a fraction of the original:
- Edge $(16,-2)$ with cost 0.4 (original cost * fraction $==\backslash(1 * 0.4 \backslash))$
- Edge $(-2,17)$ with cost 0.6 (the remaing cost)
- If more points are on the same edge, the process is repeated recursevly.
- Flip the Edges and add them to the graph:
- Edge $(17,-2)$ becomes $(-2,16)$ with cost 0.4 and is added to the graph.
- Edge $(-2,16)$ becomes $(17,-2)$ with cost 0.6 and is added to the graph.
- The total cost of the additional edges is equal to the original cost.

When driving side does not matter


- Arrival to point -2 can be achived via vertices $\mathbf{1 6}$ or $\mathbf{1 7}$.
- Affects the edges $(16,17)$ and $(17,16)$, therefore the edges are removed.
- Create four new edges:
- Work with the original edge $(16,17)$ as the fraction is a fraction of the original:
- Edge ( $16,-2$ ) with cost 0.4 (original cost * fraction $==\backslash(1 * 0.4 \backslash))$
- Edge $(-2,17)$ with cost 0.6 (the remaing cost)
- If more points are on the same edge, the process is repeated recursevly.
- Flip the Edges and add all the edges to the graph:
- Edge $(16,-2)$ is added to the graph.
- Edge $(-2,17)$ is added to the graph.
- Edge $(16,-2)$ becomes $(-2,16)$ with cost 0.4 and is added to the graph.


## See Also

- withPoints - Category


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3)

Via - Category
proposed

## Warning

Proposed functions for next mayor release.

- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.


## - pgr_dijkstraVia - Proposed

## General Information

This category intends to solve the general problem:
Given a graph and a list of vertices, find the shortest path between $\backslash($ vertex_i $)$ and $\backslash($ vertex_ $\{i+1\} \backslash)$ for all vertices In other words, find a continuos route that visits all the vertices in the order given.

## path:

represents a section of a route.
route:
is a sequence of paths

## Parameters

## Used on:

- pgr_dijkstraVia - Proposed

| Parameter | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| Edges SQL | TEXT | SQL query as described. |  |
| via vertices | ARRAY [ ANY-INTEGER ] | Array of ordered vertices identifiers that are going to be <br> visited. |  |

## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
Besides the compulsory parameters each function has, there are optional parameters that exist due to the kind of function.

Via optional parameters

## Used on all Via functions

Parameter Type Default Description

| Parameter | Type | Default | Description |  |
| :--- | :--- | :--- | :--- | :--- |
| strict | BOOLEAN | false | $\bullet$ | When true if a path is missing stops and returnsEMPTY SET |
|  |  |  |  |  |
|  |  |  |  |  |
| U_turn_on_edge | BOOLEAN | true | $\bullet$ | When true departing from a visited vertex will not try to avoid |

## Inner Queries

Depending on the function one or more inner queries are needed.

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |
|  |  |  |
|  |  |  |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Restrictions SQL

| Column | Type | Description |
| :--- | :--- | :--- |
| path | ARRAY [ANY-INTEGER] | Sequence of edge identifiers that form a path that is not allowed to be taken. - <br>  <br>  <br>  <br> Empty arrays or NULL arrays are ignored. - Arrays that have aNULL element will <br> raise an exception. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Sequential value starting from $\mathbf{1}$. |
| path_id | INTEGER | Identifier of a path. Has value $\mathbf{1}$ for the first path. |
| path_seq | INTEGER | Relative position in the path. Has value $\mathbf{1}$ for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex of the path. |
| end_vid | BIGINT | Identifier of the ending vertex of the path. |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path <br>  <br>  |
|  |  | sequence. |
|  |  | -1 for the last node of the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | Aggregate cost from start_vid to node. |  |
| route_agg_cost | FLOAT | Total cost from start_vid of seq=1 to end_vid of the current seq. |

## Note

When start_vid, end_vid and node columns have negative values, the identifier is for a Point.

## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3)
withPoints - Category
When points are added to the graph.


## Warning

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- They are not officially in the current release.
- They will likely officially be part of the next mayor release:
- The functions make use of ANY-INTEGER and ANY-NUMERICAL
- Name might not change. (But still can)
- Signature might not change. (But still can)
- Functionality might not change. (But still can)
- pgTap tests have being done. But might need more.
- Documentation might need refinement.
- withPoints - Family of functions - Functions based on Dijkstra algorithm.


## Introduction

The with points category modifies the graph on the fly by adding points on edges as required by thePoints SQL query.
The functions within this category give the ability to process between arbitrary points located outside the original graph.
This category of functions was thought for routing vehicles, but might as well work for some other application not involving vehicles.

When given a point identifier pid that its being mapped to an edge with an identifieredge_id, with a fraction from the source to the target along the edge fraction and some additional information about which side of the edge the point is orside, then processing from arbitrary points can be done on fixed networks.

All this functions consider as many traits from the "real world" as possible:

- Kind of graph:
directed graph
undirected graph
- Arriving at the point:
- Compulsory arrival on the side of the segment where the point is located.
- On either side of the segment.
- Countries with:
- Right side driving
- Left side driving
- Some points are:
- Permanent: for example the set of points of clients stored in a table in the data base.
- The graph has been modified to permanently have those points as vertices.
- There is a table on the database that describes the points
- Temporal: for example points given through a web application
- The numbering of the points are handled with negative sign.
- This sign change is to avoid confusion when there is a vertex with the same identifier as the point identifier.
- Original point identifiers are to be positive.
- Transformation to negative is done internally.
- Interpretation of the sign on the node information of the output
- positive sign is a vertex of the original graph
- negative sign is a point of the Points SQL


## Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Points SQL | TEXT | Points SQL as described below |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | BIGINT | Identifier of the starting vertex of the path. Negative value is for point's <br> identifier. |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting vertices. Negative values are for point's <br> identifiers. |
| end vid | BIGINT | Identifier of the ending vertex of the path. Negative value is for point's <br> identifier. |
| end vids | Array of identifiers of ending vertices. Negative values are for point's <br> identifiers. |  |

Optional parameters

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| driving_side | CHAR | r | Value in $[r, I]$ indicating if the driving side is: <br> - $\quad$ for right driving side <br> - \| for left driving side <br> - Any other value will be considered asr |
| details | BOOLEAN | false | - When true the results will include the points that are in the path. <br> - When false the results will not include the points that are in the path. |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Points SQL

| Parameter | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| pid | ANY-INTEGER | value | Identifier of the point. Use with positive value, as internally will be converted to negative value <br> - If column is present, it can not be NULL. <br> - If column is not present, a sequential negativevalue will be given automatically. |
| edge_id | ANY-INTEGER |  | Identifier of the "closest" edge to the point. |
| fraction | ANY-NUMERICAL |  | Value in $<0,1\rangle$ that indicates the relative postition from the first end point of the edge. |
| side | CHAR | b | Value in $[b, r, l$, NULL] indicating if the point is: <br> - In the right $r$, <br> - In the left I, <br> - In both sides b, NULL |

## Where:

## ANY-INTEGER:

## SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Combinations SQL

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- <br> INTEGER | Identifier of the departure vertex. |
| target | ANY- <br> INTEGER | Identifier of the arrival vertex. |

Where:
ANY-INTEGER:
SMALLINT, INTEGER, BIGINT

## Advanced documentation

## Contents

- About points
- Driving side
- Right driving side
- Left driving side
- Driving side does not matter
- Creating temporary vertices
- On a right hand side driving network
- On a left hand side driving network
- When driving side does not matter

About points

For this section the following city (seeSample Data) some interesing points such as restaurant, supermarket, post office, etc. will be used as example.


## - The graph is directed

- Red arrows show the (source, target) of the edge on the edge table
- Blue arrows show the (target, source) of the edge on the edge table
- Each point location shows where it is located with relation of the edge(source, target)
- On the right for points 2 and 4.
- On the left for points 1, $\mathbf{3}$ and $\mathbf{5}$.
- On both sides for point 6 .

The representation on the data base follows thePoints SQL description, and for this example:

| $1 \mid$ | $1 \mid$ | $0.4 \mid \mathrm{l}$ |
| ---: | ---: | ---: |
| $2 \mid$ | $15 \mid$ | $0.4 \mid \mathrm{r}$ |
| $3 \mid$ | $12 \mid$ | $0.6 \mid \mathrm{l}$ |
| $4 \mid$ | $6 \mid$ | $0.3 \mid \mathrm{r}$ |
| $5 \mid$ | $5 \mid$ | $0.8 \mid \mathrm{l}$ |
| $6 \mid$ | $4 \mid$ | $0.7 \mid \mathrm{b}$ |
| (6 rows) |  |  |

Driving side

In the the folowwing images:

- The squared vertices are the temporary vertices,
- The temporary vertices are added according to the driving side,
- visually showing the differences on how depending on the driving side the data is interpreted.

Right driving side


- Point 1 located on edge $(6,5)$
- Point $\mathbf{2}$ located on edge $(16,17)$
- Point $\mathbf{3}$ located on edge $(8,12)$
- Point $\mathbf{4}$ located on edge $(1,3)$
- Point 5 located on edge $(10,11)$
- Point 6 located on edges $(6,7)$ and $(7,6)$

- Point $\mathbf{1}$ located on edge $(5,6)$
- Point $\mathbf{2}$ located on edge $(17,16)$
- Point $\mathbf{3}$ located on edge $(8,12)$
- Point $\mathbf{4}$ located on edge $(3,1)$
- Point 5 located on edge $(10,11)$
- Point 6 located on edges $(6,7)$ and $(7,6)$


## Driving side does not matter

- Like having all points to be considered in both sidesb
- Prefered usage on undirected graphs

- Point 1 located on edge $(5,6)$ and $(6,5)$
- Point 2 located on edge $(17,16)$ " and " 16,17
- Point $\mathbf{3}$ located on edge $(8,12)$
- Point $\mathbf{4}$ located on edge $(3,1)$ and $(1,3)$
- Point 5 located on edge $(10,11)$
- Point 6 located on edges $(6,7)$ and $(7,6)$

Creating temporary vertices
This section will demonstrate how a temporary vertex is created internally on the graph.

## Problem

## insert point:

```
SELECT pid, edge_id, fraction, side
FROM pointsOfInterest WHERE pid = 2;
pid | edge_id | fraction | side
    2| 15| 0.4|r
(1 row)
```

On a right hand side driving network

## Right driving side



- Arrival to point-2 can be achived only via vertex $\mathbf{1 6}$.
- Does not affects edge $(17,16)$, therefore the edge is kept.
- It only affects the edge $(16,17)$, therefore the edge is removed.
- Create two new edges:
- Edge (16, -2) with cost 0.4 (original cost * fraction $==\backslash(1 * 0.4 \backslash))$
- Edge $(-2,17)$ with cost 0.6 (the remaing cost)
- The total cost of the additional edges is equal to the original cost.
- If more points are on the same edge, the process is repeated recursevly.

On a left hand side driving network

## Left driving side

- Arrival to point-2 can be achived only via vertex 17.
- Does not affects edge $(16,17)$, therefore the edge is kept.
- It only affects the edge $(17,16)$, therefore the edge is removed.
- Create two new edges:
- Work with the original edge $(16,17)$ as the fraction is a fraction of the original:
- Edge $(16,-2)$ with cost 0.4 (original cost * fraction $==\backslash(1 * 0.4 \backslash))$
- Edge $(-2,17)$ with cost 0.6 (the remaing cost)
- If more points are on the same edge, the process is repeated recursevly.
- Flip the Edges and add them to the graph:
- Edge $(17,-2)$ becomes $(-2,16)$ with cost 0.4 and is added to the graph.
- Edge $(-2,16)$ becomes $(17,-2)$ with cost 0.6 and is added to the graph.
- The total cost of the additional edges is equal to the original cost.

When driving side does not matter


- Arrival to point -2 can be achived via vertices $\mathbf{1 6}$ or $\mathbf{1 7}$.
- Affects the edges $(16,17)$ and $(17,16)$, therefore the edges are removed.
- Create four new edges:
- Work with the original edge $(16,17)$ as the fraction is a fraction of the original:
- Edge ( $16,-2$ ) with cost 0.4 (original cost * fraction $==\backslash(1 * 0.4 \backslash))$
- Edge $(-2,17)$ with cost 0.6 (the remaing cost)
- If more points are on the same edge, the process is repeated recursevly.
- Flip the Edges and add all the edges to the graph:
- Edge $(16,-2)$ is added to the graph.
- Edge $(-2,17)$ is added to the graph.
- Edge $(16,-2)$ becomes $(-2,16)$ with cost 0.4 and is added to the graph.


## See Also

- withPoints - Family of functions


## Indices and tables

- Index
- Search Page


## See Also

- Experimental Functions


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: $2.6 \mathbf{2}$.5 2.4 $2.3 \mathbf{2}$.2


## Experimental Functions

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Families

## Flow - Family of functions

- pgr_maxFlowMinCost - Experimental - Details of flow and cost on edges.
- pgr_maxFlowMinCost_Cost - Experimental - Only the Min Cost calculation.


## Chinese Postman Problem - Family of functions (Experimental)

- pgr_chinesePostman - Experimental
- pgr_chinesePostmanCost - Experimental


## Coloring - Family of functions

- pgr_bipartite -Experimental - Bipartite graph algorithm using a DFS-based coloring approach.
- pgr_edgeColoring - Experimental - Edge Coloring algorithm using Vizing's theorem.

[^1]pgr_lineGraph - Experimental - Transformation algorithm for generating a Line Graph.

- pgr_lineGraphFull - Experimental - Transformation algorithm for generating a Line Graph out of each vertex in the input graph.


## Traversal - Family of functions

- pgr_breadthFirstSearch - Experimental - Breath first search traversal of the graph.
- pgr_binaryBreadthFirstSearch - Experimental - Breath first search traversal of the graph.


## Components - Family of functions

- pgr_makeConnected - Experimental - Details of edges to make graph connected.
- Supported versions: Latest (3.3) 3.2 3.1 3.0

Chinese Postman Problem - Family of functions (Experimental)

- pgr_chinesePostman - Experimental
- pgr_chinesePostmanCost - Experimental
- Supported versions Latest (3.3) 3.23 .13 .0
pgr_chinesePostman - Experimental
pgr_chinesePostman - Calculates the shortest circuit path which contains every edge in a directed graph and starts and ends on the same vertex.


## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

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- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
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- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.0.0
- New experimental signature


## Description

## The main characteristics are:

- Process is done only on edges with positive costs.
- Running time: $\backslash(\mathrm{O}(\mathrm{E} *(\mathrm{E}+\mathrm{V} * \log \mathrm{~V})) \backslash)$
- Graph must be connected.
- Returns EMPTY SET on a disconnected graph

```
pgr_chinesePostman(Edges SQL)
RETURNS SET OF (seq, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

```
SELECT * FROM pgr_chinesePostman(
    'SELECT id, source, target, cost, reverse_cost
    FROM edges WHERE id < 17');
seq | node | edge | cost | agg_cost
1| 1| 6| 1| 0
2| 3| 7| 1| 1
    3| 7| 4| 1| 2
    4| 6| 4| 1| 3
    5| 7| 8| 1| 4
    6| 11| 8| 1| 5
    7| 10| 1| 6
    8| 12| 1| 7
    12| 13| 1| }
    17| 15| 1| 9
    16| 15| 1| 10
    17| 15| 1| 11
    16| 16| 1| 12
    15| 16| 1| 13
    16| 9| 1| 14
    11| 11| 1| 15
    12| 13| 1| 16
    17| 15| 1| 17
    16| 16| 1| 18
    | 15| 3| 1| 19
    10| 5| 1| 20
    11| 9| 1| 21
    16| 16| 1| 22
    | 15| 3| 1| 23
    |10| 2| 1| 24
    6| 1| 1| 25
    5| 1| 1| 26
    6| 4| 1| 27
    7 7| 10| 1| 28
    8| 14| 1| 29
    9| 14| 1| 30
    |\mp@code{8| 10| 1| 31}
    7 7| 7| 1| 32
    3| 6| 1| 33
    35| 1| -1| 0| 34
(35 rows)
```

Parameters

| Parameter | Type | Description |
| :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |

Inner Queries

Edges SQL

An Edges SQL that represents a directed graph with the following columns

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |

## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Returns set of (seq, node, edge, cost, agg_cost)
Column Type Description

| seq | INT | Sequential value starting from 1 |
| :--- | :--- | :--- |
| node | BIGINT | Identifier of the node in the path from start_vid to end_vid. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for the last node of the <br> path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | FLOAT | Aggregate cost from start_v to node. |

See Also

- Chinese Postman Problem - Family of functions (Experimental)
- Sample Data


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
pgr_chinesePostmanCost - Experimental
pgr_chinesePostmanCost - Calculates the minimum costs of a circuit path which contains every edge in a directed graph and starts and ends on the same vertex.


## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.0.0
- New experimental signature


## Description

## The main characteristics are:

- Process is done only on edges with positive costs.
- Running time: $\backslash(\mathrm{O}(\mathrm{E} *(\mathrm{E}+\mathrm{V} * \log \mathrm{~V})) \backslash)$
- Graph must be connected.

```
pgr_chinesePostmanCost(Edges SQL)
RETURNS FLOAT
```


## Example:

```
SELECT * FROM pgr_chinesePostmanCost(
    'SELECT id, source, target, cost, reverse_cost
    FROM edges WHERE id < 17');
pgr_chinesepostmancost
--------------------
(1 row)
```


## Parameters

| Parameter | Type | Description |  |
| :--- | :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |  |

Inner Queries

Edges SQL

An Edges SQL that represents a directed graph with the following columns

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |  |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |  |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |  |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |  |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |
|  |  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |  |
|  |  |  |  |
|  |  |  |  |
|  |  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

| Column | Type | Description |
| :--- | :--- | :--- |
| pgr_chinesepostmancost | FLOAT | Minimum costs of a circuit <br> path. |

## See Also

- Chinese Postman Problem - Family of functions (Experimental)
- Sample Data


## Indices and tables

- Index
- Search Page


## Warning

Possible server crash

- These functions might create a server crash

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Description

## The main characteristics are:

- Process is done only on edges with positive costs.
- Running time: $\backslash(\mathrm{O}(\mathrm{E} *(\mathrm{E}+\mathrm{V} * \log \mathrm{~V})) \backslash)$
- Graph must be connected.


## Parameters

| Parameter | Type | Description |
| :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |

## Inner Queries

Edges SQL

An Edges SQL that represents adirected graph with the following columns

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |
|  |  |  |
|  |  |  |

## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## See Also

## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1) 3.0
- Unsupported versions: 2.6

[^2]
## Warning

Experimental functions

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- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need c/c++ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting
- pgr_lineGraph - Experimental - Transformation algorithm for generating a Line Graph.
- pgr_lineGraphFull - Experimental - Transformation algorithm for generating a Line Graph out of each vertex in the input graph.

Supported versions: Latest (3.3) 3.2 3.1) 3.0

- Unsupported versions: $2.6 \mathbf{2 . 5}$
pgr_lineGraph - Experimental
pgr_lineGraph — Transforms the given graph into its corresponding edge-based graph.


Boost Graph Inside

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 2.5.0
- New Experimental function


## Description

Given a graph $G$, its line graph $L(G)$ is a graph such that:

- Each vertex of $\mathrm{L}(\mathrm{G})$ represents an edge of G
- Two vertices of $\mathrm{L}(\mathrm{G})$ are adjacent if and only if their corresponding edges share a common endpoint in $G$.


## Signatures

```
pgr_lineGraph(Edges SQL, [directed])
RETURNS SET OF (seq, source, target, cost, reverse_cost)
OR EMPTY SET
```


## Example:

For a directed graph

```
SELECT * FROM pgr_lineGraph(
    'SELECT id, source, target, cost, reverse_cost FROM edges'
);
seq | source | target | cost | reverse_cost
1| -18| 18| 1| 1
    2| -17| 17| 1| 1
    3| -16| -3| 1| -1
    4| -14| -10| 1| 
    6| -14| 14| 1| 1
    7| -10| -7| 1| 1
    8| -10| 
    10|-10| 10| 1| 1
    11| -9| -8| 1| 1
    12| -9| 9| 1| 1
    13|
    14| -8| -7| 1|
    15| -8| -4| 1|
    16| -8| 8| 1| 
    17| -7| -6| 1| 
    18| -7| 7| 1| 
    20| -3| -2| 1| -1
    21|
    23| -2| 
    24|
    26|
    ccc:c|c
    30| 50
    31| 8| 11| 1| 10, -1
    32| 9| -16| 1| 
    34|
    35|}1011|\mp@code{13| 1|
    30|
    40| 16| 15| 1| 1
(40 rows)
```


## Parameters

| Parameter | Type | Description |
| :--- | :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |


| Column | Type | Default | Description |  |
| :---: | :---: | :---: | :---: | :---: |
| directed | boolean | true |  | When true the graph is considered Directed |
|  |  |  |  | Whenfalse the graph is considered as Undirected. |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (source, target) |
|  |  |  |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

RETURNS SET OF (seq, source, target, cost, reverse_cost)

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. <br> - Gives a local identifier for the edge |
| source | BIGINT | Identifier of the source vertex of the current edge. <br> - When negative: the source is the reverse edge in the original graph. |
| target | BIGINT | Identifier of the target vertex of the current edge. <br> - When negative: the target is the reverse edge in the original graph. |
| cost | FLOAT | Weight of the edge (source, target). <br> - When negative: edge (source, target) does not exist, therefore it's not part of the graph. |
| reverse_cost | FLOAT | Weight of the edge (target, source). <br> - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |

See Also

- https://en.wikipedia.org/wiki/Line_graph
- The queries use the Sample Data network.


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1) 3.0
- Unsupported versions: 2.6
pgr_lineGraphFull - Experimental
pgr_lineGraphFull - Transforms a given graph into a new graph where all of the vertices from the original graph are converted to line graphs.


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need c/c++ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 2.6.0
- New Experimental function


## Description

pgr_lineGraphFull, converts original directed graph to a directed line graph by converting each vertex to a complete graph and keeping all the original edges. The new connecting edges have a cost 0 and go between the adjacent original edges, respecting the directionality.

A possible application of the resulting graph is"routing with two edge restrictions":

- Setting a cost of using the vertex when routing between edges on the connecting edge
- Forbid the routing between two edges by removing the connecting edge

This is possible because each of the intersections (vertices) in the original graph are now complete graphs that have a new edge for each possible turn across that intersection.

The main characteristics are:

- This function is for directed graphs.
- Results are undefined when a negative vertex id is used in the input graph.
- Results are undefined when a duplicated edge id is used in the input graph.
- Running time: TBD


## Signatures

## Summary

```
pgr_lineGraphFull(Edges SQL)
RETURNS SET OF (seq, source, target, cost, edge)
OR EMPTY SET
```


## Example:

Full line graph of subgraph of edges $\backslash(\backslash\{4,7,8,10 \backslash\} \backslash)$

```
SELECT * FROM pgr_lineGraphFull(
    $$SELECT id, source, target, cost, reverse_cost
    FROM edges
    WHERE id IN (4, 7, 8, 10)$$)
seq | source | target | cost | edge
    1| -1| 7| 1| 4
    6| -1 | 0| 0
    -2|
    -3| 3| 1| -7
    -4| 11| 1| 8
    -5| 8| 1| 10
    7| -2| 0)|
    7| -3| 0| 0
    7| -4| 0| 0
    7| -5| 0| 0
    -6 | -2| 0) 0
    -6| -3| 0| 0
        -6| -4| 0| 0
        -6| -5| 0| 0
        -7| -2| 0| 0
        -7| -3| 0| 0
        7
        -7| -5| 0| 0
        -8| -2| 01 0
        -8| -3| 01 0
        -8| -4| 0| 0
        -5| 0| 0
        -6| 1|
        3|
        11| --7| 10| -8
        11|-10| 0| 0
        -11| -8| 1| -10
    28| 8| -11| 0| 0
(28 rows)
```

Parameters

| Parameter | Type | Description |  |
| :--- | :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |  |
|  |  | as |  |

Inner Queries

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |
|  |  |  |
|  |  |  |

## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

RETURNS SET OF (seq, source, target, cost, edge)

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Sequential value starting from $\mathbf{1 .}$ |
|  |  | Gives a local identifier for the edge |
| source | BIGINT | Identifier of the source vertex of the current edge. |
|  |  | - When negative: the source is the reverse edge in the original graph. |


| Column | Type | Description |
| :---: | :---: | :---: |
| target | BIGINT | Identifier of the target vertex of the current edge. <br> - When negative: the target is the reverse edge in the original graph. |
| cost | FLOAT | Weight of the edge (source, target). <br> - When negative: edge (source, target) does not exist, therefore it's not part of the graph. |
| reverse_cost | FLOAT | Weight of the edge (target, source). <br> - When negative: edge (target, source) does not exist, therefore it's not part of the graph. |

## Additional Examples

```
The data
    The transformation
    Creating table that identifies transformed vertices
    - Store edge results
    - Create the mapping table
    - Filling the mapping table
- Adding a soft restriction
    - Idenifying the restriction
    - Adding a value to the restriction
- Simplifying leaf vertices
    - Using the vertex map give the leaf verices their original value.
    - Removing self loops on leaf nodes
- Complete routing graph
    - Add edges from the original graph
    - Add the newly calculated edges
- Using the routing graph
```

The examples of this section are based on theSample Data network. The examples include the subgraph including edges 4, 7,8 , and 10 with reverse_cost.

The data

This example displays how this graph transformation works to create additional edges for each possible turn in a graph.

```
SELECT id, source, target, cost, reverse_cost
FROM edges
WHERE id IN (4, 7, 8, 10);
id | source | target | cost | reverse_cost
4|
(4 rows)
```



The transformation

```
SELECT * FROM pgr_lineGraphFull(
    $$SELECT id, source, target, cost, reverse_cost
FROM edges
WHERE id IN (4, 7, 8, 10)$$)
seq | source | target | cost | edge
```




In the transformed graph, all of the edges from the original graph are still present (yellow), but we now have additional edges for every turn that could be made across vertex 7 (orange).

Creating table that identifies transformed vertices

The vertices in the transformed graph are each created by splitting up the vertices in the original graph. Unless a vertex in the original graph is a leaf vertex, it will generate more than one vertex in the transformed graph. One of the newly created vertices in the transformed graph will be given the same vertex identifier as the vertex that it was created from in the original graph, but the rest of the newly created vertices will have negative vertex ids.

Following is an example of how to generate a table that maps the ids of the newly created vertices with the original vertex that they were created from

```
Store edge results
```

The first step is to store the results of thepgr_lineGraphFull call into a table

```
SELECT seq AS id, source, target, cost, edge
INTO lineGraph_edges
FROM pgr_lineGraphFull
    $$SELECT id, source, target, cost, reverse_cost
    FROM edges
    WHERE id IN (4, 7, 8, 10)$$);
SELECT }2
```


## Create the mapping table

```
SELECT id, NULL::BIGINT original_id
INTO vertex_map
FROM vertices;
SELECT 17
```


## Add the new vertices

```
INSERT INTO vertex_map (id)
(SELECT id
FROM pgr_extractVertices(
$$SELECT id, source, target FROM lineGraph_edges$$) WHERE id < 0);
INSERT 0 11
```


## Filling the mapping table

The positive vertex identifiers are the original identifiers

```
UPDATE vertex_map
SET original_id = id
WHERE id > 0;
UPDATE 17
```


## Inspecting the vertices map

## SELECT *

FROM vertex_map ORDER BY id DESC;
id |original_id

| 17\| |
| :---: |
| 16 \| |
| 15\| |
| 14\| |
| 13 \| |
| 12 \| |
| 11\| |
| 10\| |
| 91 |
| 8\| |
| 71 |
| $6 \mid$ |
| 51 |
| 4 \| |
| 31 |
| 21 |
| 1 \| |
| -1\| |
| -2 \| |
| -3\| |
| -4 \| |
| -5 \| |
| -6\| |
| -7\| |
| -8\| |
| -9 \| |
| -10\| |
| -11\| |
| (28 rows) |

The self loops happen when there is no cost traveling to thetarget and the source has an original value.

```
SELECT *, source AS targets_original_id
    FROM lineGraph_edges
    WHERE cost = 0 and source > 0;
id | source | target | cost | edge | targets_original_id
\begin{tabular}{|c|c|c|c|c|c|}
\hline 21 & 6 & -1। & & 01 & 6 \\
\hline 71 & 71 & -2| & 01 & 01 & 7 \\
\hline 81 & 71 & -31 & & 01 & 7 \\
\hline 91 & 71 & -4| & 01 & 01 & 7 \\
\hline \(10 \mid\) & 7 & -5| & \(0 \mid\) & \(0 \mid\) & 7 \\
\hline \(24 \mid\) & 31 & -9| & \(0 \mid\) & \(0 \mid\) & 3 \\
\hline 26 & 11| & -10| & 01 & 01 & 1 \\
\hline 281 & 8| & -11| & 01 & 01 & 8 \\
\hline
\end{tabular}
(8 rows)
```

Updating values from self loops

```
WITH
self_loops AS (
SELECT DISTINCT source, target, source AS targets_original_id
FROM lineGraph_edges
WHERE cost = 0 and source >0)
UPDATE vertex_map SET original_id = targets_original_id
FROM self_loops WHERE target = id;
UPDATE }
```


## Inspecting the vertices table

```
SELECT
FROM vertex map WHERE id < 0
ORDER BY id DESC;
id | original_id
-1|
7
7
7
-6।
9| 3
-10| 11
-11| 8
(11 rows)
```

Updating from inner self loops

```
WITH
assigned_vertices
AS (SELECT id, original id
FROM vertex_map
WHERE original_id IS NOT NULL),
cross_edges
AS (SELECT DISTINCT e.source, v.original_id AS source_original_id
    FROM lineGraph edges AS e
    JOIN vertex_map AS v ON (e.target = v.id)
    WHERE source NOT IN (SELECT id FROM assigned_vertices
)
UPDATE vertex_map SET original_id = source_original_id
FROM cross_edges WHERE source = id;
UPDATE 3
```

Inspecting the vertices map
SELECT *
FROM vertex map WHERE id < 0
ORDER BY id DESC
id | original_id

| $-1 \mid$ | 6 |
| :---: | :---: |
| $-2 \mid$ | 7 |
| $-3 \mid$ | 7 |
| $-4 \mid$ | 7 |
| $-5 \mid$ | 7 |
| $-6 \mid$ | 7 |
| $-7 \mid$ | 7 |
| $-8 \mid$ | 7 |
| $-9 \mid$ | 3 |
| $-10 \mid$ | 11 |
| $-11 \mid$ | 8 |

## Adding a soft restriction

A soft restriction going from vertex 6 to vertex 3 using edges $4->7$ is wanted.

## Idenifying the restriction

Running a pgr_dijkstraNear - Proposed the edge with cost 0, edge 8, is where the cost will be increased

```
SELECT seq, path seq, start vid, end vid, node, original id, edge, cost, agg cost
FROM (SELECT * FROM pgr_dijkstraNear
$$SELECT * FROM lineGraph_edges$$,
    SELECT array_agg(id) FROM vertex_map where original_id = 6),
    (SELECT array_agg(id) FROM vertex_map where original_id = 3))) dn
JOIN vertex_map AS v1 ON (node = v1.id)
seq | path_seq | start_vid | end_vid | node | original_id | edge | cost | agg_cost
    3| -1| 3| -3| 7| 4| 1|
    1| -1| 3| -1| 6| 1| 1| 0
4| -1| 3| 3| 3| -1| 0| 2
2| -1| 3| 7| 7| 8| 0| 1
(4 rows)
```

The edge to be altered isWHERE cost $=0$ AND seq $!=1$ AND edge $!=-1$ from the previus query:

```
SELECT edge FROM pgr_dijkstraNear(
    $$SELECT * FROM lineGraph_edges$$,
    (SELECT array_agg(id) FROM vertex_map where original id = 6),
    (SELECT array_agg(id) FROM vertex_map where original_id = 3))
WHERE cost = 0 AND seq != 1 AND edge != -1;
edge
8
(1 row)
```

Adding a value to the restriction

Updating the cost to the edge:

```
UPDATE lineGraph_edges
SET cost = 100
WHERE id IN (
SELECT edge FROM pgr_dijkstraNear(
    $$SELECT * FROM lineGraph_edges$$,
    (SELECT array_agg(id) FROM vertex_map where original_id = 6),
    (SELECT array_agg(id) FROM vertex_map where original_id = 3))
WHERE cost = 0 AND seq != 1 AND edge != -1)
UPDATE }
```


## Example:

Routing from <br>(6<br>) to <br>(3<br>)

Now the route does not use edge 8 and does a U turn on a leaf vertex.

```
WITH
results AS (
SELECT * FROM pgr dijkstraNear(
    $$SELECT * FROM lineGraph edges$$,
    (SELECT array_agg(id) FROM vertex_map where original_id = 6),
    (SELECT array_agg(id) FROM vertex_map where original_id = 3)))
SELECT seq, path_seq, start_vid, end_vid, node, original_id, edge, cost, agg_cost
FROM results
LEFT JOIN vertex_map AS v1 ON (node = v1.id) ORDER BY seq
seq | path_seq | start_vid | end_vid | node | original_id | edge | cost | agg_cost
\begin{tabular}{ccccccc}
\(1 \mid\) & \(-1 \mid\) & \(3|-1|\) & \(6 \mid\) & \(1 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(-1 \mid\) & \(3 \mid\) & \(7 \mid\) & \(7 \mid\) & \(10 \mid\) & \(0 \mid\) \\
\(3 \mid\) & \(-1 \mid\) & \(3 \mid\) & \(-5 \mid\) & \(7 \mid\) & \(6 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(-1 \mid\) & \(3 \mid\) & \(8 \mid\) & \(8 \mid\) & 1 \\
\(5 \mid\) & \(-1 \mid\) & \(3|-11|\) & \(8 \mid\) & \(0 \mid\) & 2 \\
\(6 \mid\) & \(-1 \mid\) & \(3|-8|\) & \(7 \mid\) & \(20 \mid\) & \(0 \mid\) & 2 \\
\(7 \mid\) & \(-1 \mid\) & \(3|-3|\) & \(7 \mid\) & \(4 \mid\) & \(1 \mid\) & 3 \\
\(8 \mid\) & \(-1 \mid\) & \(3 \mid\) & \(3 \mid\) & \(3|-1|\) & \(0 \mid\) & 4
\end{tabular}
(8 rows)
```


## simplifying leaf vertice

In this example, there is no additional cost for traversing a leaf vertex

Using the vertex map give the leaf verices their original value.

On the source column

```
WITH
u_turns AS (
SELECT e.id AS eid, v1 .original_id
FROM linegraph_edges as e
JOIN vertex_map AS v1 ON (source = v1. .id)
AND v1 .original_id IN (3,6,8,11))
UPDATE lineGraph_edges
SET source = original_id
FROM u turns
WHERE id = eid;
UPDATE }
```


## On the target column

```
WITH
u_turns AS (
SELECT e.id AS eid, v1 .original id
FROM linegraph_edges as e
JOIN vertex_map AS v1 ON (target = v1.id)
AND v1.original_id IN (3,6,8,11))
UPDATE lineGraph_edges
SET target = original_id
FROM u_turns
WHERE id = eid;
UPDATE }
```


## Removing self loops on leaf nodes

The self loops of the leaf nodes are

```
SELECT * FROM linegraph_edges
WHERE source = target
ORDER BY id;
id | source | target | cost | edge
2| 6| 6| 0| 0
24| 3| 3| 0| 0
26| 11| 11| 0| 0
28| 8| 8| 0| 0
(4 rows)
```


## Which can be removed

```
DELETE FROM linegraph_edges
WHERE source = target;
```

DELETE

## Example:

Routing from $\backslash(6 \backslash)$ to $\backslash(3 \backslash)$
Routing can be done now using the original vertices id usingpgr_dijkstra

```
WITH
results AS (
    SELECT * FROM pgr_dijkstra(
    $$SELECT * FROM lineGraph_edges$$, 6, 3))
SELECT seq, path_seq, node, original_id, edge, cost, agg_cost
FROM results
LEFT JOIN vertex_map AS v1 ON (node = v1.id) ORDER BY seq
seq | path_seq | node | original_id | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline 1| & 1| & 61 & 61 & 1| & 1| & 0 \\
\hline 21 & 21 & 71 & 7| & 91 & 01 & 1 \\
\hline 31 & 3| & -4| & 7| & 5| & 1| & 1 \\
\hline 4 | & 4| & 11| & 11 & 25 | & 1| & 2 \\
\hline 51 & 51 & -7| & 71 & \(16 \mid\) & \(0 \mid\) & 3 \\
\hline 61 & 61 & -3| & 7| & 4| & 1| & 3 \\
\hline 71 & 71 & 31 & 3| & -1| & \(0 \mid\) & 4 \\
\hline
\end{tabular}
```


## Complete routing graph

## Add edges from the original graph

Add all the edges that are not involved in the line graph process to the new table

Some administrative tasks to get new identifiers for the edges

```
CREATE SEQUENCE new_graph_id_seq;
CREATE SEQUENCE
ALTER TABLE new_graph ALTER COLUMN id SET DEFAULT nextval('new_graph_id_seq');
ALTER TABLE
ALTER TABLE new_graph ALTER COLUMN id SET NOT NULL;
ALTER TABLE
ALTER SEQUENCE new_graph_id_seq OWNED BY new_graph.id;
ALTER SEQUENCE
SELECT setval('new_graph_id_seq', (SELECT max(id) FROM new_graph));
setval
    18
(1 row)
```

Add the newly calculated edges

INSERT INTO new_graph (source, target, cost, reverse_cost)
SELECT source, target, cost, -1 FROM lineGraph_edges;
INSERT 024

Using the routing graph

When using this method for routing with soft restrictions there will be uturns

## Example:

Routing from <br>(6<br>) to <br>(3<br>)


## Example:

Routing from $\backslash(5 \backslash)$ to $\backslash(1 \backslash)$

```
WITH
results AS (
    SELECT * FROM pgr_dijkstra(
    $$SELECT * FROM new_graph$$, 5, 1))
SELECT seq, path_seq, node, original_id, edge, cost, agg_cost
FROM results
LEFT JOIN vertex_map AS v1 ON (node = v1.id) ORDER BY seq
seq | path_seq | node | original_id | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline 1 | & 1| & 5| & 5| & 1 | & 1| & 0 \\
\hline 2 | & 21 & 6| & 6| & 35 | & 1 | & 1 \\
\hline 31 & 31 & 71 & 71 & 201 & 01 & 2 \\
\hline 4 | & 4 | & -4| & 7| & 41| & 1| & 2 \\
\hline 51 & 51 & 11| & 11 & 371 & 1| & 3 \\
\hline 61 & 61 & -7| & 7| & \(27 \mid\) & 01 & 4 \\
\hline 71 & 71 & -3| & 7| & \(40 \mid\) & \(1 \mid\) & 4 \\
\hline 81 & 81 & 31 & 31 & \(6 \mid\) & 1| & 5 \\
\hline 9 | & 91 & 1| & 1| & -1| & 01 & 6 \\
\hline
\end{tabular}
(9 rows)
```

See Also
https://en.wikipedia.org/wiki/Line_graph

- https://en.wikipedia.org/wiki/Complete_graph


## Indices and tables

- Index
- Search Page

This family of functions is used for transforming a given input graph $\backslash(\mathrm{G}(\mathrm{V}, \mathrm{E}) \backslash)$ into a new graph $\backslash\left(\mathrm{G}^{\prime}\left(\mathrm{V}^{\prime}, \mathrm{E}^{\prime}\right) \backslash\right)$.
See Also

## Indices and tables

Index

- Search Page


## categories

## Vehicle Routing Functions - Category (Experimental)

- Pickup and delivery problem
- pgr_pickDeliver - Experimental - Pickup \& Delivery using a Cost Matrix
- pgr_pickDeliverEuclidean - Experimental - Pickup \& Delivery with Euclidean distances
- Distribution problem
- pgr_vrpOneDepot - Experimental - From a single depot, distributes orders
- Supported versions: Latest (3.2 3.1) 3.0


## Vehicle Routing Functions - Category (Experimental)

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting
- Pickup and delivery problem
- pgr_pickDeliver - Experimental - Pickup \& Delivery using a Cost Matrix
- pgr_pickDeliverEuclidean - Experimental - Pickup \& Delivery with Euclidean distances
- Distribution problem
- pgr_vrpOneDepot - Experimental - From a single depot, distributes orders


## Contents

- Vehicle Routing Functions - Category (Experimental)
- Introduction
- Characteristics
- Pick \& Delivery
- Parameters
- Pick \& deliver
- Pick-Deliver optional parameters
- Inner Queries
- Orders SQL
- Vehicles SQL
- Matrix SQL
- Return columns
- Summary Row
- Handling Parameters
- Capacity and Demand Units Handling
- Locations
- Time Handling
- Factor handling
- See Also
- Supported versions: Latest (3.3) 3.2 3.1) 3.0
pgr_pickDeliver - Experimental
pgr_pickDeliver - Pickup and delivery Vehicle Routing Problem


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- Functionality might change.
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- Might need c/c++ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.0.0
- New experimental function


## Synopsis

Problem: Distribute and optimize the pickup-delivery pairs into a fleet of vehicles.

- Optimization problem is NP-hard.
- pickup and Delivery with time windows.
- All vehicles are equal.
- Same Starting location.
- Same Ending location which is the same as Starting location.
- All vehicles travel at the same speed.
- A customer is for doing a pickup or doing a deliver.
- has an open time.
- has a closing time.
- has a service time.
- has an ( $x, y$ ) location.
- There is a customer where to deliver a pickup.
- travel time between customers is distance / speed
- pickup and delivery pair is done with the same vehicle.
- A pickup is done before the delivery.


## Characteristics

- All trucks depart at time 0 .
- No multiple time windows for a location.
- Less vehicle used is considered better.
- Less total duration is better.
- Less wait time is better
- the algorithm will raise an exception when
- If there is a pickup-deliver pair than violates time window
- The speed, max_cycles, ma_capacity have illegal values
- Six different initial will be optimized - the best solution found will be result


## Signature

pgr_pickDeliver(Orders SQL, Vehicles SQL, Matrix SQL, [options])
options: [factor, max_cycles, initial_sol]
RETURNS SET OF (seq, vehicle_number, vehicle_id, stop, order_id, stop_type, cargo, travel_time, arrival_time, wait_time, service_time, departure_time)

## Example:

Solve the following problem
Given the vehicles:

```
SELECT id, capacity, start node id, start open, start close
FROM vehicles;
id | capacity | start_node_id | start_open | start_close
\begin{tabular}{ccc|}
\hline\(--+-------+-------------------+------11 \mid\) & \(0 \mid\) & 50 \\
\(1 \mid\) & \(50 \mid\) & \(11 \mid\) \\
\(2 \mid\) & \(50 \mid\) & \(11 \mid\) \\
\hline
\end{tabular}
(2 rows)
```

and the orders:

```
SELECT id, demand,
    p_node_id, p_open, p_close, p_service,
    d_node_id,d_open, d_close, d_service
FROM orders;
id | demand | p_node_id | p_open | p_close | p_service | d_node_id | d_open | d_close | d_service
\begin{tabular}{lllllllll}
\(1 \mid\) & \(10 \mid\) & \(10 \mid\) & \(2 \mid\) & \(10 \mid\) & \(3 \mid\) & \(3 \mid\) & \(6 \mid\) & \(15 \mid\) \\
\(2 \mid\) & \(20 \mid\) & \(16 \mid\) & \(4 \mid\) & \(15 \mid\) & \(2 \mid\) & \(15 \mid\) & \(6 \mid\) & \(20 \mid\) \\
\(3 \mid\) & \(30 \mid\) & \(7 \mid\) & \(2 \mid\) & \(10 \mid\) & \(3 \mid\) & \(12 \mid\) & \(3 \mid\) & \(20 \mid\) \\
\(3 \mid\)
\end{tabular}
(3 rows)
```

The query:

```
SELECT * FROM pgr_pickDeliver(
    $$SELECT id, demand,
        p_node_id, p_open, p_close, p_service,
        d_node_id, d_open, d_close, d_service
    FROMM orders$\overline{$}
    $$SELECT id, capacity, start_node_id, start_open, start_close
    FROM vehicles$$,
    $$SELECT * from pgr_dijkstraCostMatrix(
    'SELECT * FROM edges ',
    (SELECT array_agg(id) FROM (SELECT p_node_id AS id FROM orders
        UNION
        SELECT d_node_id FROM orders
        UNION
        SELECT start_node_id FROM vehicles) a))
$$);
seq | vehicle_seq | vehicle_id | stop_seq | stop_type | stop_id | order_id | cargo | travel_time | arrival_time | wait_time | service_time | departure_time
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|c|c|}
\hline 1 | & 1| & 1| & 1| & 1| & 11| & -1| & 01 & 01 & \(0 \mid\) & 01 & 01 & 0 \\
\hline \(2 \mid\) & 1| & 1| & 21 & 21 & 7| & 3| & 301 & 1 | & 1 | & 1| & 31 & 5 \\
\hline 31 & 1| & 1 | & 3| & 31 & 12 | & 31 & 01 & \(2 \mid\) & 71 & 01 & 31 & 10 \\
\hline 4| & 1| & 1| & 4| & 21 & 16| & 21 & \(20 \mid\) & 21 & 12 | & 01 & \(2 \mid\) & 14 \\
\hline 51 & 1| & 1 | & 51 & 31 & 15| & 21 & 01 & 1| & 15| & 01 & 31 & 18 \\
\hline 61 & 1| & 1| & 61 & 61 & 11| & -1| & 01 & 21 & 201 & 01 & 01 & 20 \\
\hline 71 & 21 & 21 & 1| & 1| & 11| & -1| & 01 & 01 & 01 & 01 & 01 & 0 \\
\hline 8| & 21 & 21 & \(2 \mid\) & \(2 \mid\) & 10 | & 1| & 10| & 31 & 31 & 01 & 31 & 6 \\
\hline 91 & 21 & 21 & 3| & 31 & 3| & 1| & 01 & 31 & 91 & 01 & 31 & 12 \\
\hline 10| & 21 & 21 & 4| & 61 & 11| & -1| & 01 & 21 & 14 | & 01 & 0| & 14 \\
\hline 11| & -2| & 01 & 01 & -1| & -1| & -1| & -1| & 16 | & -1| & 1| & 17 | & 34 \\
\hline
\end{tabular}
(11 rows)
```

The parameters are:

| Column | Type | Description |
| :--- | :--- | :--- |
| Orders SQL | TEXT | Orders SQL as described below. |
| Vehicles SQL | TEXT | Vehicles SQL as described <br> below. |
| Matrix SQL | TEXT | Matrix SQL as described below. |

Pick-Deliver optional parameters

Column Type Default Description

| factor | NUMERIC | 1 | Travel time multiplier. See Factor handling |
| :---: | :---: | :---: | :---: |
| max_cycles | INTEGER | 10 | Maximum number of cycles to perform on the optimization. |
| initial_sol | INTEGER | 4 | Initial solution to be used. |
|  |  |  | - 1 One order per truck |
|  |  |  | - 2 Push front order. |
|  |  |  | - 3 Push back order. |
|  |  |  | - 4 Optimize insert. |
|  |  |  | 5 Push back order that allows more orders to be inserted at the back |
|  |  |  | 6 Push front order that allows more orders to be inserted at the front |

## Orders SQL

A SELECT statement that returns the following columns:

```
id, demand
p_node_id, p_open, p_close, [p_service,]
d_node_id, d_open, d_close, [d_service,]
```

where:

| Column | Type | Description |
| :---: | :---: | :---: |
| id | ANY-INTEGER | Identifier of the pick-delivery order pair. |
| demand | ANY-NUMERICAL | Number of units in the order |
| p_open | ANY-NUMERICAL | The time, relative to 0 , when the pickup location opens. |
| p_close | ANY-NUMERICAL | The time, relative to 0 , when the pickup location closes. |
| [p_service] | ANY-NUMERICAL | The duration of the loading at the pickup location. <br> When missing: 0 time units are used |
| d_open | ANY-NUMERICAL | The time, relative to 0 , when the delivery location opens. |
| d_close | ANY-NUMERICAL | The time, relative to 0 , when the delivery location closes. |
| [d_service] | ANY-NUMERICAL | The duration of the unloading at the delivery location. <br> - When missing: 0 time units are used |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Column Type Description

| p_node_id | ANY-INTEGER | The node identifier of the pickup, must match a vertex identifier in theMatrix SQL. |
| :--- | :--- | :--- |
| d_node_id ANY-INTEGER | The node identifier of the delivery, must match a vertex identifier in theMatrix |  |
|  | SQL. |  |

Where:

## Vehicles SQL

A SELECT statement that returns the following columns:

```
id, capacity
start_node_id, start_open, start_close [, start_service,]
[end_node_id, end_open, end_close, end_service]
```

where:

| Column | Type | Description |
| :---: | :---: | :---: |
| id | ANY-NUMERICAL | Identifier of the vehicle. |
| capacity | ANY-NUMERICAL | Maiximum capacity units |
| start_open | ANY-NUMERICAL | The time, relative to 0 , when the starting location opens. |
| start_close | ANY-NUMERICAL | The time, relative to 0 , when the starting location closes. |
| [start_service] | ANY-NUMERICAL | The duration of the loading at the starting location. <br> - When missing: A duration of $\backslash(0 \backslash)$ time units is used. |
| [end_open] | ANY-NUMERICAL | The time, relative to 0 , when the ending location opens. <br> - When missing: The value ofstart_open is used |
| [end_close] | ANY-NUMERICAL | The time, relative to 0 , when the ending location closes. <br> - When missing: The value ofstart_close is used |
| [end_service] | ANY-NUMERICAL | The duration of the loading at the ending location. <br> - When missing: A duration in start_service is used. |


| Column | Type | Description |  |
| :--- | :--- | :--- | :--- |
| start_node_id | ANY-INTEGER | The node identifier of the start location, must match a vertex identifier in theMatrix <br> SQL. |  |
| [end_node_id] |  | ANY-INTEGER | The node identifier of the end location, must match a vertex identifier in theMatrix SQL. |
|  |  | When missing: end_node_id is used. |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Matrix SQL

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Return columns

```
RETURNS SET OF
(seq, vehicle_seq, vehicle_id, stop_seq, stop_type,
    travel_time, arrival_time, wait_time, service_time, departure_time)
UNION
(summary row)
```

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Sequential value starting from $\mathbf{1 .}$ |
| vehicle_seq | INTEGER | Sequential value starting from $\mathbf{1}$ for current vehicles. The $\backslash\left(\mathrm{n} \_\{\text {th }\} \backslash\right)$ vehicle in the solution. |
|  |  | - Value $\backslash(-2 \backslash)$ indicates it is the summary row. |


| Column | Type | Description |
| :--- | :--- | :--- |
| vehicle_id | BIGINT | Current vehicle identifier. |
|  |  |  |
|  |  |  |
|  |  |  |

See Also

- Vehicle Routing Functions - Category (Experimental)
- Sample Data


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
- Unsupported versions: 2.6 2.5 2.4 2.3 2.2 2.1
pgr_pickDeliverEuclidean - Experimental
pgr_pickDeliverEuclidean - Pickup and delivery Vehicle Routing Problem


## Warning

Possible server crash

- These functions might create a server crash

Experimental functions

- They are not officially of the current release.
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- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.0.0
- Replaces pgr_gsoc_vrppdtw
- New experimental function


## Synopsis

Problem: Distribute and optimize the pickup-delivery pairs into a fleet of vehicles.

- Optimization problem is NP-hard.
- Pickup and Delivery:
- capacitated
- with time windows.
- The vehicles
- have ( $x, y$ ) start and ending locations.
- have a start and ending service times.
- have opening and closing times for the start and ending locations.
- An order is for doing a pickup and a a deliver.
- has ( $x, y$ ) pickup and delivery locations.
- has opening and closing times for the pickup and delivery locations.
- has a pickup and deliver service times.
- There is a customer where to deliver a pickup.
- travel time between customers is distance / speed
- pickup and delivery pair is done with the same vehicle.
- A pickup is done before the delivery.


## Characteristics

- No multiple time windows for a location.
- Less vehicle used is considered better.
- Less total duration is better.
- Less wait time is better.
- Six different optional different initial solutions
- the best solution found will be result


## Signature

```
pgr_pickDeliverEuclidean(Orders SQL, Vehicles SQL, [options])
options: [factor, max_cycles, initial_sol]
RETURNS SET OF (seq, vehicle_number, vehicle_id, stop, order_id, stop_type, cargo, travel_time, arrival_time, wait_time, service_time,
departure_time)
```


## Example:

Solve the following problem
Given the vehicles:

```
SELECT id, capacity, start_x, start_y, start_open, start_close
```

FROM vehicles;
id | capacity | start_x | start_y | start_open | start_close

| ----+---------+-------------------------------------1 |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $1 \mid$ | $50 \mid$ | $3 \mid$ | $2 \mid$ | $0 \mid$ | 50 |
| $2 \mid$ | $50 \mid$ | $3 \mid$ | $2 \mid$ | $0 \mid$ | 50 |
| (2 rows) |  |  |  |  |  |

and the orders:

```
SELECT id, demand,
    p_x, p_y, p_open, p_close, p_service,
    d_x, d_y, d_open, d_close, d_service
FROM orders;
id | demand | p_x | p_y | p_open | p_close | p_service | d_x | d_y | d_open | d_close | d_service
\begin{tabular}{lllllllll}
\(1 \mid\) & \(10 \mid\) & \(3 \mid\) & \(1 \mid\) & \(2 \mid\) & \(10 \mid\) & \(3 \mid\) & \(1 \mid\) & \(2 \mid\) \\
\(2 \mid\) & \(20 \mid\) & \(4 \mid\) & \(2 \mid\) & \(4 \mid\) & \(15 \mid\) & \(2 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(20 \mid\) & 3 \\
\(3 \mid\) & \(30 \mid\) & \(2 \mid\) & \(2 \mid\) & \(2 \mid\) & \(10 \mid\) & \(3 \mid\) & \(3 \mid\) & \(3 \mid\) \\
3 \\
(3 rows)
\end{tabular}
```

The query:

```
SELECT * FROM pgr_pickDeliverEuclidean(
    $$SELECT id, demand,
        p_x, p_y, p_open, p_close, p_service,
        d_x, d_y, d_open, d_close, d_service
    FROM orders$$,
    $$SELECT id, capacity, start_x, start_y, start_open, start_close
    FROM vehicles$$);
seq | vehicle_seq | vehicle_id | stop_seq | stop_type | order_id | cargo | travel_time | arrival_time | wait_time | service_time | departure_time
```



## Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Orders SQL | TEXT | Orders SQL as described below. |
| Vehicles SQL | TEXT | Vehicles SQL as described <br> below. |

Pick-Deliver optional parameters
Column Type Default Description


## Orders SQL

A SELECT statement that returns the following columns:
id, demand
p_x, p_y, p_open, p_close, [p_service,]

## Where:

| Column | Type | Description |  |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the pick-delivery order pair. |  |
| demand | ANY-NUMERICAL | Number of units in the order |  |
| p_open | ANY-NUMERICAL | The time, relative to 0, when the pickup location opens. |  |
| p_close | ANY-NUMERICAL | The time, relative to 0, when the pickup location closes. |  |
| [p_service] | ANY-NUMERICAL | The duration of the loading at the pickup location. |  |
|  |  |  | When missing: 0 time units are used |

Where:

## ANY-INTEGER:

## SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

| Column | Type | Description |
| :---: | :---: | :---: |
| p_x | ANY-NUMERICAL | $\backslash(x \backslash)$ value of the pick up location |
| p_y | ANY-NUMERICAL | $\backslash(y)$ value of the pick up location |
| d_x | ANY-NUMERICAL | $\(x \mid)$ value of the delivery location |
| d_y | ANY-NUMERICAL | $\(y)$ value of the delivery location |

Where:

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Vehicles SQL

A SELECT statement that returns the following columns:

```
id, capacity
start_x, start_y, start_open, start_close [, start_service, ]
[ end_x, end_y, end_open, end_close, end_service ]
```

where:

| Column | Type | Description |
| :---: | :---: | :---: |
| id | ANY-NUMERICAL | Identifier of the vehicle. |
| capacity | ANY-NUMERICAL | Maiximum capacity units |
| start_open | ANY-NUMERICAL | The time, relative to 0 , when the starting location opens. |
| start_close | ANY-NUMERICAL | The time, relative to 0 , when the starting location closes. |
| [start_service] | ANY-NUMERICAL | The duration of the loading at the starting location. <br> - When missing: A duration of $\backslash(O \backslash)$ time units is used. |
| [end_open] | ANY-NUMERICAL | The time, relative to 0 , when the ending location opens. <br> When missing: The value ofstart_open is used |
| [end_close] | ANY-NUMERICAL | The time, relative to 0 , when the ending location closes. <br> - When missing: The value ofstart_close is used |


| Column | Type | Description |  |
| :---: | :---: | :---: | :---: |
| [end_service] | ANY-NUMERICAL | The duration of the loading at the ending loca <br> When missing: A duration in start_service is |  |
|  | Column | Type | Description |
|  | start_X | ANY-NUMERICAL | $\(x \backslash)$ value of the starting location |
|  | start_y | ANY-NUMERICAL | $\backslash(y \backslash)$ value of the starting location |
|  | [end_x] | ANY-NUMERICAL | $\backslash(x \backslash)$ value of the ending location <br> - When missing: start_ x is used. |
|  | [end_y] | ANY-NUMERICAL | $\backslash(y \backslash)$ value of the ending location <br> - When missing: start_y is used. |

Where:

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Return columns

| RETURNS SET OF <br> (seq, vehicle_seq, vehicle_id, stop_seq, stop_type, travel_time, arrival_time, wait_time, service_time, departure_time) <br> UNION <br> (summary row) |  |  |
| :---: | :---: | :---: |
| Column | Type | Description |
| seq | INTEGER | Sequential value starting from 1. |
| vehicle_seq | INTEGER | Sequential value starting from $\mathbf{1}$ for current vehicles. The $\backslash\left(n_{\_}\{t h\} \backslash\right)$ vehicle in the solution. <br> Value $\backslash(-2 \backslash)$ indicates it is the summary row. |
| vehicle_id | BIGINT | Current vehicle identifier. <br> - Sumary row has the total capacity violations. <br> - A capacity violation happens when overloading or underloading a vehicle. |
| stop_seq | INTEGER | Sequential value starting from $\mathbf{1}$ for the stops made by the current vehicle. Thel(m_\{th\} <br> ) stop of the current vehicle. <br> - Sumary row has the total time windows violations. <br> - A time window violation happens when arriving after the location has closed. |
| stop_type | INTEGER | - Kind of stop location the vehicle is at <br> - $\backslash(-1 \backslash)$ : at the solution summary row <br> - (1)): Starting location  - (2)): Pickup location  - (3): Delivery location <br> - $\backslash(6 \backslash)$ : Ending location and indicates the vehicle's summary row |
| order_id | BIGINT | Pickup-Delivery order pair identifier. <br> - Value $\backslash(-1 \backslash)$ : When no order is involved on the current stop location. |
| cargo | FLOAT | Cargo units of the vehicle when leaving the stop. <br> Value $\backslash(-1 \backslash)$ on solution summary row. |
| travel_time | FLOAT | Travel time from previous stop_seq to current stop_seq. <br> - Summary has the total traveling time: <br> - The sum of all the travel_time. |
| arrival_time | FLOAT | Time spent waiting for current location to open. $\backslash(-1 \backslash)$ : at the solution summary row. $\backslash(0)$ : at the starting location. |
| wait_time | FLOAT | Time spent waiting for current location to open. <br> Summary row has the total waiting time: <br> - The sum of all the wait_time. |


| Column | Type | Description |
| :---: | :---: | :---: |
| service_time | FLOAT | Service duration at current location. <br> - Summary row has the total service time: <br> - The sum of all the service_time. |
| departure_time | FLOAT | - The time at which the vehicle departs from the stop. <br> - (arrival\_time + wait\_time + service\_time). <br> - The ending location has the total time used by the current vehicle. <br> - Summary row has the total solution time: <br> - $\backslash$ (total $\backslash$ traveling $\backslash$ time + total $\backslash$ waiting $\backslash$ time + total service $\backslash$ time $\backslash)$. |

## Example

```
The vehicles
- The original orders
- The orders
- The query
```

This data example Ic101 is from data published at https://www.sintef.no/projectweb/top/pdptw/li-lim-benchmark/

The vehicles

There are 25 vehciles in the problem all with the same characteristics.

```
CREATE TABLE v_lc101(
    id BIGINT NOT NULLL primary key,
    capacity BIGINT DEFAULT 200,
    start_x FLOAT DEFAULT 30,
    start y FLOAT DEFAULT 50,
    start_open INTEGER DEFAULT 0,
    start_close INTEGER DEFAULT 1236);
CREATE TABLE
/* create 25 vehciles */
INSERT INTO v_Ic101 (id)
(SELECT * FROM generate_series(1, 25));
INSERT 0 25
```

The original orders
The data comes in different rows for the pickup and the delivery of the same order.

```
CREATE table Ic101_c(
id BIGINT not null primary key,
    x DOUBLE PRECISION
    y DOUBLE PRECISION,
    demand INTEGER,
    open INTEGER,
    close INTEGER,
    service INTEGER,
    pindex BIGINT
    dindex BIGINT
);
CREATE TABLE
/* the original data */
INSERT INTO Ic101 c(
    d, x, y, demand, open, close, service, pindex, dindex) VALUES
( 1, 45, 68, -10, 912, 967, 90, 11, 0),
( 2, 45, 70, -20, 825, 870, 90, 6, 0),
( 3, 42, 66, 10, 65, 146, 90, 0, 75),
( 4, 42, 68, -10, 727, 782, 90, 9, 0),
( 5, 42, 65, 10, 15, 67, 90, 0, 7),
( 6, 40, 69, 20, 621, 702, 90, 0, 2),
( 7, 40, 66, -10, 170, 225, 90, 5, 0),
( 8, 38, 68, 20, 255, 324, 90, 0, 10),
( 9, 38, 70, 10, 534, 605, 90, 0, 4),
(10, 35, 66, -20, 357, 410, 90, 8, 0),
(11, 35, 69, 10, 448, 505, 90, 0, 1),
(12, 25, 85, -20, 652, 721, 90, 18, 0),
(13, 22, 75, 30, 30, 92, 90, 0, 17),
(14, 22, 85, -40, 567, 620, 90, 16, 0),
(15, 20, 80, -10, 384, 429, 90, 19, 0),
(16, 20, 85, 40, 475, 528, 90, 0, 14),
(17, 18, 75, -30, 99, 148, 90, 13, 0),
(18, 15, 75, 20, 179, 254, 90, 0, 12),
(19, 15, 80, 10, 278, 345, 90, 0, 15),
(20, 30, 50, 10, 10, 73, 90, 0, 24),
(21, 30, 52, -10, 914, 965, 90, 30, 0),
(22, 28, 52, -20, 812, 883, 90, 28, 0),
(23, 28, 55, 10, 732, 777, 0, 0, 103),
(24, 25, 50, -10, 65, 144, 90, 20, 0),
(25, 25, 52, 40, 169, 224, 90, 0, 27),
(26, 25, 55, -10, 622, 701, 90, 29, 0),
```



## The orders

The original data needs to be converted to an appropiate table:

```
WITH deliveries AS (SELECT * FROM Ic101_c WHERE dindex = 0)
SELECT
row_number() over() AS id, p.demand,
p.id as p_node_id, p.x AS p_x, p.y AS p_y, p.open AS p_open, p.close as p_close, p.service as p_service,
d.id as d_node_id, d.x AS d_x, d.y AS d_y,d.open AS d_open, d.close as d_close, d.service as d_service
INTO c_Ic101
FROM deliveries as d JOIN Ic101_c as p ON (d.pindex = p.id);
SELECT 53
SELECT * FROM c_Ic101 LIMIT 1;
id | demand | p_node_id | p_x | p_y | p_open | p_close | p_service | d_node_id | d_x d_y |_open | d_close | d_service
1| 10| 3| 42| 66| 65| 146| 90| 75| 45| 65| 997| 1068| 90
(1 row)
```

The query

Showing only the relevant information to compare with the best solution information published on https://www.sintef.no/projectweb/top/pdptw/100-customers/

- The best solution found forlc101 is a travel time: 828.94
- This implementation's travel time: 854.54

```
SELECT travel_time, 828.94 AS best
FROM pgr_pickDeliverEuclidean(
    $$SELECT * FROM c_Ic101 $$,
    $$SELECT * FROM v_lc101 $$,
    max_cycles => 2, initial_sol => 4) WHERE vehicle_seq = -2;
    travel_time | best
854.5412705652799 | 828.94
(1 row)
```


## See Also

- Vehicle Routing Functions - Category (Experimental)
- The queries use the Sample Data network.


## Indices and tables

## Index

- Search Page


## - Supported versions: Latest (3.3) 3.2 3.1 3.0 <br> - Unsupported versions: $2.6 \mathbf{2 . 5} 2.4 \mathbf{2 . 3} \mathbf{2 . 2} \mathbf{2 . 1}$

pgr_vrpOneDepot - Experimental

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## No documentation available

## Availability

- Version 2.1.0

New experimental function

- TBD


## Description

- TBD


## Signatures

- TBD


## Parameters

- TBD


## Inner Queries

- TBD

Result Columns

- TBD


## Additional Example:

```
BEGIN;
BEGIN
SET client_min_messages TO NOTICE;
SET
SELECT * FROM pgr_vrpOneDepot(
    'SELECT * FROM solomon_100_RC_101',
    'SELECT * FROM vrp_vehicles',
    'SELECT * FROM vrp_distance',
    1);
oid | opos | vid | tarrival | tdepart
\begin{tabular}{|c|c|c|c|c|}
\hline -1| & 1| & 1| & 01 & 0 \\
\hline 71 & \(2 \mid\) & 1| & 01 & 0 \\
\hline 91 & 3 | & 1| & 01 & 0 \\
\hline 8| & 4 | & 1| & 01 & 0 \\
\hline 6| & 5 & 1| & 01 & 0 \\
\hline 51 & \(6 \mid\) & 1| & 01 & 0 \\
\hline \(4 \mid\) & 71 & 1| & 01 & 0 \\
\hline 21 & 8 | & 1| & 01 & 0 \\
\hline \(6 \mid\) & 9| & 1| & \(40 \mid\) & 51 \\
\hline 8| & \(10 \mid\) & \(1 \mid\) & 62| & 89 \\
\hline 91 & 11| & \(1 \mid\) & 94| & 104 \\
\hline 7| & 12| & \(1 \mid\) & 110| & 120 \\
\hline 4| & 13| & \(1 \mid\) & 131 | & 141 \\
\hline 2| & 14| & \(1 \mid\) & 144 | & 155 \\
\hline 51 & \(15 \mid\) & \(1 \mid\) & 162 | & 172 \\
\hline -1| & 16| & 1 | & 208 | & 208 \\
\hline -1| & & \(2 \mid\) & \(0 \mid\) & 0 \\
\hline 10| & 21 & 21 & 01 & 0 \\
\hline 11| & 31 & 21 & \(0 \mid\) & 0 \\
\hline 10| & 41 & 21 & 34| & 101 \\
\hline 11| & 51 & 21 & 106| & 129 \\
\hline -1| & & 21 & 161 | & 161 \\
\hline -1| & 1| & 31 & 01 & 0 \\
\hline 3| & \(2 \mid\) & 31 & \(0 \mid\) & 0 \\
\hline 3| & 3| & 31 & 31| & 60 \\
\hline -1| & 4| & 31 & 91| & 91 \\
\hline -1| & & 01 & -1| & 460 \\
\hline
\end{tabular}
(27 rows)
ROLLBACK
ROLLBACK
```

```
DROP TABLE IF EXISTS solomon_100_RC_101 cascade;
CREATE TABLE solomon_100_RC_101(
    id integer NOT NULL PRIMARY KEY,
    order_unit integer,
    open_time integer,
    close_time integer,
    service_time integer
    x float8,
    y float8
);
INSERT INTO solomon_100_RC_101 (id, x, y, order_unit, open_time, close_time, service_time) VALUES
(1, 40.000000, 50.000000, 0, 0, 240, 0),
(2, 25.000000, 85.000000, 20, 145, 175, 10),
(3, 22.000000, 75.000000, 30, 50, 80, 10),
(4, 22.000000, 85.000000, 10, 109, 139, 10),
(5, 20.000000, 80.000000, 40, 141, 171, 10),
(6, 20.000000, 85.000000, 20, 41, 71, 10),
(7, 18.000000, 75.000000, 20, 95, 125, 10),
(8, 15.000000, 75.000000, 20, 79, 109, 10),
(9, 15.000000, 80.000000, 10, 91, 121, 10),
(10, 10.000000, 35.000000, 20, 91, 121, 10),
(11, 10.000000, 40.000000, 30, 119, 149, 10);
DROP TABLE IF EXISTS vrp_vehicles cascade;
CREATE TABLE vrp_vehicles (
    vehicle_id integer not null primary key,
    capacity integer,
    case_no integer
);
INSERT INTO vrp_vehicles (vehicle_id, capacity, case_no) VALUES
(1, 200, 5),
(2, 200,5),
(3, 200, 5);
DROP TABLE IF EXISTS vrp_distance cascade;
WITH
the_matrix_info AS (
    SELECT A.id AS src_id, B.id AS dest_id, sqrt( (a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y)) AS cost
    FROM solomon_100_rc_101 AS A, solomon_100_rc_101 AS B WHERE A.id != B.id
)
SELECT src_id, dest_id, cost, cost AS distance, cost AS traveltime
INTO vrp_distance
FROM the_matrix_info;
```


## See Also

## - https://en.wikipedia.org/wiki/Vehicle_routing_problem

## Indices and tables

## Index

- Search Page


## Introduction

Vehicle Routing Problems VRP are NP-hard optimization problem, it generalises the travelling salesman problem (TSP).

- The objective of the VRP is to minimize the total route cost.
- There are several variants of the VRP problem,


## pgRouting does not try to implement all variants.

## Characteristics

- Capacitated Vehicle Routing Problem CVRP where The vehicles have limited carrying capacity of the goods.
- Vehicle Routing Problem with Time Windows VRPTW where the locations have time windows within which the vehicle's visits must be made.
- Vehicle Routing Problem with Pickup and Delivery VRPPD where a number of goods need to be moved from certain pickup locations to other delivery locations.


## Limitations

- No multiple time windows for a location.
- Less vehicle used is considered better.
- Less total duration is better.
- Less wait time is better.
- Times are relative to 0
- The vehicles
- have start and ending service duration times.
- have opening and closing times for the start and ending locations.
- have a capacity.
- The orders
- Have pick up and delivery locations.
- Have opening and closing times for the pickup and delivery locations.
- Have pickup and delivery duration service times.
- have a demand request for moving goods from the pickup location to the delivery location.
- Time based calculations:
- Travel time between customers is $\backslash$ (distance / speed $\backslash$ )
- Pickup and delivery order pair is done by the same vehicle.
- A pickup is done before the delivery.

Parameters

Pick \& deliver

Used in pgr_pickDeliverEuclidean - Experimental

| Column | Type | Description |
| :--- | :--- | :--- |
| Orders SQL | TEXT | Orders SQL as described below. |
| Vehicles SQL | TEXT | Vehicles SQL as described <br> below. |

Used in pgr_pickDeliver - Experimental

| Column | Type | Description |
| :--- | :--- | :--- |
| Orders SQL | TEXT | Orders SQL as described below. |
| Vehicles SQL | TEXT | Vehicles SQL as described <br> below. |
| Matrix SQL | TEXT | Matrix SQL as described below. |

Pick-Deliver optional parameters

| Column | Type | Default | Description |
| :---: | :---: | :---: | :---: |
| factor | NUMERIC | 1 | Travel time multiplier. See Factor handling |
| max_cycles | INTEGER | 10 | Maximum number of cycles to perform on the optimization. |
| initial_sol | INTEGER | 4 | Initial solution to be used. |
|  |  |  | - 1 One order per truck |
|  |  |  | - 2 Push front order. |
|  |  |  | - 3 Push back order. |
|  |  |  | - 4 Optimize insert. |
|  |  |  | 5 Push back order that allows more orders to be inserted at the back |
|  |  |  | 6 Push front order that allows more orders to be inserted at the front |

## Inner Queries

Orders SQL

Common columns for the orders SQL in both implementations:

| Column | Type | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the pick-delivery order pair. |
| demand | ANY-NUMERICAL | Number of units in the order |
| p_open | ANY-NUMERICAL | The time, relative to 0, when the pickup location opens. |
| p_close | ANY-NUMERICAL | The time, relative to 0, when the pickup location closes. |
| [p_service] | ANY-NUMERICAL | The duration of the loading at the pickup location. |
|  |  | When missing: 0 time units are used |
| d_open | ANY-NUMERICAL | The time, relative to 0, when the delivery location <br> opens. |


| Column | Type | Description |
| :--- | :--- | :--- |
| d_close | ANY-NUMERICAL | The time, relative to 0, when the delivery location <br> closes. |
| [d_service] ANY-NUMERICAL | The duration of the unloading at the delivery location. |  |
|  |  | When missing: 0 time units are used |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

For pgr_pickDeliver - Experimental the pickup and delivery identifiers of the locations are needed:
Column Type Description
p_node_id ANY-INTEGER The node identifier of the pickup, must match a vertex identifier in theMatrix SQL.
d_node_id ANY-INTEGER The node identifier of the delivery, must match a vertex identifier in theMatrix SQL.

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
For pgr_pickDeliverEuclidean - Experimental the $\backslash((x, y) \backslash)$ values of the locations are needed:

| Column | Type | Description |
| :--- | :--- | :--- |
| $p \_x$ | ANY-NUMERICAL | $\backslash(x \backslash)$ value of the pick up location |
| $p-y$ | ANY-NUMERICAL | $\backslash(y \backslash)$ value of the pick up location |
| $d \_x$ | ANY-NUMERICAL | $\backslash(x \backslash)$ value of the delivery <br> location |
| $d-y$ | ANY-NUMERICAL | $\backslash(y \backslash)$ value of the delivery <br> location |

Where:

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Vehicles SQL

Common columns for the vehicles SQL in both implementations:

| Column | Type | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-NUMERICAL | Identifier of the vehicle. |
| capacity | ANY-NUMERICAL | Maiximum capacity units |
| start_open | ANY-NUMERICAL | The time, relative to 0, when the starting location <br> opens. |
| start_close | ANY-NUMERICAL | The time, relative to 0, when the starting location <br> closes. |
| [start_service] | ANY-NUMERICAL | The duration of the loading at the starting location. |
|  |  | $0 \quad$ When missing: A duration of $\backslash(0 \backslash)$ time units is used. |

[end_open] ANY-NUMERICAL The time, relative to 0 , when the ending location opens.

- When missing: The value ofstart_open is used

|  | When missing: The value ofstart_open is used |
| :--- | :--- |
| [end_close] ANY-NUMERICAL The time, relative to 0 , when the ending location closes. |  |

- When missing: The value ofstart_close is used
[end_service] ANY-NUMERICAL The duration of the loading at the ending location.
- When missing: A duration in start_service is used.

For pgr_pickDeliver - Experimental the starting and ending identifiers of the locations are needed:
Column Type Description
start_node_id ANY-INTEGER The node identifier of the start location, must match a vertex identifier in theMatrix SQL.
[end_node_id] ANY-INTEGER The node identifier of the end location, must match a vertex identifier in theMatrix SQL.

- When missing: end_node_id is used.

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
For pgr_pickDeliverEuclidean - Experimental the $\backslash((x, y) \backslash)$ values of the locations are needed:

| Column | Type | Description |
| :---: | :---: | :---: |
| start_x | ANY-NUMERICAL | $\backslash(x \backslash)$ value of the starting location |
| start_y | ANY-NUMERICAL | $\backslash(y \backslash)$ value of the starting location |
| [end_x] | ANY-NUMERICAL | $\backslash(x \backslash)$ value of the ending location <br> - When missing: start_x is used. |
| [end_y] | ANY-NUMERICAL | $\backslash(y \backslash)$ value of the ending location <br> - When missing: start_y is used. |

Where:

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Matrix SQL

Set of (start_vid, end_vid, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| agg_cost | FLOAT | Aggregate cost from start_vid to end_vid. |

## Return columns

```
RETURNS SET OF
(seq, vehicle_seq, vehicle_id, stop_seq, stop_type,
    travel_time, arrival_time, wait_time, service_time, departure_time)
UNION
(summary row)
```

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. |
| vehicle_seq | INTEGER | Sequential value starting from $\mathbf{1}$ for current vehicles. The $\backslash\left(n_{\_}\{t h\} \backslash\right)$ vehicle in the solution. - Value $\backslash(-2 \backslash)$ indicates it is the summary row. |
| vehicle_id | BIGINT | Current vehicle identifier. <br> - Sumary row has the total capacity violations. <br> - A capacity violation happens when overloading or underloading a vehicle. |
| stop_seq | INTEGER | Sequential value starting from 1 for the stops made by the current vehicle. Thel(m_\{th\} <br> ) stop of the current vehicle. <br> - Sumary row has the total time windows violations. <br> - A time window violation happens when arriving after the location has closed. |
| stop_type | INTEGER | - Kind of stop location the vehicle is at <br> - $\backslash(-1 \backslash)$ : at the solution summary row <br> - (1)): Starting location  - (2): Pickup location <br> - (3): Delivery location <br> - $\backslash(6 \backslash)$ : Ending location and indicates the vehicle's summary row |


| Column | Type | Description |
| :---: | :---: | :---: |
| order_id | BIGINT | Pickup-Delivery order pair identifier. <br> - Value $\backslash(-1 \backslash)$ : When no order is involved on the current stop location. |
| cargo | FLOAT | Cargo units of the vehicle when leaving the stop. <br> - Value $\backslash(-1 \backslash)$ on solution summary row. |
| travel_time | FLOAT | Travel time from previous stop_seq to current stop_seq. <br> - Summary has the total traveling time: <br> - The sum of all the travel_time. |
| arrival_time | FLOAT | Time spent waiting for current location to open. <br> - $\backslash(-1 \backslash)$ : at the solution summary row. <br> - <br> (0)): at the starting location. |
| wait_time | FLOAT | Time spent waiting for current location to open. <br> - Summary row has the total waiting time: <br> - The sum of all the wait_time. |
| service_time | FLOAT | Service duration at current location. <br> - Summary row has the total service time: <br> - The sum of all the service_time. |
| departure_time | FLOAT | - The time at which the vehicle departs from the stop. <br> - (arrival\_time + wait\_time + service\_time). <br> - The ending location has the total time used by the current vehicle. <br> - Summary row has the total solution time: <br> - (total\ traveling $\backslash$ time + total $\backslash$ waiting $\backslash$ time + total $\backslash$ service $\backslash$ time $\$ ). |

Summary Row

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Continues the sequence |
| vehicle_seq | INTEGER | Value $\backslash(-2 \backslash)$ indicates it is the summary row. |
| vehicle_id | BIGINT | total capacity violations: |
|  |  |  |
| stop_seq | INTEGER | total time windows violations: |
|  |  |  |

## Handling Parameters

To define a problem, several considerations have to be done, to get consistent results. This section gives an insight of how parameters are to be considered.

- Capacity and Demand Units Handling
- Locations
- Time Handling
- Factor Handling

The capacity of a vehicle, can be measured in:

- Volume units like $\backslash\left(m^{\wedge} 3 \backslash\right)$.
- Area units like $\backslash\left(m^{\wedge} 2 \backslash\right)$ (when no stacking is allowed).
- Weight units like $\backslash(\mathrm{kg} \backslash)$.
- Number of boxes that fit in the vehicle.
- Number of seats in the vehicle

The demand request of the pickup-deliver orders must use the same units as the units used in the vehicle'sapacity.
To handle problems like: 10 (equal dimension) boxes of apples and 5 kg of feathers that are to be transported (not packed in boxes).

- If the vehicle's capacity is measured in boxes, a conversion of kg of feathers to number of boxes is needed.
- If the vehicle's capacity is measured in kg , a conversion of box of apples to kg is needed.

Showing how the 2 possible conversions can be done
Let: - $\backslash(f \backslash$ boxes $\backslash)$ : number of boxes needed for 1 kg of feathers. $-\backslash(a \backslash$ weight $\backslash)$ : weight of 1 box of apples.

| Capacity Units | apples | feathers |  |
| :--- | :--- | :--- | :--- |
| boxes | 10 | $\backslash(5$ | $*$ |
|  |  | f_boxes $\backslash)$ |  |
| kg | $\backslash(10 * a \backslash$ weight $\backslash$ | 5 |  |

## Locations

- When using pgr_pickDeliverEuclidean - Experimental:
- The vehicles have $\backslash((x, y) \backslash)$ pairs for start and ending locations.
- The orders Have $\backslash((x, y) \backslash)$ pairs for pickup and delivery locations.
- When using pgr_pickDeliver - Experimental:
- The vehicles have identifiers for the start and ending locations.
- The orders have identifiers for the pickup and delivery locations.
- All the identifiers are indices to the given matrix.

Time Handling

The times are relative to $\mathbf{0}$. All time units have to be converted to a0 reference and the same time units.
Suppose that a vehicle's driver starts the shift at 9:00 am and ends the shift at 4:30 pm and the service time duration is 10 minutes with 30 seconds.

| Meaning of $\mathbf{0}$ | time units | $\mathbf{9 : 0 0} \mathbf{a m}$ | $\mathbf{4 : 3 0} \mathbf{~ p m}$ | $\mathbf{1 0} \mathbf{~ m i n} \mathbf{3 0} \mathbf{~ s e c s ~}$ |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $0: 00 \mathrm{am}$ | hours | 9 | 16.5 | $\backslash(10.5$ | $/$ | $60=$ |
|  |  |  |  | $0.175 \backslash)$ |  |  |
| $0: 00 \mathrm{am}$ | minutes | $\backslash(9 * 60=54 \backslash)$ | $\backslash(16.5 * 60=990 \backslash)$ | 10.5 |  |  |
| $9: 00 \mathrm{am}$ | hours | 0 | 7.5 | $\backslash(10.5$ | $/$ | $60=$ |
|  |  |  |  | $0.175 \backslash)$ |  |  |
| $9: 00 \mathrm{am}$ | minutes | 0 | $\backslash(7.5 * 60=540 \backslash)$ | 10.5 |  |  |

Factor handling
factor acts as a multiplier to convert from distance values to time units the matrix values or the euclidean values.

- When the values are already in the desired time units
- factor should be $\mathbf{1}$
- When factor > 1 the travel times are faster
- When factor $<1$ the travel times are slower

For the pgr_pickDeliverEuclidean - Experimental:
Working with time units in seconds, and $x / y$ in lat/lon: Factor: would depend on the location of the points and on the average velocity say $25 \mathrm{~m} / \mathrm{s}$ is the velocity.

| Latitude | Conversion | Factor |  |  |
| :--- | :--- | :--- | :--- | :--- |
| 45 | 1 longitude degree is $(78846.81 \mathrm{~m}) /(25 \mathrm{~m} / \mathrm{s})$ | 3153 s |  |  |
| 0 | $1 \quad$ longitude degree is $(111319.46$ | 4452 s |  |  |
|  | $\mathrm{~m}) /(25 \mathrm{~m} / \mathrm{s})$ |  |  |  |
|  |  |  |  |  |
|  |  |  |  |  |

For the pgr_pickDeliver - Experimental:
Given $\backslash(v=d / t \backslash)$ therefore $\backslash(t=d / v \backslash)$ And the factor becomes $\backslash(1 / v \backslash)$

Where:
v:
Velocity
d:
Distance
t:
Time
For the following equivalences $\backslash(10 \mathrm{~m} / \mathrm{s}$ \approx $600 \mathrm{~m} / \mathrm{min}$ \approx $36 \mathrm{~km} / \mathrm{hr} \backslash$ )
Working with time units in seconds and the matrix been in meters: For a 1000 m lenght value on the matrix:

| Units | velocity | Conversion | Factor | Result |
| :--- | :--- | :--- | :--- | :--- |
| seconds | $\backslash(10 \mathrm{~m} / \mathrm{s} \backslash)$ | $\backslash(\backslash$ frac $\{1\}\{10 \mathrm{~m} / \mathrm{s}\} \backslash)$ | $\backslash(0.1 \mathrm{~s} / \mathrm{m} \backslash)$ | $\backslash(1000 \mathrm{~m} * 0.1 \mathrm{~s} / \mathrm{m}=100 \mathrm{~s} \backslash)$ |
| minutes | $\backslash(600 \mathrm{~m} / \mathrm{min} \backslash)$ | $\backslash(\backslash$ frac $\{1\}\{600 \mathrm{~m} / \mathrm{min}\} \backslash)$ | $\backslash(0.0016 \mathrm{~min} / \mathrm{m} \backslash)$ | $\backslash(1000 \mathrm{~m} * 0.0016 \mathrm{~min} / \mathrm{m}=1.6 \mathrm{~min} \backslash)$ |
| Hours | $\backslash(36 \mathrm{~km} / \mathrm{hr} \backslash)$ | $\backslash(\backslash$ frac $\{1\}\{36 \mathrm{~km} / \mathrm{hr}\} \backslash)$ | $\backslash(0.0277 \mathrm{hr} / \mathrm{km} \backslash)$ | $\backslash(1 \mathrm{~km} * 0.0277 \mathrm{hr} / \mathrm{km}=0.0277 \mathrm{hr} \backslash)$ |

See Also

- https://en.wikipedia.org/wiki/Vehicle_routing_problem
- The queries use the Sample Data network.


## Indices and tables

- Index
- Search Page


## Not classified

- pgr_bellmanFord - Experimental
- pgr_dagShortestPath - Experimental
- pgr_edwardMoore - Experimental
- pgr_isPlanar - Experimental
- pgr_stoerWagner - Experimental
- pgr_topologicalSort - Experimental
- pgr_transitiveClosure - Experimental
- pgr_turnRestrictedPath - Experimental
- pgr_lengauerTarjanDominatorTree -Experimental
- Supported versions: Latest (3.3) 3.2 3.1 3.0


## pgr_bellmanFord - Experimental

pgr_bellmanFord — Shortest path(s) using Bellman-Ford algorithm.

## boost

Boost Graph Inside

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.2.0
- New experimental signature:
- pgr_bellmanFord (Combinations)
- Version 3.0.0
- New experimental signatures:
- pgr_bellmanFord (One to One)
- pgr_bellmanFord (One to Many)
- pgr_bellmanFord (Many to One)
- pgr_bellmanFord (Many to Many)


## Description

Bellman-Ford's algorithm, is named after Richard Bellman and Lester Ford, who first published it in 1958 and 1956, respectively.It is a graph search algorithm that computes shortest paths from a starting vertex (start_vid) to an ending vertex (end_vid) in a graph where some of the edge weights may be negative. Though it is more versatile, it is slower than Dijkstra's algorithm. This implementation can be used with a directed graph and an undirected graph.

## The main characteristics are:

- Process is valid for edges with both positive and negative edge weights.
- Values are returned when there is a path.
- When the start vertex and the end vertex are the same, there is no path. The agg_cost would be(0<br>).
- When the start vertex and the end vertex are different, and there exists a path between them without having a negative cycle. The agg_cost would be some finite value denoting the shortest distance between them.
- When the start vertex and the end vertex are different, and there exists a path between them, but it contains a negative cycle. In such case, agg_cost for those vertices keep on decreasing furthermore, Hence agg_cost can't be defined for them.
- When the start vertex and the end vertex are different, and there is no path. The agg_cost i\$(\infty<br>).
- For optimization purposes, any duplicated value in the start_vids or end_vids are ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending
- Running time: $\backslash(\mathrm{O}(\mid$ start $\backslash$ _vids $\mid *(\mathrm{~V} * \mathrm{E})) \backslash)$


## Signatures

## Summary

```
pgr_bellmanFord(Edges SQL, start vid, end vid, [directed])
pgr_bellmanFord(Edges SQL, start vid, end vids, [directed])
pgr_bellmanFord(Edges SQL, start vids, end vid, [directed])
pgr_bellmanFord(Edges SQL, start vids, end vids, [directed])
pgr_bellmanFord(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (seq, path_seq, [start_vid], [end_vid], node, edge, cost, agg_cost)
OR EMPTY SET
```


## One to One

```
pgr_bellmanFord(Edges SQL, start vid, end vid, [directed])
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

## SELECT * FROM pgr_bellmanFord(

'SELECT id, source, target, cost, reverse_cost FROM edges',
6, 10, true);
seq | path_seq | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ |
| $3 \mid$ | $3 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ |
| $4 \mid$ | $4 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ |
| $5 \mid$ | $5 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ |
| $6 \mid$ | $6 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ |
| 6 rows) |  |  |  |  |

(6 rows)

## One to Many

```
pgr_bellmanFord(Edges SQL, start vid, end vids, [directed])
RETURNS SET OF (seq, path_seq, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(6 \backslash)$ to vertices $\backslash(\backslash\{10,17 \backslash\} \backslash)$ on a directed graph

```
SELECT * FROM pgr bellmanFord(
    SELECT id, source, target, cost, reverse cost FROM edges',
6, ARRAY[10, 17]);
seq | path_seq | end_vid | node | edge | cost | agg_cost
    10| 6| 4| 1| 0
    2| 10| 7| 8| 1| 1
    3| 10| 11| 9| 1| 2
    | 10| 16| 16| 1| 3
    10| 15| 3| 1| 4
    10| 10| -1| 0| 5
    17| 7 1 0 1 1 1 
    17| 11| 11| 1| 2
    9|
11| 5| 17| 17| -1| 0| 4
(11 rows)
```


## Many to One

pgr_bellmanFord(Edges SQL, start vids, end vid, [directed])
RETURNS SET OF (seq, path_seq, start_vid, node, edge, cost, agg_cost OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,1 \backslash\} \backslash)$ to vertex $\backslash(17 \backslash)$ on a directed graph

```
SELECT * FROM pgr_bellmanFord(
    SELECT id, source, target, cost, reverse_cost FROM edges',
    ARRAY[6, 1], 17);
seq | path_seq | start_vid | node | edge | cost | agg_cost
    1| 1| 1| 1| 6| 1| 0
    2| 2| 1| 3| 7| 1| 1
    4| 4| 1| 11| 11| 1| - 
    5| 5| 1| 12| 13| 1| 
    6| 6| 1| 17| -1| 0| 5
    |lllllllll
    8| 2| 6| 7 < 8| 1| 1
    9| 3| 6| 11| 11| 1| 2
    10| 4| 6| 12| 13| 1| 3
11| 5| 6| 17| -1| 0| 4
(11 rows)
```


## Many to Many

RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{6,1 \backslash\} \backslash)$ to vertices $\backslash(\backslash\{10,17 \backslash\} \backslash)$ on an undirected graph

```
SELECT * FROM pgr_bellmanFord(
    'SELECT id, source, target, cost, reverse_cost FROM edges',
    ARRAY[6, 1], ARRAY[10, 17]
    directed => false)
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|c|}
\hline 1 | & 1| & 1| & \(10 \mid\) & 1| & 61 & 1| & 0 \\
\hline \(2 \mid\) & \(2 \mid\) & 1| & 10| & 31 & 7| & 1| & 1 \\
\hline 31 & 31 & 1| & 10| & 7 & 4 | & 1| & 2 \\
\hline \(4 \mid\) & \(4 \mid\) & 1| & 10| & 61 & 21 & 1| & 3 \\
\hline 51 & 51 & 1| & 10| & \(10 \mid\) & -1| & 0| & 4 \\
\hline \(6 \mid\) & 1| & 1| & 17| & 1| & \(6 \mid\) & 1| & 0 \\
\hline 71 & 21 & 1| & 17| & 31 & 71 & 1| & 1 \\
\hline 8| & 31 & 1| & 17| & 71 & 81 & 1| & 2 \\
\hline 91 & \(4 \mid\) & 1| & 17| & 11| & 11| & \(1 \mid\) & 3 \\
\hline \(10 \mid\) & 51 & 1 | & \(17 \mid\) & \(12 \mid\) & \(13 \mid\) & 1| & \\
\hline 11| & 61 & 1 | & \(17 \mid\) & 17| & -1| & 01 & 5 \\
\hline \(12 \mid\) & \(1 \mid\) & 61 & \(10 \mid\) & \(6 \mid\) & 21 & 1| & 0 \\
\hline 131 & 21 & 61 & \(10 \mid\) & \(10 \mid\) & -1| & 01 & 1 \\
\hline \(14 \mid\) & \(1 \mid\) & 61 & \(17 \mid\) & 6| & 4| & 1 | & 0 \\
\hline 151 & 21 & 61 & \(17 \mid\) & 7| & 8| & 1 | & 1 \\
\hline \(16 \mid\) & 31 & 61 & \(17 \mid\) & 11| & 11| & 1| & \\
\hline \(17 \mid\) & 4 | & 61 & \(17 \mid\) & 12| & 13| & 1| & 3 \\
\hline 18| & 51 & 61 & \(17 \mid\) & \(17 \mid\) & -1| & 01 & 4 \\
\hline
\end{tabular}
```


## Combinations

```
pgr_bellmanFord(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

Using a combinations table on anundirected graph.

The combinations table:

```
SELECT source, target FROM combinations;
source | target
    5| 6
    5| 10
    6| 5
    6| 15
    6| 14
(5 rows)
```

The query:

```
SELECT * FROM pgr_bellmanFord(
    SELECT id, source, target, cost, reverse_cost FROM edges',
    'SELECT source, target FROM combinations',
false);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
\begin{tabular}{lllllcll}
\(1 \mid\) & \(1 \mid\) & \(5 \mid\) & \(6 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(5 \mid\) & \(6 \mid\) & \(6 \mid\) & \(-1 \mid\) & \(0 \mid\) & 1 \\
\(3 \mid\) & \(1 \mid\) & \(5 \mid\) & \(10 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) & 0 \\
\(4 \mid\) & \(2 \mid\) & \(5 \mid\) & \(10 \mid\) & \(6 \mid\) & \(2 \mid\) & \(1 \mid\) & 1 \\
\(5 \mid\) & \(3 \mid\) & \(5 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 2 \\
\(6 \mid\) & \(1 \mid\) & \(6 \mid\) & \(5 \mid\) & \(6 \mid\) & \(1 \mid\) & \(1 \mid\) & 0 \\
\(7 \mid\) & \(2 \mid\) & \(6 \mid\) & \(5 \mid\) & \(5 \mid\) & \(-1 \mid\) & \(0 \mid\) & 1 \\
\(8 \mid\) & \(1 \mid\) & \(6 \mid\) & \(15 \mid\) & \(6 \mid\) & \(2 \mid\) & \(1 \mid\) & 0 \\
\(9 \mid\) & \(2 \mid\) & \(6 \mid\) & \(15 \mid\) & \(10 \mid\) & \(3 \mid\) & \(1 \mid\) & 1 \\
\(10 \mid\) & \(3 \mid\) & \(6 \mid\) & \(15 \mid\) & \(15 \mid\) & \(-1 \mid\) & \(0 \mid\) & 2 \\
\((10\) rows \()\) & & & & & & &
\end{tabular}
```


## Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |


| Column | Type | Description |
| :--- | :--- | :--- |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | BIGINT | Identifier of the starting vertex of the path. |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting vertices. |
| end vid | BIGINT | Identifier of the ending vertex of the path. |
| end vids | ARRAY[BIGINT] | Array of identifiers of ending vertices. |

Optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| directed | BOOLEAN true | $\ddots$ | When true the graph is considered Directed |
|  |  |  When false the graph is considered as <br>   <br>   <br>   |  |

Inner Queries

Edges SQL

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER | Identifier of the edge. |  |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |  |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |  |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |

- When negative: edge (target, source) does not exist, therefore it's not part of the graph.

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Combinations SQL

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- | Identifier of the departure vertex. |
|  | INTEGER |  |
| target | ANY- | Identifier of the arrival vertex. |
|  | INTEGER |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Return columns

Returns set of (seq, path_seq [, start_vid] [, end_vid], node, edge, cost, agg_cost)

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. |
| path_seq | INTEGER | Relative position in the path. Has value1 for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex. Returned when multiple starting vetrices are in the query. |
|  |  | - Many to One <br> - Many to Many |
| end_vid | BIGINT | Identifier of the ending vertex. Returned when multiple ending vertices are in the query. |
|  |  | One to Many <br> Many to Many |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |


| Column | Type | Description |
| :--- | :--- | :--- |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for <br> the last node of the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | FLOAT | Aggregate cost from start_vid to node. |

## Additional Examples

## Example 1:

Demonstration of repeated values are ignored, and result is sorted.

```
SELECT * FROM pgr_bellmanFord(
    'SELECT id, source, target, cost, reverse_cost FROM edges',
    ARRAY[7, 10, 15, 10, 10, 15], ARRAY[10, 7, 10, 15]);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|c|}
\hline 1 | & \(1 \mid\) & 7| & \(10 \mid\) & 71 & 81 & 1| & 0 \\
\hline \(2 \mid\) & \(2 \mid\) & 7| & 10| & 11| & & 1| & 1 \\
\hline 3| & 3| & 71 & 10| & \(16 \mid\) & 16| & 1| & 2 \\
\hline 41 & 4 | & 71 & \(10 \mid\) & 15| & 31 & \(1 \mid\) & 3 \\
\hline 51 & 5| & 7| & \(10 \mid\) & 10| & -1| & 01 & 4 \\
\hline 61 & 1| & 7| & \(15 \mid\) & 7| & 8| & 1| & 0 \\
\hline 7| & \(2 \mid\) & 71 & 15| & 11| & 9| & 1| & 1 \\
\hline 8। & 31 & 71 & 15| & 16| & 16| & 11 & 2 \\
\hline 9| & 4 | & 71 & 15| & 15| & -1| & 01 & 3 \\
\hline \(10 \mid\) & \(1 \mid\) & \(10 \mid\) & 71 & 10| & 5| & 1| & 0 \\
\hline 11| & 2 | & \(10 \mid\) & 71 & 11| & 8। & 1| & 1 \\
\hline 12 | & 31 & \(10 \mid\) & 7| & 71 & -1| & 01 & 2 \\
\hline 13 | & \(1 \mid\) & \(10 \mid\) & 15 & \(10 \mid\) & 51 & \(1 \mid\) & 0 \\
\hline 14 | & 2 | & \(10 \mid\) & \(15 \mid\) & 11| & 9| & \(1 \mid\) & \\
\hline 15 | & 31 & \(10 \mid\) & \(15 \mid\) & 16| & 16 & | 1 & \\
\hline \(16 \mid\) & 4 | & \(10 \mid\) & 15 | & 15| & -1| & | \(0 \mid\) & \\
\hline \(17 \mid\) & \(1 \mid\) & 151 & 71 & 15| & 31 & \(1 \mid\) & 0 \\
\hline 18| & \(2 \mid\) & \(15 \mid\) & 71 & \(10 \mid\) & 2। & \(1 \mid\) & 1 \\
\hline 19| & 31 & 151 & 71 & 61 & 4। & 1| & 2 \\
\hline \(20 \mid\) & 4 | & \(15 \mid\) & 71 & & -1| & 01 & 3 \\
\hline 21| & \(1 \mid\) & 151 & \(10 \mid\) & 151 & 31 & 11 & 0 \\
\hline 22 | & 2 | & 15 | & \(10 \mid\) & 10| & -1| & 0 & \\
\hline \multicolumn{8}{|l|}{(22 rows)} \\
\hline
\end{tabular}
```


## Example 2:

Making start vids the same as end vids.

```
SELECT * FROM pgr_bellmanFord(
    'SELECT id, source, target, cost, reverse_cost FROM edges',
    ARRAY[7, 10, 15], ARRAY[7, 10, 15]);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|c|}
\hline 1| & 1| & 7| & 10| & 7| & 8। & & 0 \\
\hline 21 & \(2 \mid\) & 71 & \(10 \mid\) & 11| & 9| & 1| & 1 \\
\hline 31 & 31 & 71 & \(10 \mid\) & 16| & 16| & \(1 \mid\) & 2 \\
\hline 4 | & 41 & 7| & \(10 \mid\) & 15| & 31 & 1| & 3 \\
\hline 51 & 51 & 71 & \(10 \mid\) & 10| & -1| & \(0 \mid\) & 4 \\
\hline 61 & 1| & 71 & 15| & 71 & 8| & 1| & 0 \\
\hline 71 & 21 & 71 & \(15 \mid\) & 11| & 91 & 1| & 1 \\
\hline 8| & 31 & 71 & 15| & 16| & 16| & 1| & 2 \\
\hline 9| & 4 | & 71 & \(15 \mid\) & 15| & -1| & 01 & 3 \\
\hline 10| & 1| & \(10 \mid\) & 71 & 10| & 51 & 1| & 0 \\
\hline 11| & \(2 \mid\) & \(10 \mid\) & 71 & 11| & 8। & 1| & 1 \\
\hline 12 | & 31 & \(10 \mid\) & 71 & 71 & -1| & 01 & 2 \\
\hline 13 | & \(1 \mid\) & \(10 \mid\) & 15| & \(10 \mid\) & 5। & \(1 \mid\) & 0 \\
\hline 14 | & 21 & \(10 \mid\) & 15 & 11| & 9 & \(1 \mid\) & \\
\hline 15 & 31 & \(10 \mid\) & 15| & 16 & 16| & 1 & \\
\hline \(16 \mid\) & 4 | & \(10 \mid\) & \(15 \mid\) & 15 & -1| & 0 & 3 \\
\hline \(17 \mid\) & 1 | & \(15 \mid\) & 7| & 15| & 3| & 1| & 0 \\
\hline 18 | & \(2 \mid\) & \(15 \mid\) & 71 & \(10 \mid\) & 21 & 1| & 1 \\
\hline 19 & 31 & \(15 \mid\) & 71 & 61 & 4| & 1| & 2 \\
\hline \(20 \mid\) & 41 & 151 & 71 & 71 & -1| & 01 & 3 \\
\hline \(21 \mid\) & 1 | & \(15 \mid\) & 10| & & 31 & \(1 \mid\) & 0 \\
\hline 22| & 21 & \(15 \mid\) & \(10 \mid\) & 10 & -1| & 0 & \\
\hline \multicolumn{8}{|l|}{(22 rows)} \\
\hline
\end{tabular}
```


## Example 3:

Manually assigned vertex combinations.

## SELECT * FROM pgr bellmanFord(

'SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT * FROM (VALUES $(6,10),(6,7),(12,10))$ AS combinations (source, target)');
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $6 \mid$ | $7 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $6 \mid$ | $7 \mid$ | $7 \mid$ | $-1 \mid$ | $0 \mid$ |
| $3 \mid$ | $1 \mid$ | $6 \mid$ | $10 \mid$ | $6 \mid$ | $4 \mid$ | $1 \mid$ |
| $4 \mid$ | $2 \mid$ | $6 \mid$ | $10 \mid$ | $7 \mid$ | $8 \mid$ | $1 \mid$ |
| $5 \mid$ | $3 \mid$ | $6 \mid$ | $10 \mid$ | $11 \mid$ | $9 \mid$ | $1 \mid$ |
| $6 \mid$ | $4 \mid$ | $6 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ |
| $7 \mid$ | $5 \mid$ | $6 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ |
| $8 \mid$ | $6 \mid$ | $6 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ |
| $9 \mid$ | $1 \mid$ | $12 \mid$ | $10 \mid$ | $12 \mid$ | $13 \mid$ | $1 \mid$ |
| $10 \mid$ | $2 \mid$ | $12 \mid$ | $10 \mid$ | $17 \mid$ | $15 \mid$ | $1 \mid$ |
| 10 |  |  |  |  |  |  |
| $11 \mid$ | $3 \mid$ | $12 \mid$ | $10 \mid$ | $16 \mid$ | $16 \mid$ | $1 \mid$ |
| $12 \mid$ | $4 \mid$ | $12 \mid$ | $10 \mid$ | $15 \mid$ | $3 \mid$ | $1 \mid$ |
| $13 \mid$ | $5 \mid$ | $12 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ |
| 13 rows) |  |  |  |  | 3 |  |

## See Also

- https://en.wikipedia.org/wiki/Bellman\�\�\�Ford_algorithm
- Sample Data


## Indices and tables

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## pgr_dagShortestPath - Experimental

pgr_dagShortestPath - Returns the shortest path(s) for weighted directed acyclic graphs(DAG). In particular, the DAG shortest paths algorithm implemented by Boost.Graph.
boost

Boost Graph Inside

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need c/c++ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting
- Version 3.2.0
- New experimental function:
- pgr_dagShortestPath(Combinations)
- Version 3.0.0
- New experimental function


## Description

Shortest Path for Directed Acyclic Graph(DAG) is a graph search algorithm that solves the shortest path problem for weighted directed acyclic graph, producing a shortest path from a starting vertex (start_vid) to an ending vertex énd_vid).

This implementation can only be used with a directed graph with no cycles i.e. directed acyclic graph.
The algorithm relies on topological sorting the dag to impose a linear ordering on the vertices, and thus is more efficient for DAG's than either the Dijkstra or Bellman-Ford algorithm.

The main characteristics are:

- Process is valid for weighted directed acyclic graphs only. otherwise it will throw warnings.
- Values are returned when there is a path.
- When the starting vertex and ending vertex are the same, there is no path.
- The agg_cost the non included values $(v, v)$ is 0
- When the starting vertex and ending vertex are the different and there is no path:
- The agg_cost the non included values $(u, v)$ is $\backslash(\backslash i n f t y \backslash)$
- For optimization purposes, any duplicated value in the start_vids or end_vids are ignored.
- The returned values are ordered:
- start_vid ascending
- end_vid ascending
- Running time: $\backslash\left(\mathrm{O}\left(\mid\right.\right.$ start $\backslash$ vids $\left.\left.\left.\right|^{*}(\mathrm{~V}+\mathrm{E})\right) \backslash\right)$


## Signatures

## Summary

```
pgr_dagShortestPath(Edges SQL, start vid, end vid)
pgr_dagShortestPath(Edges SQL, start vid, end vids)
pgr_dagShortestPath(Edges SQL, start vids, end vid)
pgr_dagShortestPath(Edges SQL, start vids, end vids)
pgr_dagShortestPath(Edges SQL, Combinations SQL)
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET
```


## One to One

```
pgr_dagShortestPath(Edges SQL, start vid, end vid)
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(5 \backslash)$ to vertex $\backslash(11 \backslash)$ on a directed graph

```
SELECT * FROM pgr_dagShortestPath(
    'SELECT id, source, target, cost FROM edges',
5,11);
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{cccc|c}
\(1 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 2 \\
\(4 \mid\) & \(11 \mid\) & \(-1 \mid\) & \(0 \mid\) & 3
\end{tabular}
(4 rows)
```


## One to Many

```
pgr_dagShortestPath(Edges SQL, start vid, end vids)
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
```


## Example:

From vertex $\backslash(5 \backslash)$ to vertices $\backslash(\backslash\{7,11 \backslash\} \backslash)$

```
SELECT * FROM pgr_dagShortestPath(
    'SELECT id, source, target, cost FROM edges',
5, ARRAY[7, 11]);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 5| 1| 1| 0
2| 2| 6| 4| 1| 1
3| 3| 7|-1| 0| 2
4| 1| 5| 1| 1| 0
5| 2| 6| 4| 1| 1
6| 3| 7| 8| 1| 2
7| 4| 11|-1| 0| 3
(7 rows)
```


## Many to One

pgr_dagShortestPath(Edges SQL, start vids, end vid)
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET

## Example:

From vertices $\backslash(\backslash\{5,10 \backslash\} \backslash)$ to vertex $\backslash(11 \backslash)$

```
SELECT * FROM pgr_dagShortestPath(
    'SELECT id, source, target, cost FROM edges',
ARRAY[5, 10], 11);
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{cccccc}
\(1 \mid\) & \(1 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(3 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 2 \\
\(4 \mid\) & \(4 \mid\) & \(11 \mid\) & \(-1 \mid\) & \(0 \mid\) & 3 \\
\(5 \mid\) & \(1 \mid\) & \(10 \mid\) & \(5 \mid\) & \(1 \mid\) & 0 \\
\(6 \mid\) & \(2 \mid\) & \(11 \mid\) & \(-1 \mid\) & \(0 \mid\) & 1 \\
(6 rows)
\end{tabular}
```


## Many to Many

```
pgr_dagShortestPath(Edges SQL, start vids, end vids)
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From vertices $\backslash(\backslash\{5,15 \backslash\} \backslash)$ to vertices $\backslash(\backslash\{11,17 \backslash\} \backslash)$ on an undirected graph

```
SELECT * FROM pgr_dagShortestPath(
    'SELECT id, source, target, cost FROM edges',
    ARRAY[5, 15], ARRAY[11, 17]);
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{ccccc}
\(1 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 2 \\
\(4 \mid\) & \(11 \mid\) & \(-1 \mid\) & \(0 \mid\) & 3 \\
\(1 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 2 \\
\(4 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) & 3 \\
\(5 \mid\) & \(16 \mid\) & \(15 \mid\) & \(1 \mid\) & 4 \\
\(6 \mid\) & \(17 \mid\) & \(-1 \mid\) & \(0 \mid\) & 5 \\
\(1 \mid\) & \(15 \mid\) & \(16 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(16 \mid\) & \(15 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(17 \mid\) & \(-1 \mid\) & \(0 \mid\) & 2
\end{tabular}
(13 rows)
```

```
pgr_dagShortestPath(Edges SQL, Combinations)
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
```

OR EMPTY SET

## Example:

Using a combinations table on anundirected graph
The combinations table:

```
SELECT source, target FROM combinations;
source | target
    5| 6
    5| 10
    6| 5
    6| 15
    6| 14
(5 rows)
```

The query:

```
SELECT * FROM pgr_dagShortestPath(
    'SELECT id, source, target, cost FROM edges',
    'SELECT source, target FROM combinations')
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{llllll}
\(1 \mid\) & \(1 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(6 \mid\) & \(-1 \mid\) & \(0 \mid\) & 1
\end{tabular}
(2 rows)
```


## Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | Edges SQL as described below |
| Combinations SQL | TEXT | Combinations SQL as described below |
| start vid | BIGINT | Identifier of the starting vertex of the path. |
| start vids | ARRAY[BIGINT] | Array of identifiers of starting vertices. |
| end vid | BGGINT | Identifier of the ending vertex of the path. |
| end vids | ARRAY[BIGINT] | Array of identifiers of ending vertices. |

Inner Queries
Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (source, target) |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |

Where:

## ANY-INTEGER:

## SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Combinations SQL

| Parameter | Type | Description |
| :--- | :--- | :--- |
| source | ANY- <br> INTEGER | Identifier of the departure vertex. |
| target | ANY- <br> INTEGER | Identifier of the arrival vertex. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## Resturn Columns

Returns set of (seq, path_seq [, start_vid] [, end_vid], node, edge, cost, agg_cost)

| Column | Type | Description |
| :---: | :---: | :---: |
| seq | INTEGER | Sequential value starting from 1. |
| path_seq | INTEGER | Relative position in the path. Has value1 for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex. Returned when multiple starting vetrices are in the query. Many to One Many to Many |
| end_vid | BIGINT | Identifier of the ending vertex. Returned when multiple ending vertices are in the query. One to Many Many to Many |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |
| edge | BIGINT | Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for the last node of the path. |
| cost | FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |
| agg_cost | FLOAT | Aggregate cost from start_vid to node. |

## Additional Examples

## Example 1:

Demonstration of repeated values are ignored, and result is sorted.

```
SELECT * FROM pgr_dagShortestPath(
    'SELECT id, source, target, cost FROM edges',
    ARRAY[5, 10, 5, 10, 10, 5], ARRAY[11, 17, 17, 11]);
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|}
\hline 1 & 1 & 5 & 1 & 1 & 0 \\
\hline \(2 \mid\) & \(2 \mid\) & 61 & 4| & 1| & 1 \\
\hline 31 & 3| & 71 & 8| & 1| & 2 \\
\hline 4 & \(4 \mid\) & 11| & -1| & \(0 \mid\) & 3 \\
\hline 51 & 1| & 51 & 1| & 1| & 0 \\
\hline 61 & \(2 \mid\) & 61 & 4| & 1| & 1 \\
\hline 71 & 3| & 71 & 8। & 1| & 2 \\
\hline 8 | & \(4 \mid\) & 11| & 91 & 1| & 3 \\
\hline 91 & 51 & 16| & 15| & 1| & 4 \\
\hline \(10 \mid\) & 61 & 17| & -1| & \(0 \mid\) & 5 \\
\hline 11| & \(1 \mid\) & \(10 \mid\) & 51 & \(1 \mid\) & 0 \\
\hline 12| & \(2 \mid\) & 11| & -1| & 01 & 1 \\
\hline 13 | & 1 & 10 & 5 & 1 & 0 \\
\hline 14| & \(2 \mid\) & 11| & 91 & \(1 \mid\) & 1 \\
\hline 15| & \(3 \mid\) & 16| & 15 & 1| & 2 \\
\hline 16| & 4 | & 17| & -1| & \(0 \mid\) & 3 \\
\hline (16 rows) & & & & & \\
\hline
\end{tabular}
```


## Example 2:

Making start_vids the same as end_vids

```
SELECT * FROM pgr_dagShortestPath(
    'SELECT id, source, target, cost FROM edges',
    ARRAY[5, 10, 11], ARRAY[5, 10, 11]);
seq | path_seq | node | edge | cost | agg_cost
\begin{tabular}{ccccc}
\(1 \mid\) & \(1 \mid\) & \(5 \mid\) & \(1 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(2 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(3 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(4 \mid\) & \(11 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(5 \mid\) & \(1 \mid\) & 3 \\
\(5 \mid\) & \(10 \mid\) & \(5 \mid\) & \(1 \mid\) & 0 \\
\(6 \mid\) & \(2 \mid\) & \(11 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\((6\) rows) & & &
\end{tabular}
(6 rows)
```


## Example 3:

Manually assigned vertex combinations.

```
SELECT * FROM pgr_dagShortestPath(
    'SELECT id, source, target, cost FROM edges',
    'SELECT * FROM (VALUES (6, 10), (6, 7), (12, 10)) AS combinations (source, target)');
seq | path_seq | node | edge | cost | agg_cost
```

```
1| 1| 6| 4| 1| 0
```

1| 1| 6| 4| 1| 0
2| 2| 7| -1| 0| 1
2| 2| 7| -1| 0| 1
(2 rows)

```

\section*{See Also}

\section*{- Sample Data}
- https://en.wikipedia.org/wiki/Topological_sorting

\section*{Indices and tables}

\section*{Index}
- Search Page

\section*{- Supported versions: Latest (3.3) 3.2}

\section*{pgr_edwardMoore - Experimental}
pgr_edwardMoore - Returns the shortest path using Edward-Moore algorithm.

\section*{Warning}

Possible server crash
- These functions might create a server crash

\section*{Warning}

Experimental functions
- They are not officially of the current release.
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- pgTap tests might be missing.
- Might need c/c++ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting

\section*{Availability}
- Version 3.2.0
- New experimental signature:
- pgr_edwardMoore (Combinations)
- Version 3.0.0
- New experimental signatures:
- pgr_edwardMoore (One to One)
- pgr_edwardMoore (One to Many)
- pgr_edwardMoore (Many to One)
- pgr_edwardMoore (Many to Many)

\section*{Description}

Edward Moore's Algorithm is an improvement of the Bellman-Ford Algorithm. It can compute the shortest paths from a single source vertex to all other vertices in a weighted directed graph. The main difference between Edward Moore's Algorithm and Bellman Ford's Algorithm lies in the run time.

The worst-case running time of the algorithm is \(\backslash(O(|V| *|E|) \backslash)\) similar to the time complexity of Bellman-Ford algorithm. However, experiments suggest that this algorithm has an average running time complexity of \(\backslash(O(|E|) \backslash)\) for random graphs. This is significantly faster in terms of computation speed.

Thus, the algorithm is at-best, significantly faster than Bellman-Ford algorithm and is at-worst,as good as Bellman-Ford algorithm

\section*{The main characteristics are:}
- Values are returned when there is a path.
- When the starting vertex and ending vertex are the same, there is no path.
- The agg_cost the non included values \((v, v)\) is \(\backslash(0 \backslash)\)
- When the starting vertex and ending vertex are the different and there is no path: - The agg_cost the non included values ( \(u, v\) ) is \\(\infty \(\backslash\) )
- For optimization purposes, any duplicated value in the start vids or end vids are ignored.
- The returned values are ordered:
- start vid ascending
- end vid ascending
- Running time:
- Worst case: \(\backslash(\mathrm{O}(|\mathrm{V}| *|\mathrm{E}|) \backslash)\)
- Average case: \(\backslash(\mathrm{O}(|\mathrm{E}|) \backslash)\)

\section*{Signatures}

\section*{Summary}
pgr_edwardMoore(Edges SQL, start vid, end vid, [directed])
pgr_edwardMoore(Edges SQL, start vid, end vids, [directed])
pgr_edwardMoore(Edges SQL, start vids, end vid, [directed])
pgr_edwardMoore(Edges SQL, start vids, end vids, [directed])
pgr_edwardMoore(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (seq, path_seq, [start_vid], [end_vid], node, edge, cost, agg_cost)
OR EMPTY SET

\section*{One to One}
```

pgr_edwardMoore(Edges SQL, start vid, end vid, [directed])
RETURNS SET OF (seq, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertex \\(6) to vertex \(\backslash(10 \backslash)\) on a directed graph
```

SELECT * FROM pgr_edwardMoore(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6,10, true);
seq | path_seq | node | edge | cost | agg_cost
1| 1| 6| 4| 1| 0
2| 7| 8| 1| 1
3| 11| 9| 1| 2
4| 16| 16| 1| 3
5| 15| 3| 1| 4
6| 6| 10| -1| 0| 5
(6 rows)

```

\section*{One to Many}
```

pgr_edwardMoore(Edges SQL, start vid, end vids, [directed])
RETURNS SET OF (seq, path_seq, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

From vertex \(\backslash(6 \backslash)\) to vertices \(\backslash(\backslash\{10,17 \backslash\} \backslash)\) on a directed graph
SELECT * FROM pgr_edwardMoore(
'SELECT id, source, target, cost, reverse_cost FROM edges',
6 , ARRAY[10, 17]);
seq | path_seq | end_vid | node | edge | cost | agg_cost
\begin{tabular}{rccccc}
\(1 \mid\) & \(1 \mid\) & \(10 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(2 \mid\) & \(10 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(3 \mid\) & \(10 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) \\
\(4 \mid\) & \(4 \mid\) & \(10 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) \\
\(5 \mid\) & \(5 \mid\) & \(10 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(6 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\(7 \mid\) & \(1 \mid\) & \(17 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(8 \mid\) & \(2 \mid\) & \(17 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) \\
\(9 \mid\) & \(3 \mid\) & \(17 \mid\) & \(11 \mid\) & \(11 \mid\) & \(1 \mid\) \\
10 \\
\(10 \mid\) & \(4 \mid\) & \(17 \mid\) & \(12 \mid\) & \(13 \mid\) & \(1 \mid\) \\
\(11 \mid\) & \(5 \mid\) & \(17 \mid\) & \(17 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\((11\) rows)
\end{tabular}

Many to One
```

pgr_edwardMoore(Edges SQL, start vids, end vid, [directed])
RETURNS SET OF (seq, path_seq, start_vid, node, edge, cost, agg_cost)

```
OR EMPTY SET

\section*{Example:}

From vertices \(\backslash(\backslash\{6,1 \backslash\} \backslash)\) to vertex \(\backslash(17 \backslash)\) on a directed graph
```

SELECT * FROM pgr_edwardMoore
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[6, 1], 17);
seq path_seq start_vid node | edge cost | agg_cost

```


Many to Many
```

pgr_edwardMoore(Edges SQL, start vids, end vids, [directed])
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)

``` OR EMPTY SET

\section*{Example:}

From vertices \(\backslash(\backslash\{6,1 \backslash\} \backslash)\) to vertices \(\backslash(\backslash\{10,17 \backslash\} \backslash)\) on an undirected graph
\begin{tabular}{ccccc|c|c|c}
\(1 \mid\) & \(1 \mid\) & \(1 \mid\) & \(10 \mid\) & \(1 \mid\) & \(6 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(1 \mid\) & \(10 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(3 \mid\) & \(1 \mid\) & \(10 \mid\) & \(7 \mid\) & \(4 \mid\) & \(1 \mid\) & 2 \\
\(4 \mid\) & \(4 \mid\) & \(1 \mid\) & \(10 \mid\) & \(6 \mid\) & \(2 \mid\) & \(1 \mid\) & 3 \\
\(5 \mid\) & \(5 \mid\) & \(1 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 4 \\
\(6 \mid\) & \(1 \mid\) & \(1 \mid\) & \(17 \mid\) & \(1 \mid\) & \(6 \mid\) & \(1 \mid\) & 0 \\
\(7 \mid\) & \(2 \mid\) & \(1 \mid\) & \(17 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) & 1 \\
\(8 \mid\) & \(3 \mid\) & \(1 \mid\) & \(17 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 2 \\
\(9 \mid\) & \(4 \mid\) & \(1 \mid\) & \(17 \mid\) & \(11 \mid\) & \(11 \mid\) & \(1 \mid\) & 3 \\
\(10 \mid\) & \(5 \mid\) & \(1 \mid\) & \(17|12|\) & \(13 \mid\) & \(1 \mid\) & 4 \\
\(11 \mid\) & \(6 \mid\) & \(1 \mid\) & \(17 \mid\) & \(17 \mid\) & \(-1 \mid\) & \(0 \mid\) & 5 \\
\(12 \mid\) & \(1 \mid\) & \(6 \mid\) & \(10 \mid\) & \(6 \mid\) & \(2 \mid\) & \(1 \mid\) & 0 \\
\(13 \mid\) & \(2 \mid\) & \(6 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 1 \\
\(14 \mid\) & \(1 \mid\) & \(6 \mid\) & \(17 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) & 0 \\
\(15 \mid\) & \(2 \mid\) & \(6 \mid\) & \(17 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 1 \\
\(16 \mid\) & \(3 \mid\) & \(6 \mid\) & \(17 \mid\) & \(11 \mid\) & \(11 \mid\) & \(1 \mid\) & 2 \\
\(17 \mid\) & \(4 \mid\) & \(6 \mid\) & \(17 \mid\) & \(12 \mid\) & \(13 \mid\) & \(1 \mid\) & 3 \\
\(18 \mid\) & \(5 \mid\) & \(6 \mid\) & \(17 \mid\) & \(17 \mid\) & \(-1 \mid\) & \(0 \mid\) & 4
\end{tabular}

\section*{Combinations}
```

pgr_edwardMoore(Edges SQL, Combinations SQL, [directed])
RETURNS SET OF (seq, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)
OR EMPTY SET

```

\section*{Example:}

Using a combinations table on anundirected graph.

The combinations table:
```

SELECT source, target FROM combinations;
source | target
5| 6
5| 10
6| 5
6| 15
6| 14
(5 rows)

```

The query:
```

SELECT * FROM pgr_edwardMoore(
'SELECT id, source, target, cost, reverse_cost FROM edges',
SELECT source, target FROM combinations',
false);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| $1 \mid$ | $1 \mid$ | $5 \mid$ | $6 \mid$ | $5 \mid$ | $1 \mid$ | $1 \mid$ | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \mid$ | $2 \mid$ | $5 \mid$ | $6 \mid$ | $6 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |
| $3 \mid$ | $1 \mid$ | $5 \mid$ | $10 \mid$ | $5 \mid$ | $1 \mid$ | $1 \mid$ | 0 |
| $4 \mid$ | $2 \mid$ | $5 \mid$ | $10 \mid$ | $6 \mid$ | $2 \mid$ | $1 \mid$ | 1 |
| $5 \mid$ | $3 \mid$ | $5 \mid$ | $10 \mid$ | $10 \mid$ | $-1 \mid$ | $0 \mid$ | 2 |
| $6 \mid$ | $1 \mid$ | $6 \mid$ | $5 \mid$ | $6 \mid$ | $1 \mid$ | $1 \mid$ | 0 |
| $7 \mid$ | $2 \mid$ | $6 \mid$ | $5 \mid$ | $5 \mid$ | $-1 \mid$ | $0 \mid$ | 1 |
| $8 \mid$ | $1 \mid$ | $6 \mid$ | $15 \mid$ | $6 \mid$ | $2 \mid$ | $1 \mid$ | 0 |
| $9 \mid$ | $2 \mid$ | $6 \mid$ | $15\|10\|$ | $3 \mid$ | $1 \mid$ | 1 |  |
| $10 \mid$ | $3 \mid$ | $6 \mid$ | $15 \mid$ | $15 \mid$ | $-1 \mid$ | $0 \mid$ | 2 |

(10 rows)

```

\section*{Parameters}
\begin{tabular}{lll} 
Column & Type & Description \\
\hline Edges SQL & TEXT & Edges SQL as described below \\
\hline Combinations SQL & TEXT & Combinations SQL as described below \\
& & \\
\hline start vid & BIGINT & Identifier of the starting vertex of the path. \\
\hline start vids & ARRAY[BIGINT] & Array of identifiers of starting vertices. \\
\hline end vid & BIGINT & Identifier of the ending vertex of the path. \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|c|}
\hline Column & \multicolumn{2}{|c|}{Type} & \multicolumn{2}{|r|}{Description} & \\
\hline end vids & \multicolumn{2}{|r|}{ARRAY[BIGINT]} & \multicolumn{2}{|r|}{Array of identifiers of ending vertices.} & \\
\hline \multicolumn{6}{|l|}{onal parameters} \\
\hline & Column & Type & Default & Description & \\
\hline & directed & BOOLEAN & true & \begin{tabular}{l}
- When true the graph is considered Directed \\
- Whenfalse the graph is considered Undirected.
\end{tabular} & as \\
\hline
\end{tabular}

\section*{Inner Queries}

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL -1 & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Combinations SQL}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline source & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the departure vertex. \\
\hline target & \begin{tabular}{l} 
ANY- \\
INTEGER
\end{tabular} & Identifier of the arrival vertex. \\
\hline
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{Return columns}

Returns set of (seq, path_seq [, start_vid] [, end_vid], node, edge, cost, agg_cost)
\begin{tabular}{|c|c|c|}
\hline Column & Type & Description \\
\hline seq & INTEGER & Sequential value starting from 1. \\
\hline path_seq & INTEGER & Relative position in the path. Has value \(\mathbf{1}\) for the beginning of a path. \\
\hline start_vid & BIGINT & \begin{tabular}{l}
Identifier of the starting vertex. Returned when multiple starting vetrices are in the query.
Many to One \\
- Many to Many
\end{tabular} \\
\hline end_vid & BIGINT & \begin{tabular}{l}
Identifier of the ending vertex. Returned when multiple ending vertices are in the query. \\
One to Many \\
Many to Many
\end{tabular} \\
\hline node & BIGINT & Identifier of the node in the path fromstart_vid to end_vid. \\
\hline edge & BIGINT & Identifier of the edge used to go fromnode to the next node in the path sequence.-1 for the last node of the path. \\
\hline cost & FLOAT & Cost to traverse from node using edge to the next node in the path sequence. \\
\hline agg_cost & FLOAT & Aggregate cost from start_vid to node. \\
\hline
\end{tabular}

\section*{Example 1:}

Demonstration of repeated values are ignored, and result is sorted.
```

SELECT * FROM pgr_edwardMoore(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[7, 10, 15, 10, 10, 15], ARRAY[10, 7, 10, 15]);
seq| path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1\| | 1 \| | 71 | $10 \mid$ | 71 | 8। | 1\| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 21 | $2 \mid$ | 71 | 10\| | 11\| | 9\| | 1 \| | 1 |
| 31 | 31 | 71 | 10\| | $16 \mid$ | 16\| | 1\| | 2 |
| 4 \| | $4 \mid$ | 71 | $10 \mid$ | 15 | 3\| | 1 \| | 3 |
| 51 | 51 | 71 | 10\| | 10\| | -1\| | $0 \mid$ | 4 |
| 61 | 1\| | 71 | 15 \| | 71 | 8\| | 1\| | 0 |
| 71 | 21 | 71 | 15\| | 11\| | 9\| | 1 \| | 1 |
| 81 | 31 | 71 | 15 \| | 16\| | 16\| | 1\| | 2 |
| 91 | $4 \mid$ | 71 | 15\| | 15 | -1\| | 01 | 3 |
| $10 \mid$ | 1 \| | $10 \mid$ | 71 | 10\| | 51 | 1\| | 0 |
| 11\| | $2 \mid$ | $10 \mid$ | 71 | 11\| | 8। | 1 \| | 1 |
| $12 \mid$ | 31 | $10 \mid$ | 71 | 71 | -1\| | 01 | 2 |
| 131 | 1 \| | $10 \mid$ | 15\| | $10 \mid$ | 51 | $1 \mid$ | 0 |
| $14 \mid$ | $2 \mid$ | $10 \mid$ | 15\| | 11\| | 91 | 1\| | 1 |
| 151 | 31 | $10 \mid$ | 15\| | $16 \mid$ | 16\| | \| 1 | 2 |
| $16 \mid$ | $4 \mid$ | $10 \mid$ | 15\| | 15\| | -1\| | $0 \mid$ | 3 |
| $17 \mid$ | 1 \| | 151 | 71 | 15 | 16\| | 1\| | 0 |
| 18\| | 21 | 151 | 71 | $16 \mid$ | 9\| | 1\| | 1 |
| 191 | 31 | 151 | 71 | 11\| | 8। | 1\| | 2 |
| 201 | 4 \| | 151 | 71 | 71 | -1\| | $0 \mid$ | 3 |
| $21 \mid$ | 1 \| | 151 | $10 \mid$ | $15 \mid$ | 31 | 1\| | 0 |
| 221 | 21 | 151 | 10\| | $10 \mid$ | -1\| | $0 \mid$ | 1 |

(22 rows)

```

Example 2:
Making start vids the same as end vids.
```

SELECT * FROM pgr_edwardMoore(
'SELECT id, source, target, cost, reverse_cost FROM edges',
ARRAY[7, 10, 15], ARRAY[7, 10, 15]);
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost

| 1\| | $1 \mid$ | 71 | 10\| | 7\| | 8 \| | 1\| | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 21 | 21 | 71 | 10\| | 11\| | 91 | 1\| | 1 |
| 3 | 31 | 71 | $10 \mid$ | $16 \mid$ | 16\| | $1 \mid$ | 2 |
| $4 \mid$ | 4 \| | 71 | $10 \mid$ | 15\| | 31 | 1\| | 3 |
| 51 | 51 | 71 | 10\| | 10\| | -1\| | 01 | 4 |
| 61 | 1 \| | 71 | 15\| | 71 | 81 | $1 \mid$ | 0 |
| 71 | $2 \mid$ | 71 | 15 \| | 11\| | 91 | 1\| | 1 |
| 81 | 31 | 71 | 15 \| | 16\| | 16\| | $1 \mid$ | 2 |
| 91 | 4 \| | 7 \| | 15 \| | 15\| | -1\| | 01 | 3 |
| $10 \mid$ | 1 \| | $10 \mid$ | 71 | $10 \mid$ | 51 | 1 \| | 0 |
| 11\| | $2 \mid$ | $10 \mid$ | 71 | 11\| | 8\| | 1 \| | 1 |
| $12 \mid$ | 31 | $10 \mid$ | 71 | 71 | -1\| | $0 \mid$ | 2 |
| 131 | 1 \| | $10 \mid$ | 15 \| | 10 | 51 | 1 \| | 0 |
| $14 \mid$ | $2 \mid$ | $10 \mid$ | 15 \| | 11\| | 91 | 1 \| | 1 |
| 151 | 3 | $10 \mid$ | 15 \| | 16 | 16\| | 1 | 2 |
| $16 \mid$ | 4 \| | $10 \mid$ | 15\| | 15 | -1\| | 0 | 3 |
| $17 \mid$ | 1 \| | 151 | 71 | $15 \mid$ | $16 \mid$ | 1 \| | 0 |
| 181 | $2 \mid$ | 15 \| | 71 | $16 \mid$ | 91 | $1 \mid$ | 1 |
| 191 | 31 | 151 | 71 | 11\| | 8। | $1 \mid$ | 2 |
| 201 | 4 \| | 151 | 71 | 71 | -1\| | 01 | 3 |
| 21\| | 1 \| | 15 \| | $10 \mid$ | 15 | 31 | 1\| | 0 |
| 221 | 21 | 15 \| | 10 \| | 10 | -1\| | 0 | 1 |

```

Example 3:
Manually assigned vertex combinations.

\section*{SELECT * FROM pgr_edwardMoore}
'SELECT id, source, target, cost, reverse_cost FROM edges',
'SELECT * FROM (VALUES \((6,10),(6,7),(12,10))\) AS combinations (source, target)');
seq | path_seq | start_vid | end_vid | node | edge | cost | agg_cost
\begin{tabular}{rllllll|l}
\(1 \mid\) & \(1 \mid\) & \(6 \mid\) & \(7 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) & 0 \\
\(2 \mid\) & \(2 \mid\) & \(6 \mid\) & \(7 \mid\) & \(7 \mid\) & \(-1 \mid\) & \(0 \mid\) & 1 \\
\(3 \mid\) & \(1 \mid\) & \(6 \mid\) & \(10 \mid\) & \(6 \mid\) & \(4 \mid\) & \(1 \mid\) & 0 \\
\(4 \mid\) & \(2 \mid\) & \(6 \mid\) & \(10 \mid\) & \(7 \mid\) & \(8 \mid\) & \(1 \mid\) & 1 \\
\(5 \mid\) & \(3 \mid\) & \(6 \mid\) & \(10 \mid\) & \(11 \mid\) & \(9 \mid\) & \(1 \mid\) & 2 \\
\(6 \mid\) & \(4 \mid\) & \(6 \mid\) & \(10 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) & 3 \\
\(7 \mid\) & \(5 \mid\) & \(6 \mid\) & \(10 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) & 4 \\
\(8 \mid\) & \(6 \mid\) & \(6 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 5 \\
\(9 \mid\) & \(1 \mid\) & \(12 \mid\) & \(10 \mid\) & \(12 \mid\) & \(13 \mid\) & \(1 \mid\) & 0 \\
\(10 \mid\) & \(2 \mid\) & \(12 \mid\) & \(10 \mid\) & \(17 \mid\) & \(15 \mid\) & \(1 \mid\) & 1 \\
\(11 \mid\) & \(3 \mid\) & \(12 \mid\) & \(10 \mid\) & \(16 \mid\) & \(16 \mid\) & \(1 \mid\) & 2 \\
\(12 \mid\) & \(4 \mid\) & \(12 \mid\) & \(10 \mid\) & \(15 \mid\) & \(3 \mid\) & \(1 \mid\) & 3 \\
\(13 \mid\) & \(5 \mid\) & \(12 \mid\) & \(10 \mid\) & \(10 \mid\) & \(-1 \mid\) & \(0 \mid\) & 4 \\
\((13\) rows) & & & & &
\end{tabular}

\section*{See Also}
- Sample Data
- https://en.wikipedia.org/wiki/Shortest_Path_Faster_Algorithm

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2
pgr_isPlanar - Experimental
pgr_isPlanar - Returns a boolean depending upon the planarity of the graph.

Boost Graph Inside

\section*{Warning}

Possible server crash
- These functions might create a server crash

\section*{Warning}

Experimental functions
- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need \(c / c++\) coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting
- Version 3.2.0
- New experimental function

\section*{Description}

A graph is planar if it can be drawn in two-dimensional space with no two of its edges crossing. Such a drawing of a planar graph is called a plane drawing. Every planar graph also admits a straight-line drawing, which is a plane drawing where each edge is represented by a line segment. When a graph has \(\backslash\left(K \_5 \backslash\right)\) or \(\backslash\left(K \_\{3,3\} \backslash\right)\) as subgraph then the graph is not planar.

The main characteristics are:
- This implementation use the Boyer-Myrvold Planarity Testing.
- It will return a boolean value depending upon the planarity of the graph.
- Applicable only for undirected graphs.
- The algorithm does not considers traversal costs in the calculations.
- Running time: \(\backslash(\mathrm{O}(|\mathrm{V}|) \backslash)\)

\section*{Signatures}

\section*{Summary}
```

pgr_isPlanar(Edges SQL)

```
RETURNS BOOLEAN
```

SELECT * FROM pgr_isPlanar(
'SELECT id, source, target, cost, reverse_cost
FROM edges
);
pgr_isplanar
t
(1 row)

```

\section*{Parameters}
\begin{tabular}{lll} 
Parameter & Type & Description \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL as described \\
below.
\end{tabular}
\end{tabular}

\section*{Inner Queries}

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Description \\
\hline source & ANY-INTEGER & Identifier of the edge. \\
\hline target & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Identifier of the second end point vertex of the edge. \\
\hline reverse_cost & ANY-NUMERICAL & -1 \\
& & Weight of the edge (source, target) \\
& & Weight of the edge (target, source) \\
& & \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& &
\end{tabular}

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns a boolean (pgr_isplanar)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline pgr_isplanar & BOOLEAN & \begin{tabular}{l} 
true when the graph is planar. \\
false when the graph is not \\
planar.
\end{tabular}
\end{tabular}

\section*{Additional Examples}

The following edges will make the subgraph with vertices \(\{10,15,11,16,13\}\) al(K_1 \()\) graph.
```

INSERT INTO edges (source, target, cost, reverse_cost) VALUES
(10, 16, 1, 1), (10, 13, 1, 1),
(15, 11, 1, 1), (15, 13, 1, 1),
(11, 13, 1, 1), (16, 13, 1, 1)
INSERT 0 6

```

The new graph is not planar because it has \(a \backslash(\mathrm{~K}, 5 \backslash)\) subgraph. Edges in blue represent \(\backslash\left(\mathrm{K} \_5 \backslash\right)\) subgraph.

```

SELECT * FROM pgr_isPlanar(
'SELECT id, source, target, cost, reverse_cost
FROM edges');
pgr_isplanar
f
(1 row)

```

\section*{See Also}
- Sample Data
- https://www.boost.org/libs/graph/doc/boyer_myrvold.html

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1) 3.0
pgr_stoerWagner - Experimental
pgr_stoerWagner - The min-cut of graph using stoerWagner algorithm.

\section*{8 boost}

Boost Graph Inside

\section*{Warning}

\section*{Possible server crash}
- These functions might create a server crash

\section*{Experimental functions}
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- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
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- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting

\section*{Availability}
- Version 3.0
- New Experimental function

\section*{Description}

In graph theory, the Stoer-Wagner algorithm is a recursive algorithm to solve the minimum cut problem in undirected weighted graphs with non-negative weights. The essential idea of this algorithm is to shrink the graph by merging the most intensive vertices, until the graph only contains two combined vertex sets. At each phase, the algorithm finds the minimum s-t cut for two vertices \(s\) and \(t\) chosen as its will. Then the algorithm shrinks the edge between \(s\) and \(t\) to search for non s-t cuts. The minimum cut found in all phases will be the minimum weighted cut of the graph.

A cut is a partition of the vertices of a graph into two disjoint subsets. A minimum cut is a cut for which the size or weight of the cut is not larger than the size of any other cut. For an unweighted graph, the minimum cut would simply be the cut with the least edges. For a weighted graph, the sum of all edges' weight on the cut determines whether it is a minimum cut.

\section*{The main characteristics are:}

Process is done only on edges with positive costs.
- It's implementation is only on undirected graph.
- Sum of the weights of all edges between the two sets is mincut.
- A mincut is a cut having the least weight.
- Values are returned when graph is connected.
- When there is no edge in graph then EMPTY SET is return.
- When the graph is unconnected then EMPTY SET is return.
- Sometimes a graph has multiple min-cuts, but all have the same weight. The this function determines exactly one of the min-cuts as well as its weight.
- Running time: \(\backslash\left(\mathrm{O}\left(\mathrm{V}^{*} \mathrm{E}+\mathrm{V}^{\wedge} 2^{*} \log \mathrm{~V}\right) \backslash\right)\).

\section*{Signatures}
```

pgr_stoerWagner(Edges SQL)

```
```

RETURNS SET OF (seq, edge, cost, mincut)

```
OR EMPTY SET

\section*{Example:}
min cut of the main subgraph
```

SELECT * FROM pgr_stoerWagner(
'SELECT id, source, target, cost, reverse_cost
FROM edges WHERE id < 17');
seq | edge | cost | mincut
1| 6| 1| 1
(1 row)

```

\section*{Parameter Type Description}

Edges SQL TEXT Edges SQL as described below.

\section*{Inner Queries}

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL & -1
\end{tabular}
- When negative: edge (target, source) does not exist, therefore it's not part of the graph.

\section*{Where:}

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT
ANY-NUMERICAL:
SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns set of (seq, edge, cost, mincut)
\begin{tabular}{lll} 
Column & Type & Description \\
\hline seq & INT & Sequential value starting from \(\mathbf{1 .}\) \\
\hline edge & BIGINT & \begin{tabular}{l} 
Edges which divides the set of vertices into \\
two.
\end{tabular} \\
\hline cost & FLOAT & Cost to traverse of edge. \\
\hline mincut & FLOAT & Min-cut weight of a undirected graph. \\
\hline
\end{tabular}

\section*{Additional Example:}

\section*{Example:}
min cut of an edge
```

SELECT * FROM pgr_stoerWagner(
'SELECT id, source, target, cost, reverse_cost
FROM edges WHERE id = 18');
seq | edge | cost | mincut
1| 18| 1|-----------+------
(1 row)

```

\section*{Example:}

Using pgr_connectedComponents
```

SELECT * FROM pgr_stoerWagner(

$$
SELECT id, source, target, cost, reverse_cost FROM edges
WHERE source IN (
    SELECT node FROM pgr_connectedComponents(
    'SELECT id, source, target, cost, reverse_cost FROM edges ')
    WHERE component = 2)
$$

seq | edge | cost | mincut
------------------------
(1 row)

```

\section*{See Also}

\section*{Indices and tables}
- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0

\section*{pgr_topologicalSort - Experimental}
pgr_topologicalSort - Linear ordering of the vertices for directed acyclic graphs (DAG).

\section*{boost}

Boost Graph Inside

\section*{Warning}

Possible server crash
- These functions might create a server crash

\section*{Warning}

Experimental functions
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- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need c/c++ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
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- Might depend on a deprecated function of pgRouting

\section*{Availability}
- Version 3.0.0
- New experimental function

\section*{Description}

The topological sort algorithm creates a linear ordering of the vertices such that if edgel((u,v)\\) appears in the graph, then \(\backslash(v \backslash)\) comes before \(\backslash(u \backslash)\) in the ordering.

The main characteristics are:
- Process is valid for directed acyclic graphs only. otherwise it will throw warnings.
- For optimization purposes, if there are more than one answer, the function will return one of them.
- The returned values are ordered in topological order:
- Running time: \(\backslash(O(V+E) \backslash)\)

\section*{Signatures}

\section*{Summary}
```

RETURNS SET OF (seq, sorted_v)
OR EMPTY SET

```

\section*{Example:}

Topologically sorting the graph
```

SELECT * FROM pgr_topologicalsort(

$$
SELECT id, source, target, cost
    FROM edges WHERE cost >= 0
    UNION
    SELECT id, target, source, reverse_cost
    FROM edges WHERE cost < 0
$$);

seq| sorted_v
1| 1
2| 5
2
4
3
13
14
15
|
10|
11|}
12|}
13|}
14 11
|}1
16| 12
17| 17
(17 rows)

```

\section*{Parameters}
\begin{tabular}{lllll} 
Parameter & Type & Description \\
\hline Edges SQL TEXT & \begin{tabular}{l} 
Edges SQL \\
below.
\end{tabular} & as described
\end{tabular}

Inner Queries

Edges SQL
\begin{tabular}{lll} 
Column & Type & Default \\
\hline id & ANY-INTEGER & Identifier of the edge. \\
\hline source & ANY-INTEGER & Identifier of the first end point vertex of the edge. \\
\hline target & ANY-INTEGER & Identifier of the second end point vertex of the edge. \\
\hline cost & ANY-NUMERICAL & Weight of the edge (source, target) \\
\hline reverse_cost & ANY-NUMERICAL -1 & Weight of the edge (target, source) \\
& & \\
& & \\
& & When negative: edge (target, source) does not exist, therefore it's \\
& & \\
& &
\end{tabular}

Where:

\section*{ANY-INTEGER:}

SMALLINT, INTEGER, BIGINT

\section*{ANY-NUMERICAL:}

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

\section*{Result Columns}

Returns set of (seq, sorted_v)
\begin{tabular}{lllll} 
Column & Type & Description & \\
\hline seq & INTEGER & Sequential value starting from \\
(1)) & \\
\hline sorted_v & BIGINT & \begin{tabular}{l} 
Linear \\
vertices
\end{tabular} & &
\end{tabular}

\section*{Additional examples}

\section*{Example:}

Topologically sorting the one way segments
```

SELECT * FROM pgr_topologicalsort(

$$
SELECT id, source, target, cost, -1 AS reverse_cost
FROM edges WHERE cost >= 0
UNION
SELECT id, source, target, -1, reverse_cost
FROM edges WHERE cost < 0
$$);

seq| sorted_v
1| 5
2| 2
4| 13
5| 14
6| 1
8| 15
9| 10
10| 6
11|}
12| 8
13|}
14| 11
15| 12
16| 16
17| 17
(17 rows)

```

\section*{Example:}

Graph is not a DAG
```

SELECT * FROM pgr_topologicalsort(

$$
SELECT id, source, target, cost, reverse_cost FROM edges
$$);
ERROR: The graph must be a DAG
HINT: Working with Directed Graph
CONTEXT: SQL function "pgr_topologicalsort" statement 1
```

## See Also

- Sample Data
- https://en.wikipedia.org/wiki/Topological_sorting


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.23 .13 .0
pgr_transitiveClosure - Experimental
pgr_transitiveClosure - Transitive closure graph of a directed graph.


## boost

Boost Graph Inside

## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

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- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.0.0
- New experimental function


## Description

Transforms the input directed graph into the transitive closure of the graph.
The main characteristics are:

- Process is valid for directed graphs.
- The transitive closure of an undirected graph produces a cluster graph
- Reachability between vertices on an undirected graph happens when they belong to the same connected component. (see pgr_connectedComponents)
- The returned values are not ordered
- The returned graph is compresed
- Running time: $\backslash(\mathrm{O}(|\mathrm{V}||E|) \backslash)$


## Signatures

## Summary

The pgr_transitiveClosure function has the following signature:

```
pgr_transitiveClosure(Edges SQL)
RETURNS SET OF (seq, vid, target_array)
```


## Example:

Rechability of a subgraph

```
SELECT * FROM pgr_transitiveclosure(
    'SELECT id, source, target, cost, reverse_cost
    FROM edges WHERE id IN (2, 3, 5, 11, 12, 13, 15)')
ORDER BY vid;
seq|vid | target_array
    1| 6|{}
    8|{12,17,16}
    2| 10|{12,17,16,11,6}
    4| 11 |{12,17,16}
    5| 12|{17,16}
    3| 15|{12,17,16,10,11,6}
    8| 16|{17,16}
    7| 17|{17,16}
(8 rows)
```


## Parameters

| Parameter | Type | Description |
| :--- | :--- | :--- | :--- | :--- |
| Edges SQL TEXT | Edges SQL as described <br> below. |  |

## Inner Queries

Edges SQL

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER |  | Identifier of the edge. |


| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| source | ANY-INTEGER |  | Identifier of the first end point vertex of the edge. |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |  |
| cost | ANY-NUMERICAL | Weight of the edge (source, target) |  |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |
|  |  |  | When negative: edge (target, source) does not exist, therefore it's <br>  |
|  |  |  |  |
|  |  |  |  |

Where:

## ANY-INTEGER: <br> SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

RETURNS SET OF (seq, vid, target_array)

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Sequential value starting from $\backslash(1 \backslash)$ |
| vid | BIGINT | Identifier of the source of the edges |
| target_array | BIGINT | Identifiers of the targets of the edges |
|  |  | a  <br>   <br>   <br>   |

## See Also

- Sample Data
- https://en.wikipedia.org/wiki/Transitive_closure


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2 3.1 3.0
pgr_turnRestrictedPath - Experimental
pgr_turnRestrictedPath Using Yen's algorithm Vertex -Vertex routing with restrictions


## Warning

Possible server crash

- These functions might create a server crash


## Warning

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting


## Availability

Version 3.0.0

- New experimental function


## Description

Using Yen's algorithm to obtain K shortest paths and analyze the paths to select the paths that do not use the restrictions

## Signatures

```
pgr_turnRestrictedPath(Edges SQL, Restrictions SQL, start vid, end vid, K, [options])
options: [directed, heap_paths, stop_on_first, strict]
RETURNS SET OF (seq, path_id, path_seq, node, edge, cost, agg_cost)
OR EMPTY SET
```


## Example:

From vertex $\backslash(3 \backslash)$ to vertex $\backslash(8 \backslash)$ on a directed graph

```
SELECT * FROM pgr_turnRestrictedPath(
$$SELECT id, source, target, cost, reverse_cost FROM edges$$,
$$SELECT path, cost FROM new_restrictions$$,
3, 8, 3);
seq | path_id | path_seq | node | edge | cost | agg_cost
\begin{tabular}{rrrrrrr}
\(1 \mid\) & \(1 \mid\) & \(1 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) Infinity \\
\(2 \mid\) & \(1 \mid\) & \(2 \mid\) & \(7 \mid\) & \(10 \mid\) & \(1 \mid\) & 1 \\
\(3 \mid\) & \(1 \mid\) & \(3 \mid\) & \(8 \mid\) & \(-1 \mid\) & \(0 \mid\) & 2
\end{tabular}
```


## Parameters

| Column | Type | Description |
| :--- | :--- | :--- |
| Edges SQL | TEXT | SQL query as described. |
| start vid | ANY-INTEGER | Identifier of the departure vertex. |
| end vid | ANY-INTEGER | Identifier of the departure vertex. |
| K | ANY-INTEGER | Number of required paths |

Where:

## ANY-INTEGER:

```
SMALLINT, INTEGER, BIGINT
```

Optional parameters


## KSP Optional parameters

Column Type Default Description
heap_paths BOOLEAN false - When false Returns at most K paths

- When true all the calculated paths while processing are returned.
- Roughly, when the shortest path has N edges, the heap will contain about thanN * K paths for small value of $K$ and $K>5$.


## Special optional parameters

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| stop_on_first | BOOLEAN true | When true stops on first path found that dos not violate <br> restrictions <br> - |  |
|  |  | When false returns at most K paths |  |

Edges SQL

| Column | Type | Default |
| :--- | :--- | :--- |
| id | ANY-INTEGER | Description |
| source | ANY-INTEGER | Identifier of the edge. |
| target | ANY-INTEGER | Identifier of the first end point vertex of the edge. |
| cost | ANY-NUMERICAL | Identifier of the second end point vertex of the edge. |
| reverse_cost | ANY-NUMERICAL | -1 |
|  |  | Weight of the edge (source, target) |
|  |  |  |
|  |  | Weight of the edge (target, source) |
|  |  |  |
|  |  | When negative: edge (target, source) does not exist, therefore it's |
|  |  |  |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

Restrictions SQL

| Column | Type | Description |
| :--- | :--- | :--- |
| path | ARRAY [ANY-INTEGER] | Sequence of edge identifiers that form a path that is not allowed to be taken. - <br>  <br>  <br> Empty arrays or NULL arrays are ignored. - Arrays that have aNULL element will <br> raise an exception. |
| Cost | ANY-NUMERICAL | Cost of taking the forbidden path. |

Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

Returns set of (seq, path_id, path_seq, start_vid, end_vid, node, edge, cost, agg_cost)

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Sequential value starting from $\mathbf{1 .}$ |
| path_id | INTEGER | Path identifier. |
|  |  | Has value $\mathbf{1}$ for the first of a path fromstart_vid to end_vid. |
| path_seq | INTEGER | Relative position in the path. Has value $\mathbf{1}$ for the beginning of a path. |
| start_vid | BIGINT | Identifier of the starting vertex. |
| end_vid | BIGINT | Identifier of the ending vertex. |
| node | BIGINT | Identifier of the node in the path fromstart_vid to end_vid. |
| edge | FLOAT | Identifier of the edge used to go fromnode to the next node in the path sequence. <br>  <br> the last node of the path. <br> cost |
| FLOAT | Cost to traverse from node using edge to the next node in the path sequence. |  |

## Additional Examples

## Example:

From vertex $\backslash(3 \backslash)$ to $\backslash(8 \backslash)$ with strict flag on.
No results because the only path available follows a restriction.

```
SELECT * FROM pgr_turnRestrictedPath(
    $$SELECT id, source, target, cost, reverse_cost FROM edges$$,
    $$SELECT path, cost FROM new_restrictions$$,
    3, 8, 3,
    strict => true);
seq | path_id | path_seq | node | edge | cost | agg_cost
(0 rows)
```


## Example:

From vertex $\backslash(3 \backslash)$ to vertex $\backslash(8 \backslash)$ on an undirected graph

```
SELECT * FROM pgr_turnRestrictedPath(
$$SELECT id, source, target, cost, reverse cost FROM edges$$,
$$SELECT path, cost FROM new_restrictions$$,
3, 8, 3,
directed => false)
seq | path_id | path_seq | node | edge | cost | agg_cost
\begin{tabular}{cccccc}
\(1 \mid\) & \(1 \mid\) & \(1 \mid\) & \(3 \mid\) & \(7 \mid\) & \(1 \mid\) \\
\(2 \mid\) & \(1 \mid\) & \(2 \mid\) & \(7 \mid\) & \(4 \mid\) & \(1 \mid\) \\
\(3 \mid\) & \(1 \mid\) & \(3 \mid\) & \(6 \mid\) & \(2 \mid\) & \(1 \mid\) \\
3 & 2 \\
\(4 \mid\) & \(1 \mid\) & \(4 \mid\) & \(10 \mid\) & \(5 \mid\) & \(1 \mid\) \\
\(5 \mid\) & \(1 \mid\) & \(5 \mid\) & \(11 \mid\) & \(11 \mid\) & \(1 \mid\) \\
\(6 \mid\) & \(1 \mid\) & \(6 \mid\) & \(12 \mid\) & \(12 \mid\) & \(1 \mid\) \\
\(7 \mid\) & \(1 \mid\) & \(7 \mid\) & \(8 \mid\) & \(-1 \mid\) & \(0 \mid\) \\
\((7\) rows \()\)
\end{tabular}
```


## Example:

From vertex $\backslash(3 \backslash)$ to vertex $\backslash(8 \backslash)$ with more alternatives

```
SELECT * FROM pgr_turnRestrictedPath(
$$SELECT id, source, target, cost, reverse_cost FROM edges$$,
$$SELECT path, cost FROM new restrictions$$,
3, 8, 3,
directed => false,
heap_paths => true,
stop_on_first => false);
seq | path_id | path_seq | node | edge | cost | agg_cost
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline 1| & 1| & 1| & 31 & 7| & 1| & 0 \\
\hline \(2 \mid\) & 1| & 21 & 71 & 4| & 1| & 1 \\
\hline 31 & 1| & 31 & 61 & 21 & 1| & 2 \\
\hline 4 | & 1 | & 4 | & 10 & 5| & 1| & 3 \\
\hline 51 & \(1 \mid\) & 51 & 11| & 11| & 1| & 4 \\
\hline 61 & 1| & 61 & 12| & 12| & \(1 \mid\) & 5 \\
\hline 71 & 1| & 71 & 8| & -1| & 01 & 6 \\
\hline 81 & 21 & 1| & 31 & 71 & 1| & 0 \\
\hline 91 & 21 & 21 & 71 & 8। & 1| & 1 \\
\hline 10| & \(2 \mid\) & \(3 \mid\) & 11 & 9| & 1| & 2 \\
\hline 11| & \(2 \mid\) & 4 | & 16 & 15 & \(1 \mid\) & 3 \\
\hline 12 | & \(2 \mid\) & 5 & 17 & 13 & 1| & 4 \\
\hline 13| & \(2 \mid\) & \(6 \mid\) & 12 & \(12 \mid\) & 1| & 5 \\
\hline 14| & 21 & \(7 \mid\) & 8। & -1| & 0| & 6 \\
\hline
\end{tabular}
(14 rows)
```


## See Also

```
- K shortest paths - Category
- Sample Data
```


## Indices and tables

- Index
- Search Page
- Supported versions: Latest (3.3) 3.2
pgr_lengauerTarjanDominatorTree -Experimental
pgr_lengauerTarjanDominatorTree - Returns the immediate dominator of all vertices.


## boost

Boost Graph Inside

## Warning

Possible server crash

These functions might create a server crash

Experimental functions

- They are not officially of the current release.
- They likely will not be officially be part of the next release:
- The functions might not make use of ANY-INTEGER and ANY-NUMERICAL
- Name might change.
- Signature might change.
- Functionality might change.
- pgTap tests might be missing.
- Might need $c / c++$ coding.
- May lack documentation.
- Documentation if any might need to be rewritten.
- Documentation examples might need to be automatically generated.
- Might need a lot of feedback from the comunity.
- Might depend on a proposed function of pgRouting
- Might depend on a deprecated function of pgRouting


## Availability

- Version 3.2.0
- New experimental function


## Description

The algorithm calculates the immidiate dominator of each vertex calledidom, once idom of each vertex is calculated then by making every idom of each vertex as its parent, the dominator tree can be built.

## The main Characteristics are:

- The algorithm works in directed graph only.
- The returned values are not ordered.
- The algorithm returns idom of each vertex.
- If the root vertex not present in the graph then it returns empty set.
- Running time: $\backslash(\mathrm{O}((\mathrm{V}+\mathrm{E}) \log (\mathrm{V}+\mathrm{E})) \backslash)$


## Signatures

## Summary

```
pgr_lengauerTarjanDominatorTree(Edges SQL, root vertex)
RETURNS SET OF (seq, vertex_id, idom)
OR EMPTY SET
```


## Example:

The dominator tree with root vertex $\backslash(5 \backslash)$

```
SELECT * FROM pgr_lengauertarjandominatortree(
    $$SELECT id,source,target,cost,reverse_cost FROM edges$$,
    5) ORDER BY vertex_id;
seq | vertex_id | idom
    1| 2
    2| 0
```



```
    4| 0
    5| 0
    6| 17
    7| 4
    8| 3
    9| 7
    10| 16
    11| 3
    12| 3
    13| 0
    14| 0
    5| 15
    16|}
14| 17 | 3
(17 rows)
```

| Column | Type | Description |
| :--- | :--- | :--- | :--- |
| Edges SQL | TEXT | SQL query as described above. |
| root vertex | BIGINT | Identifier of the starting <br> vertex. |

## Inner Queries

Edges SQL

| Column | Type | Default | Description |
| :--- | :--- | :--- | :--- |
| id | ANY-INTEGER |  | Identifier of the edge. |
| source | ANY-INTEGER | Identifier of the first end point vertex of the edge. |  |
| target | ANY-INTEGER | Identifier of the second end point vertex of the edge. |  |
| cost | ANY-NUMERICAL |  | Weight of the edge (source, target) |
| reverse_cost | ANY-NUMERICAL | -1 | Weight of the edge (target, source) |

- When negative: edge (target, source) does not exist, therefore it's not part of the graph.


## Where:

## ANY-INTEGER:

SMALLINT, INTEGER, BIGINT

## ANY-NUMERICAL:

SMALLINT, INTEGER, BIGINT, REAL, FLOAT

## Result Columns

Returns set of (seq, vertex_id, idom)

| Column | Type | Description |
| :--- | :--- | :--- |
| seq | INTEGER | Sequential value starting from |
|  |  | $\mathbf{1 .}$ |
| vertex_id | BIGINT | Identifier of vertex . |
| idom | BIGINT | Immediate dominator of vertex. |

## Additional Examples

## Example:

Dominator tree of another component.

```
SELECT * FROM pgr_lengauertarjandominatortree(
    $$SELECT id,source,target,cost,reverse_cost FROM edges$$,
    13) ORDER BY vertex id;
seq | vertex id | idom
    1| 1| 0
    9| 2| 0
    2| 3| 0
10| 4| 0
17| 5| 0
4| 6| 0
    7| 0
    8| 0
    9| 0
    10| 0
    11| 0
    12| 0
    13| 0
    14| 12
    15| 0
    | 16| 0
    14| 17| 0
(17 rows)
```


## See Also

[^3]- Index
- Search Page

See Also

Indices and tables

- Index
- Search Page


## Release Notes

- Supported versions: Latest (3.3) 3.2 3.1 3.0
- Unsupported versions: $2.6 \mathbf{2 . 5} \mathbf{2 . 4} \mathbf{2 . 3} \mathbf{2 . 2} 2.12 .0$


## Release Notes

To see the full list of changes check the list of Git commits on Github.

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```
pgRouting 2.0.1 Release Notes
pgRouting 2.0.0 Release Notes
pgRouting 1.x Release Notes
- Changes for release 1.05
- Changes for release 1.03
- Changes for release }\mathbf{1.02
- Changes for release 1.01
- Changes for release 1.0
- Changes for release 1.0.0b
- Changes for release 1.0.0a
- Changes for release 0.9.9
- Changes for release 0.9.8
```


## pgRouting 3.3.4 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.3.4

## Issue fixes

- \#2400: pgRouting 3.3.3 does not build in focal


## pgRouting 3.3.3 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.3.3

## Issue fixes

- \#1891: pgr_ksp doesn't give all correct shortest path


## Official functions changes

## - Flow functions

- pgr_maxCardinalityMatch(text,boolean)
- Ignoring optional boolean parameter, as the algorithm works only for undirected graphs.


## pgRouting 3.3.2 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.3.2

- Revised documentation
- Simplifying table names and table columns, for example:
- edges instead of edge_table
- Removing unused columns category_id and reverse_category_id.
- combinations instead of combinations_table
- Using PostGIS standard for geometry column.
- geom instead of the_geom
- Avoiding usage of functions that modify indexes, columns etc on tables.
- Using pgr_extractVertices to create a routing topology
- Restructure of the pgRouting concepts page.


## Issue fixes

- \#2276: edgeDisjointPaths issues with start_vid and combinations
- \#2312: pgr_extractVertices error when target is not BIGINT
- \#2357: Apply clang-tidy performance-*


## pgRouting 3.3.1 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.3.1 on Github.

## Issue fixes

- \#2216: Warnings when using clang


## pgRouting 3.3.0 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.3.0 on Github.

## Issue fixes

- \#2057: trspViaEdges columns in different order
- \#2087: pgr_extractVertices to proposed
- \#2201: pgr_depthFirstSearch to proposed
- \#2202: pgr_sequentialVertexColoring to proposed
- \#2203: pgr_dijkstraNear and pgr_dijkstraNearCost to proposed


## New experimental functions

- Coloring
- pgr_edgeColoring


## Experimental promoted to Proposed

- Dijkstra
- pgr_dijkstraNear
- pgr_dijkstraNear(Combinations)
- pgr_dijkstraNear(Many to Many)
- pgr_dijkstraNear(Many to One)
- pgr_dijkstraNear(One to Many)
- pgr_dijkstraNearCost
- pgr_dijkstraNearCost(Combinations)
- pgr_dijkstraNearCost(Many to Many)
- pgr_dijkstraNearCost(Many to One)
- pgr_dijkstraNearCost(One to Many)
- Coloring
- pgr_sequentialVertexColoring
- Topology
- pgr_extractVertices
- Traversal
- pgr_depthFirstSearch(Multiple vertices)
- pgr_depthFirstSearch(Single vertex)


## pgRouting 3.2.2 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.2.2 on Github.

## Issues

- \#2093: Compilation on Visual Studio
- \#2189: Build error on RHEL 7


## pgRouting 3.2.1 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.2.1 on Github.

## Issue fixes

- \#1883: pgr_TSPEuclidean crashes connection on Windows
- The solution is to use Boost::graph::metric_tsp_approx
- To not break user's code the optional parameters related to the TSP Annaeling are ignored
- The function with the annaeling optional parameters is deprecated


## pgRouting 3.2.0 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.2.0 on Github.

## Build

- \#1850: Change Boost min version to 1.56


## New experimental functions

```
- pgr_bellmanFord(Combinations)
p pgr_binaryBreadthFirstSearch(Combinations)
 pgr_bipartite
- pgr_dagShortestPath(Combinations)
- pgr_depthFirstSearch
D Dijkstra Near
- pgr dijkstraNear
    - pgr_dijkstraNear(One to Many)
    - pgr_dijkstraNear(Many to One)
    - pgr_dijkstraNear(Many to Many)
    - pgr_dijkstraNear(Combinations)
- pgr dijkstraNearCost
    - pgr_dijkstraNearCost(One to Many)
    - pgr_dijkstraNearCost(Many to One)
    - pgr_dijkstraNearCost(Many to Many)
    - pgr dijkstraNearCost(Combinations)
- pgr_edwardMoore(Combinations)
- pgr_isPlanar
- pgr_lengauerTarjanDominatorTree
pgr_makeConnected
Flow
    - pgr_maxFlowMinCost(Combinations)
    - pgr maxFlowMinCost Cost(Combinations)
- pgr sequentialVertexColoring
```


## New proposed functions

- Astar
- pgr_aStar(Combinations)
- pgr_aStarCost(Combinations)
- Bidirectional Astar
- pgr_bdAstar(Combinations)
- pgr_bdAstarCost(Combinations)
- Bidirectional Dijkstra
- pgr_bdDijkstra(Combinations)
- pgr_bdDijkstraCost(Combinations)
- Flow
- pgr_boykovKolmogorov(Combinations)
- pgr_edgeDisjointPaths(Combinations)
- pgr_edmondsKarp(Combinations)
- pgr_maxFlow(Combinations)
- pgr_pushRelabel(Combinations
- pgr_withPoints(Combinations)
- pgr_withPointsCost(Combinations)


## pgRouting 3.1.4 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.1.4 on Github.

## Issues fixes

- 2189: Build error on RHEL 7


## pgRouting 3.1.3 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.1.3 on Github.

## Issues fixes

- \#1825: Boost versions are not honored
- \#1849: Boost 1.75 .0 geometry "point_xy.hpp" build error on macOS environment
- \#1861: vrp functions crash server

To see all issues \& pull requests closed by this release see theGit closed milestone for $\mathbf{3 . 1 . 2}$ on Github.

## Issues fixes

- \#1304: FreeBSD 12 64-bit crashes on pgr_vrOneDepot tests Experimental Function
- \#1356: tools/testers/pg_prove_tests.sh fails when PostgreSQL port is not passed
- \#1725: Server crash on pgr_pickDeliver and pgr_vrpOneDepot on openbsd
- \#1760: TSP server crash on ubuntu 20.04 \#1760
- \#1770: Remove warnings when using clang compiler


## pgRouting 3.1.1 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.1.1 on Github.

## Issues fixes

- \#1733: pgr_bdAstar fails when source or target vertex does not exist in the graph
- \#1647: Linear Contraction contracts self loops
- \#1640: pgr_withPoints fails when points_sql is empty
- \#1616: Path evaluation on C++ not updated before the results go back to C
- \#1300: pgr_chinesePostman crash on test data


## pgRouting 3.1.0 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for $\mathbf{3 . 1 . 0}$ on Github.

## New proposed functions

- pgr_dijkstra(combinations)
- pgr_dijkstraCost(combinations)


## Build changes

- Minimal requirement for Sphinx: version 1.8
pgRouting 3.0.6 Release Notes
To see all issues \& pull requests closed by this release see theGit closed milestone for 3.0.6 on Github.


## Issues fixes

- \#2189: Build error on RHEL 7


## pgRouting 3.0.5 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.0.5 on Github.

## Backport issues fixes

- \#1825: Boost versions are not honored
- \#1849: Boost 1.75 .0 geometry "point_xy.hpp" build error on macOS environment
- \#1861: vrp functions crash server


## pgRouting 3.0.4 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.0.4 on Github.

## Backport issues fixes

\#1304: FreeBSD 12 64-bit crashes on pgr_vrOneDepot tests Experimental Function

- \#1356: tools/testers/pg_prove_tests.sh fails when PostgreSQL port is not passed
- \#1725: Server crash on pgr_pickDeliver and pgr_vrpOneDepot on openbsd
- \#1760: TSP server crash on ubuntu 20.04 \#1760
- \#1770: Remove warnings when using clang compiler


## pgRouting 3.0.3 Release Notes

## Backport issues fixes

- \#1733: pgr_bdAstar fails when source or target vertex does not exist in the graph
- \#1647: Linear Contraction contracts self loops
- \#1640: pgr_withPoints fails when points_sql is empty
- \#1616: Path evaluation on $\mathrm{C}++$ not updated before the results go back to C
- \#1300: pgr_chinesePostman crash on test data


## pgRouting 3.0.2 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.0.2 on Github.

## Issues fixes

- \#1378: Visual Studio build failing


## pgRouting 3.0.1 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for 3.0.1 on Github.

## Issues fixes

- \#232: Honor client cancel requests in C /C++ code


## pgRouting 3.0.0 Release Notes

To see all issues \& pull requests closed by this release see theGit closed milestone for $\mathbf{3 . 0 . 0}$ on Github.

## Fixed Issues

\#1153: Renamed pgr_eucledianTSP to pgr_TSPeuclidean

- \#1188: Removed CGAL dependency
- \#1002: Fixed contraction issues:
- \#1004: Contracts when forbidden vertices do not belong to graph
- \#1005: Intermideate results eliminated
- \#1006: No loss of information


## New functions

- Kruskal family
- pgr_kruskal
- pgr_kruskalBFS
- pgr_kruskalDD
- pgr_kruskalDFS
- Prim family
- pgr_prim
- pgr_primDD
- pgr_primDFS
- pgr_primBFS


## Proposed moved to official on pgRouting

- aStar Family
- pgr_aStar(one to many)
- pgr_aStar(many to one)
- pgr_aStar(many to many)
- pgr_aStarCost(one to one)
- pgr_aStarCost(one to many)
- pgr_aStarCost(many to one)
- pgr_aStarCost(many to many)
- pgr_aStarCostMatrix(one to one)
- pgr_aStarCostMatrix(one to many)
- pgr_aStarCostMatrix(many to one)
- pgr_aStarCostMatrix(many to many)
- bdAstar Family

```
pgr_bdAstar(one to many)
pgr_bdAstar(many to one)
pgr_bdAstar(many to many)
pgr_bdAstarCost(one to one)
pgr_bdAstarCost(one to many)
pgr_bdAstarCost(many to one)
pgr_bdAstarCost(many to many)
pgr_bdAstarCostMatrix(one to one)
pgr_bdAstarCostMatrix(one to many)
pgr_bdAstarCostMatrix(many to one)
pgr_bdAstarCostMatrix(many to many)
- bdDijkstra Family
    - pgr_bdDijkstra(one to many)
    pgr_bdDijkstra(many to one)
    pgr_bdDijkstra(many to many)
    pgr_bdDijkstraCost(one to one)
    pgr_bdDijkstraCost(one to many)
    pgr_bdDijkstraCost(many to one)
    pgr_bdDijkstraCost(many to many)
    pgr_bdDijkstraCostMatrix(one to one)
    pgr_bdDijkstraCostMatrix(one to many)
    pgr_bdDijkstraCostMatrix(many to one)
    pgr_bdDijkstraCostMatrix(many to many)
- Flow Family
    pgr_pushRelabel(one to one)
    pgr_pushRelabel(one to many)
    pgr_pushRelabel(many to one)
    pgr_pushRelabel(many to many)
    pgr_edmondsKarp(one to one)
    pgr_edmondsKarp(one to many)
    pgr_edmondsKarp(many to one)
    pgr_edmondsKarp(many to many)
    pgr_boykovKolmogorov (one to one)
    pgr_boykovKolmogorov (one to many)
    pgr_boykovKolmogorov (many to one)
    pgr_boykovKolmogorov (many to many)
    pgr_maxCardinalityMatching
    pgr_maxFlow
    pgr_edgeDisjointPaths(one to one)
    pgr_edgeDisjointPaths(one to many)
    pgr_edgeDisjointPaths(many to one)
    pgr_edgeDisjointPaths(many to many)
- Components family
    pgr_connectedComponents
    pgr_strongComponents
    pgr_biconnectedComponents
    pgr_articulationPoints
    pgr_bridges
- Contraction:
    - Removed unnecessary column seq
    Bug Fixes
```


## New Experimental functions

- pgr_maxFlowMinCost
- pgr_maxFlowMinCost_Cost
- pgr_extractVertices
- pgr_turnRestrictedPath
- pgr_stoerWagner
- pgr_dagShortestpath
- pgr_topologicalSort
- pgr_transitiveClosure
- VRP category
- pgr_pickDeliverEuclidean
pgr_pickDeliver
- Chinese Postman family
- pgr_chinesePostman
- pgr_chinesePostmanCost
- pgr_breadthFirstSearch
- pgr_binaryBreadthFirstSearch
- Bellman Ford family
- pgr_bellmanFord
- pgr_edwardMoore


## Moved to legacy

- Experimental functions
- pgr_labelGraph - Use the components family of functions instead.
- Max flow - functions were renamed on v2.5.0
- pgr_maxFlowPushRelabel
- pgr_maxFlowBoykovKolmogorov
- pgr_maxFlowEdmondsKarp - pgr_maximumcardinalitymatching
- VRP
- pgr_gsoc_vrppdtw
- TSP old signatures
- pgr_pointsAsPolygon
- pgr_alphaShape old signature


## pgRouting 2.6.3 Release Notes

To see the issues closed by this release see the Git closed milestone for 2.6.3 on Github.

## Bug fixes

- \#1219 Implicit cast for via_path integer to text
- \#1193 Fixed pgr_pointsAsPolygon breaking when comparing strings in WHERE clause
- \#1185 Improve FindPostgreSQL.cmake


## pgRouting 2.6.2 Release Notes

To see the issues closed by this release see theGit closed milestone for 2.6.2 on Github.

## Bug fixes

[^4]
## pgRouting 2.6.1 Release Notes

To see the issues closed by this release see the Git closed milestone for 2.6.1 on Github.

- Fixes server crash on several functions.
- pgr_floydWarshall
- pgr_johnson
- pgr_astar
- pgr_bdAstar
- pgr_bdDijstra
- pgr_alphashape
- pgr_dijkstraCostMatrix
- pgr_dijkstra
- pgr_dijkstraCost
- pgr_drivingDistance
- pgr_KSP
- pgr_dijkstraVia (proposed)
- pgr_boykovKolmogorov (proposed)
- pgr_edgeDisjointPaths (proposed)
- pgr_edmondsKarp (proposed)
- pgr_maxCardinalityMatch (proposed)
- pgr_maxFlow (proposed)
- pgr_withPoints (proposed)
- pgr_withPointsCost (proposed)
- pgr_withPointsKSP (proposed)
- pgr_withPointsDD (proposed)
- pgr_withPointsCostMatrix (proposed)
- pgr_contractGraph (experimental)
- pgr_pushRelabel (experimental)
- pgr_vrpOneDepot (experimental)
- pgr_gsoc_vrppdtw (experimental)
- Fixes for deprecated functions where also applied but not tested
- Removed compilation warning for g++8
- Fixed a fallthrugh on Astar and bdAstar.


## pgRouting 2.6.0 Release Notes

To see the issues closed by this release see the Git closed milestone for $\mathbf{2 . 6 . 0}$ on Github.

## New experimental functions

- pgr_lineGraphFull


## Bug fixes

- Fix pgr_trsp(text,integer,double precision,integer,double precision,boolean,boolean[,text])
- without restrictions
- calls pgr_dijkstra when both end points have a fraction $\operatorname{IN}(0,1)$
- calls pgr_withPoints when at least one fraction NOT IN $(0,1)$
- with restrictions
- calls original trsp code


## Internal code

- Cleaned the internal code of trsp(text,integer,integer,boolean,boolean [, text])
- Removed the use of pointers
- Internal code can accept BIGINT
- Cleaned the internal code of withPoints


## pgRouting 2.5.5 Release Notes

To see the issues closed by this release see theGit closed milestone for $\mathbf{2 . 5 . 5}$ on Github.

## Bug fixes

- Fixes driving distance when vertex is not part of the graph
- Fixes windows test
- Fixes build for python3 and perl5


## pgRouting 2.5.4 Release Notes

To see the issues closed by this release see the Git closed milestone for 2.5.4 on Github.

- Fixes server crash on several functions.
- pgr_floydWarshall
- pgr_johnson
- pgr_astar
- pgr_bdAstar
- pgr_bdDijstra
- pgr_alphashape
- pgr_dijkstraCostMatrix
- pgr_dijkstra
- pgr_dijkstraCost
- pgr_drivingDistance
- pgr_KSP
- pgr_dijkstraVia (proposed)
- pgr_boykovKolmogorov (proposed)
- pgr_edgeDisjointPaths (proposed)
- pgr_edmondsKarp (proposed)
- pgr_maxCardinalityMatch (proposed)
- pgr_maxFlow (proposed)
- pgr_withPoints (proposed)
- pgr_withPointsCost (proposed)
- pgr_withPointsKSP (proposed)
- pgr_withPointsDD (proposed)
- pgr_withPointsCostMatrix (proposed)
- pgr_contractGraph (experimental)
- pgr_pushRelabel (experimental)
- pgr_vrpOneDepot (experimental)
- pgr_gsoc_vrppdtw (experimental)
- Fixes for deprecated functions where also applied but not tested
- Removed compilation warning for g++8
- Fixed a fallthrugh on Astar and bdAstar.


## pgRouting 2.5.3 Release Notes

To see the issues closed by this release see the Git closed milestone for 2.5.3 on Github.

## Bug fixes

- Fix for postgresql 11: Removed a compilation error when compiling with postgreSQL


## pgRouting 2.5.2 Release Notes

To see the issues closed by this release see the Git closed milestone for 2.5.2 on Github.

## Bug fixes

- Fix for postgresql 10.1: Removed a compiler condition


## pgRouting 2.5.1 Release Notes

To see the issues closed by this release see the Git closed milestone for 2.5.1 on Github.

## Bug fixes

- Fixed prerequisite minimum version of: cmake


## pgRouting 2.5.0 Release Notes

To see the issues closed by this release see the Git closed issues for 2.5.0 on Github.

## enhancement:

- pgr_version is now on SQL language


## Breaking change on:

- pgr_edgeDisjointPaths:
- Added path_id, cost and agg_cost columns on the result
- Parameter names changed
- The many version results are the union of the one to one version


## New Signatures:

pgr_bdAstar(one to one)

## New Proposed functions

```
pgr_bdAstar(one to many)
pgr_bdAstar(many to one)
pgr_bdAstar(many to many)
- pgr_bdAstarCost(one to one)
- pgr_bdAstarCost(one to many)
pgr_bdAstarCost(many to one)
- pgr_bdAstarCost(many to many)
- pgr_bdAstarCostMatrix
- pgr_bdDijkstra(one to many)
- pgr_bdDijkstra(many to one)
```

pgr_bdDijkstra(many to many)

- pgr_bdDijkstraCost(one to one)
- pgr_bdDijkstraCost(one to many)
- pgr_bdDijkstraCost(many to one)
- pgr_bdDijkstraCost(many to many)
- pgr_bdDijkstraCostMatrix
. pgr_lineGraph
- pgr_lineGraphFull
- pgr_connectedComponents
- pgr_strongComponents
- pgr_biconnectedComponents
- pgr_articulationPoints
- pgr_bridges


## Deprecated Signatures

- pgr_bdastar - use pgr_bdAstar instead


## Renamed Functions

- pgr_maxFlowPushRelabel - use pgr_pushRelabel instead
- pgr_maxFlowEdmondsKarp -use pgr_edmondsKarp instead
- pgr_maxFlowBoykovKolmogorov - use pgr_boykovKolmogorov instead
- pgr_maximumCardinalityMatching - use pgr_maxCardinalityMatch instead


## Deprecated function

- pgr_pointToEdgeNode


## pgRouting 2.4.2 Release Notes

To see the issues closed by this release see theGit closed milestone for 2.4.2 on Github.

## Improvement

- Works for postgreSQL 10


## Bug fixes

- Fixed: Unexpected error column "cname"
- Replace __linux__ with __GLIBC__ for glibc-specific headers and functions


## pgRouting 2.4.1 Release Notes

To see the issues closed by this release see the Git closed milestone for 2.4.1 on Github.

## Bug fixes

- Fixed compiling error on macOS
- Condition error on pgr_withPoints


## pgRouting 2.4.0 Release Notes

To see the issues closed by this release see the Git closed issues for 2.4.0 on Github.

## New Signatures

- pgr_bdDijkstra


## New Proposed Signatures

- pgr_maxFlow
- pgr_astar(one to many)
- pgr_astar(many to one)
- pgr_astar(many to many)
- pgr_astarCost(one to one)
- pgr_astarCost(one to many)
- pgr_astarCost(many to one)
- pgr_astarCost(many to many)
- pgr_astarCostMatrix


## Deprecated Signatures

- pgr_bddijkstra - use pgr_bdDijkstra instead


## Deprecated Functions

- pgr_pointsToVids


## Bug fixes

- Bug fixes on proposed functions
- pgr_withPointsKSP: fixed ordering
- TRSP original code is used with no changes on the compilation warnings


## pgRouting 2.3.2 Release Notes

To see the issues closed by this release see the Git closed issues for 2.3.2 on Github.

## Bug Fixes

- Fixed pgr_gsoc_vrppdtw crash when all orders fit on one truck.
- Fixed pgr_trsp:
- Alternate code is not executed when the point is in reality a vertex
- Fixed ambiguity on seq


## pgRouting 2.3.1 Release Notes

To see the issues closed by this release see the Git closed issues for 2.3.1 on Github.

## Bug Fixes

- Leaks on proposed max_flow functions
- Regression error on pgr_trsp
- Types discrepancy on pgr_createVerticesTable


## pgRouting 2.3.0 Release Notes

To see the issues closed by this release see the Git closed issues for 2.3.0 on Github.

## New Signatures

- pgr_TSP
- pgr_aStar


## New Functions

- pgr_eucledianTSP


## New Proposed functions

pgr_dijkstraCostMatrix

- pgr_withPointsCostMatrix
- pgr_maxFlowPushRelabel(one to one)
- pgr_maxFlowPushRelabel(one to many)
- pgr_maxFlowPushRelabel(many to one)
- pgr_maxFlowPushRelabel(many to many)
- pgr_maxFlowEdmondsKarp(one to one)
- pgr_maxFlowEdmondsKarp(one to many)
- pgr_maxFlowEdmondsKarp(many to one)
- pgr_maxFlowEdmondsKarp(many to many)
- pgr_maxFlowBoykovKolmogorov (one to one)
- pgr_maxFlowBoykovKolmogorov (one to many)
- pgr_maxFlowBoykovKolmogorov (many to one)
- pgr_maxFlowBoykovKolmogorov (many to many)
- pgr_maximumCardinalityMatching
- pgr_edgeDisjointPaths(one to one)
- pgr_edgeDisjointPaths(one to many)
- pgr_edgeDisjointPaths(many to one)
- pgr_edgeDisjointPaths(many to many)
- pgr_contractGraph


## Deprecated Signatures

- pgr_tsp - use pgr_TSP or pgr_eucledianTSP instead
- pgr_astar - use pgr_aStar instead


## Deprecated Functions

- pgr_flip_edges
- pgr_vidsToDmatrix
- pgr_pointsToDMatrix
- pgr_textToPoints


## pgRouting 2.2.4 Release Notes

To see the issues closed by this release see the Git closed issues for 2.2.4 on Github.

## Bug Fixes

- Bogus uses of extern "C"
- Build error on Fedora 24 + GCC 6.0
- Regression error pgr_nodeNetwork


## pgRouting 2.2.3 Release Notes

To see the issues closed by this release see the Git closed issues for 2.2.3 on Github.

## Bug Fixes

- Fixed compatibility issues with PostgreSQL 9.6.


## pgRouting 2.2.2 Release Notes

To see the issues closed by this release see the Git closed issues for 2.2.2 on Github.

## Bug Fixes

- Fixed regression error on pgr_drivingDistance


## pgRouting 2.2.1 Release Notes

To see the issues closed by this release see the Git closed issues for 2.2.1 on Github.

## Bug Fixes

- Server crash fix on pgr_alphaShape
- Bug fix on With Points family of functions


## pgRouting 2.2.0 Release Notes

To see the issues closed by this release see the Git closed issues for 2.2.0 on Github.

## Improvements

- pgr_nodeNetwork
- Adding a row_where and outall optional parameters
- Signature fix
- pgr_dijkstra - to match what is documented


## New Functions

- pgr_floydWarshall
- pgr」ohnson
- pgr_dijkstraCost(one to one)
- pgr_dijkstraCost(one to many)
- pgr_dijkstraCost(many to one)
- pgr_dijkstraCost(many to many)


## Proposed functionality

- pgr_withPoints(one to one)
- pgr_withPoints(one to many)
- pgr_withPoints(many to one)
- pgr_withPoints(many to many)
- pgr_withPointsCost(one to one)
- pgr_withPointsCost(one to many)
- pgr_withPointsCost(many to one)
- pgr_withPointsCost(many to many)
- pgr_withPointsDD(single vertex)
- pgr_withPointsDD(multiple vertices)
- pgr_withPointsKSP
- pgr_dijkstraVia


## Deprecated functions:

pgr_apspWarshall use pgr_floydWarshall instead

- pgr_apspJohnson use pgr_Johnson instead
pgr_kDijkstraCost use pgr_dijkstraCost instead
- pgr_kDijkstraPath use pgr_dijkstra instead


## Renamed and deprecated function

- pgr_makeDistanceMatrix renamed to _pgr_makeDistanceMatrix


## pgRouting 2.1.0 Release Notes

To see the issues closed by this release see the Git closed issues for 2.1.0 on Github.

## New Signatures

- pgr_dijkstra(one to many)
- pgr_dijkstra(many to one)
- pgr_dijkstra(many to many)
- pgr_drivingDistance(multiple vertices)


## Refactored

- pgr_dijkstra(one to one)
- pgr_ksp
- pgr_drivingDistance(single vertex)


## Improvements

- pgr_alphaShape function now can generate better (multi)polygon with holes and alpha parameter.


## Proposed functionality

- Proposed functions from Steve Woodbridge, (Classified as Convenience by the author.)
- pgr_pointToEdgeNode - convert a point geometry to a vertex_id based on closest edge.
- pgr_flipEdges - flip the edges in an array of geometries so the connect end to end.
- pgr_textToPoints - convert a string of $x, y ; x, y ; \ldots$ locations into point geometries.
- pgr_pointsToVids - convert an array of point geometries into vertex ids.
- pgr_pointsToDMatrix - Create a distance matrix from an array of points.
- pgr_vidsToDMatrix - Create a distance matrix from an array of vertix_id.
- pgr_vidsToDMatrix - Create a distance matrix from an array of vertix_id.
- Added proposed functions from GSoc Projects:
- pgr_vrppdtw


## Deprecated functions

- pgr_getColumnName
- pgr_getTableName
- pgr_isColumnCndexed
- pgr_isColumnInTable
- pgr_quote_ident
- pgr_versionless
- pgr_startPoint
- pgr_endPoint
- pgr_pointTold


## No longer supported

- Removed the 1.x legacy functions


## Bug Fixes

- Some bug fixes in other functions


## Refactoring Internal Code

- A C and C++ library for developer was created
- encapsulates postgreSQL related functions
- encapsulates Boost.Graph graphs
- Directed Boost.Graph
- Undirected Boost.graph.
- allow any-integer in the id's
- allow any-numerical on the cost/reverse_cost columns
- Instead of generating many libraries: - All functions are encapsulated in one library - The library has the prefix 2-1-0


## pgRouting 2.0.1 Release Notes

Minor bug fixes.

## Bug Fixes

- No track of the bug fixes were kept.


## pgRouting 2.0.0 Release Notes

To see the issues closed by this release see the Git closed issues for $\mathbf{2 . 0 . 0}$ on Github.
With the release of pgRouting 2.0.0 the library has abandoned backwards compatibility topgRouting 1.x Release Notes releases. The main Goals for this release are:

- Major restructuring of pgRouting.
- Standardization of the function naming
- Preparation of the project for future development.

As a result of this effort:

- pgRouting has a simplified structure
- Significant new functionality has being added
- Documentation has being integrated
- Testing has being integrated
- And made it easier for multiple developers to make contributions.


## Important Changes

- Graph Analytics - tools for detecting and fixing connection some problems in a graph
- A collection of useful utility functions
- Two new All Pairs Short Path algorithms (pgr_apspJohnson, pgr_apspWarshall)
- Bi-directional Dijkstra and A-star search algorithms (pgr_bdAstar, pgr_bdDijkstra)
- One to many nodes search (pgr_kDijkstra)
- K alternate paths shortest path (pgr_ksp)
- New TSP solver that simplifies the code and the build process (pgr_tsp), dropped "Gaul Library" dependency
- Turn Restricted shortest path (pgr_trsp) that replaces Shooting Star
- Dropped support for Shooting Star
- Built a test infrastructure that is run before major code changes are checked in
- Tested and fixed most all of the outstanding bugs reported against 1.x that existing in the 2.0-dev code base.
- Improved build process for Windows
- Automated testing on Linux and Windows platforms trigger by every commit
- Modular library design
- Compatibility with PostgreSQL 9.1 or newer
- Compatibility with PostGIS 2.0 or newer
- Installs as PostgreSQL EXTENSION
- Return types re factored and unified
- Support for table SCHEMA in function parameters
- Support for st_ PostGIS function prefix
- Added pgr_ prefix to functions and types
- Better documentation: https://docs.pgrouting.org
- shooting_star is discontinued


## pgRouting 1.x Release Notes

To see the issues closed by this release see the Git closed issues for 1.x on Github. The following release notes have been copied from the previous RELEASE_NOTES file and are kept as a reference.

Changes for release 1.05

- Bug fixes

Changes for release 1.03

- Much faster topology creation
- Bug fixes

Changes for release 1.02

- Shooting* bug fixes
- Compilation problems solved

Changes for release 1.01

- Shooting* bug fixes

Changes for release 1.0

- Core and extra functions are separated
- Cmake build process
- Bug fixes

Changes for release 1.0.0b

- Additional SQL file with more simple names for wrapper functions
- Bug fixes

Changes for release 1.0.0a

- Shooting* shortest path algorithm for real road networks
- Several SQL bugs were fixed

Changes for release 0.9.9

- PostgreSQL 8.2 support
- Shortest path functions return empty result if they could not find any path

Changes for release 0.9.8

- Renumbering scheme was added to shortest path functions
- Directed shortest path functions were added
- routing_postgis.sql was modified to use dijkstra in TSP search


## Indices and tables

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[^0]:    - Supported versions: Latest (3.3) 3.2 3.1) 3.0
    - Unsupported versions: 2.6 2.5 2.4

    ```
    Driving Distance - Category
    ```

[^1]:    Transformation - Family of functions (Experimental)

[^2]:    Transformation - Family of functions (Experimental)

[^3]:    - Sample Data
    - Boost: Lengauer-Tarjan dominator tree algorithm
    - Wikipedia: dominator tree

[^4]:    \#1152 Fixes driving distance when vertex is not part of the graph

    - \#1098 Fixes windows test
    - \#1165 Fixes build for python3 and perl5

